

## Last session

- Time-domain filters
- Fourier series & transform
- Frequency-domain filters

## This session

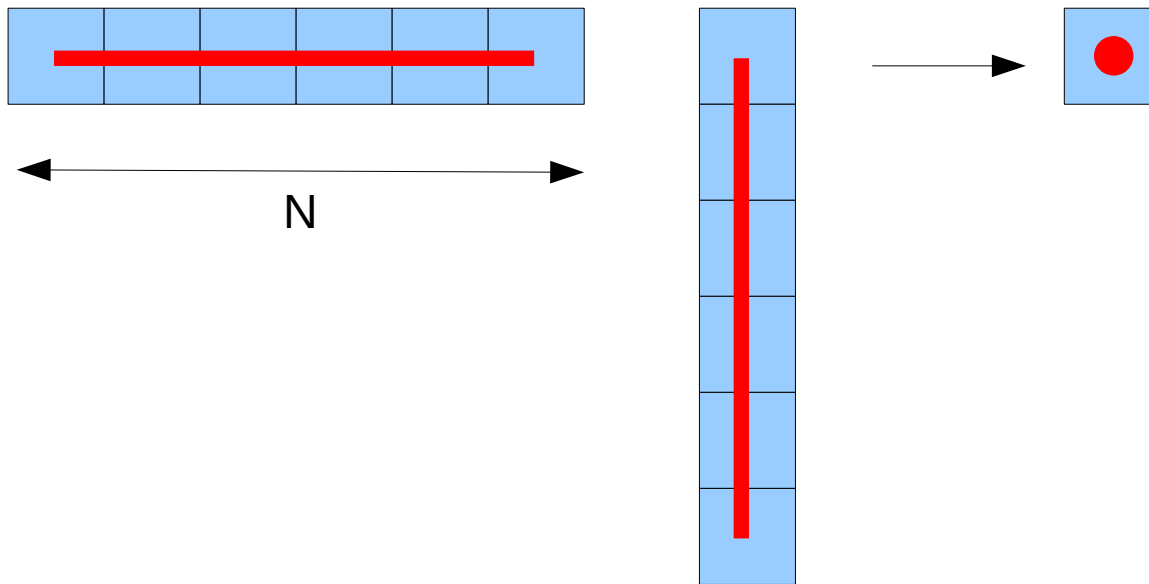
- Complexity and big-O
- The FFT
- Fourier transform pairs

# Concept: “Time Complexity”

- Question: How does program execution time scale with size of input?
- Examples:
  - Dot product:  $\vec{u} \cdot \vec{v}$  Execution time grows as length of vector. (Explanation on board)
  - Matrix multiplication:  $A B$   
Execution time grows as length of vector cubed. (Explanation on board.)
- Finding “performant” algorithms is a main goal of numerical analysis.

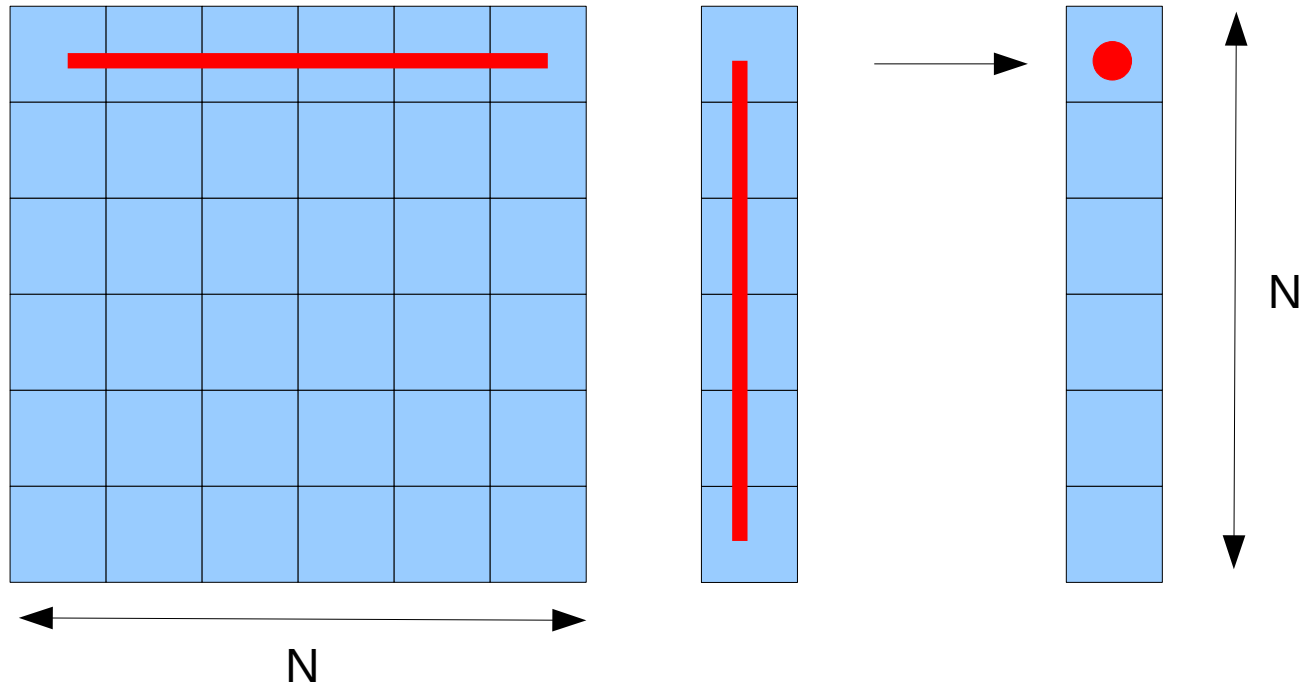
# Dot product time complexity

- Two vectors, each length  $N$



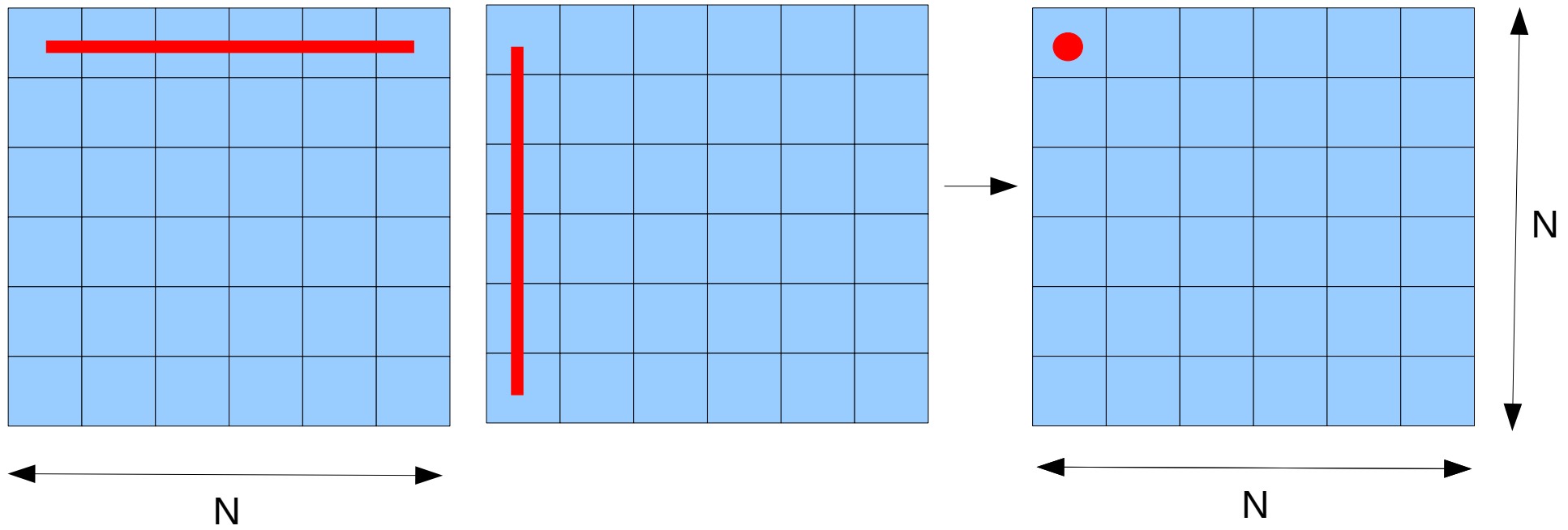
- Complexity  $O(N)$

# Complexity: matrix vector multiply



- Complexity  $O(N^2)$

# Complexity: matrix matrix multiply



- Complexity  $O(N^3)$

# Complexity: Determinant (naive)

- 2x2 matrix

a	b
c	d

$$\text{Det} = ad - bc$$

2 mults

- 3x3 matrix


+


+


3x2 mults

- 4x4 matrix





4x3x2 mults

# Complexity: Determinant (naive)

- In general case, for  $N \times N$  matrix, the naive computation of determinant has  $O(N!)$  complexity!
- This is terrible.
- Avoid computing determinants, particularly the naive way.

# Big-O notation

- Consider growth of execution time as input size goes to infinity.
- Only keep the fastest growing part.  
Example:  $O(N^2 + N) = O(N^2)$
- Ignore any constant multiplicative factors – we are only interested in *scaling*.
- $O(N)$  – Execution time grows linearly with input data size
- $O(\log(N))$  – Execution time grows as log of input data size
- $O(N^2)$  – Execution time grows as square of input data size.



# Examples of different algorithms

- $O(1)$  – Return first element of vector.
- $O(N)$  -- Sum elements of a vector. Search an unordered list.
- $O(N^2)$  – Multiply two  $N$  digit numbers. Bubble sort.
- $O(\log N)$  – Binary search on ordered list.  
Example: finding a name in a phone book of  $N$  names.
- $O(N^3)$  – matrix multiply (Naive).
- $O(N!)$  -- Computing determinant of matrix using permutations (what you learned as undergraduate).

# Example $O(1)$

- Function which computes a number and returns.
- Independent of input data size
  - This fcn works only on scalars.

```
function y = derivative( f, x, h )  
    % This fcn returns the forward-difference  
    % derivative of f at x using step h  
    y = (f(x+h) - f(x)) / (h);  
end
```

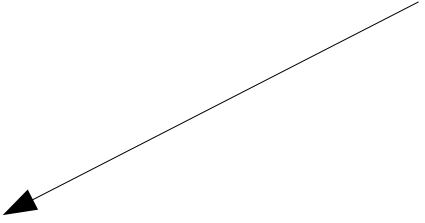
# Example $O(N^3)$ – matrix multiplication

```
function z = mymatmul(x, y)
    % Matrix multiplication the naive way, using loops.
    % This algorithm is  $O(n^3)$ 
    %  $z = x*y$ 
    % size(x) = [n, m]
    % size(y) = [m, p]
    % size(z) = [n, p]

    [n, m] = size(x);
    [m, p] = size(y);

    z = zeros(n, p);
    for row = 1:n
        for col = 1:p
            for idx = 1:m
                z(row, col) = z(row, col) + x(row, idx)*y(idx, col);
            end
        end
    end
end
```

Three nested loops



- In this class we generally identify the complexity of an implementation by the number of nested loops.
- Performant algorithms have low-order complexity.
  - $O(N)$  is good.
  - $O(N \log N)$  is good.
  - $O(N^p)$  where  $p \gg 1$  not so good.
- Consider (naive) matrix multiply --  $O(N^3)$ 
  - Suppose  $N=10$  takes 1 sec.
  - Then  $N = 100$  takes 1000 sec.
  - $N = 1000$  takes 1e6 seconds.

# Numerical algorithm: Fast Fourier Transform

- Algorithm implementing Fourier Transform for sampled signal.
- “The most important numerical algorithm of our lifetime”, Gilbert Strang, American Scientist 82 (3): 253 (1994).
- Operates on sampled (discrete) signal.
- **Complexity:  $O(N \log N)$**
- Matlab: `fft()`, `ifft()`

Continuous Fourier transform

$$Y(\omega) = \int_{-\infty}^{\infty} dt y(t) e^{-i\omega t}$$

$$y(t) = \frac{1}{2\pi} \int_{-\infty}^{\infty} d\omega Y(\omega) e^{i\omega t}$$

# Fourier transform – from continuous to discrete

## Continuous Fourier transform

$$Y(\omega) = \int_{-\infty}^{\infty} dt y(t) e^{-i\omega t}$$

$$y(t) = \frac{1}{2\pi} \int_{-\infty}^{\infty} d\omega Y(\omega) e^{i\omega t}$$

## Discrete Fourier transform

$$Y_k = \sum_{n=0}^{N-1} y_n e^{-i2\pi nk/N}$$

$$y_n = \frac{1}{N} \sum_{k=0}^{N-1} Y_k e^{i2\pi nk/N}$$

$$\omega \Leftrightarrow 2\pi k/N \quad \text{“Frequency”}$$

$$Y(\omega) \Leftrightarrow Y_k \quad \text{“Frequency domain”}$$

$$y(t) \Leftrightarrow y_n \quad \text{“Time domain”}$$

# Computation of DFT using FFT

- $Y_f = \text{fft}(y)$  takes vector of length  $N$  and returns vector of length  $N$ .

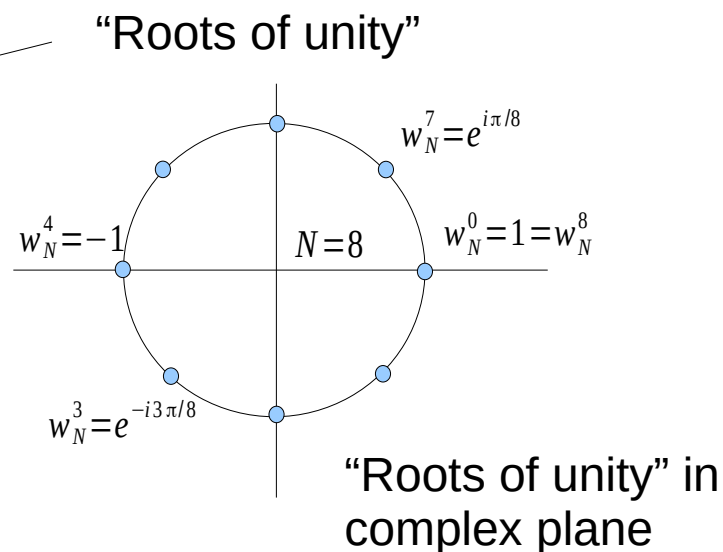
$$Y_k = \sum_{n=0}^{N-1} y_n e^{-i2\pi nk/N}$$

Discrete Fourier transform

- Define:  $w_N^k = e^{-i2\pi k/N}$

- So:

$$Y_k = \sum_{n=0}^{N-1} y_n (w_N^k)^n$$



# Consider 4 point example

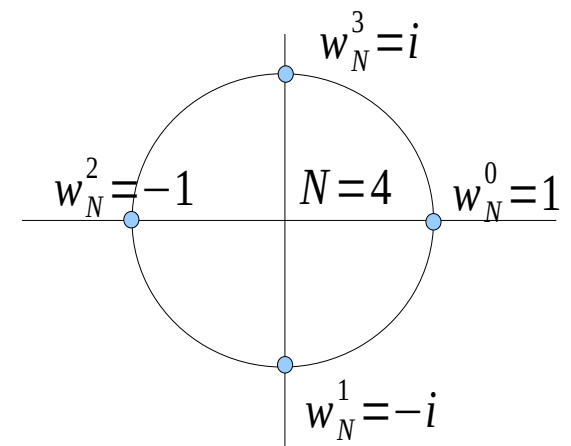
$$Y_k = \sum_{n=0}^{N-1} y_n (w_N^k)^n$$

$$Y_0 = y_0 (w_N^0)^0 + y_1 (w_N^0)^1 + y_2 (w_N^0)^2 + y_3 (w_N^0)^3$$

$$Y_1 = y_0 (w_N^1)^0 + y_1 (w_N^1)^1 + y_2 (w_N^1)^2 + y_3 (w_N^1)^3$$

$$Y_2 = y_0 (w_N^2)^0 + y_1 (w_N^2)^1 + y_2 (w_N^2)^2 + y_3 (w_N^2)^3$$

$$Y_3 = y_0 (w_N^3)^0 + y_1 (w_N^3)^1 + y_2 (w_N^3)^2 + y_3 (w_N^3)^3$$





# Evaluate w coefficients

$$Y_0 = y_0(w_N^0)^0 + y_1(w_N^0)^1 + y_2(w_N^0)^2 + y_3(w_N^0)^3$$

$$Y_1 = y_0(w_N^1)^0 + y_1(w_N^1)^1 + y_2(w_N^1)^2 + y_3(w_N^1)^3$$

$$Y_2 = y_0(w_N^2)^0 + y_1(w_N^2)^1 + y_2(w_N^2)^2 + y_3(w_N^2)^3$$

$$Y_3 = y_0(w_N^3)^0 + y_1(w_N^3)^1 + y_2(w_N^3)^2 + y_3(w_N^3)^3$$

$$Y_0 = y_0 + y_1 + y_2 + y_3$$

$$Y_1 = y_0 - iy_1 - y_2 + iy_3$$

$$Y_2 = y_0 - y_1 + y_2 - y_3$$

$$Y_3 = y_0 + iy_1 - y_2 - iy_3$$

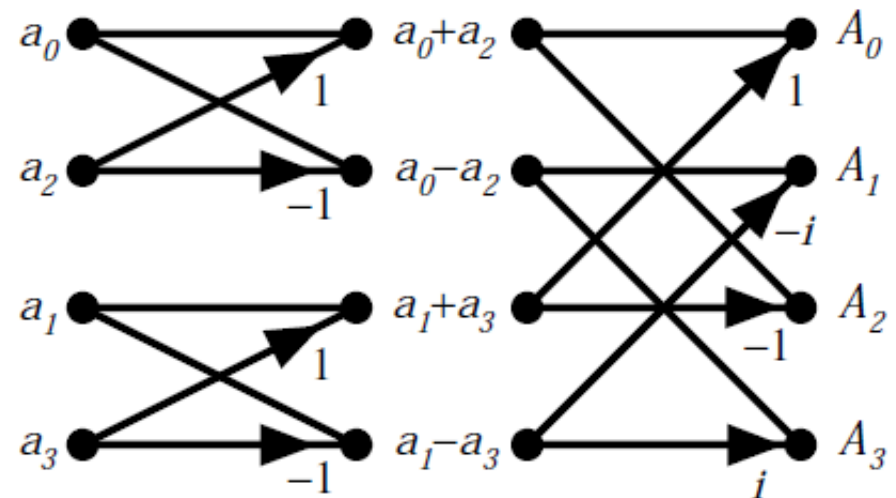
$$Y_0 = (y_0 + y_2) + (y_1 + y_3)$$

$$Y_1 = (y_0 - y_2) - i(y_1 - y_3)$$

$$Y_2 = (y_0 + y_2) - (y_1 + y_3)$$

$$Y_3 = (y_0 - y_2) + i(y_1 - y_3)$$

Think of this as two step computation.....

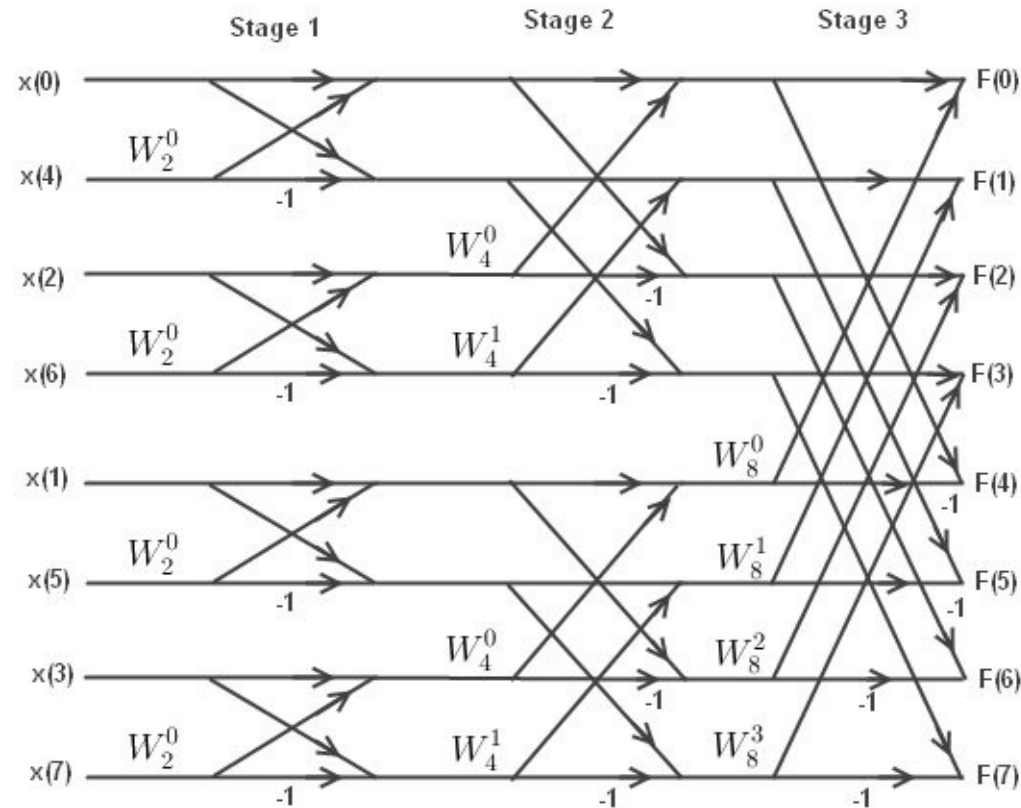


"Butterfly diagram"

*I have computed 4 output values in 4\*2 multiplications*

# “Butterfly diagram”

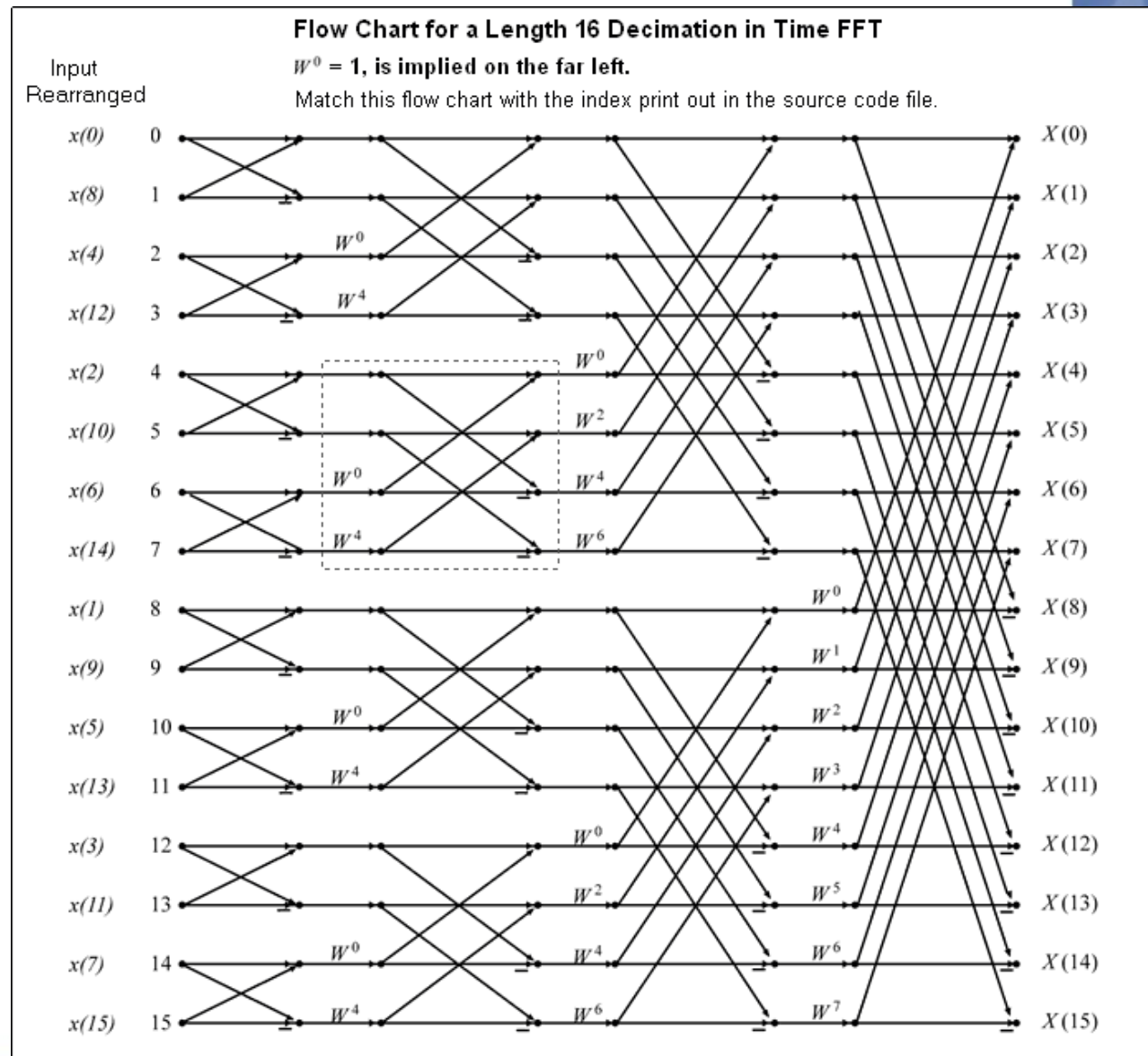
An 8 Input Butterfly. Note, you double a 4 input butterfly, extend output lines, then connect the upper and lower butterflies together with diagonal lines.



- 8 input values, 3 layers of multiplication.
- $\text{Log}_2(8) = 3$
- Number of multiplications =  $n \log n$

# FFT is $O(n \log n)$

- 4 input points  
– 4\*2 operations
- 8 input points  
– 8\*3 operations
- 16 input points  
– 16\*4 operations



# Different versions of FFT

- FFT (fast Fourier transform)
  - $\exp(-i\omega t)$  basis functions.
  - Input & output complex
- DCT (discrete cosine transform)
  - $\cos(\omega t)$  basis functions
  - Input & output real.
- DST (discrete sine transform)
  - $\sin(\omega t)$  basis functions
  - Input & output real

# FFTW

- Matlab uses FFTW to compute FFTs.



## Introduction

FFTW is a C subroutine library for computing the discrete Fourier transform (DFT) in one or more dimensions, of arbitrary input size, and of both real and complex data (as well as of even/odd data, i.e. the discrete cosine/sine transforms or DCT/DST). We believe that FFTW, which is [free software](#), should become the [FFT](#) library of choice for most applications.

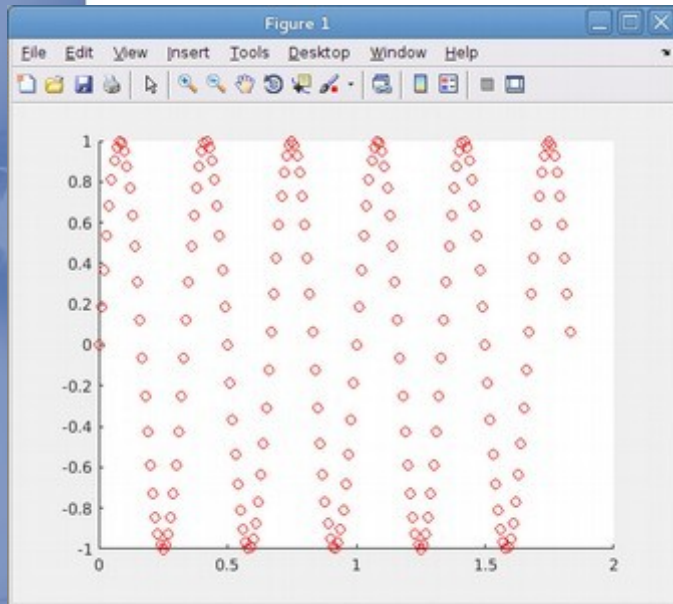
The latest official release of FFTW is version **3.3.6**, available from [our download page](#). Version 3.3 introduced support for the AVX x86 extensions, a distributed-memory implementation on top of MPI, and a Fortran 2003 API. Version 3.3.1 introduced support for the ARM Neon extensions. See the [release notes](#) for more information.

The FFTW package was developed at [MIT](#) by [Matteo Frigo](#) and [Steven G. Johnson](#).

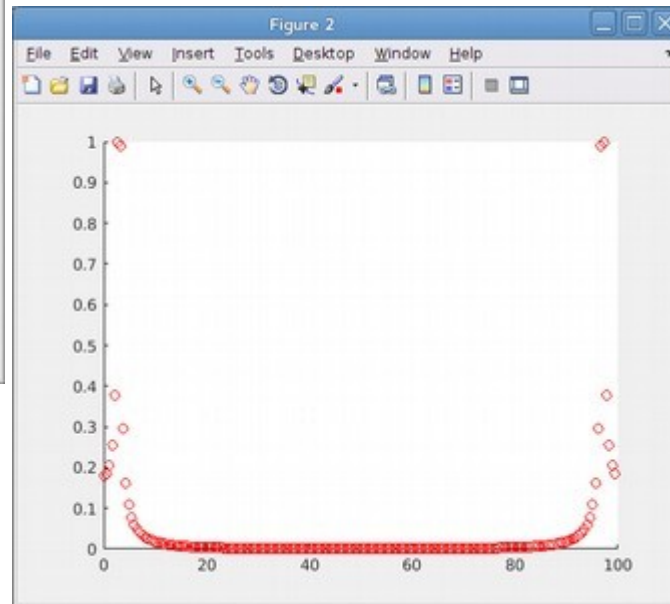
Our [benchmarks](#), performed on a variety of platforms, show that FFTW's performance is typically superior to that of other publicly available FFT software, and is even competitive with vendor-tuned codes. In contrast to



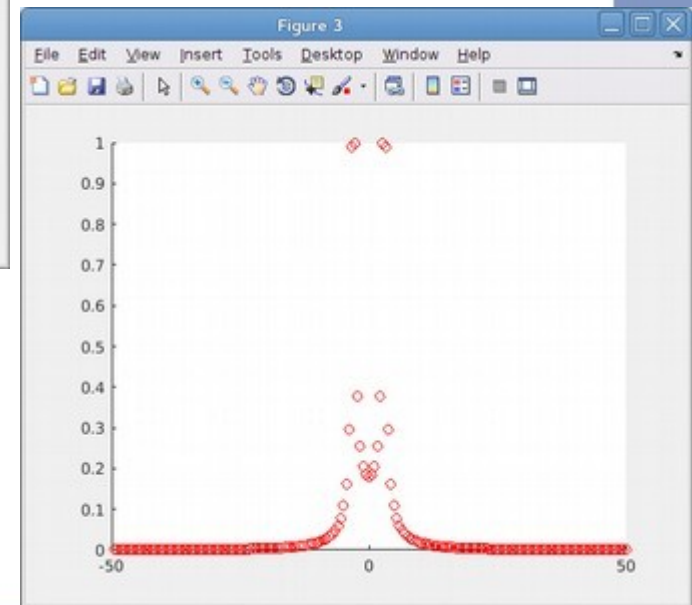
# Fourier transform of sine wave



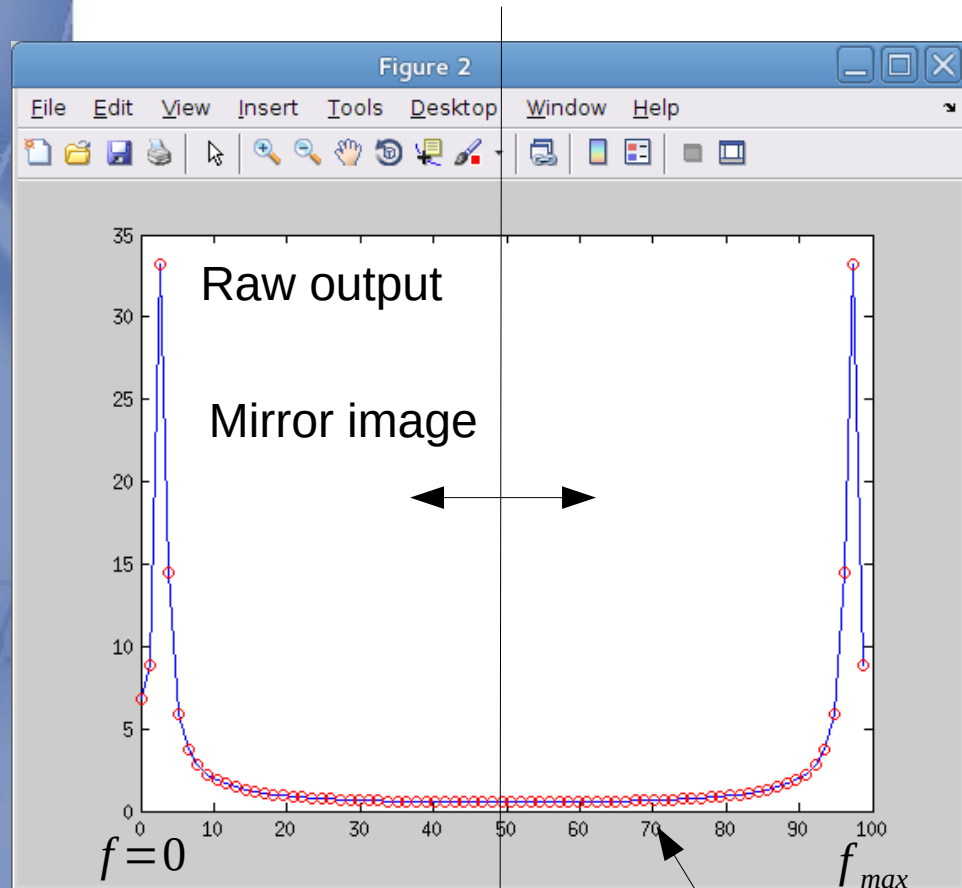
`fft()`



`fftshift()`

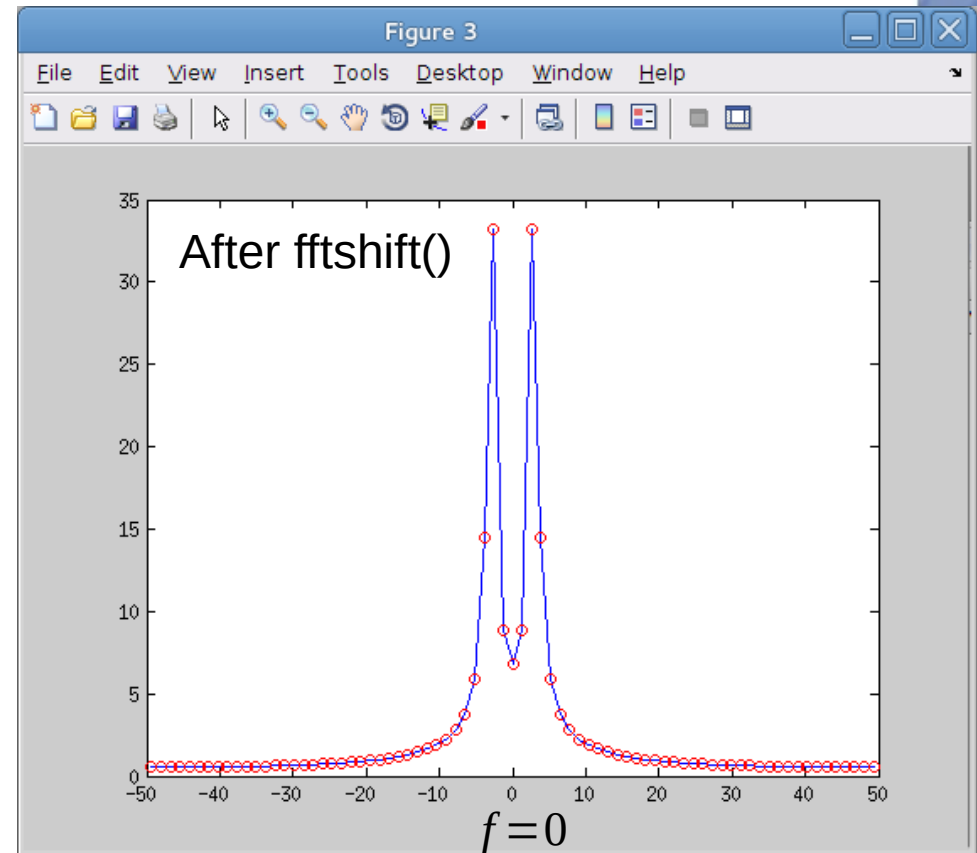


# Raw output of FFT



$$f = \frac{f_{max}}{2}$$

No new information in upper 1/2 of spectrum



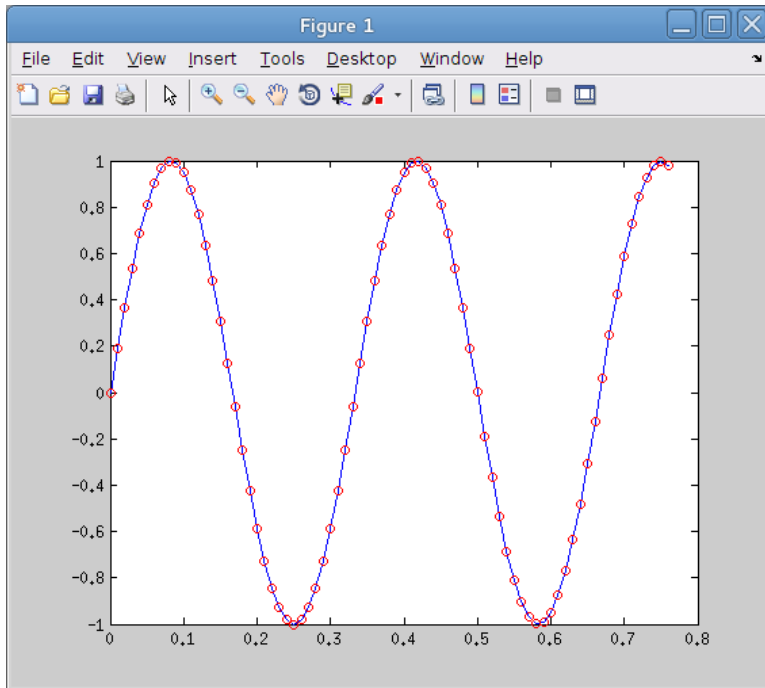
$$f = -\frac{f_{max}}{2}$$

$$f = \frac{f_{max}}{2}$$

`fftshift()` folds upper 1/2 of spectrum down to give two-sided spectrum

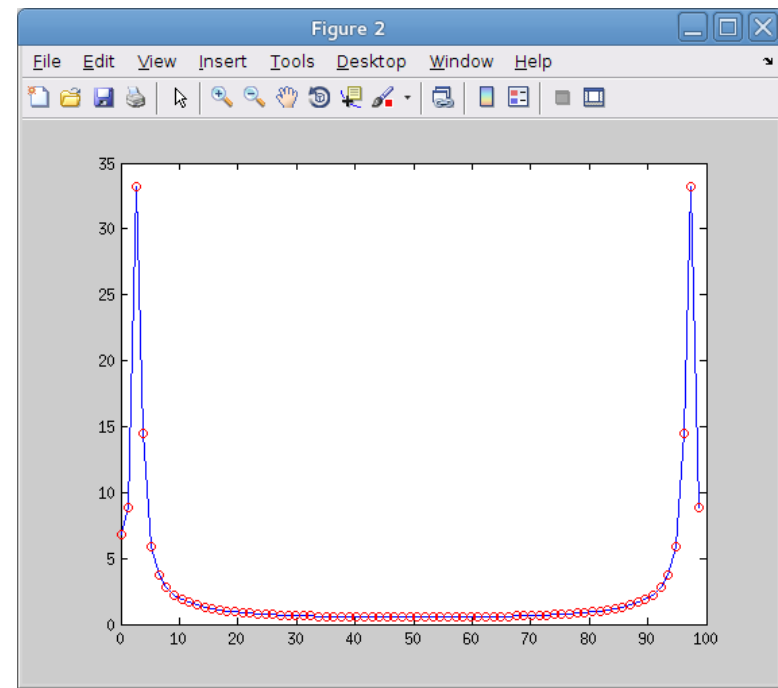
# time $\Leftrightarrow$ frequency units

Input signal



Number of samples  $N$   
Sample period  $\Delta t$   
Sample length  $T = (N - 1) \Delta t$   
Sample frequency  $f_s = \frac{1}{\Delta t}$

Output FFT

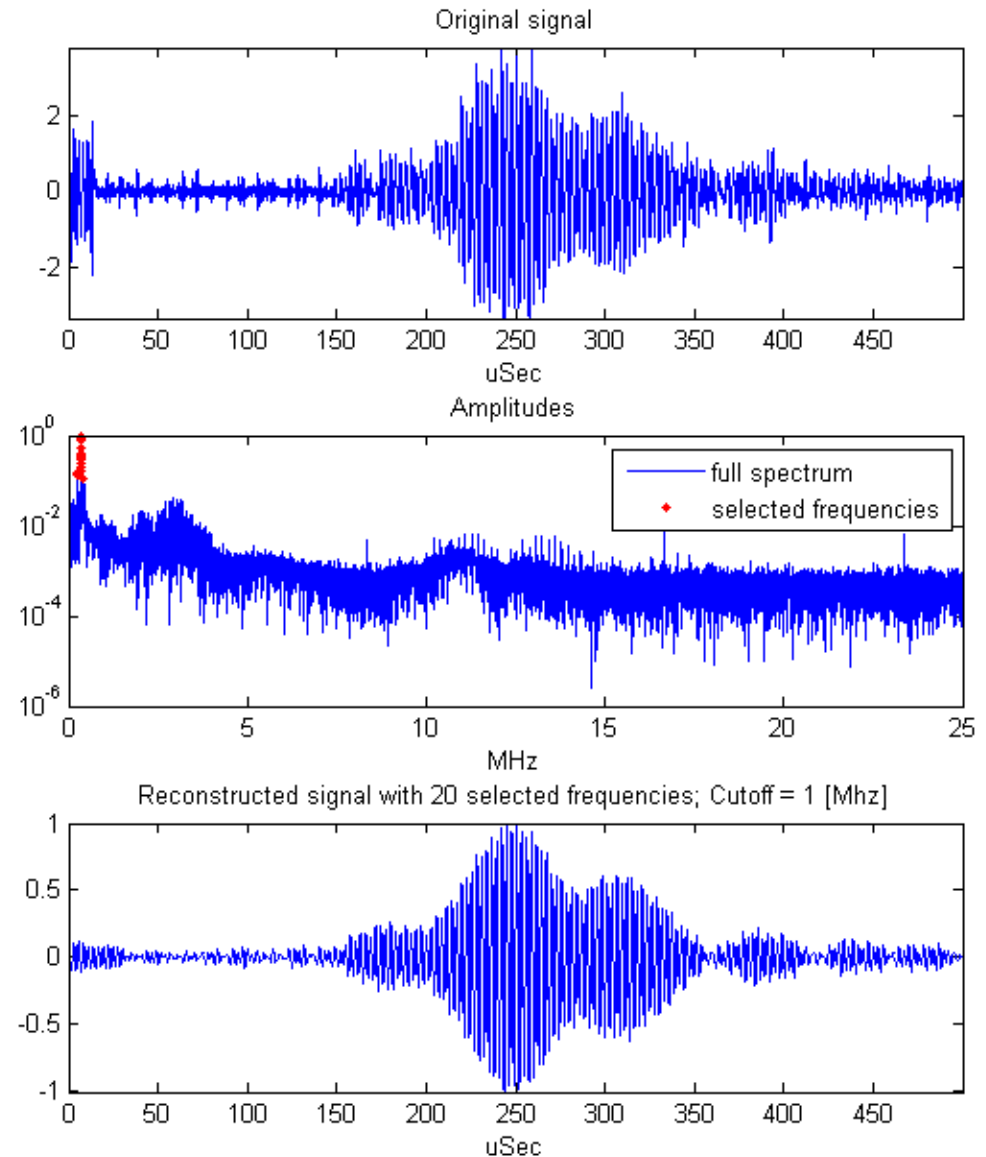


Number of samples  $M = N$   
Maximum frequency  $\frac{M-1}{M} f_s$   
Frequency step  $\Delta f = \frac{f_s}{M}$



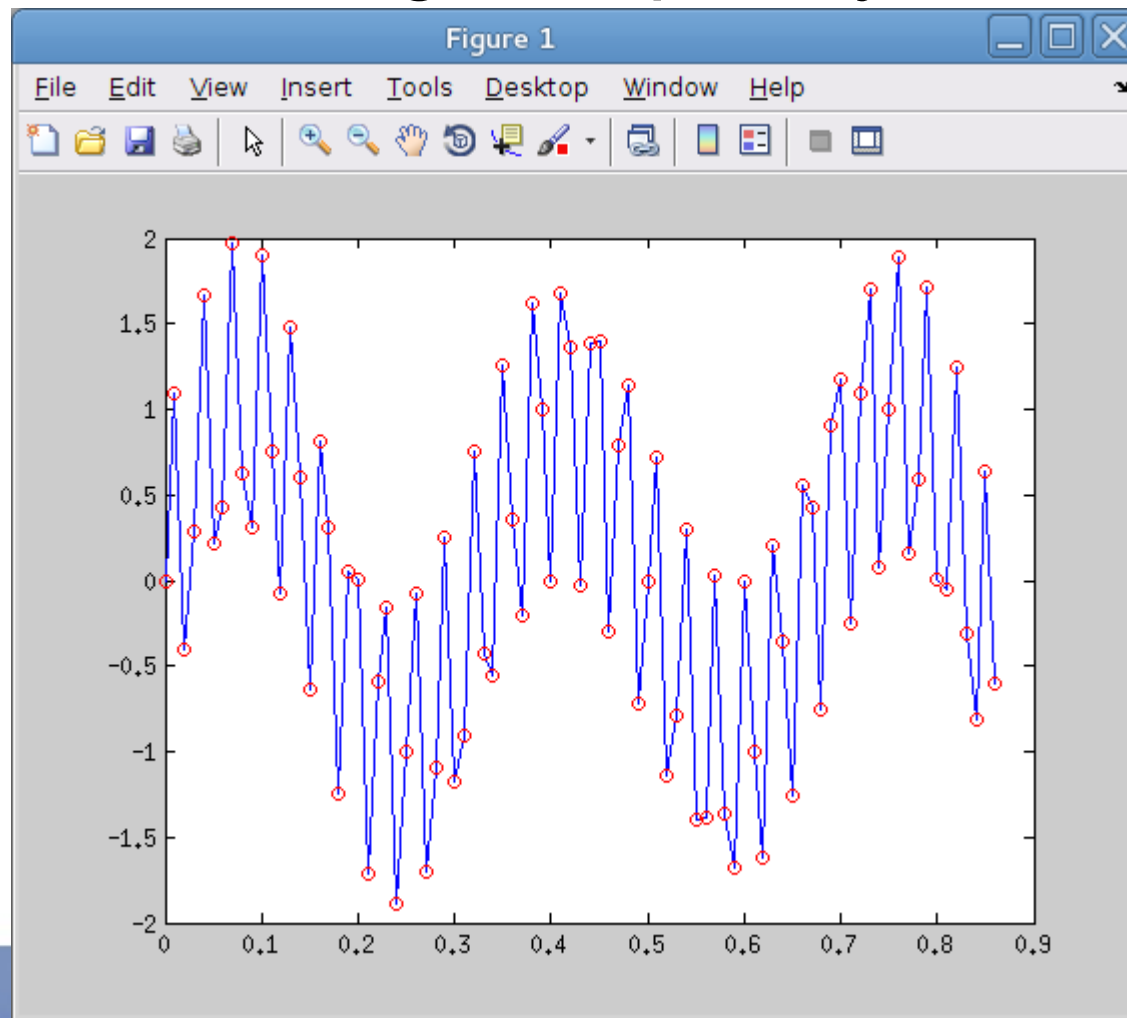
# Our goal: filter signals using FFT

1. FFT your input time-domain signal.
2. Select frequencies you want to keep.
3. Inverse FFT to get filtered time-domain signal

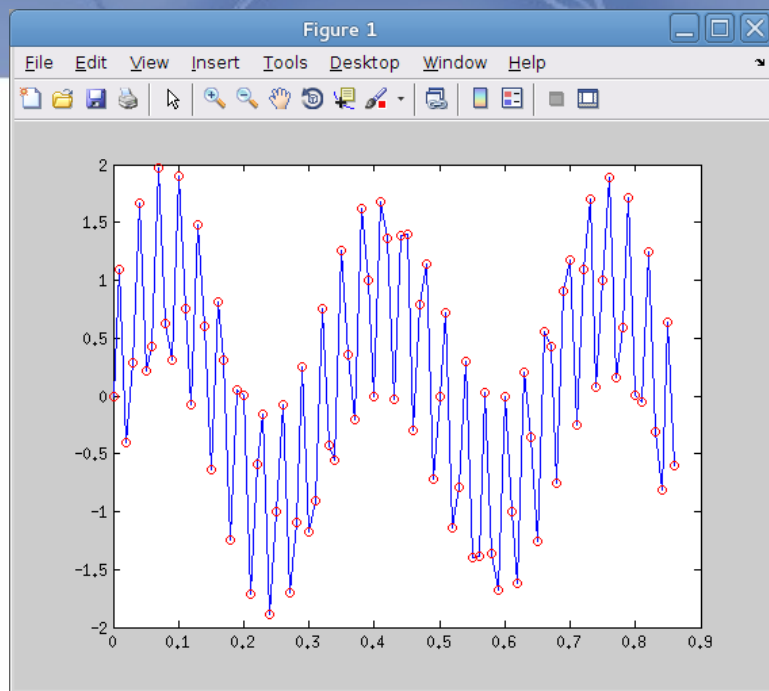


# Example: low-pass filter

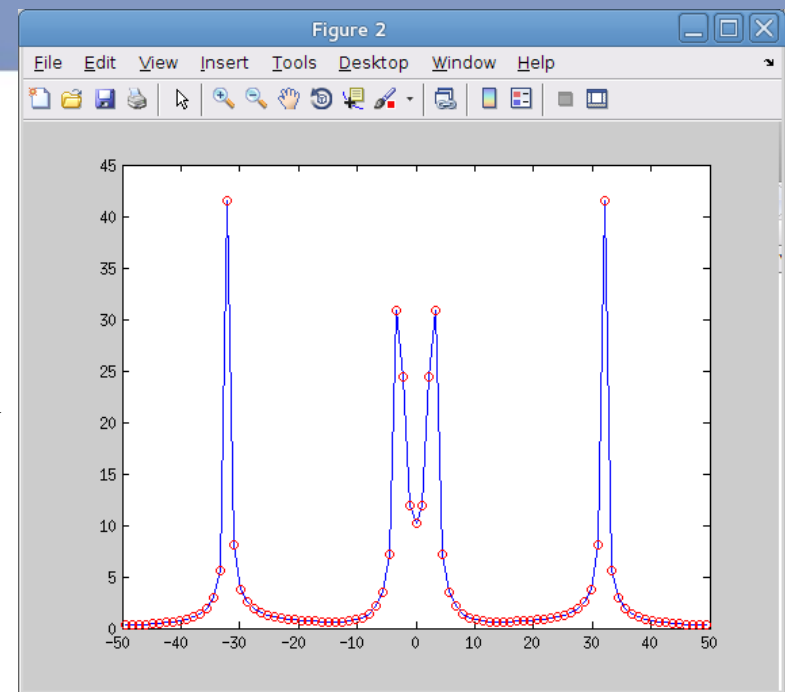
- Two sine waves: 3Hz and 32Hz.
- Want to cut out high-frequency sine wave.



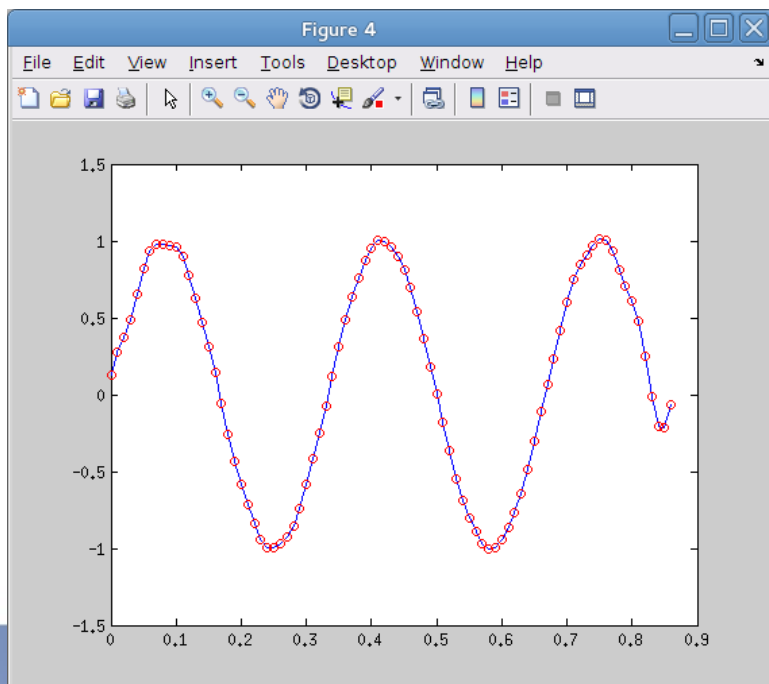
Look in  
fft\_playpen  
for code



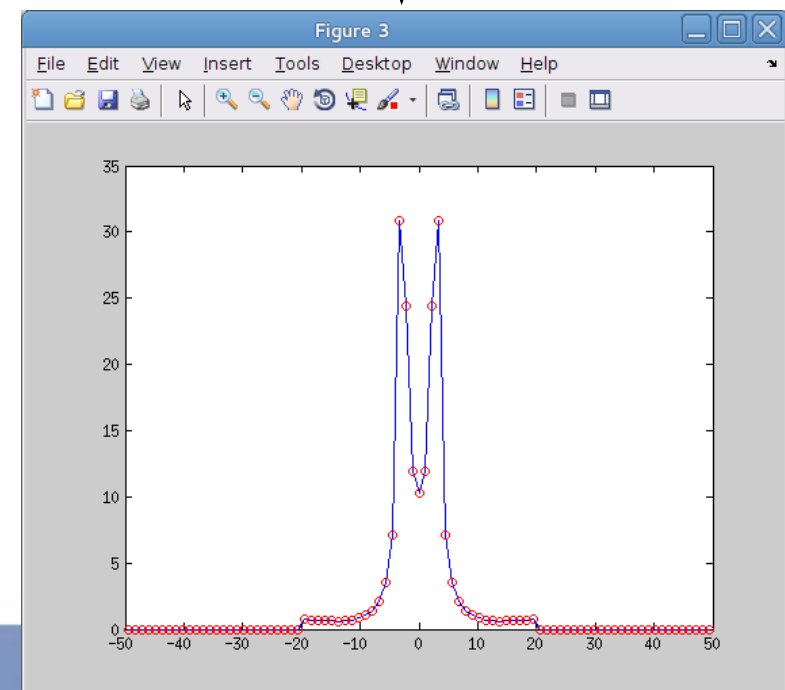
fft



Filter



ifft



```

function filter_sines()
    % This fcn creates two sine waves of different frequencies.
    % It then FFTs the time-domain signal and zeros out
    % the high frequency stuff (low-pass filter), then
    % does ifft.

    M = 87;      % Number of sample points
    dt = .01;    % 10mS sample period.
    fs = 1/dt;   % Sample freq
    t = linspace(0, (M-1)*dt, M); % Vector of timestamps

    % Create sine waves at 3 and 32 Hz
    w1 = 3;      % 3 Hz signal
    w2 = 32;
    x = sin(2*pi*w1*t) + sin(2*pi*w2*t);      samples
    figure(1)
    plot(t, x)
    hold on
    plot(t, x, 'ro')

    % Now do FFT of input signal
    Xf = fft(x); % Fast Fourier Transform
    w = linspace(0, (M-1)*(fs/M), M);

```

```

% Shift over negative freqs on frequency axis in prep for fftshift
w1 = w;
w1(w >= fs/2) = w(w >= fs/2) - fs;
ws = fftshift(w1);
Xfs = fftshift(Xf);
figure(2)
plot(ws, abs(Xfs))
hold on
plot(ws, abs(Xfs), 'ro')

% Now zero out all frequency components above 20 Hz.
idx = find(abs(ws) > 20);
Xfs(idx) = 0;
figure(3)
plot(ws, abs(Xfs))
hold on
plot(ws, abs(Xfs), 'ro')

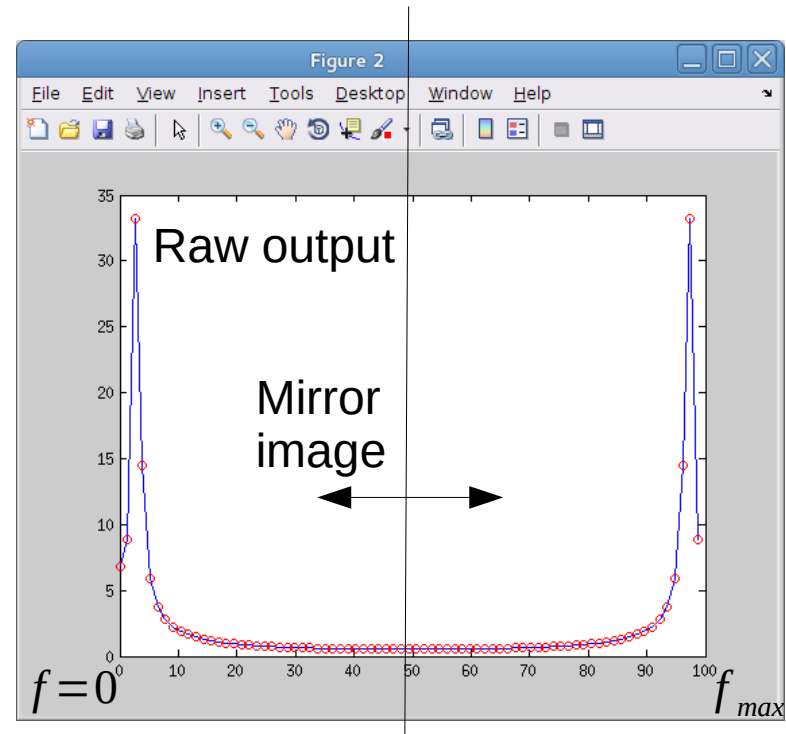
% Now ifft signal back to time domain
Xf = ifftshift(Xfs);
xnew = ifft(Xf);
figure(4)
plot(t, real(xnew))
hold on
plot(t, real(xnew), 'ro')

```

end

# Nyquist frequency

- Upper 1/2 of spectrum from FFT is redundant (no new information).
- Nyquist frequency is maximum frequency of signal which can be correctly captured and reconstructed using FFT.



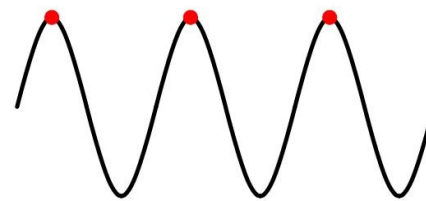
$$f_{\text{Nyquist}} = \frac{f_{\text{max}}}{2}$$

# Sampling theorem

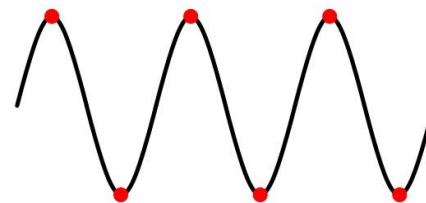
- Our usual goal in using the FFT:
  - Capture some time-domain signal
  - FFT the signal
  - Do some operations in frequency domain (filter the signal)
  - Inverse FFT to reconstruct time-domain signal.
- Sampling theorem: You must sample the signal at  $f_s \geq 2 \times$  maximum frequency in signal to avoid distortion in reconstructed signal

# Aliasing

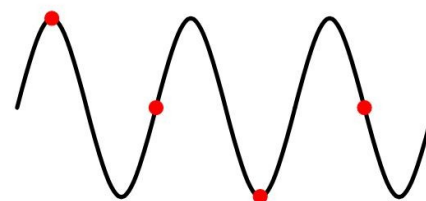
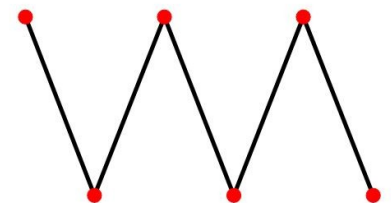
- If signal is not sampled fast enough, the sampled signal does not represent the actual (continuous) signal.
- For correct reconstruction of the actual signal, you must have  
 $f_s \geq$  Maximum frequency component in signal
- If you don't, the effect is called “aliasing”.



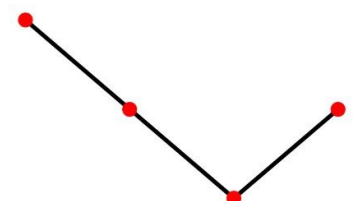
A  
Sampled at  $f$



B  
Sampled at  $2f$



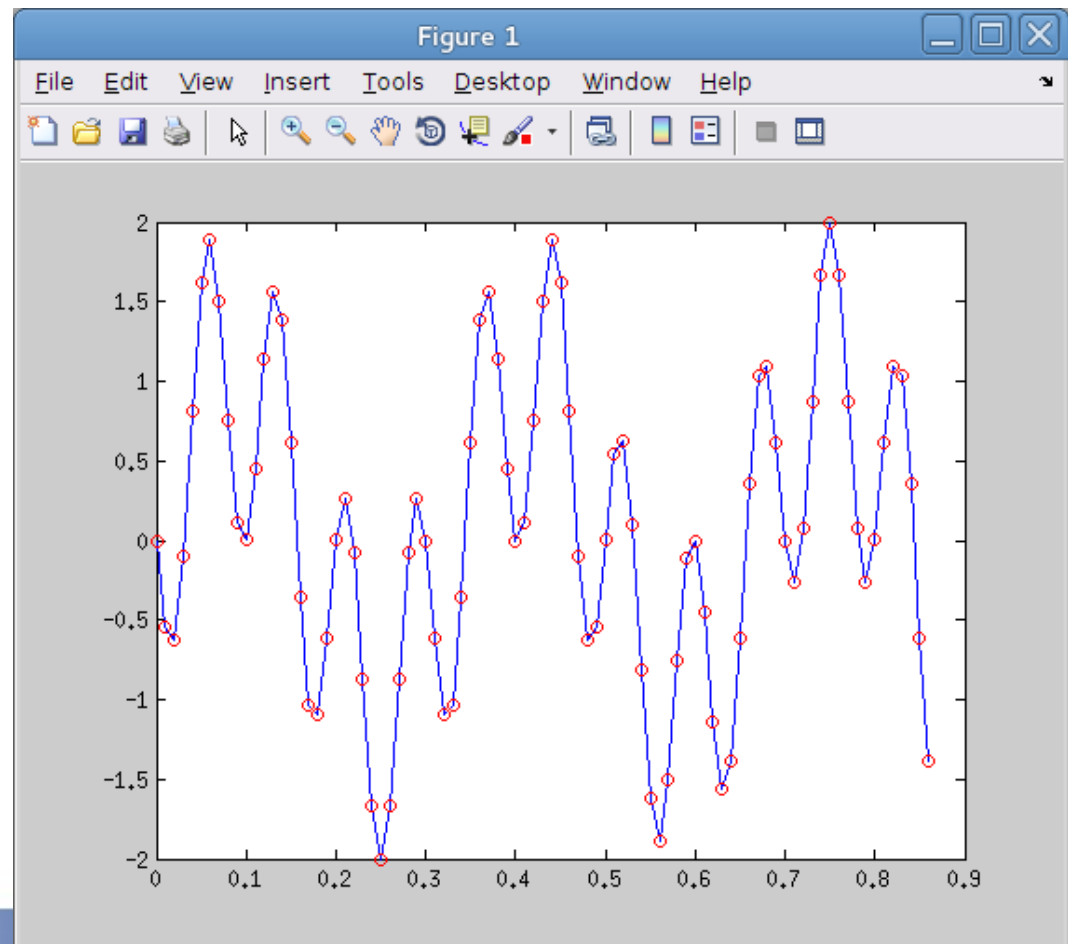
C  
Sampled at  $4f/3$



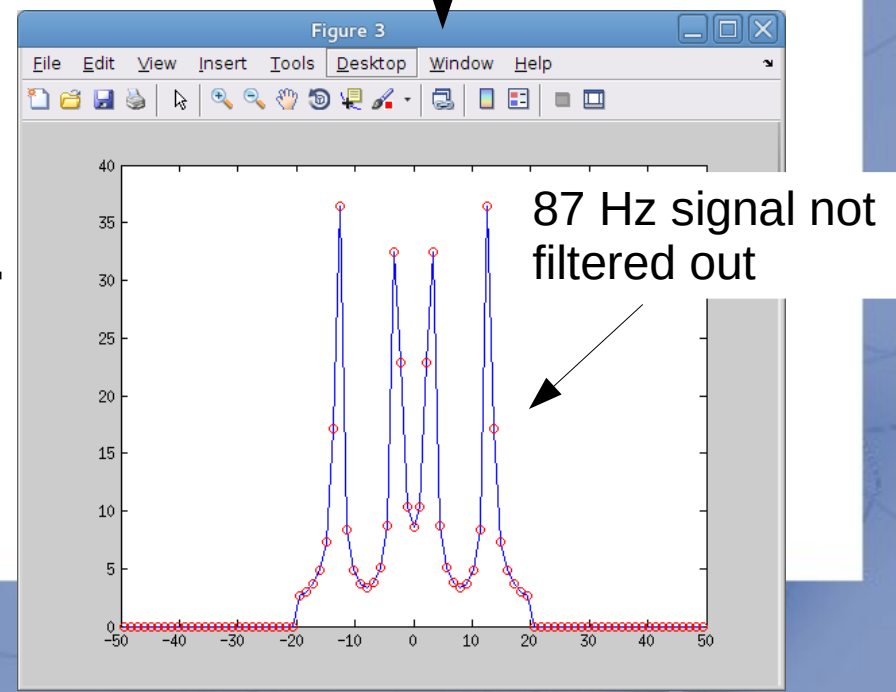
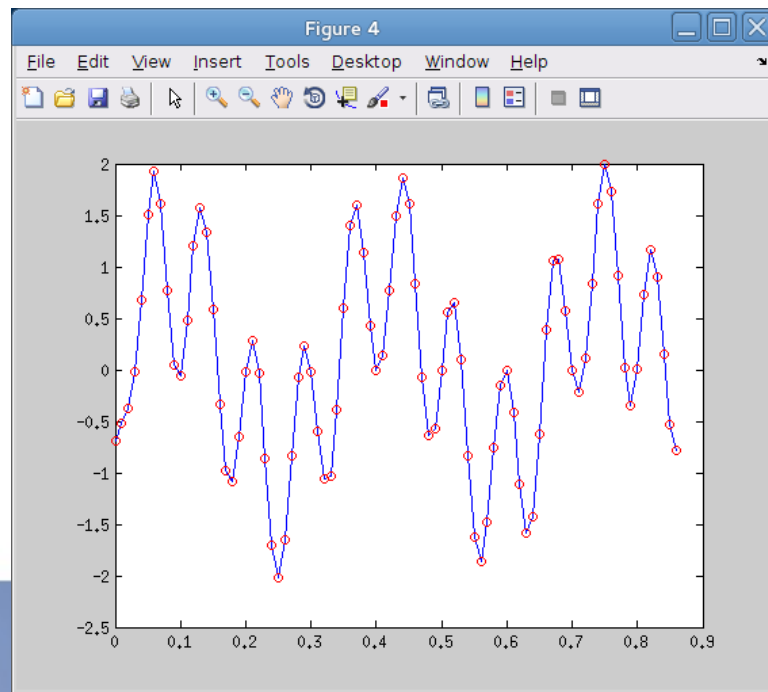
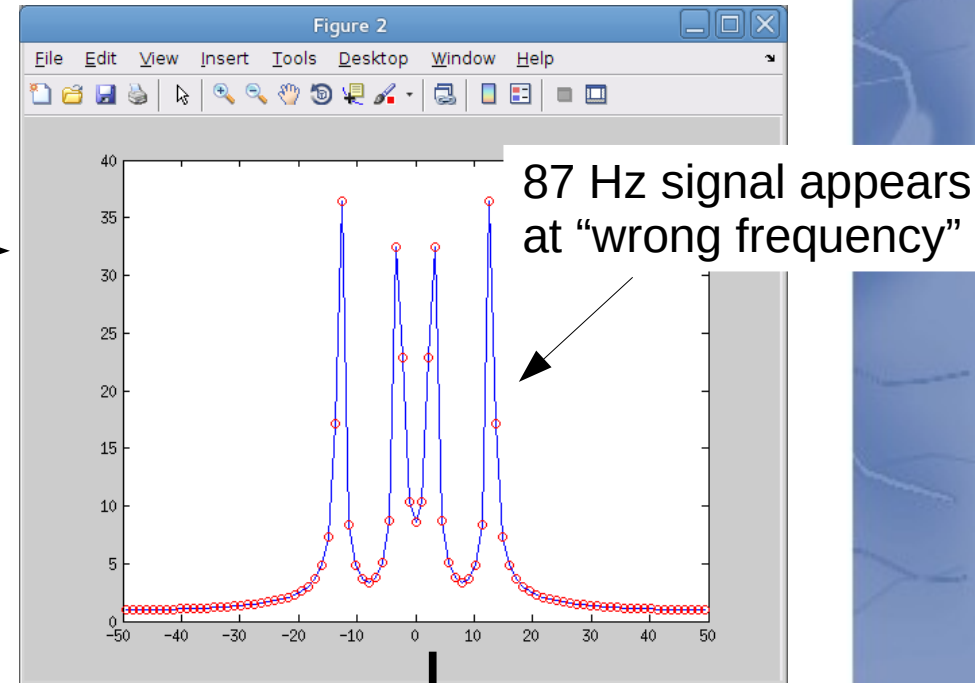
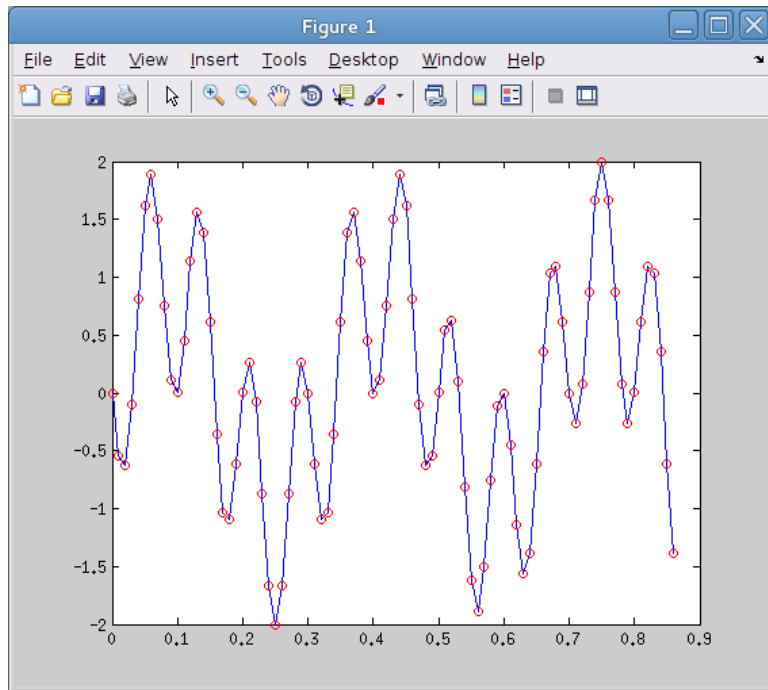


# Effect of undersampling

- Two sines: 3Hz and 87Hz
- Sample frequency: 100Hz (0.01sec sample period)
- This input signal has frequency components larger than Nyquist frequency.



# Effect of undersampling



## Takeaway points (so far)

- Algorithm to compute FFT is fast ( $O(N \log N)$ ).
- Must pay close attention to time and frequency axis.
- Must sample signal with  $f_s \geq 2 \times \text{highest frequency component of signal}$ .
- If you don't, aliasing will introduce spurious components into your signal.

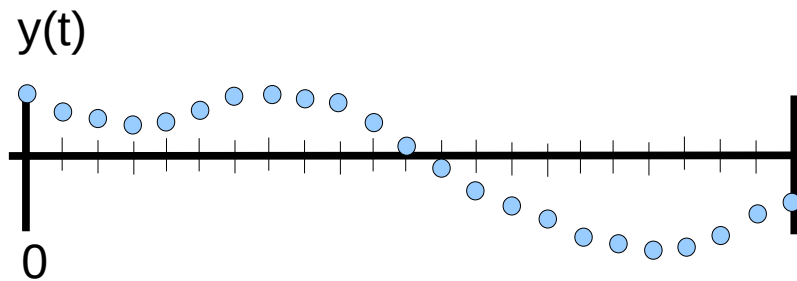
# Next topic: Fourier transform pairs

- Fourier transform pair: Every time-domain function has a frequency-domain dual.

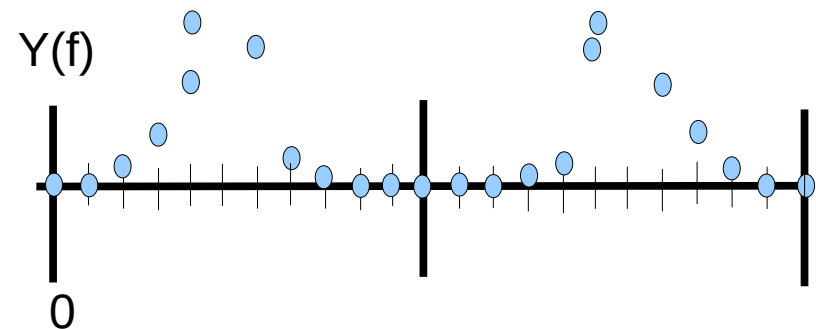
$$y(t) = \frac{1}{2\pi} \int_{-\infty}^{\infty} d\omega Y(\omega) e^{i\omega t}$$



$$Y(\omega) = \int_{-\infty}^{\infty} dt y(t) e^{-i\omega t}$$



Time domain



Frequency domain

# Fourier transform pairs

Time domain

Frequency domain

$$y(t) = \frac{1}{2\pi} \int_{-\infty}^{\infty} d\omega Y(\omega) e^{i\omega t}$$



$$Y(\omega) = \int_{-\infty}^{\infty} dt y(t) e^{-i\omega t}$$

Gaussian

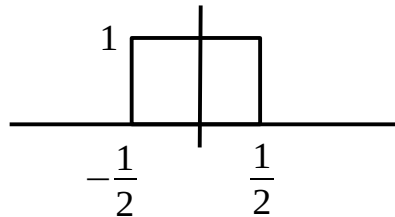
$$e^{-t^2/(2\sigma^2)}$$



$$\sigma \sqrt{2\pi} e^{-\sigma^2 \omega^2/2}$$

Gaussian

Box function



$$\frac{1}{\sqrt{2\pi}} \frac{\sin(\omega/2)}{\omega/2}$$

sinc

Delta function

$$\delta(t)$$



$$1$$

Constant over  
all frequencies

Sine

$$\sin(\omega_0 t)$$



$$-i\pi (\delta(\omega - \omega_0) - \delta(\omega + \omega_0))$$

Two delta functions at +/- input  
frequency

# More Fourier transform pairs...

Shift  $y(t - t_0) \longleftrightarrow Y(\omega) e^{-i\omega t_0}$

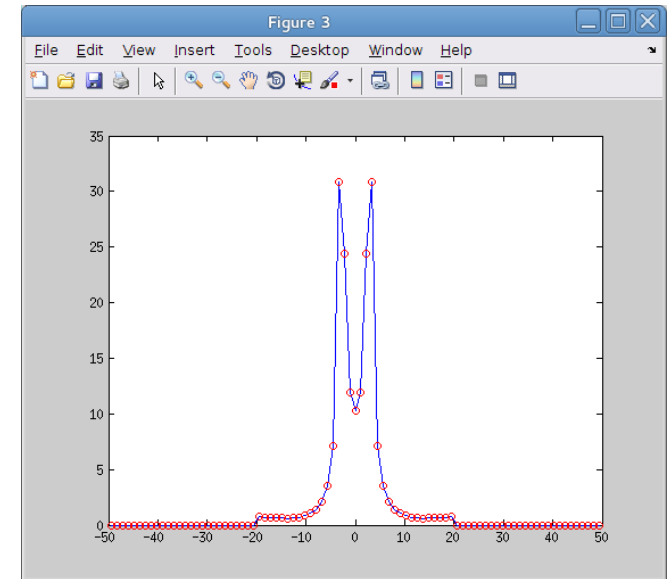
Time derivative  $\left(\frac{d}{dt}\right)y(t) \longleftrightarrow (-i\omega)Y(\omega)$

Convolution  $\int_{-\infty}^{\infty} d\tau y(\tau)z(t - \tau) \longleftrightarrow Y(\omega)Z(\omega)$  Multiplication

- Note that these pairs apply to operations, not just to functions

# Convolution and filtering

- Recall our low-pass filter from earlier this session.
- That filter was equivalent to multiplying by a box function in frequency domain.  $Y(\omega)Z(\omega)$



- This is equivalent to convolution by  $\text{sinc}()$  in time domain.
- Recall this filter kernel from session 1?

$$\frac{\sin(\omega/2)}{\omega/2}$$

# Filtering and convolution

- Recall this weighted-average filter from Session 1?

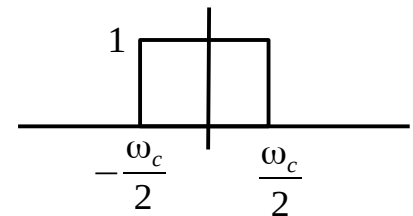
$$y_n = \sum_{i=-m}^m w_{n-i} x_i$$

- That filter was performing a convolution in the time domain.

$$\int_{-\infty}^{\infty} d\tau y(\tau) z(t-\tau)$$

- Hard-wall filter in frequency is same as convolution with sinc() in time domain.

$$\frac{\sin(t/2)}{t/2}$$



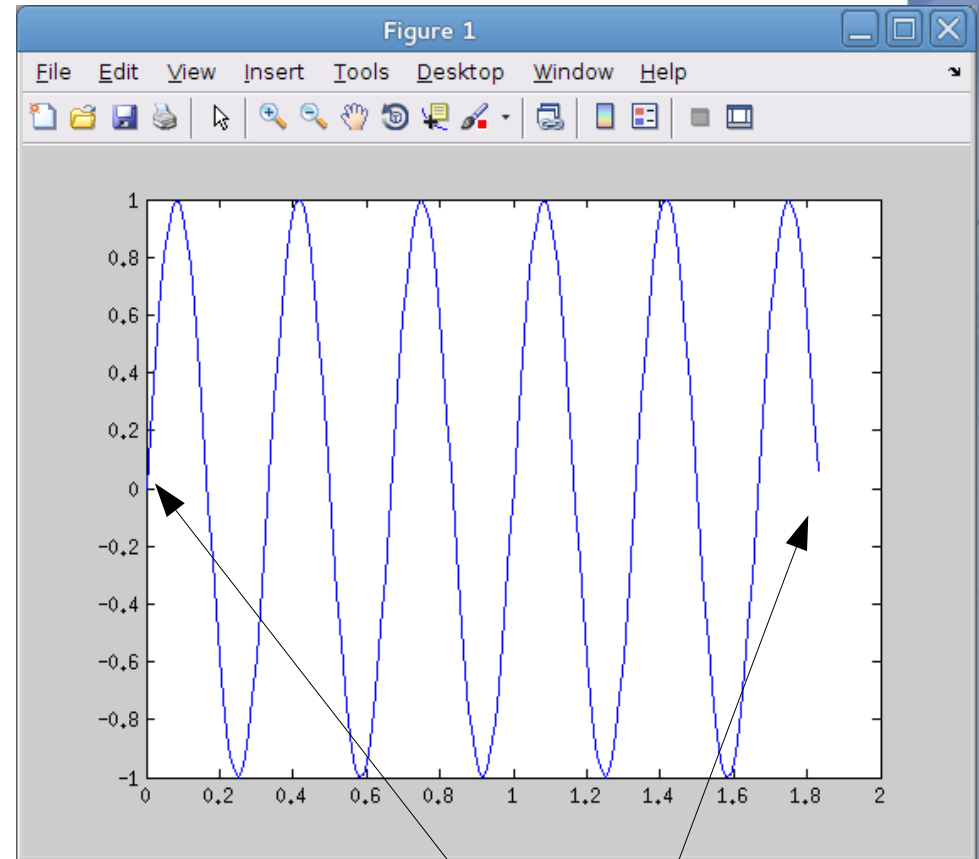


# Remarks

- Fourier transforms allow you to create a desired frequency response for your filter.
- Time-domain filter kernel may be analyzed in frequency domain.
- Two ways to filter your signal:
  - FFT into frequency domain and manipulate its spectrum. Fast
  - Create desired filter kernel in frequency domain, then inverse FFT into time domain and use convolution to filter your signal slower

# Final topic: Windowing

- FFT is sensitive to discontinuities at ends of signal.
- FFT “thinks” signal is periodic.
- Therefore, if the ends don't match, the FFT senses a discontinuity.
  - This causes effects in the frequency spectrum.



Abrupt change of signal at ends because FFT thinks signal is periodic.

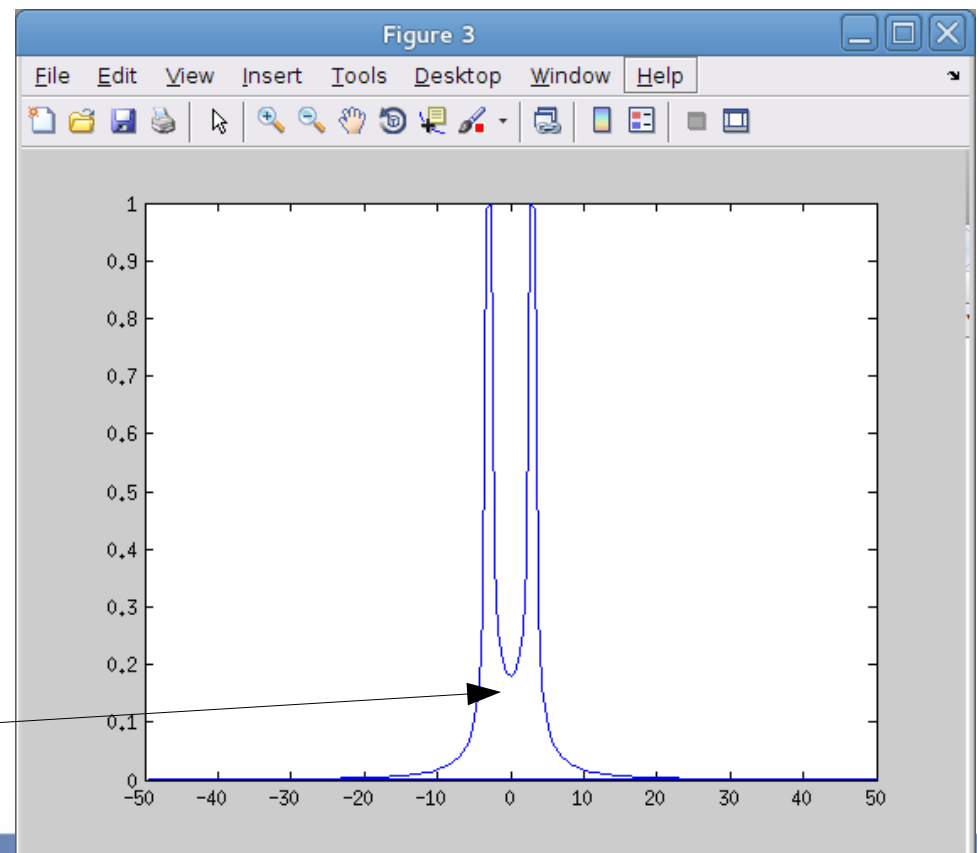
# Effect on computed spectrum

- Recall Fourier transform pair:

$$\sin(\omega_0 t) \longleftrightarrow -i\pi(\delta(\omega - \omega_0) - \delta(\omega + \omega_0))$$

- The FFT should give two delta functions at  $\pm 3\text{Hz}$ .
- But the deltas sit on a pedestal!

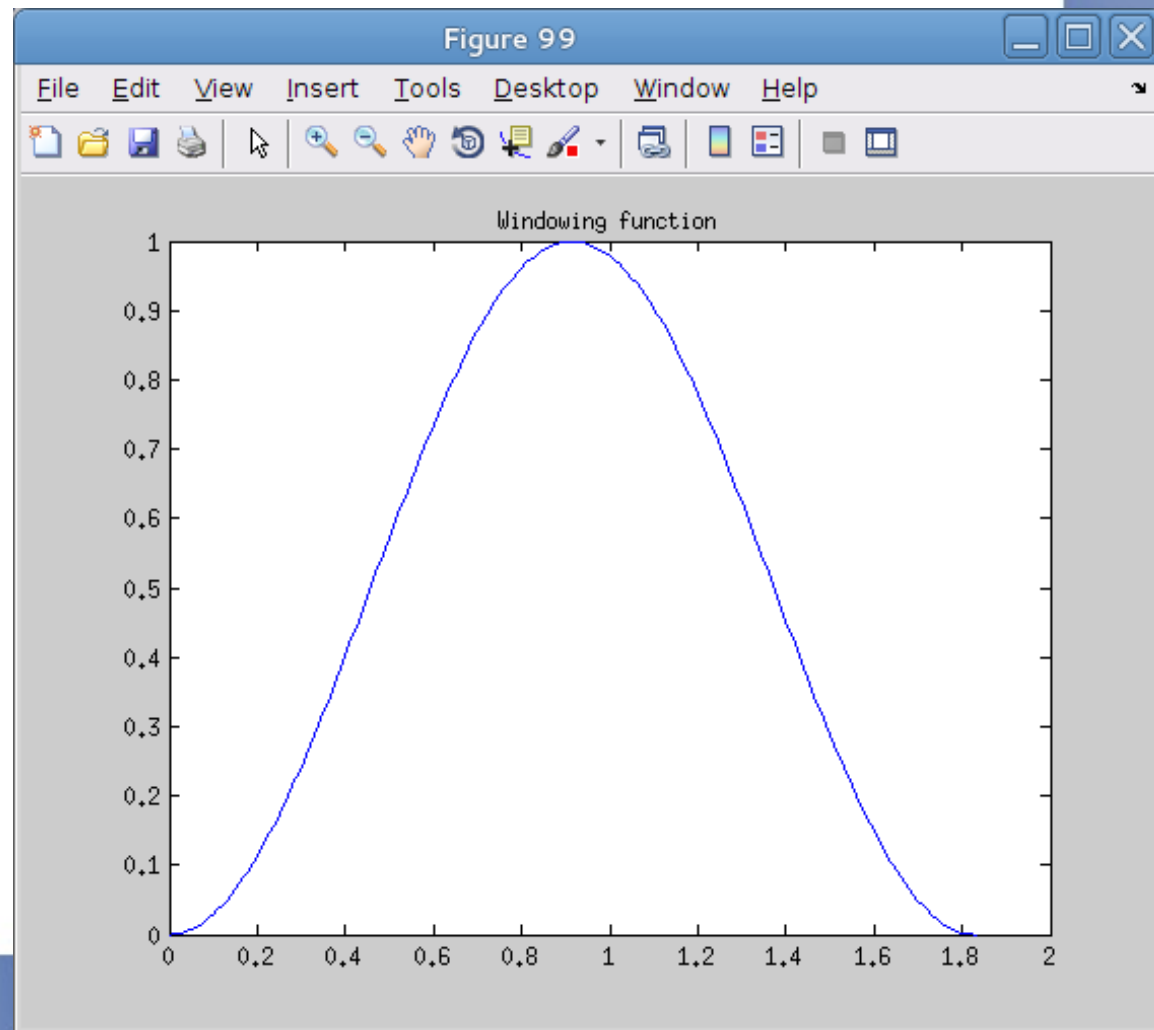
Pedestal



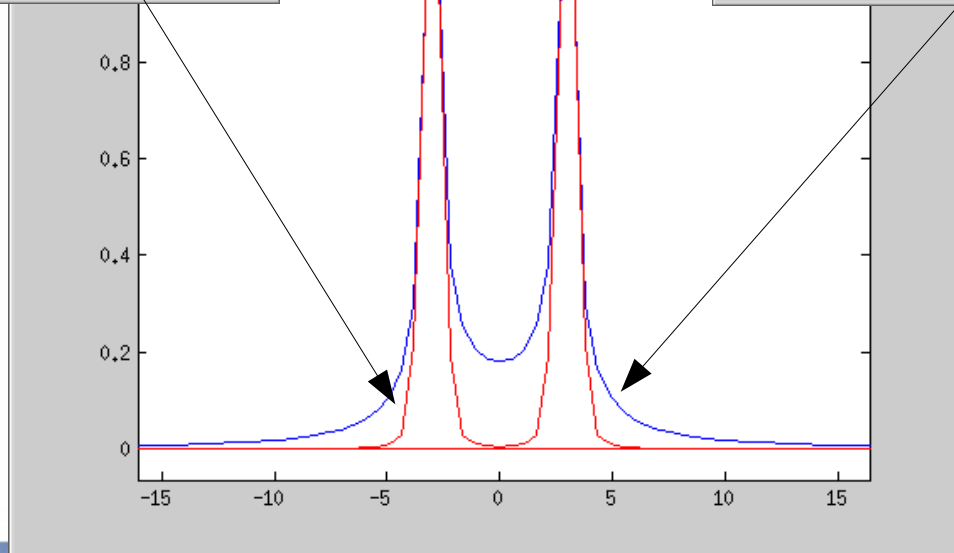
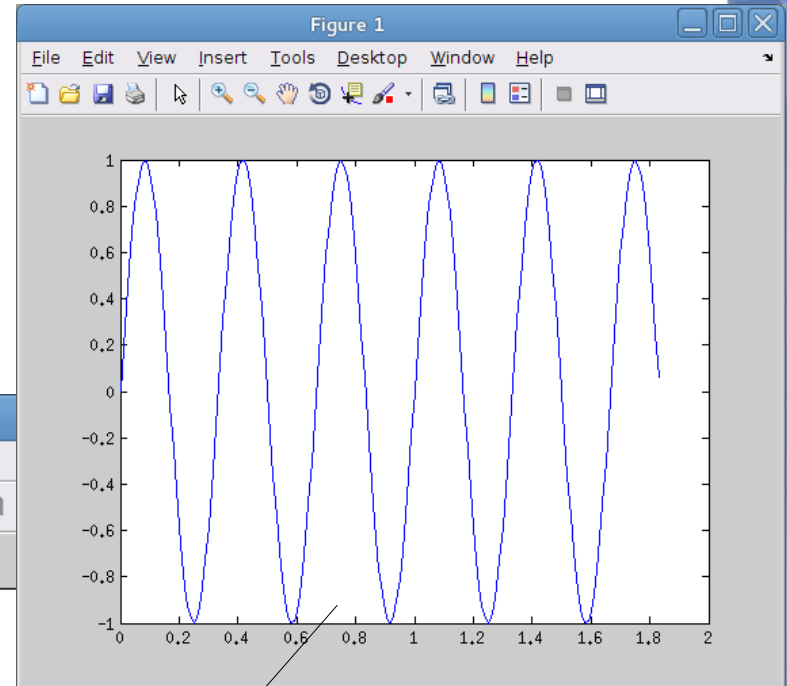
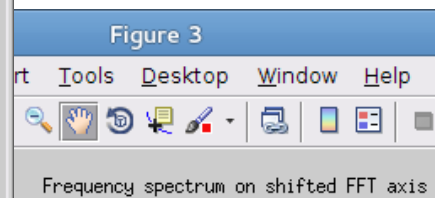
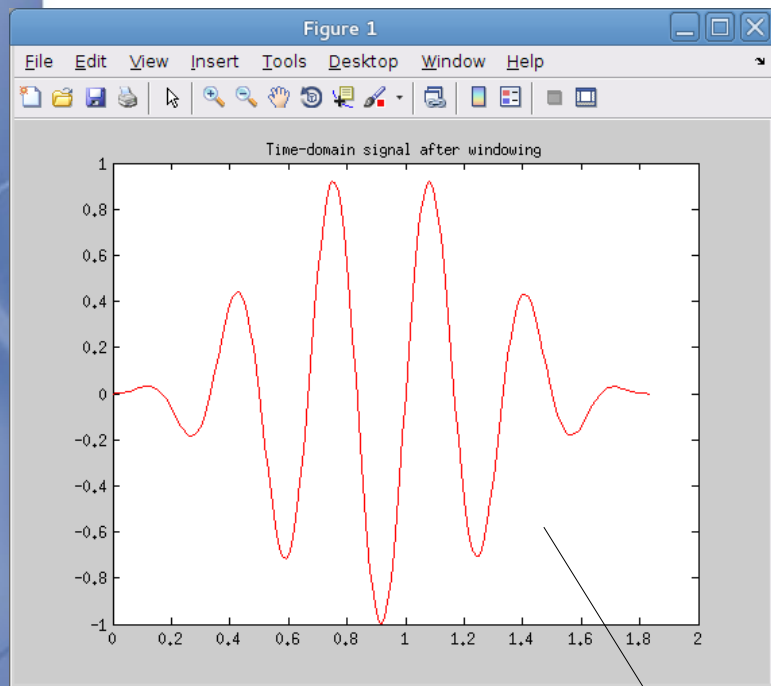
# Solution: Use window

- Multiply input signal by smooth function which goes to zero at ends.
- Example:

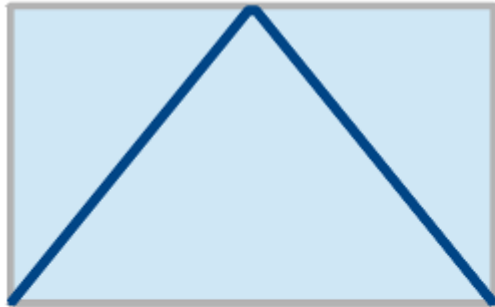
$$u(t) = \frac{1 - \cos(2\pi t/T)}{2}$$



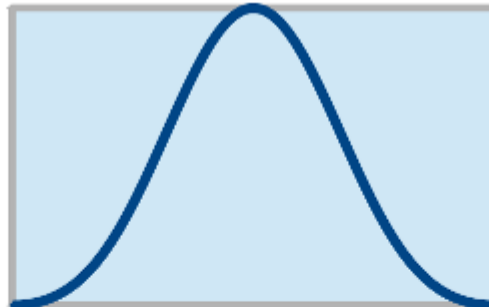
# FFT: Window vs. No window



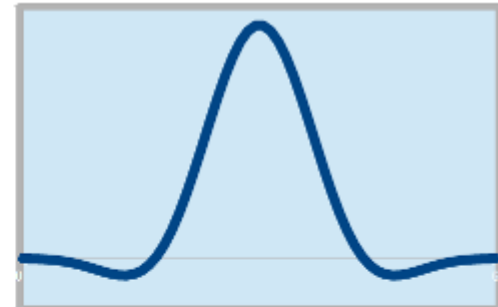
# Common window functions



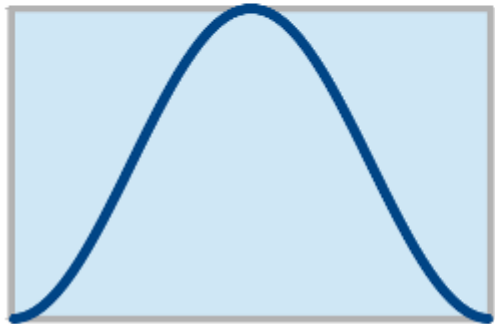
**Bartlett**



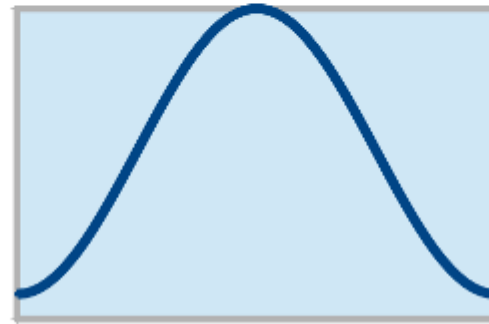
**Blackman**



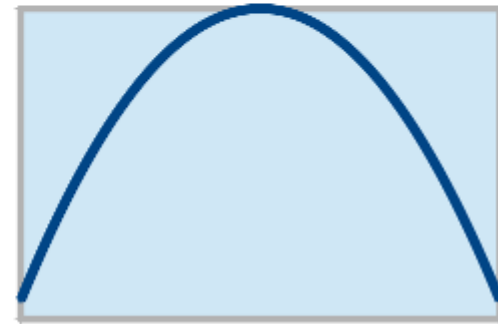
**Flat top**



**Hanning**



**Hamming**



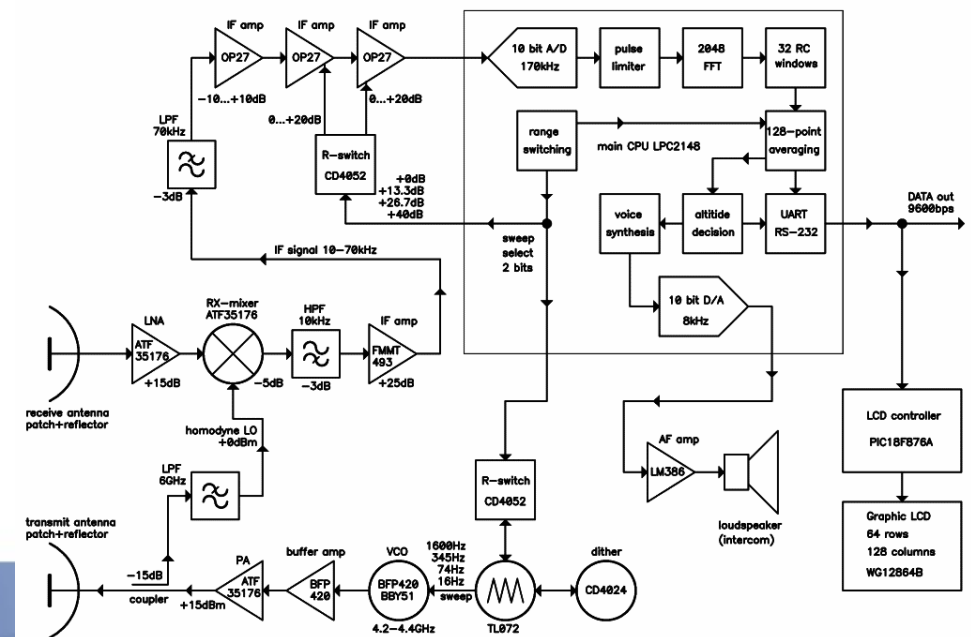
**Welch**

- Window to use depends upon details of your signal, your application, etc.

# Last slide: DSP

- DSP = Digital signal processing
- Very important, very large branch of electrical engineering, very math intensive.

- Radar
- Sonar
- Audio processing
- etc.





# Session summary

- Complexity and big-O.
- The FFT – why it is important and why it is fast.
- Filtering using the FFT.
- Nyquist frequency and aliasing.
- Filtering and convolution.
- Windowing.