## DS & AI

Database Management System

Super 1500+

Lecture No. 08



#### **Recap of Previous Lecture**









File organization and Indexing



### **Topics to be Covered**



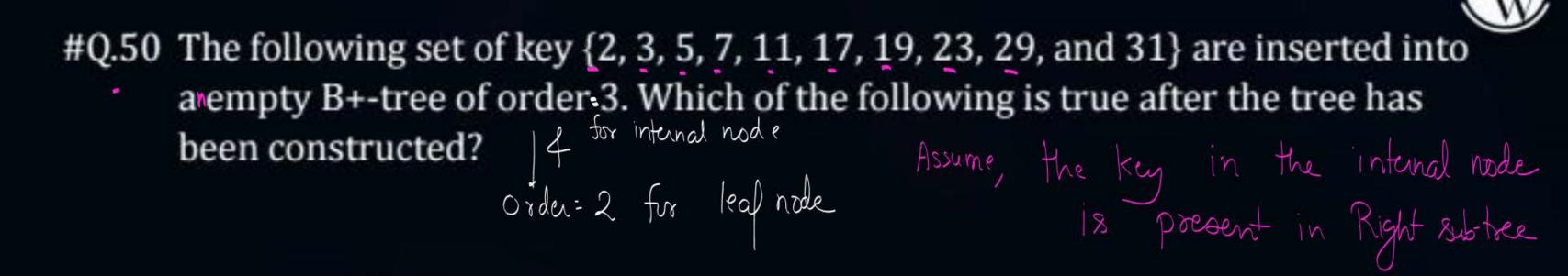






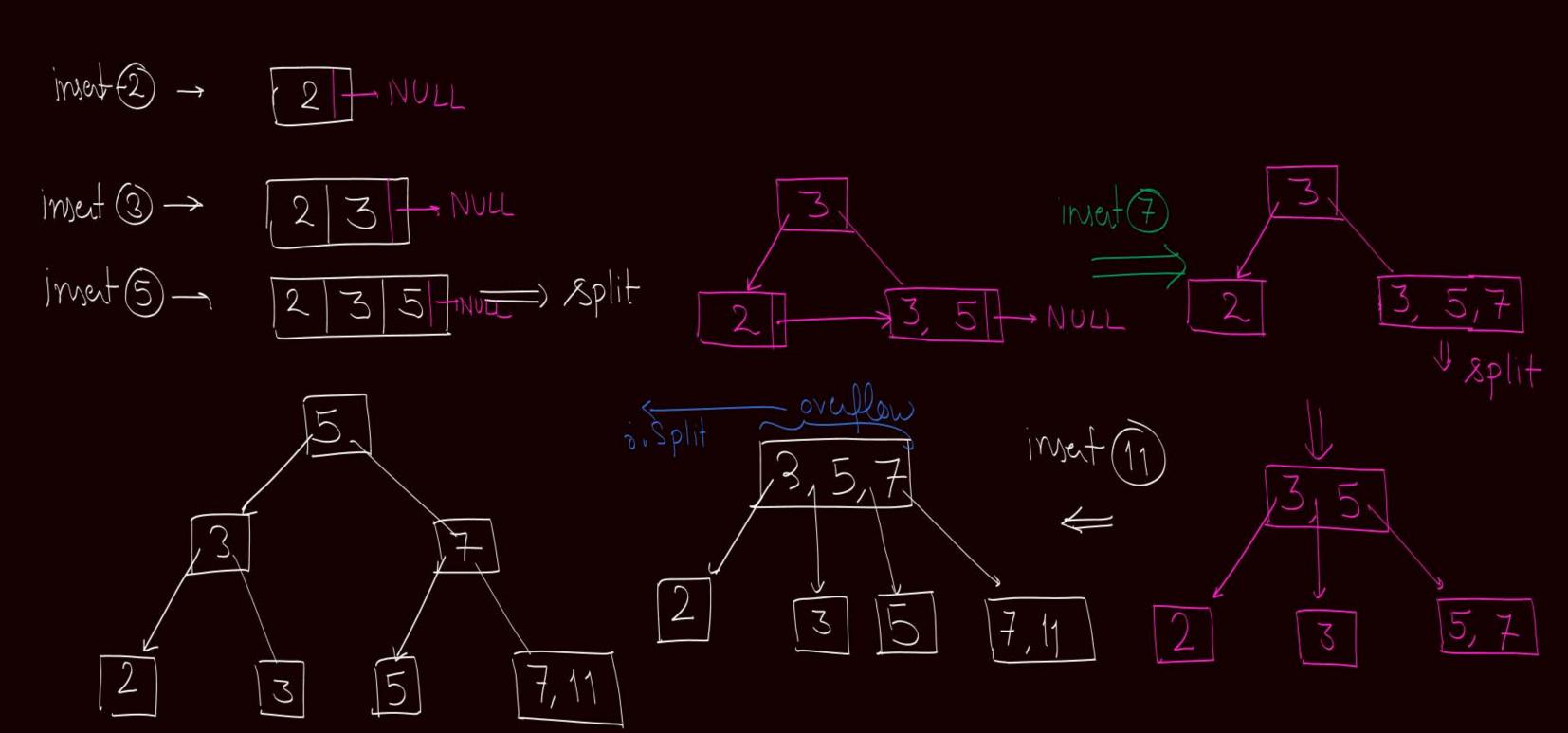
File organization and Indexing

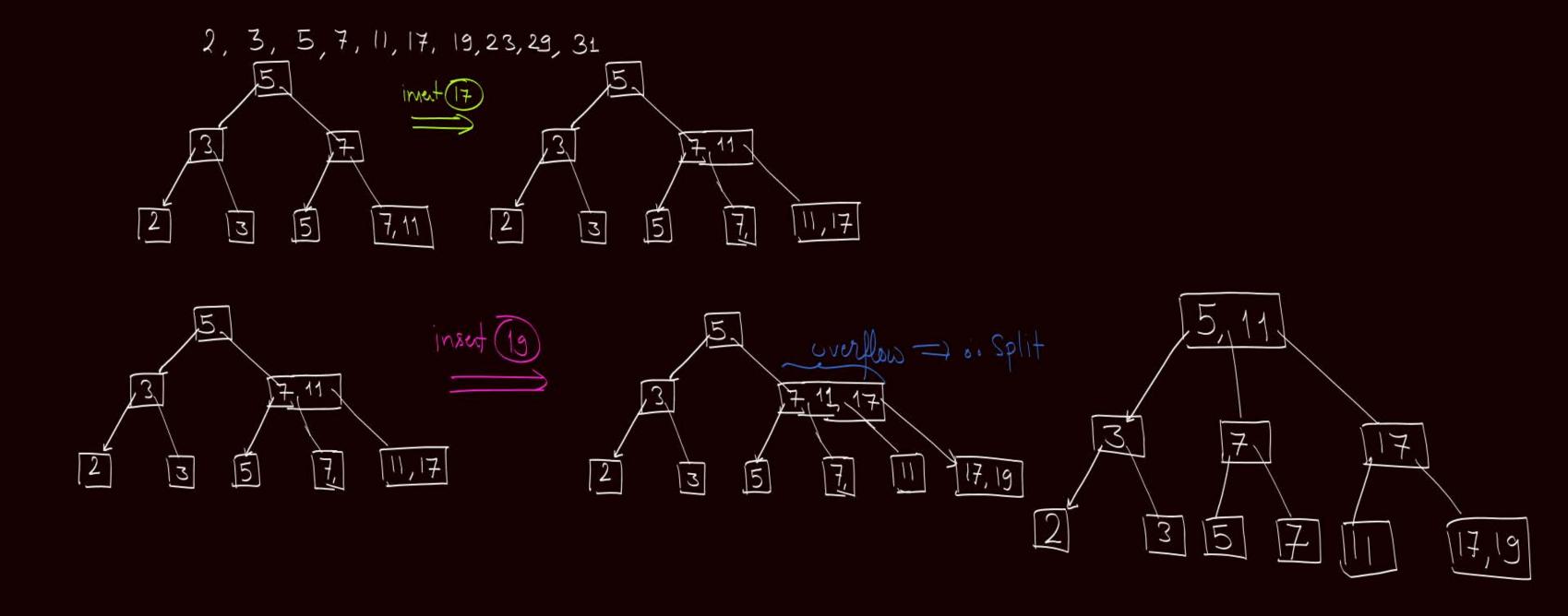




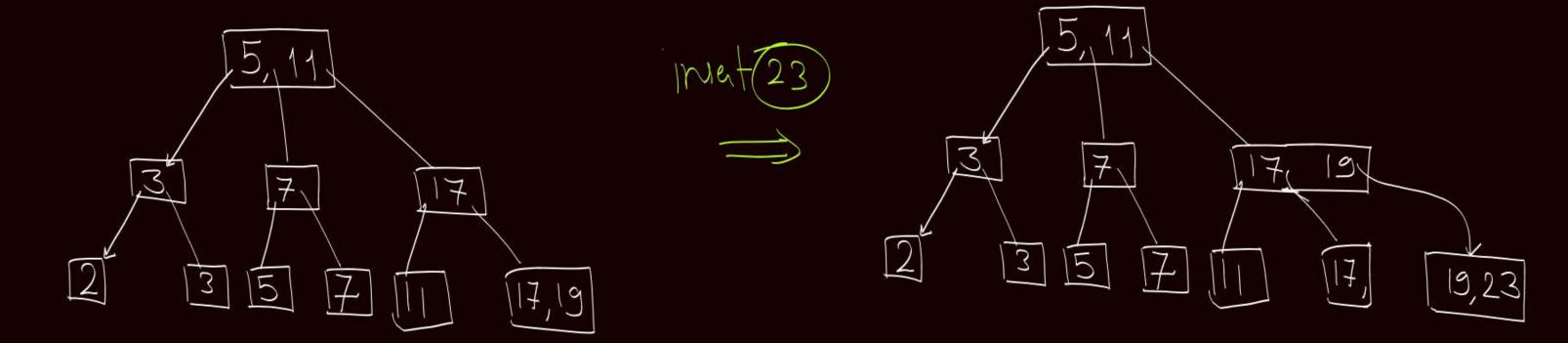
- (a) The root has only one key value 19.
- (b) The root has two key values 17 and 19.
- (c) The root has only one key value 11.
- (d) Key value 31 is alone in a leaf node

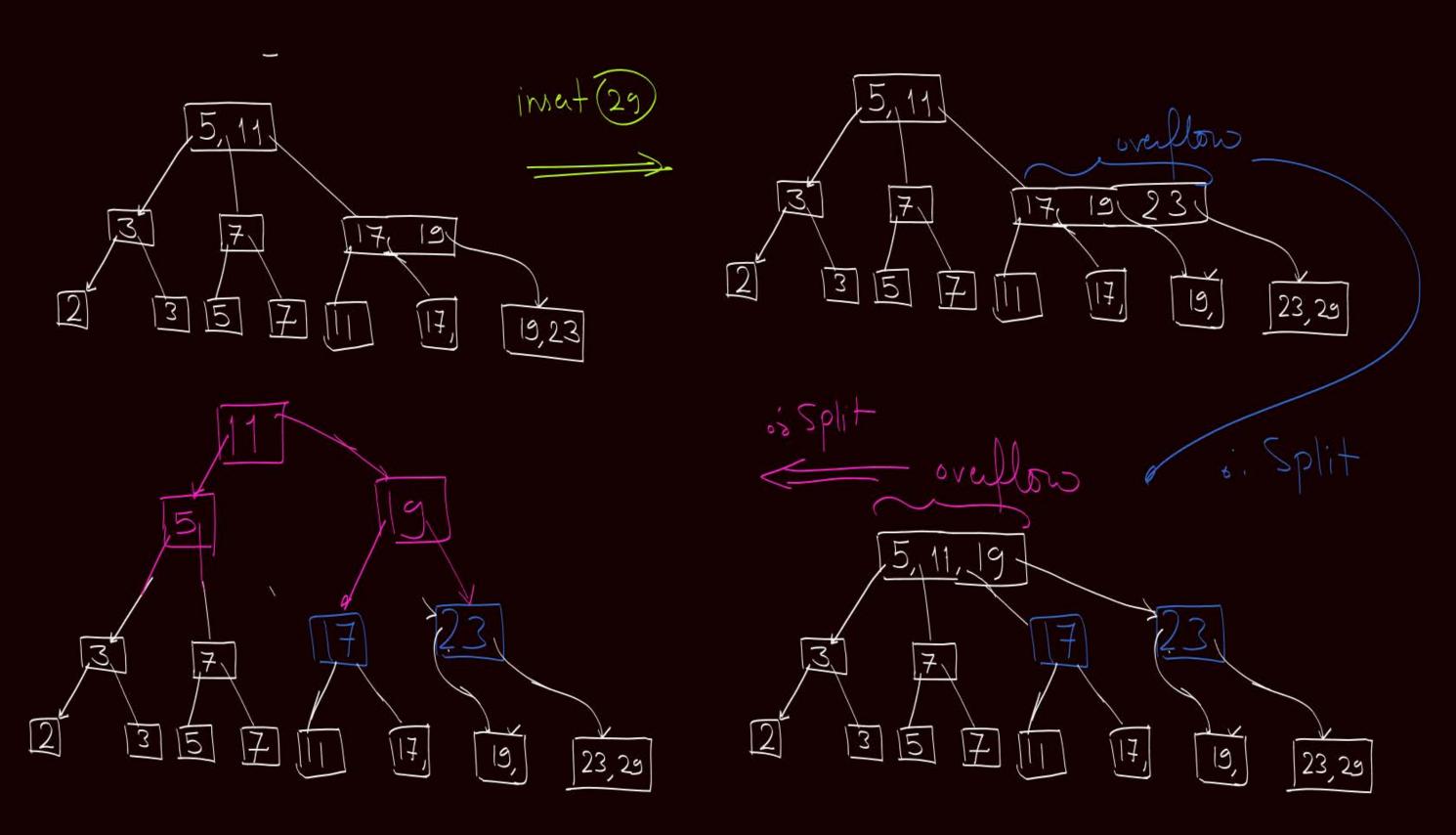
2, 3, 5, 7, 11, 17, 19, 23, 29, 31



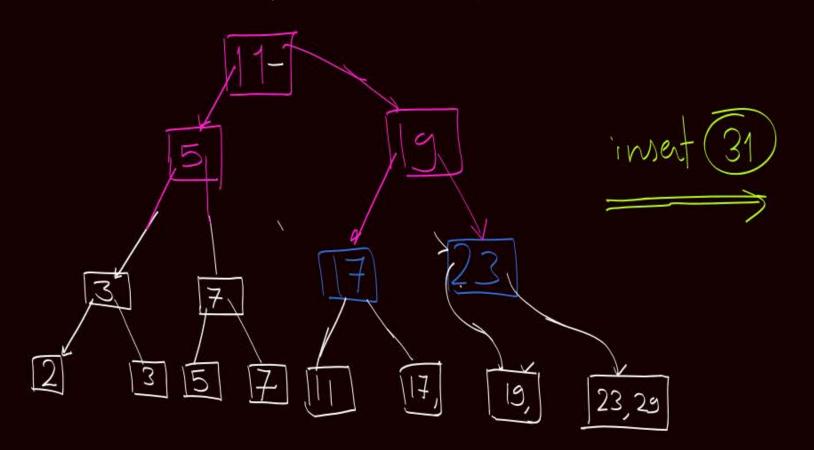


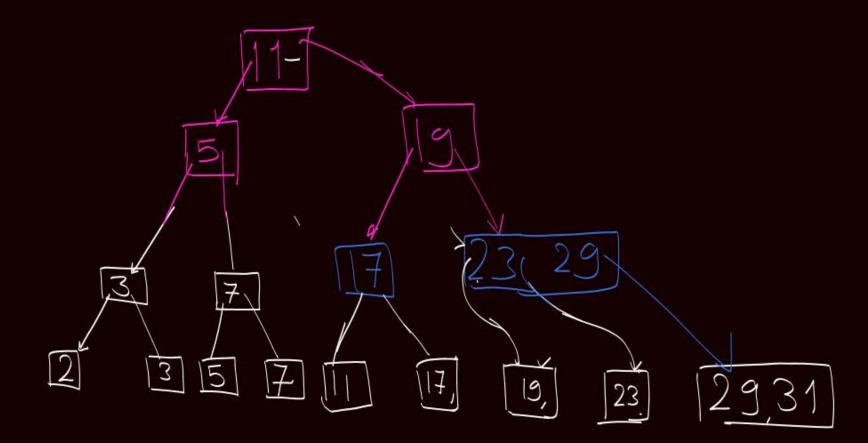
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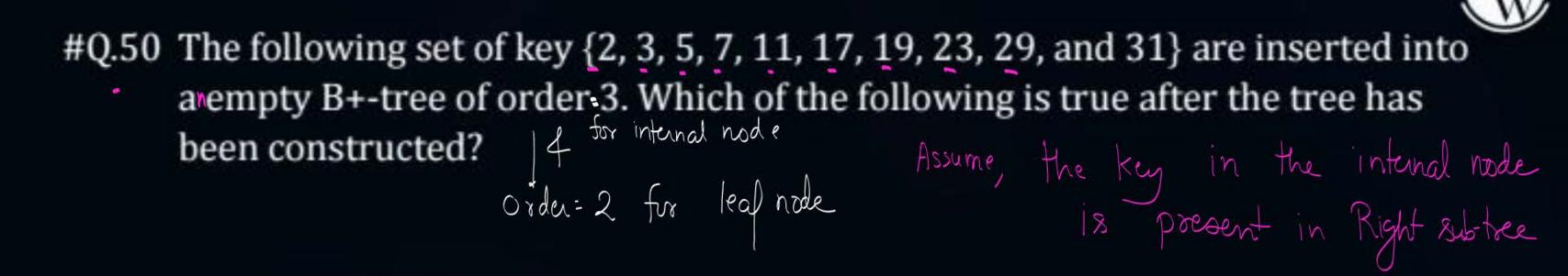




2, 3, 5, 7, 11, 17, 19, 23, 29, 31







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#Q.51 values 23, 65, 37, 60, 46, 92, 48, 71, 56, 59, 18, 21, 10, 74, and 78 are inserted in the given order in a B-tree of order = 4. The number of time the node of B tree splits during insertion is (assume that left biasing is used)\_\_\_\_\_

More Kysin left hand side Node after splitting Max Keys = 3

23, 65, 37, 60, 46, 92, 48, 71, 56, 59, 18, 21, 10, 74, 78

inset 23, 65, 37 =>

23, 37, 65 Invert 60

23,37,60,65 overflow

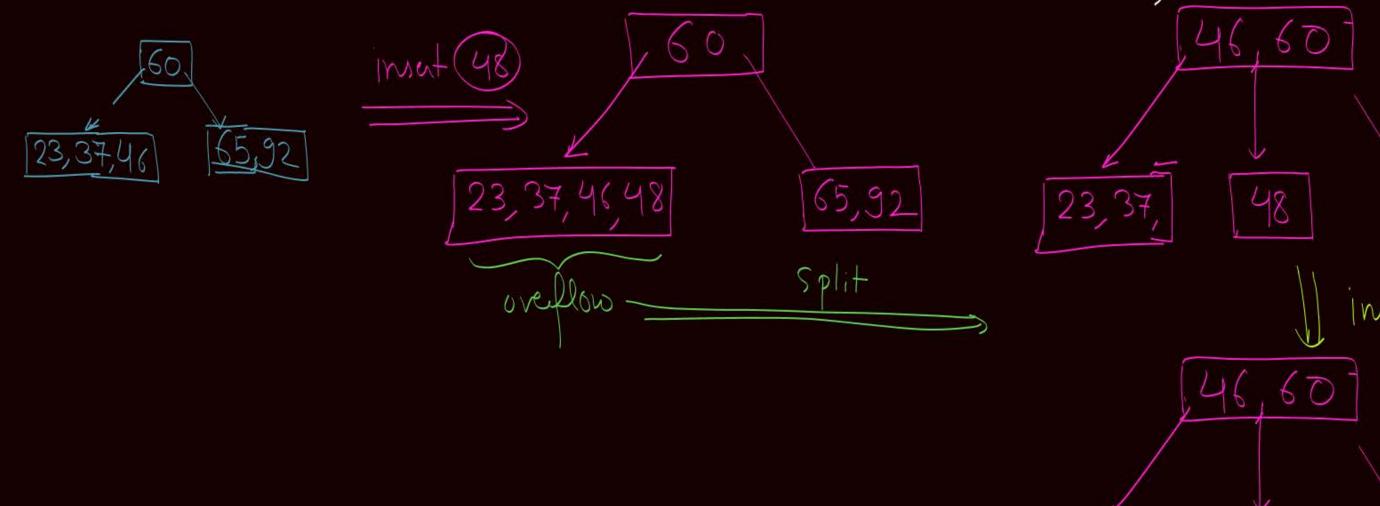
23,37

ment 46 ,60,

inset (92)

Max keys = 3

23, 65, 37, 60, 46, 92, 48, 71, 56, 59, 18, 21, 10, 74, 78



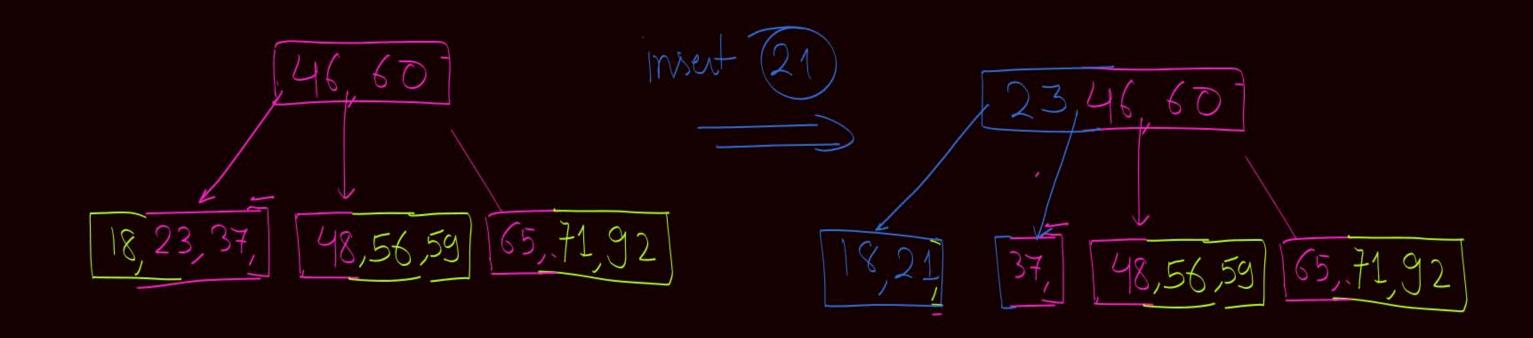
65,92

48

23,37,

Max keys = 3

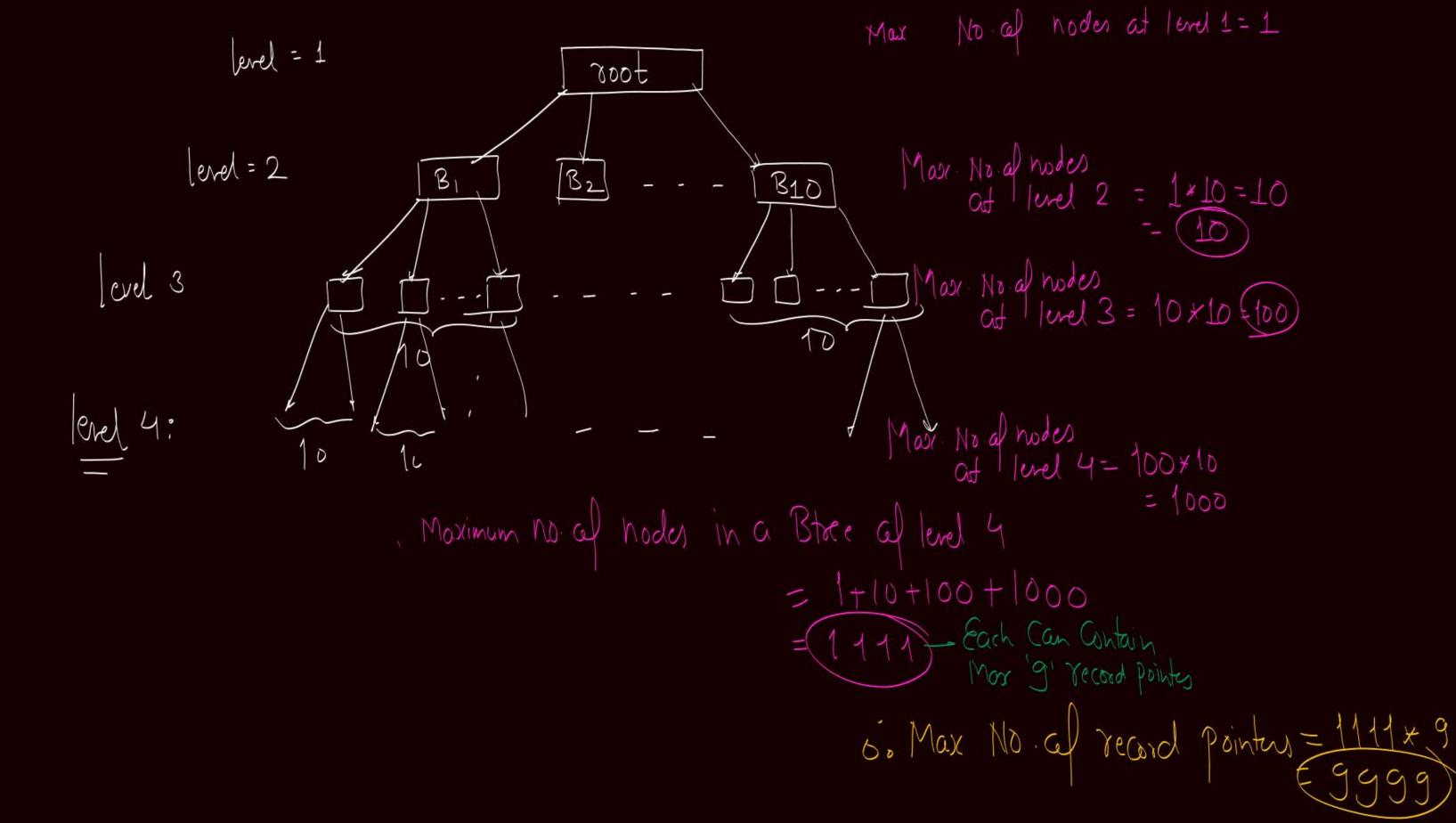
23, 65, 37, 60, 46, 92, 48, 71, 56, 59, 18, 21, 10, 74, 78



Max keys = 3 23, 65, 37, 60, 46, 92, 48, 71, 56, 59, 18, 21, 10, 74, 78 23,46,60 23,46,60 48,56,59 65,71,92 60 mout (78 23,46 More 23,46 74 741 Split on insution (5,71 78,92 92 6x In total 5' 10,18,21 (5,71 48,56,59 48,56,59

Max No. of Child pointer = 10

Max No. of keys = Max No of occord pointer = 10-1 R #Q.52 Consider the B tree of order 10. Assume that the number of levels in the tree is four including root. The maximum number of record pointers that can be stored in B tree is Am - 1999 Maximum ho of record pointers Will be there if each node is full. Le When Node occupancy is maximum





Mai no af dild pointer necessary = \left[\frac{10}{2}\right] = 5

Thin no af Child pointer necessary = \left[\frac{10}{2}\right] = 5

Thin no af Child pto necessary for root = "2"

No col record pointer will be
minimum in the tree, when

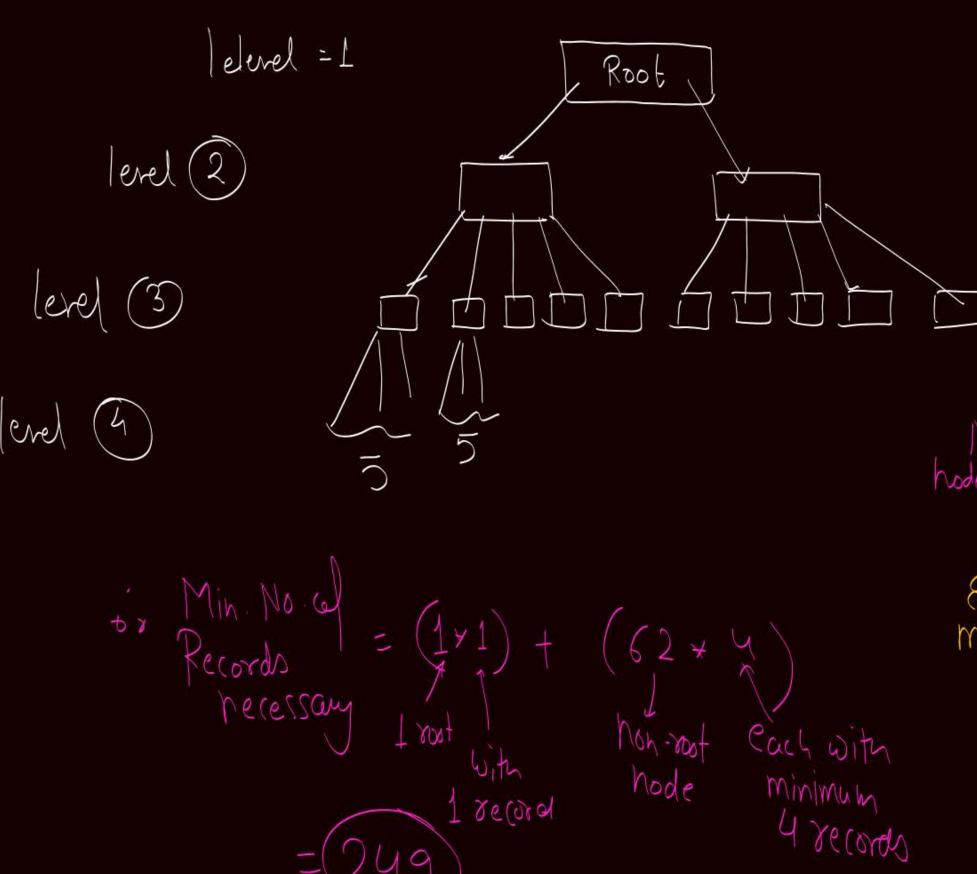
row = \left[\frac{10}{2}\right] = 5

for non-root node \left[\frac{8}{2}\right] \text{Minimum}

Minimum

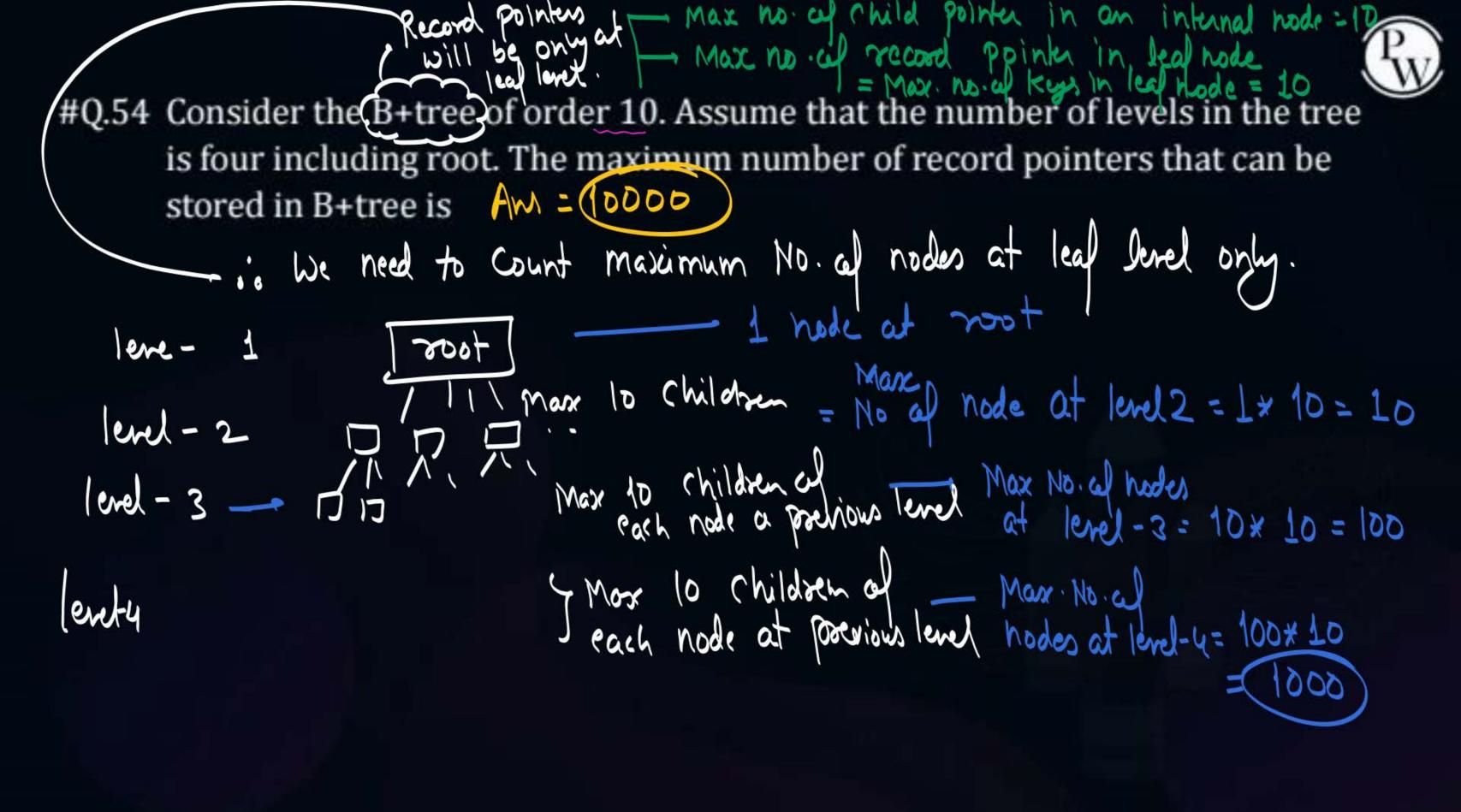
The tree, when

Rode occupancy of each node



Min No of node at level Root Min. No al node at level 2 must Contain Minimum Each one al them Must Contain minimum  $\lceil \frac{p}{2} \rceil - 1 = \lceil \frac{10}{2} \rceil - 1 = 4$  decords

Scord



In a Bt tree of level 4 and order=10, We can have maximum '1000' nodex at leaf level. And each node at leaf level can have maximum '10' record
Pointer ob Moximum No. a) = (1000) x (10) = (10,000)

Max No al nodex Maximum No. a)

at leal node

in each leal node

min no al child ptr at non-root intered node = [0] #Q.55 Consider the B+tree of order 10. Assume that the number of levels in the tree is four including root. The min mum number of record pointers that can be stored in B+tree is \_\_\_\_\_ Min. No. of keys at non-voot

Min No of nodes at lead 1 = 1 leaf node - order = [7] = 5 level-1 Min. no. of nodes at level-2 = 1+2 = 2 lend-2 Each node at porvious Min. No all hoder 10 \* 5- 50 knd will have at level-4 = 10 \* 5- 50 Min '5' child pto

- + Minimum '50' nodes at leaf level.

  + and each node at leaf level must have minimum [10] = 5 record Pointur.

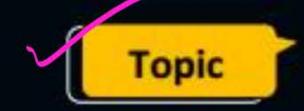
is Min. no. cel accord pointers in a B+ ltree of Oroda = 10 and level = 4

Mh 50 each with hodes min'5' record Pointry



#### 2 mins Summary





File organization and Indexing



# THANK - YOU