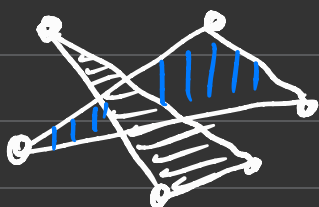


# Rank Reduction for SDPs

THEOREM. [Radon's Theorem] Let  $S$  be a set of  $d+2$  points in  $\mathbb{R}^d$ . Then  $\exists P, Q$  such that  
 $P \cap Q = \emptyset$ ;  $P \cup Q = S$ ;  $\text{conv}(P) \cap \text{conv}(Q) \neq \emptyset$ .



HOMOGENEOUS

$$\sum \alpha_i x_i = 0 \quad d+1 \text{ equations.}$$

$$\sum \alpha_i = 0 \quad d+2 \text{ variables.}$$



$\exists x \neq 0$  satisfying this.

$$P = \{i: \alpha_i > 0\}, \quad Q = [d] - P.$$

$$\sum_P \alpha_i = 1.$$

$$\frac{1}{Z} \sum_P \alpha_i x_i = \frac{1}{Z} \sum_Q \alpha_i x_i.$$

$S_P, S_Q$  witness.

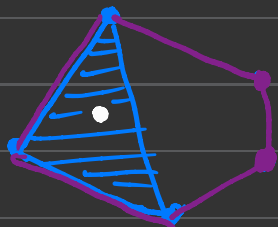
APPLICATION [VC Dimen] Take linear classifiers in  $d$  dimensions.  $VC = \max_n \left\{ n\text{-set } S \text{ is perfectly } \begin{matrix} \text{classifiable} \\ \text{w/ labellings} \end{matrix} \right\}$

Linear classifiers can shatter  $d+1$  points (simplex).

But given  $d+2$ ,  $\exists$  a labelling such that polytope distance between  $P, N = 0$ .

$$VC(\text{Linear Classifier}) = d+1.$$

THEOREM. [CARATHÉODORY] If  $x \in \text{conv}(P)$ , then  
 $\exists Q \subseteq P, |Q| = d+1, x \in \text{conv}(Q)$ .



$$x = \sum_{i=1}^m \alpha_i x_i ; x_i \in P. \sum \alpha_i = 1, \alpha_i \geq 0.$$

← finitely supported.

minimal representation.

$$\text{mlog } m \geq d+2.$$

Claim that  $\exists B \neq 0, \sum B_i x_i = 0 ; \sum B_i = 0$  (RADON-like)  
 $(d+1 \text{ equations} ; d+2 \text{ variables}).$

Define  $\gamma = \max \{ \gamma \geq 0 : \alpha_i - \gamma B_i \geq 0 \forall i \}.$

exists because  $\exists i, B_i > 0.$

$$\Delta_i = \alpha_i - \gamma B_i. \quad x = \sum_{i=1}^m \Delta_i x_i, \quad \sum \Delta_i = 1, \quad \Delta_i \geq 0.$$

Also it is  $m-1$  supported i.e.  $\Delta_i = 0$  for  $\exists i$ .

## SPARSIFYING LPs

In general, sparsest solution to a linear system is NP-hard.

Let  $x = \text{solve}(\min_x C^T x ; Ax = b ; x \geq 0)$ . Posit  $x^+ = x + \alpha \delta, \alpha > 0$ .

$x^+$  is OPT.  $C^T \delta = 0$ .

← Subsume. O/w  $x$  is NOT OPT.

$x^+$  is feasible.  $A\delta = 0$

why? Add  $\pm(\text{small}) \delta$ .

$x^+ \geq 0$  or  $\delta_j = 0 \Leftrightarrow x_j = 0$ .

\*  $m$  equations,  $|x|$  variables. ALA  $m < |x|, \exists \delta \neq 0, |x^+| < |x|$ .

$\therefore$  can reduce untill  $|x| \leq m$ .

## SDP EQUIVALENT

Let  $X = \text{SOLVE}(\min_x \langle C, x \rangle; A(x) = b; x \geq 0)$ . Posit  $X^+ = X + VSV^T$ .

where  $X = V\Sigma V^T$ ,  $S$  is symmetric.

$X^+$  is OPT.  $\langle C, VSV^T \rangle = 0$  SUBSUMED.

$X^+$  is feasible.  $A(VSV^T) = 0$  \*  $m$  constraints,  $r(r+1)/2$  variables.

CAN find  $S \neq 0$  to construct  $r(X^+) < r(X)$

whenever  $m < r(r+1)/2$ .

∴ can reduce  $r(r+1)/2 \leq m$ . [PATAKI, BARVINOK].

$$\text{RANK}(X) \leq \sqrt{2m + \frac{1}{4}} - \frac{1}{2}.$$