

Q-1) - Discrete Cosine Transform

P-1

on (FFT)  $DCT(X) = e^{\frac{i\omega K}{2}} \overline{F_{2N}(X)}$

$$F_{2N} = \sum_{n=0}^{N-1} x(n) e^{\frac{2\pi i kn}{2N}} + e^{\frac{i\omega K}{2}} F_{2N}(X)$$

for  $K = \{0, 1, \dots, N-1\}$

Discrete Cosine Transform is special case of Fast Fourier Transform.

$$\begin{aligned} F_{2N} = C_K &= \sum_{n=0}^{N-1} x_n e^{\left(-i \frac{2\pi K n}{N}\right)} \\ &= \sum_{n=0}^{\frac{1}{2}N} x_n e^{\left(-i \frac{2\pi K n}{N}\right)} + \sum_{n=\frac{1}{2}N+1}^{N-1} x_n e^{\left(-i \frac{2\pi K n}{N}\right)} \\ &= \sum_{n=0}^{\frac{1}{2}N} x_n e^{\left(-i \frac{2\pi K n}{N}\right)} + \sum_{n=\frac{1}{2}N+1}^{N-1} x_{N-n} e^{\frac{i 2\pi K (N-n)}{N}} \end{aligned}$$

function is symmetric  $x_0 = x_N$ ,  
 $x_1 = x_{N-1}$  and  
 $e^{i 2\pi K} = 1 \quad \forall K \in \{0, 1, 2, \dots, N-1\}$

change  $(N-n) \rightarrow n$  in right hand expr

$$\begin{aligned} C_K &= \sum_{n=0}^{\frac{1}{2}N} x_n e^{-i \left(\frac{2\pi K n}{N}\right)} + \sum_{n=1}^{\frac{1}{2}N-1} x_n e^{\left(i \frac{2\pi K n}{N}\right)} \\ &= x_0 + x_{N/2} \cos\left(\frac{2\pi K (N/2)}{N}\right) + 2 \sum_{n=1}^{\frac{1}{2}N-1} x_n \cos\left(\frac{2\pi K n}{N}\right) \end{aligned}$$

Cosine transform is applied to real samples, which implies that coefficients  $C_K$ , will all be real as the sum of all real terms.

As  $x_n$  and  $C_k$  are real

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$$C_{N-k} = C_k^* = C_k \quad \text{inverse function}$$

$$x_n = \frac{1}{N} \sum_{k=0}^{N-1} C_k e^{i\left(\frac{2\pi kn}{N}\right)}$$

$$x_n = \frac{1}{N} \left\{ \sum_{k=0}^{\frac{N}{2}-1} C_k e^{i\left(\frac{2\pi kn}{N}\right)} + \sum_{k=\frac{N}{2}}^{N-1} C_k e^{i\left(\frac{2\pi kn}{N}\right)} \right\}$$

for each  $C_k$  one has to  $N$  complex multiplication and  $(N-1)$  addition i.e.,  
i.e.,  $(2N-1)$  complex operations.

Since  $N$   $C_k$ 's the total number of operations is  $O(N^2)$

Fast Fourier Transform algorithm is simple to understand when no. of sample  $N = 2^m$  divide the group into two as even and odd groups.

$$\equiv E_k + e^{-i\frac{2\pi k}{N}} O_k$$

$$\text{and } C_k = \begin{cases} \frac{1}{\sqrt{2}} & k=0 \\ 1 & \text{otherwise} \end{cases}$$

$$\text{so } x_n = \frac{1}{N} \sum_{k=0}^{N-1} \frac{1}{\sqrt{2}} e^{i\left(\frac{2\pi kn}{N}\right)}$$

$$x_n = \frac{1}{N} \sum_{k=0}^{N-1} \frac{1}{\sqrt{2}} e^{i(n\omega_k)} \quad \omega_k = \left(\frac{2\pi k}{N}\right)$$

$$= \frac{1}{\sqrt{2}N} \sum_{k=0}^{N-1} e^{i(n\omega_k)}$$



row vector of size  $1 \times N$ 

$$C_N^k(n) = e^{\frac{\pi i n (n-N)k}{N}}$$

$$j, k \in (0 \rightarrow N-1)$$

$$V_{j,k} = \frac{1}{N} e^{\frac{2\pi i j n}{N}} C_N^k(n) = \frac{1}{N} e^{\frac{\pi i n}{N} (2j + (n-N)k)}$$

Form  $(N \times N^2)$  matrix  $V$ 

Columns are

$$\{V_{j,k}\}_{j,k=0}^{N-1}$$

row vectors

$$N=1, \quad C_1^k(n) = e^{\frac{\pi (1-N)k}{N}}$$

$$C_1^k(n) = e^{\pi (n-1)k}$$

$$C_2^k(n) = e^{\frac{\pi (n-2)k}{2}}$$

$$V_{j,k} = \frac{1}{N} e^{\frac{2\pi i j n}{N}}$$

$$j = 0, 1, 2, \dots, N-1$$

$$k = 0, 1, 2, \dots, N-1$$

$$V_{(0,0)} = \frac{1}{N} e^0 = \frac{1}{N}$$

$$V_{(j,k)} = \frac{1}{N} e^{\frac{2\pi i j n}{N}} C_N^k(n) = \frac{1}{N} e^{\frac{\pi i n}{N} (2j + (n-N)k)}$$

$$j=0, k=0$$

$$V_{(0,0)} = \frac{1}{N} e^{\frac{\pi i n (0)}{N}} = \frac{1}{N} e^0 = \frac{1}{N}$$

$$V_{(1,1)} = \frac{1}{N} e^{\frac{\pi i n (2 + n - N)}{N}}$$

$$V_{(2,2)} = \frac{1}{N} e^{\frac{\pi i n (4 + 2n - 2N)}{N}}$$

$$V_{(N-1, N-1)} = \frac{1}{N-1} e^{\frac{\pi i n (2N-2 + (n-N)(N-1))}{N}}$$

$$2N-2 + nN - n - N^2 + N$$

$$(3N + nN - N^2 - n - 2)$$

$$U_{j,K} = \frac{1}{N} \left[ \begin{array}{c} e^{\frac{\pi 2jn}{N}} (n+2-N) \\ e^{\frac{2\pi n}{N}} (2n+4-2K) \end{array} \right]$$

$$U_{j,K} = \frac{1}{N} e^{\frac{2\pi n}{N}} \left[ \begin{array}{c} 1 \quad (n+2-N) \quad (2n+4-2K) \text{ --- } \\ 2 \quad (n+4-N) \quad (2n+8-4K) \text{ --- } \\ \vdots \\ n \quad (n+2N) \quad (2n+2N-2K) \text{ --- } \end{array} \right]$$

$$\forall \quad j = 0, 1, 2, \dots, N-1$$

$$K = 0, 1, 2, \dots, N-1$$

$$C_{N-1}^K(n) = e^{\frac{\pi 2jn}{N-1}} \left( \frac{n-(N-1)}{N-1} \right)^K$$