

## ITSRUNTYM

Dynamic Programming (DP) is a powerful algorithmic technique used to **solve problems by breaking them down into smaller subproblems, solving each subproblem only once, and storing their results (memoization)** to avoid redundant work.

### WHY USE DYNAMIC PROGRAMMING?

- **When to use DP:**

Use DP when a problem has:

1. **Overlapping Subproblems**

→ Subproblems are solved multiple times (e.g., Fibonacci).

2. **Optimal Substructure**

→ Optimal solution of the problem depends on the optimal solutions of its subproblems.

- **Why it helps:**

Instead of recalculating the result for the same input multiple times (like in plain recursion), **DP saves the result in memory and reuses it**, leading to a significant performance boost.

### TYPES OF DP APPROACHES

Type	Description	Example Problem
<b>Top-Down (Memoization)</b>	Use recursion + store results in a table	Fibonacci using recursion and dp[]
<b>Bottom-Up (Tabulation)</b>	Solve subproblems iteratively, build solution from base	Fibonacci using a loop and dp[]
<b>Space Optimization</b>	Optimize space by storing only required results	Fibonacci using two variables

1. **Top-Down (Memoization)**

2. **Bottom-Up (Tabulation)**

3. **Space Optimization**

We will explain:

1. Why it's named like that
2. The **difference between the three**
3. When to use each
4. With diagrams and code for each

## 1. Top-Down Approach (Memoization)

Why this name?

- **"Top-Down"**: We start from the **main problem** and break it down into **subproblems recursively**.
- **"Memoization"**: We **store answers to subproblems** in a table (dp[]) to avoid recomputation.

Characteristics:

- Uses **recursion**
- Uses a **cache (dp[])** to save results of subproblems
- Useful when you're comfortable thinking recursively

### General Template (Recursive + Cache)

```
int solve(int n, int[] dp) {  
    if (n == 0 || n == 1) return n;  
    if (dp[n] != -1) return dp[n]; // already solved  
    return dp[n] = solve(n - 1, dp) + solve(n - 2, dp);  
}
```

### Diagram (Top-Down Flow)

Start from fib(5)

|

fib(5) = fib(4) + fib(3)

|

recurse on fib(4), fib(3), ...

Store in dp[] during the way down

## 2. Bottom-Up Approach (Tabulation)

Why this name?

- **"Bottom-Up"**: We solve **smaller subproblems first**, and use them to solve **bigger problems**.
- **"Tabulation"**: We use a **table (array)** to store the result in a **bottom-up fashion**.

Characteristics:

- Uses **loops (iteration)**, no recursion
- Builds the solution from **base cases upward**

- More memory efficient than top-down (no call stack overhead)

### General Template

```
int fib(int n) {
    int[] dp = new int[n + 1];
    dp[0] = 0;
    dp[1] = 1;
    for (int i = 2; i <= n; i++) {
        dp[i] = dp[i - 1] + dp[i - 2];
    }
    return dp[n];
}
```

### Diagram (Bottom-Up Flow)

Build from  $dp[0] \rightarrow dp[1] \rightarrow dp[2] \rightarrow dp[3] \dots$  up to  $dp[n]$

$dp[0] = 0$

$dp[1] = 1$

$dp[2] = dp[1] + dp[0] = 1$

...

$dp[n] = \text{result}$

## 3. Space Optimization

### Why this name?

- We **optimize space** usage by noticing that we don't need the whole  $dp[]$  table – only a few previous values.
















### Characteristics:

- Applies only when **current state depends on few previous states**
- Reduces space from  $O(n) \rightarrow O(1)$  in some cases

### General Template

```
int fib(int n) {
    if (n <= 1) return n;
    int prev2 = 0, prev1 = 1;
    for (int i = 2; i <= n; i++) {
        int curr = prev1 + prev2;
        prev2 = prev1;
        prev1 = curr;
    }
    return prev1;
}
```

Approach	Time Complexity	Space Complexity	Notes
Top-Down (Memoization)	$O(n)$	$O(n)$	Uses recursion + dp[]
Bottom-Up (Tabulation)	$O(n)$	$O(n)$	Iterative, no recursion
Space Optimization	$O(n)$	$O(1)$	Best for linear recurrence problems

Problem	DP Type	LeetCode Link
Fibonacci Number	1D DP	<a href="#"> Link</a>
Climbing Stairs	1D DP	<a href="#"> Link</a>
Min Cost Climbing Stairs	1D DP	<a href="#"> Link</a>
House Robber	1D DP	<a href="#"> Link</a>
Unique Paths	2D Grid DP	<a href="#"> Link</a>
Minimum Path Sum	2D Grid DP	<a href="#"> Link</a>
Longest Common Subsequence	2D DP on Strings	<a href="#"> Link</a>
Longest Palindromic Subsequence	2D DP on Strings	<a href="#"> Link</a>
Partition Equal Subset Sum	0/1 Knapsack (1D DP)	<a href="#"> Link</a>
Combination Sum IV	Subset Sum / Count Ways	<a href="#"> Link</a>
Target Sum	0/1 Knapsack / Subset Sum	<a href="#"> Link</a>
Edit Distance	2D DP on Strings	<a href="#"> Link</a>
Dungeon Game	2D DP (Reverse Build)	<a href="#"> Link</a>
Jump Game II	1D DP + Greedy	<a href="#"> Link</a>
Cherry Pickup II	3D DP (Advanced Grid)	<a href="#"> Link</a>