

Basic Feasible solution :- (BFS)

A feasible solution to a $(m \times n)$ transportation problem that contains no more than $(m+n-1)$ non-negative allocations is called basic feasible solution to the transportation pbm.

NORTH - WEST CORNER RULE:-

Determine basic feasible solution to the following transportation pbm using NWCR :-

	A	B	C	D	E	Supply
P	2	11	10	3	7	4
origin	Q	1	4	7	2	1
R	3	9	4	8	12	9
Demand	3	3	4	5	6	21

Sol:-

Since $\sum a_i = \sum b_j$ the given transportation pbm is balance

$$\sum a_i = 21 \quad \text{and} \quad \sum b_j = 21$$

Step 1:-

	A	B	C	D	E	Supply
P	2	11	10	3	7	4
origin	Q	1	4	7	2	1
R	3	9	4	8	12	9
Demand	3	3	4	5	6	

Step : 2

	B	C	D	E	Supply
P	11	10	3	7	X
Q	4	7	2	1	8
R	9	4	8	12	9
Demand	3	4	5	6	
	2				

Step 3 :-

B	B	C	D	E	Supply
Q	4	7	2	1	8
R	9	4	8	12	9
Demand	2	4	5	6	

Step 4 :

	C	D	E	Supply
Q	7	2	1	8
R	4	8	12	9
Demand	A	5	6	

Step 5:

	D	E	Supply
Q	2	1	X
R	8	12	8

Demand: 3 6

Step 6:

	D	E	Supply
R	8	12	X 6

Demand 3 6

Step 7:

	E	Supply
R	12	X

Demand X

2)

	A	B	C	D	E	
P	3 2	1 11	10	3	7	
Q	1 4	2 7	4 2	2	1	
R	3	9	4 3	8	6 12	

$$Z = 3 \times 2 + 1 \times 11 + 2 \times 4 + 4 \times 7 + 2 \times 2 + 3 \times 8 + 6 \times 12$$

$$Z = \text{Rs. } 153$$

→ Least cost Method (or) Matrix Minimum method:-

- Pbm: Obtain an initial basic feasible solution to the following T.P using the matrix minimum method.

	D ₁	D ₂	D ₃	D ₄	Capacity
O ₁	1	2	3	4	6
O ₂	4	3	2	0	8
O ₃	0	2	2	1	10
Demand	4	6	8	6	

Where O_i and D_j denote ith origin and jth destination respectively.

$$\sum a_{ij} = 24, \quad \sum b_j = 24$$

So. Given pbm is balanced.

Therefore given pbm has feasible solution.

Step 1:-

	D1	D2	D3	D4	capacity
O1	1	2	3	4	6
O2	4	3	2	0	8
O3	4	0	2	1	6
Demand	4	6	8	6	

Step 2:-

	D2	D3	D4	Capacity
O1	2	3	4	6
O2	3	2	6	8
O3	2	2	1	6
Demand	6	8	6	

2

Step 3 :-

	D_2	D_3	C
01	2	3	6
02	3	<u>2</u>	2
03	2	2	6
	6	86	

Step 4 :

C

†

+

-

	D_2	D_3	C
01	<u>b</u> 2	3	X
03	2	2	6
	b	b	

Step 5 :-

	D_2	D_3	C
03	2	<u>b</u> 2	b
	b	X	

	D ₁	D ₂	D ₃	D ₄
O ₁	1	6	3	4
O ₂	4	3	2	0
O ₃	4	0	2	1

$$Z = 4x_0 + 6x_2 + 2x_2 + 6x_2 + 6x_0$$

$$Z = \text{Rs. } 28.$$

Vogel's Approximation Method (VAM)

prob: Use Vogel's approximation method to obtain initial basic feasible solution of the transportation prbm.

	D	E	F	G ₁	Available
A	11	13	17	14	250
B	16	18	14	10	300
C	21	24	15	10	400
Demand	200	225	275	250	

$$\sum a_{ij} = 950, \sum b_{ij} = 950$$

Step 1

11	13	17	14	250	(2) $\leftarrow 13-11$
16	18	14	10	300	(4) $\leftarrow 14-10$
21	24	15	10	400	(3) $\leftarrow 13-10$

260 225 275 250

(5) (5) (1) (0)

2)

Step 2 :-

10 13	17	14	561	(1)
18	14	10	300	(4)
24	13	10	400	(3)

175
~~225~~ 275 250

(5) (1) (0)

+

-1 Step 3 .

18	14	10	125	(4)
24	13	10	300	(3)

1 175 275 250

(6) (1) (0)

Step 4 :-

14	10	125	125	(4)
13	10		400	(3)

275 280 125

(1) (0)

Step 5:-

275	10
125	400 (3)
275	10

275	10
275	125
(0)	(0)

Step 6

~~275~~ 275

275

200	50	17	14
11	13		
16	175	14	125 10
21	24	275 13	125 10

$$Z = 200 \times 11 + 50 \times 13 + 175 \times 18 + 125 \times 10 + 275 \times 13 + 125 \times 10.$$

$$Z = 12075.$$

2) MODI METHOD (Modified distribution method)

Test for optimal solution

Pbm :-

Solve the following pbm using MODI Method

	1	2	3	4	Supply
I	21	16	25	13	11
II	17	18	14	23	13
III	32	27	18	41	19
Demand	6	10	12	15	43

$$\sum a_{ij} = \sum b_j$$

First we should solve this pbm by Vogel's approximation method or W-W corner rule.

	1	2	3	4		
I	21	16	25	13	11	(3)
II	17	18	14	23	13	(3)
III	32	27	18	41	19	(9)

$$\begin{array}{cccc}
 6 & 10 & 12 & 15 \\
 (4) & (2) & (4) & (10)
 \end{array}$$

	1	2	3	4	9	(3)
<u>II</u>	17	18	14	23	18	
<u>III</u>	32	27	18	41	19	(9)
	6	10	12	A		
	(15)	(9)	(4)	(18)		

	1	2	3	93	(3)
<u>II</u>	17	18	14	93	(3)
<u>III</u>	32	27	18	19	(9)
	6	10	12		
	(15)	(9)	(4)		

	2	3	3	(4)
<u>II</u>	18	14	3	(4)
<u>III</u>	27	18	19	(9)
	162	12		
	(9)	(4)		

	2	3	12	3
<u>II</u>	27	18	12	(9)
	162	12		
	0	0		

2)

	1	2	3	4	v_1
I	21	16	25	11 13	$\rightarrow v_1$
II	6 17	3 18	14	4 23	v_2
III	32	7 27	12 18	41	v_3

 v_2

From this table, we see that the no. of non negative independent allocation is $(m+n-1) = 6$. Hence the solution is non degenerate basic feasible. \therefore The initial transportation cost is 796/-.

To find the optimal solution by modi method.

Consider the above transportation table.

since $m+n-1 = 6$, we can apply modi method.

Now we determine the set of values.

u_i and v_j for each occupied cell (i, j)

by using the relation $c_{ij} = u_i + v_j$

since the second row contains the maximum no. of allocations, then we choose $u_2 = 0$.

$$c_{21} = u_2 + v_1 \Rightarrow 17 = 0 + v_1 = v_1 = 17$$

$$c_{22} = u_2 + v_2 \Rightarrow 18 = 0 + v_2 = v_2 = 18$$

$$c_{24} = u_2 + v_4 \Rightarrow 23 = 0 + v_4 = v_4 = 23$$

$$c_{14} = u_1 + v_4 \Rightarrow 13 = u_1 + 23 \Rightarrow u_1 = -10.$$

$$c_{22} = u_2 + v_2 \Rightarrow 27 = u_2 + 18 \Rightarrow u_2 = 9.$$

$$c_{33} = u_3 + v_3 \Rightarrow 18 = u_3 + v_3 \Rightarrow v_3 = 9.$$

	1	2	3	4	
I	$\boxed{u_i+v_j}$ 21 \cancel{x}	14	25 \cancel{x}	13 \cancel{x}	$u_1 = -10$
II	6 \cancel{x}	17 \cancel{x}	18 \cancel{x}	14 \cancel{x}	$u_2 = 0$
III	22 \cancel{x}	7 \cancel{x}	12 \cancel{x}	41 \cancel{x}	$u_3 = 9$

$v_1 = 17 \quad v_2 = 18 \quad v_3 = 9 \quad v_4 = 23$

Now we find $u_i + v_j$ for each unoccupied cell (i, j) and enter at the corresponding unoccupied cell (i, j) then we find the cell evaluation $d_{ij} = c_{ij} - (u_i + v_j)$ for each unoccupied cell (i, j) and enter the corresponding unoccupied cell (i, j) . Thus we get the following table.

	1	2	3	4	
I	$\boxed{+}$ 21 \cancel{x}	14 \cancel{x}	8 \cancel{x}	8 \cancel{x}	$u_1 = -10$
II	6 \cancel{x}	3 \cancel{x}	9 \cancel{x}	5 \cancel{x}	$u_2 = 0$
III	26 \cancel{x}	6 \cancel{x}	7 \cancel{x}	12 \cancel{x}	$u_3 = 9$

$v_1 = 17 \quad v_2 = 18 \quad v_3 = 9 \quad v_4 = 23$

$$2) d_{11} = c_{11} - (v_1 + v_1) = 21 - 7 = 14$$

Since all $d_{ij} > 0$, the solution under the test is optimal and unique.

The optimum allocation schedule is given by

$$x_{14} = 11, x_{21} = 6, x_{22} = 3, x_{24} = 4, x_{32} = 7$$

$x_{33} = 12$ and the optimum (minimum).

transportation cost $Z = \text{Rs. } 796/-$

Pbm:-

Find the optimal transportation cost of the following matrix using least cost method for finding the critical solution.

	A	B	C	D	E	Available
R	7	1	2	6	9	100
Q	6	4	3	5	7	120
R	5	2	6	4	8	120
Demand	10	50	70	90	90	340

$$\sum a_i = \sum b_j$$

	A	B	C	D	E		
P	4	1	2	6	9	100	(1)
Q	6	4	3	5	7	120	(1)
R	5	2	6	4	8	120 80	(2)
	16	50	70	90	90		
(1)	(1)	(1)	(1)	(1)	(1)		

	B	C	D	E		
P	1	2	6	9	100	(1)
Q	4	3	5	7	120	(1)
R	2	6	4	8	80 30	(2)
	50	70	90	90		
(1)	(1)	(1)	(1)	(1)		

	C	D	E		
P	2	6	9	100	(4)
Q	3	5	7	120	(2)
R	6	4	8	30	(2)
	70	90	90		

2)

	A	B	C	D	E	
P	4	1	2	6	9	190 50
Q	6	4	3	5	7	120
R	5	2	6	4	8	120
	40	80	70	90	90	

	A	C	D	E	
P	4	2	6	9	50
Q	6	3	5	7	120
R	5	6	4	8	120
	40	70	90	90	
		20			

	A	C	D	E	
Q	6	3	5	7	120 100
R	5	6	4	8	120
	40	70	90	90	

	A	D	E	
Q	6	5	7	100
R	5	4	8	120 30
	40	90	90	

	A	E	
Q	6	7	100
R	5	8	36
	16	90	
	10		

	A	E	
Q	6	7	100 90
	16	90	

	E	
Q	7	90
	90	

	A	B	C	D	E	
P	4	<u>50</u> ₁	<u>50</u> ₂	6	9	<u>90</u>
Q	<u>10</u>	6	<u>20</u> ₃	5	<u>7</u>	
R	<u>30</u> ₅	2	6	<u>90</u> ₄	8	
	v ₁	v ₂	v ₃	v ₄	v ₅	

$$Z = 1410 \text{ l-}$$

2) For optimality since the no. of non-negative independent allocations is 7 which is equal to $m+n-1$, then we can apply MODI method.

$$U_2 = 0$$

$$U_2 + V_1 = 6 \Rightarrow V_1 = 6$$

$$U_2 + V_3 = 3 \Rightarrow V_3 = 3$$

$$U_2 + V_5 = 7 \Rightarrow V_5 = 7$$

$$V_3 + U_1 = 2 \Rightarrow 3 + U_1 = 2 \Rightarrow U_1 = -1$$

$$V_2 + U_1 = 1 \Rightarrow V_2 - 1 = 1 \Rightarrow V_2 = 2$$

$$V_1 + U_3 = 5 \Rightarrow 6 + U_3 = 5 \Rightarrow U_3 = -1$$

$$V_4 + U_3 = 4 \Rightarrow V_4 - 1 = 4 \Rightarrow V_4 = 5$$

	A	B	C	D	E	
P	5 4 x	6 1 x	50 1 2 x	4 6 9 x	2 6 9 x	$U_1 = -1$
Q	10 6 x	2 4 x	2 20 3 x	5 5 0 90 7 x	5 0 90 7 x	$U_2 = 0$
R	30 5 x	1 2 x	12 6 x	4 90 4 x	6 8 x	$U_3 = -1$
	$V_1 = 6$	$V_2 = 2$	$V_3 = 3$	$V_4 = 5$	$V_5 = 7$	

since $d_{ii} = -1 < 0$, the current solution is not optimal. Now let us form a new basic feasible solution by given maximum allocation to the cell (i, j) for which d_{ij} is most negative by making an unoccupied cell empty here the cell $(1, 1)$ having the negative value $d_{11} = -1$. we draw a closed loop consisting of horizontal and vertical lines beginning and ending at the cell $(1, 1)$ and having other corners at some occupied cells. Along the closed loop indicate +o and -o alternatively at the corners. we have.

	A	B	C	D	E
P	+o 4	50 1	50 2	-o 6	9
Q	10 -o 6		20 3	+o 5	90 7
R	30 5	2	6	10 4	8

From the two cell $(1, 3)$ and $(2, 1)$ having -o, we find that the minimum of the allocations that is $\min \{10, 50\} = 10$. Add this 10 to the cells with +o and subtract this 10 to the cells o

10	50	40		
4	1	2	6	9
		20		90
6	4	3	5	7

30				
5	2	6	90 4	8

Hence the new basic feasible solution is displayed in the following table.

2)

$$U_1 = 0$$

$$U_1 + V_1 = 4 \Rightarrow V_1 = 4$$

$$U_1 + V_2 = 1 \Rightarrow V_2 = 1$$

$$U_1 + V_3 = 2 \Rightarrow V_3 = 2$$

$$U_2 + V_3 = 3 \Rightarrow U_2 + 2 = 3 \Rightarrow U_2 = 1$$

$$U_2 + V_5 = 7 \Rightarrow 1 + V_5 = 7 \Rightarrow V_5 = 6$$

$$U_3 + V_1 = 5 \Rightarrow U_3 + 4 = 5 \Rightarrow U_3 = 1$$

$$U_3 + V_4 = 4 \Rightarrow 1 + V_4 = 4 \Rightarrow V_4 = 3$$

10	4	50	1	40	2	3	6	2	6	9	3
5	6	1	2	2	30	4	5	1	90	7	8
6	4	x	x	3	x	x	x	x	x	x	x
30	5	2	0	3	3	90	4	7	8	1	x

$$V_1 = 4 \quad V_2 = 1 \quad V_3 = 2 \quad V_4 = 3 \quad V_5 = 6$$

$$U_1 = 0$$

$$U_2 = 1$$

$$U_3 = 1$$

All $d_{ij} > 0$ the solution under the test

is optimal & unique.

$$Z = 4 \times 10 + 1 \times 50 + 2 \times 40 + 3 \times 30 + 7 \times 90 + 30 \times 5 + 4 \times 90$$

$$Z = RS \cdot 140/-$$

$$Z = 1160$$

Pbm :-

6	1	9	3
11	5	2	8
10	12	4	7

Available units	
70	(2)
55	(2)
90	(2)

Required units

85	35	50	45
15			
(4)	(4)	(2)	(4)

11	5	2	8
10	12	4	7

85	20	(3)
90		(3)

15	35	50	45
(1)	(7)	(2)	(1)

11	2	8
10	4	7

20		(6)
90		(3)

15	50	45
(1)	(2)	(1)

10	4	7	60	(3)
15	30	45		

(0)	(0)	(0)
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2)

$$\begin{array}{|c|c|} \hline 10 & 7 \\ \hline 15 & 45 \\ \hline 0 & 10 \\ \hline \end{array} \quad 60^{15} (3)$$

$$\begin{array}{|c|} \hline 10 \\ \hline \end{array} \quad 15 \\ 15$$

6	1	9	3
11	35	20	8
15	10	30	45

$$m+n-1 \geq 6$$

$$Z = 70 \times 6 + 35 \times 5 + 20 \times 20 + 15 \times 10 + 30 \times 4 + 45 \times 7$$

$$Z = Rs - 1580/-$$

$$U_3 = 0$$

$$V_1 + U_3 = 10 \Rightarrow V_1 = 10$$

$$V_3 + U_3 = 4 \Rightarrow V_3 = 4$$

$$V_4 + U_3 = 7 \Rightarrow V_4 = 7$$

$$U_2 + V_2 = 5 \Rightarrow -2 + V_2 = 5 \Rightarrow V_2 = 7$$

$$U_2 + V_3 = 2 \Rightarrow U_2 + 4 = 2 \Rightarrow U_2 = -2$$

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$$U_1 + V_1 = 6 \Rightarrow U_1 + 10 = 6 \Rightarrow U_1 = -4$$

70	3	-2	0	9	3	0
-0		+0				
8	3	35	20	5	8	3
11		5	2	+0		
15	7	5	20	45	7	
10		12	4	-0		
+0						

$V_1 = 10$ $V_2 = 7$ $V_3 = 4$ $V_4 = 7$
 $U_1 = -4$ $U_2 = -2$ $U_3 = 0$.

cell $(2, 2)$, $(3, 3)$, $(1, 1)$

$$\min \{70, 35, 30\} = 30$$

70				
60		+0		
11	25	05	20	8
15	10	12	30	45

40				
6	30	1	9	3
11	85	5	20	8
10	45	12	30	45

$V_1 = 10$ $V_2 = 7$ $V_3 = 4$ $V_4 = 7$
 $U_1 = -4$ $U_2 = -2$ $U_3 = 0$.

$$U_3 = 0$$

$$U_3 + V_1 = 10 \Rightarrow V_1 = 10$$

$$U_3 + V_4 = 7 \Rightarrow V_4 = 7$$

$$U_3 + V_2 = 7 \Rightarrow V_2 = 7$$

2)

$$U_2 + V_2 = 5 \Rightarrow -2 + V_2 = 5 \Rightarrow V_2 = 7$$

$$U_2 + V_3 = 2 \Rightarrow U_2 + 4 = 2 \Rightarrow U_2 = -2$$

$$U_1 + V_1 = 6 \Rightarrow U_1 + 10 = 6 \Rightarrow U_1 = -4$$

$$X_1 + X_2 = 11 \Rightarrow X_2 = 11 - X_1$$

6	4	0	30	1	0	9	3	3	0	$U_1 = -4$
8	1	3	5	1	50	2	5	8	3	$U_2 = -2$
11	x	5	x	2	x	8	x	3	x	$U_3 = 0$
45	7	1	5	9	4	0	45	7	x	$V_1 = 10$
10	12	x	4	x	7	4	45	x	$V_2 = 7$	$V_3 = 4$
										$V_4 = 7$

$$Z = 6 \times 40 + 30 \times 1 + 5 \times 5 + 50 \times 2 + 45 \times 10 + 45 \times 7$$

$$Z = \text{Rs} - 1160/-$$

Degeneracy in Transportation pbm: X

In a trans pbm, whenever the no. of non-negative independent allocation is less than $m+n-1$, the Tp is said to be degenerate. Degeneracy may occur either at the initial stage (or) at an intermediate stage at some subsequent iteration.

To resolve degeneracy, we allocate an extremely small amount (which is close to zero), to one (or) more empty cells of the T table (generally minimum cost cells if possible), so that the total number of occupied cells becomes $(m+n-1)$ at independent positions. We denote this small amount by ϵ (epsilon) satisfying the following condition

- i) $0 < \epsilon < x_{ij}$ for all $x_{ij} > 0$
- ii) $x_{ij} + \epsilon = x_{ij}$ for all $x_{ij} > 0$.

Pbm :-

NNC rule

					Supply
10	20	5	7		10
13	9	12	8		20
4	5	7	9		30
14	7	1	0		40
3	12	5	19		50.

Demand 60 60 20 10
 50 .

13	9	12	8	20
4	5	7	9	30
14	7	1	0	40
3	12	5	19	50.

60 60 20 10
30

2)

4	5	7	9
14	7	1	0
3	12	5	19
36	60	20	10

36

40

50

14	7	1	0	40
3	12	5	19	50
60	20	10		

7	1	0	40
12	5	19	50
60	20	10	
20			

5	19	36	10
26	10		

19 10 -

10

~~max~~

10	0	19	4	12	5	7	26	7	19
20		22	-13	15	12	3	29	8	-19
13		9	x	x	x	x	x	x	
30	4	13	5	-8	6	7	20	9	-11
-2	14	16	40	+	0	1	18	-14	
14	+0	x	x	-0		x	x	x	
5	20	12		20	5		10	19	
-0	x	x	+0						
v ₁ = 3		v ₂ = 12		v ₃ = 5		v ₄ = 19			

$$v_1 = 7$$

$$v_2 = 10$$

$$v_3 = 1$$

$$v_4 = -5$$

$$v_5 = 0$$

$$Z = 12 \cdot 30 + 3 \cdot 5$$

$$\min Z = 12 \cdot 30$$

$$v_5 = 0$$

$$v_5 + v_1 = 3 \Rightarrow v_1 = 3$$

$$v_5 + v_2 = 12 \Rightarrow v_2 = 12$$

$$v_5 + v_3 = 5 \Rightarrow v_3 = 5$$

$$v_5 + v_4 = 19 \Rightarrow v_4 = 19.$$

$$v_4 + v_2 = 7 \Rightarrow v_4 + 12 = 7 \Rightarrow v_4 = -5$$

$$v_3 + v_1 = 4 \Rightarrow v_3 + 3 = 4 \Rightarrow v_3 = 1$$

$$v_2 + v_1 = 13 \Rightarrow v_2 + 3 = 13 \Rightarrow v_2 = 10.$$

$$v_1 + v_1 = 10 \Rightarrow v_1 + 3 = 10 \Rightarrow v_1 = 7$$

$$\min \left\{ (4, 2) \quad (5, 1) \right\} = \{ 7, 3 \} = 3.$$

2)

$\frac{10}{20}$	10	20	5	7
$\frac{20}{15}$	15	9	12	8
$\frac{30}{4}$	4	5	7	9
$\frac{40}{14}$	14	$\frac{7}{7}$	1	0
$\frac{\Sigma}{30}$	30	$\frac{20}{12}$	$\frac{20}{5}$	$\frac{10}{19}$

CB

M	$\frac{10}{20}$	10	20	5	7	$v_1 = 10$
0	$\frac{20}{13}$	13	9	12	8	$v_2 = 13$
-M	$\frac{30}{4}$	4	5	7	9	$v_3 = 4$
	$\frac{3}{14}$	14	$\frac{7}{7}$	1	0	$v_4 = 14$
	$\frac{3-\Sigma}{30}$	30	$\frac{23}{12}$	$\frac{20}{5}$	$\frac{10}{19}$	$v_5 = 30$
		$v_1 = 0$	$v_2 = -7$	$\sqrt{3} = -25$	$v_4 = -11$	

$$V_1 = 0$$

$$V_1 + v_1 = 10 \Rightarrow v_1 = 10$$

$$V_1 + v_2 = 13 \Rightarrow v_2 = 13$$

$$V_1 + v_3 = 4 \Rightarrow v_3 = 4$$

$$V_1 + v_4 = 14 \Rightarrow v_4 = 14$$

$$V_1 + v_5 = 30 \Rightarrow v_5 = 30$$

$$V_2 + V_4 = 7 \Rightarrow V_2 + 14 = 7 \Rightarrow V_2 = -7$$

$$V_3 + V_5 = 5 \Rightarrow V_3 + 30 = 5 \Rightarrow V_3 = -25$$

$$V_4 + V_5 = 19 \Rightarrow V_4 + 30 = 19 \Rightarrow V_4 = -11$$

$$Z = 10 \times 10 + 20 \times 13 + 30 \times 4 + 3 \times 14 + 37 \times 7 \\ + (3 - \varepsilon) \times 30 + 23 \times 12 + 20 \times 5 + 10 \times 19$$

$$Z = 1347 + 90 - 30\varepsilon$$

$$\varepsilon = 0$$

$$Z = 1347 + 90 \Rightarrow \text{RS } 1437/-$$

2) Vogel's

	(1)	(2)	(3)	(4)	
Supply	10	20	5	7	
	13	9	12	8	
	4	5	7	9	
	14	7	1	0	
	3	12	5	19	

Demand
60 (1)
60 (2)
20 (4)
10 (1)

2)

10	20	5	10	(5)
13	9	12	20	(3)
4	5	7	30	(1)
14	7	1	60	(6)
3	12	5	50	(2)

60 60 20
 (1) (2) (4)

CP

M

O

-T

10	20	16	(10)
13	9	20	(4)
4	5	30	(1)
14	7	10	(7)
3	12	50	(9)

60 50 60
 (1) (2)

13	9	20	(4)
4	5	30	(1)
14	7	10	(7)
3	12	50	(9)

80 60
 (1) (2)

9	20
5	30
7	10
12	0

9	20
X	10
12	0

3620.

6030

209	20
12	0
	20

10	20	5	7
13	9	12	8
4	5	7	9
14	7	1	0
3	12	5	12

I = Rs. 670/-

2) 3) Least count:

10	20	5	7	10
13	9	12	8	20
4	5	7	9	30
14	7	1	0	46 30
3	12	5	19	50
60	60	20	16	

CE

M

10	20	5	10	10
13	9	12	20	20
4	5	7	30	30
14	7	1	10	30 10
3	12	5	50	50
60	60	20		

10	20	10
13	9	20
4	5	30
14	7	10
3	12	50
60	60	

10	20
18	9
4	5
14	7
18	60

10
20
30 20
10

20
9
5
7

10
20
20
10

60
40

20
9
7

10
20
10

40 30

20
9

10
20

30 10

2)

$$\boxed{20} \quad 10$$

	$\boxed{10}$			
$\boxed{10}$	20	5	7	
13	$\overset{20}{\cancel{9}}$	12	8	
$\boxed{10}$	$\overset{20}{\cancel{5}}$	$\cancel{7}$	$\overset{2}{\cancel{9}}$	
14	$\boxed{10}$	$\overset{20}{\cancel{1}}$	$\overset{10}{\cancel{0}}$	
$\cancel{50}$	3	12	5	19

C
M
C
+

Assignment Method (or) Hungarian Method.

Pbm:- A department has five emp with 5 jobs to be performed. The time (in hours) each man will take to perform each job is given in the effectiveness matrix.

	I	II	III	IV	V
A	10	5	13	15	16
B	3	9	18	13	6
C	10	7	2	2	2
D	7	11	9	7	12
E	7	9	10	4	12

Select minimum in row and sub

Step 1: Row wise subtraction

5	0	8	10	11
0	6	15	10	3
8	5	0	0	0
0	4	2	0	5
3	5	6	0	8

2 Step 2 : Column wise

5	<input type="checkbox"/>	8	10	11
<input type="checkbox"/>	6	15	10	3
8	5	<input type="checkbox"/>	X	X
X	4	2	<input type="checkbox"/>	5
3	5	6	X	8

Sub min value of the table zero, and mark it as and mark first we have to see row wise

5	<input type="checkbox"/>	3	10	6
<input type="checkbox"/>	6	15	10	3
8	0	<input type="checkbox"/>	X	X
X	4	2	<input type="checkbox"/>	5
3	5	6	X	8

X in corresponding row after see column wise zero same procedure

Entha row the assigned 0 "" illaiyo verum X mattum irundha anthe row ah 1 tick pannikanum next col la and X enga iruko anga tick pannikanum tick agirukura col la enga iruko anga anthe row la tick pannikanum.

Step 2 :- Column wise.

A hand-drawn coordinate grid on a whiteboard. The horizontal axis (x-axis) has tick marks and labels: 5, 0, 8, 10, 4, 6, 15, 10, 3, 8, 5, 0, 2, 0, 5, 3, 5, 6, -8, and 3. The vertical axis (y-axis) has tick marks and labels: 3, 5, 6, -8, and 3. Points are marked with small squares or circles. A vertical line at x=0 is labeled "unmarked".

unmarked rows
and
marked column.

Sub min value of the table.

2)

I II III IV V

A 7 0 8 12 11B 0 4 15 10 1C 10 5 ~~X~~ 2 0D ~~X~~ 42 0 ~~X~~ 3E 3 3 4 6 $A \rightarrow \underline{\text{II}}$ $C \rightarrow \underline{\text{IV}}$ $E \rightarrow \underline{\text{IV}}$ $B \rightarrow \underline{\text{I}}$ $D \rightarrow \underline{\text{III}}$

$$Z = 5 + 3 + 2 + 9 + 4 \Rightarrow 23 \text{ hours}$$

Maximization case in the Assignment problem :-

	1	2	3	4
A	(16)	10	14	11
B	14	11	15	15
C	15	15	13	12
D	13	12	14	15

Select max no and sub it with all value
in table

	1	2	3	4	5
A	0	6	2	5	
B	2	5	1	1	
C	1	0	3	0	4
D	3	4	2	0	1

Step 1 & 2 :-

	1	2	3	4
A	0	6	2	5
B	1	4	0	3
C	0	0	2	1
D	2	3	1	0

2)

	1	2	3	4
A	5	1	5	
B	4	0	1	
C	0	2	4	
D	2	2	0	0

$$A \rightarrow 1 \quad B \rightarrow 3 \quad C \rightarrow 2 \quad D \rightarrow 4$$

Rs: - $16 + 15 + 15 + 15 = 61.$

	1	2	3	4	5	6	7
A	60	50	40	30			
B	40	30	20	15			
C	40	20	35	10			
D	30	20	25	20			

	1	2	3	4
A	0	10	20	30
B	20	30	40	45
C	20	40	25	50
D	30	30	35	40

	1	2	3	4
A	0	10	20	30
B	0	10	20	25
C	0	20	10	30
D	10	10	15	20

	1	2	3	4
A	0	0	5	10
B	0	0	5	5
C	0	10	0	10
D	10	0	0	0

	1	2	3	4
A	0	0	15	10
B	0	0	15	5
C	0	10	0	10
D	10	0	10	0

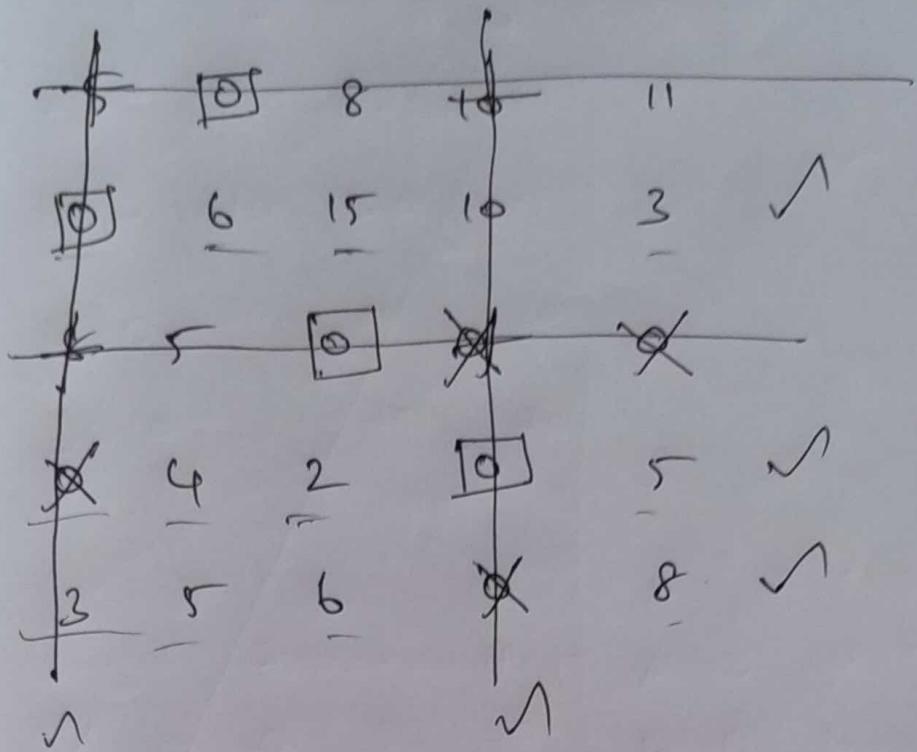
	1	2	3	4
A	0	10	20	30
B	0	10	20	25
C	0	20	5	30
D	0	0	5	10

	1	2	3	4
A	0	10	15	20
B	0	10	15	15
C	0	20	0	20
D	0	0	0	0

	1	2	3	4
A	0	10	15	20
B	0	10	15	15
C	0	20	0	20
D	0	0	0	0

	1	2	3	4
A	0	0	15	10
B	0	0	15	5
C	0	10	0	10
D	10	0	10	0

✓



	I	II	III	IV	V
A	+	<input type="checkbox"/>	8	12	11
B	<input type="checkbox"/>	4	13	10	1
C	<input checked="" type="checkbox"/>	5	<input checked="" type="checkbox"/>	2	<input type="checkbox"/>
D	<input checked="" type="checkbox"/>	2	<input type="checkbox"/>	<input checked="" type="checkbox"/>	3
E	3	3	4	<input type="checkbox"/>	6.

A → I
 B → II
 C → III
 D → IV
 E → V

$$5 + 3 + 2 + 9 + 4 = 23$$