

TWO PASS STRATEGY

Algorithm 1: TWO PASS STRATEGY

```
Activated  $\leftarrow \phi$ ;  
for Every event  $(k, v_k)$  do  
     $v(k) = v_k$ ;  
    for  $\forall j \in \text{fanout of } k$  do  
        Update the value of  $j$ ;  
        Activated  $\leftarrow$  Activated  $\cup \{j\}$ ;  
    end  
end  
for  $j \in \text{Activated}$  do  
    begin  
         $v' = \text{evaluate}(j)$ ;  
        if  $v' \neq \text{lsv}(j)$  then  
            add event  $(j, v')$  to the event list at time  $t + d(j)$ ;  
             $\text{lsv}(j) = v'$ ;  
        end  
    end  
end
```

ONE PASS STRATEGY

Algorithm 2: ONE PASS STRATEGY

```

for Every event  $(k, v_k)$  do
   $v(k) = v_k$ ;
  for  $\forall j \in \text{fanout of } k$  do
    Update the value of  $j$ ;
    Instead of putting  $j$  to activated set, process it here;
    begin
       $v' = \text{evaluate}(j)$ ;
      if  $v' \neq lsv(j)$  then
        add event  $(j, v')$  to the event list at time  $t + d(j)$ ;
         $lsv(j) = v'$ ;
      end
    end
  end
end

```

ONE PASS STRATEGY WITH ZERO-WIDTH SPIKES

Algorithm 3: ONE PASS STRATEGY WITH ZERO-WIDTH SPIKES

```

for Every event  $(k, v_k)$  do
   $v(k) = v_k$ ;
  for  $\forall j \in \text{fanout of } k$  do
    Update the value of  $j$ ;
    Instead of putting  $j$  to activated set, process it here;
    begin
       $v' = \text{evaluate}(j)$ ;
      if  $v' \neq lsv(j)$  then
         $t' = t + d(j)$ ;
        if  $t' = lst(j)$  then
          | cancel event  $(j, lsv(j))$  at time  $t'$ ;
        end
        add event  $(j, v')$  to the event list at time  $t + d(j)$ ;
         $lsv(j) = v'$ ;
         $lst(j) = t'$ ;
      end
    end
  end
end

```
