Computability and Complexity theory Assignment # 5

April 18, 2023

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- 1. Which of the following statements are true, which are false? Give a short proof of your answers.
 - (a) If L and L' are P-complete, so is $L \cup L'$.

Ans:

(b) If L = P, then every nontrivial language in L is P-complete.

Ans:

(c) $\mathsf{PolyL} \stackrel{\triangle}{=} \bigcup_{i \in \mathbb{N}} \mathsf{DSpace}(\log^i n)$ has complete problems (under log space many-one reductions).

Ans:

2. Show that the following problem is hard for NL under log-space many-one reductions. Input: A set A with an associative binary operation '*', a subset $S \subseteq A$, and an element $w \in A$.

Output: YES if there is a set of elements $a_1, \ldots, a_k \in S$ such that $a_1 * a_2 * \ldots * a_k = w$. Assume that A is closed under *. Is the problem complete for NL? Prove your answer.

Ans:

3. Let G = (V, E) be a directed graph. Consider two persons playing the following game: The first player occupies some given vertex v_0 . Then the second player chooses a non-occupied neighbor v_1 of v_0 (i.e. $(v_0, v_1) \in E$) and occupies it. Then the first player chooses some non-occupied neighbor v_2 of v_1 and occupies it (in addition to v_0). The game is continued in this way until one play is not able to extend the path v_0, v_1, \ldots by a non-occupied neighbor of the current vertex. The player who played the last move wins the game.

GEO is the following problem: Given a directed graph G and a vertex v_0 of G decide whether the first player has a winning strategy for the game above.

Prove that GEO is PSPACE-complete.

Ans:

4. Let G = (V, E) be a directed graph. Two vertices u and v are said to be strongly connected if there is a directed path from u to v in G as well as there is a directed path from v to u in G. Let

 $StrConn = \{\langle G, s, t \rangle \mid s \text{ and } t \text{ are strongly connected in } G\}.$

Show that StrConn is NL-complete (under logarithmic space many-one reductions.)

Ans:

- 5. A language A is called downward self-reducible, if there is a polynomial time oracle deterministic Turing machine M that on input x, only queries oracle strings with length <|x| such that $A=L(M^A)$.
 - (a) Show that SAT is downward self reducible.
 - (b) Show that if A is downward self reducible, then $A \in \mathsf{PSPACE}$.

Ans: