Binary Search

↳ Divide & conquer

i = 0
J = 7

1) Sorted array

0	1	2	3	4	5	6	7
2	4	8	12	20	25	50	70

arr = [2, 4, 8, 12, 20, 25, 50, 70]

x = 50

Small Problem → Single element

x = 50 ⟷ [50]

only one condition

```
if arr(i) == x :
    return i
return -1
```

Big Problem →

mid ⇒ (i + J) // 2 = (0 + 7) // 2 — overflow

= 3

✓

⇒ i + (J − i) // 2

(0 + 7) // 2

= 3

$x_2 = 4$, $x_1 = 50$, $x = 12$

0	1	2	3	4	5	6	7
2	4	8	12	20	25	50	70

→ Sorted

BinarySearch (arr, 0, 7, x)

1. mid = 3

i <= J

2. arr(mid) == x :

BinarySearch (arr, i, J, x)

Psuedocode

return mid

12 < 50

arr(mid) < x — Recursion

BinarySearch(arr, mid+1, J, x)

OR

i = mid + 1

12 > 4

arr(mid) > x — Recursion

BinarySearch (arr, i, mid - 1, x)

OR

J = mid - 1

return -1