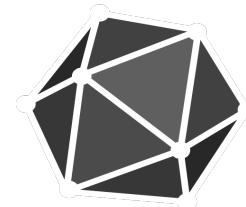
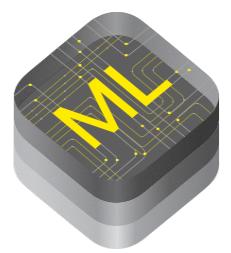


VTA: Open & Flexible DL Acceleration

Thierry Moreau
TVM Conference, Dec 12th 2018



TVM Stack



High-Level Differentiable IR

Tensor Expression IR

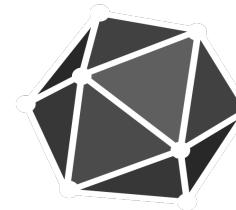
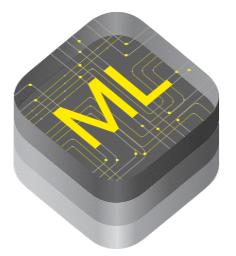
LLVM

CUDA

Metal



TVM Stack



High-Level Differentiable IR

Tensor Expression IR

LLVM

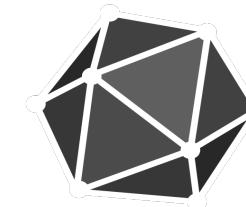
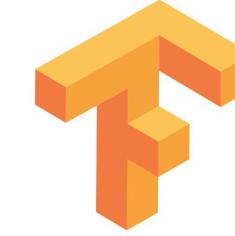
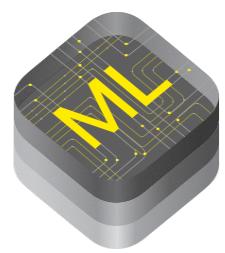
CUDA

Metal

VTA: Open Hardware Accelerator



TVM Stack



High-Level Differentiable IR

Tensor Expression IR

LLVM

CUDA

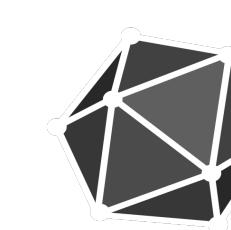
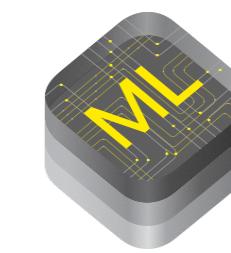
Metal

VTA: Open Hardware Accelerator



Edge FPGA

TVM Stack



High-Level Differentiable IR

Tensor Expression IR

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VTA: Open Hardware Accelerator



Edge FPGA

Cloud FPGA

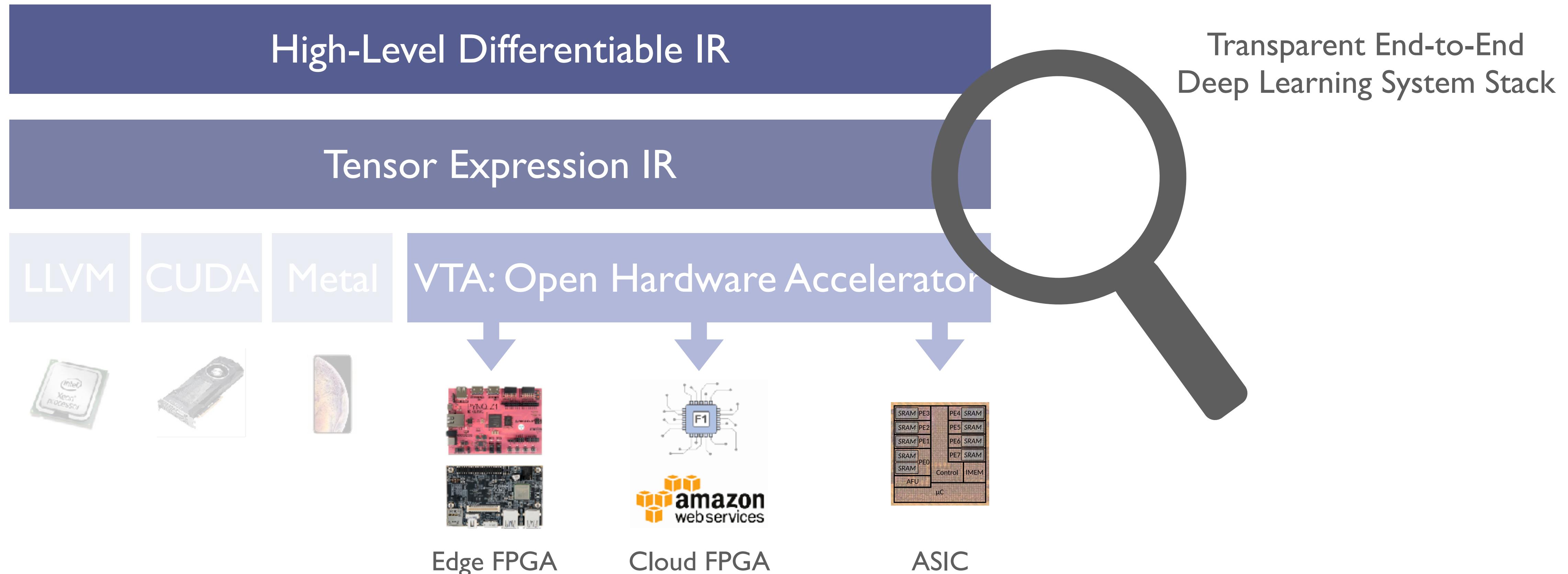
TVM Stack



TVM Stack

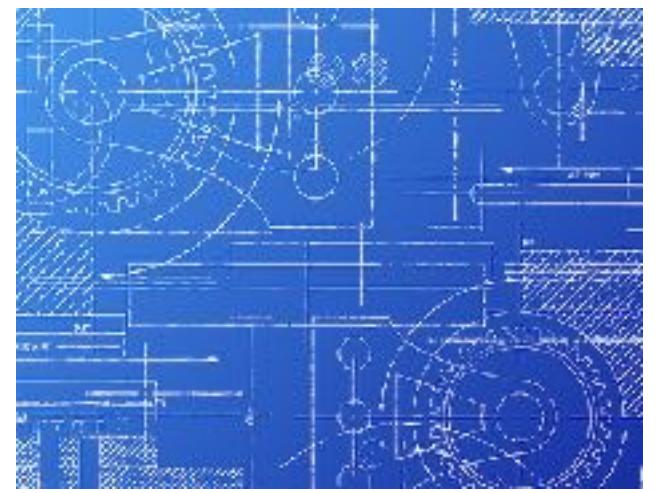


TVM Stack



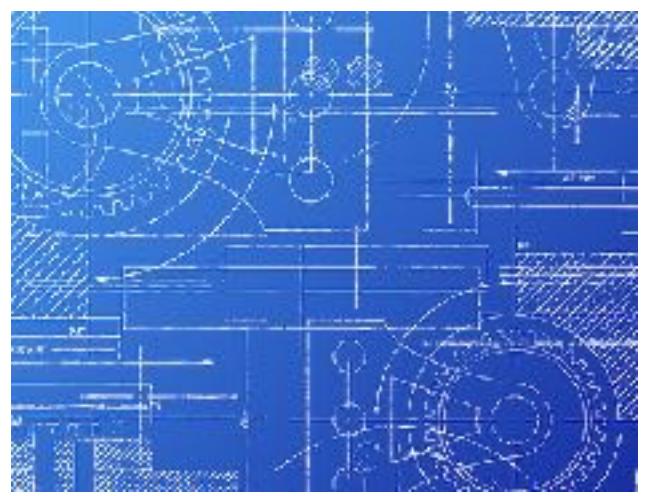
TVM+VTA Stack Goals

TVM+VTA Stack Goals



- Blue-print for a complete deep learning acceleration stack

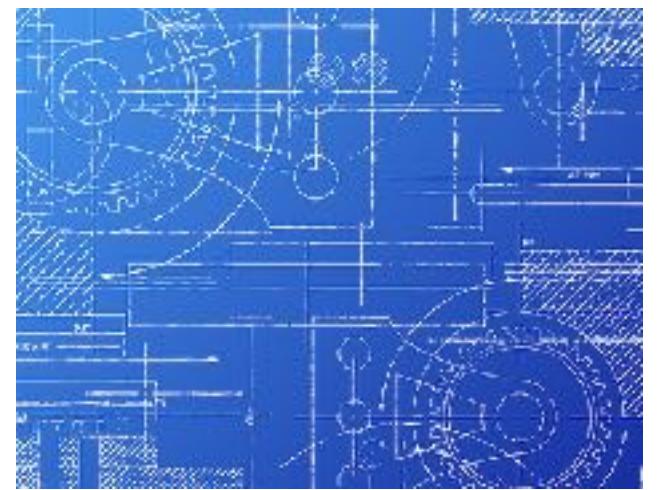
TVM+VTA Stack Goals



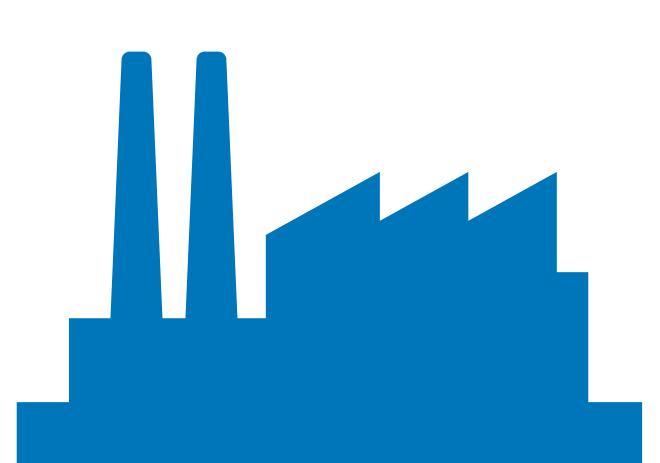
- Blue-print for a complete deep learning acceleration stack
- Experimentation framework for cross-stack deep learning optimizations



TVM+VTA Stack Goals



- Blue-print for a complete deep learning acceleration stack
- Experimentation framework for cross-stack deep learning optimizations
- Open-source community for industrial-strength deep learning acceleration



VTA Overview

Extensible Hardware Architecture

Programmability Across the Stack

Facilitates HW-SW Co-Design

VTA Overview

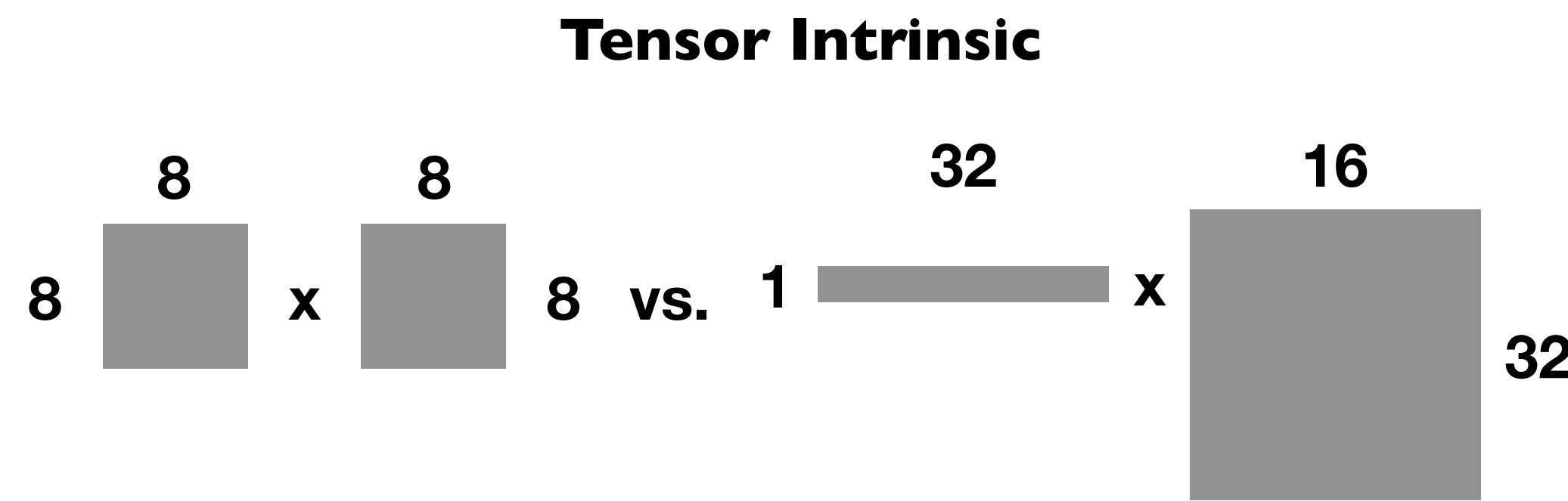
Extensible Hardware Architecture

Programmability Across the Stack

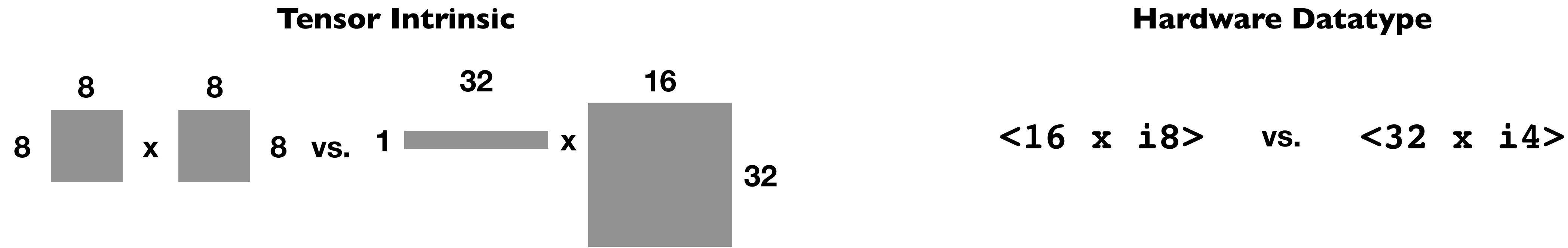
Facilitates HW-SW Co-Design

VTA: General DL Architecture

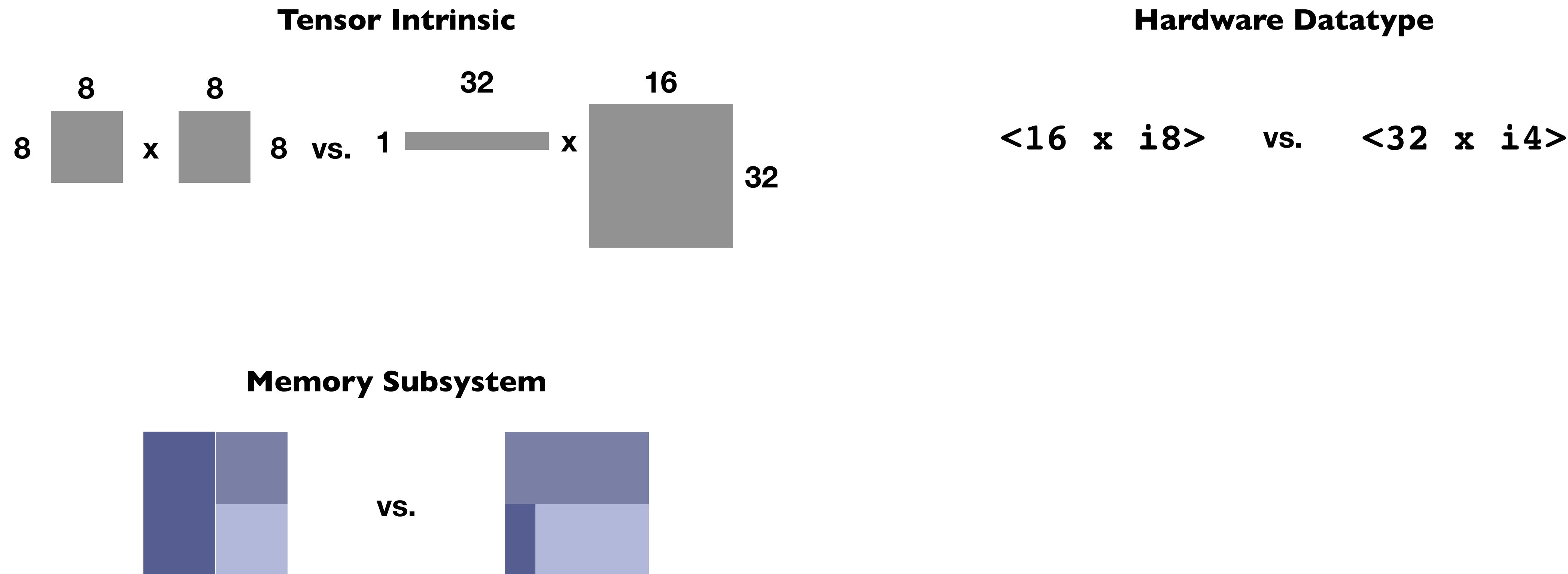
VTA: General DL Architecture



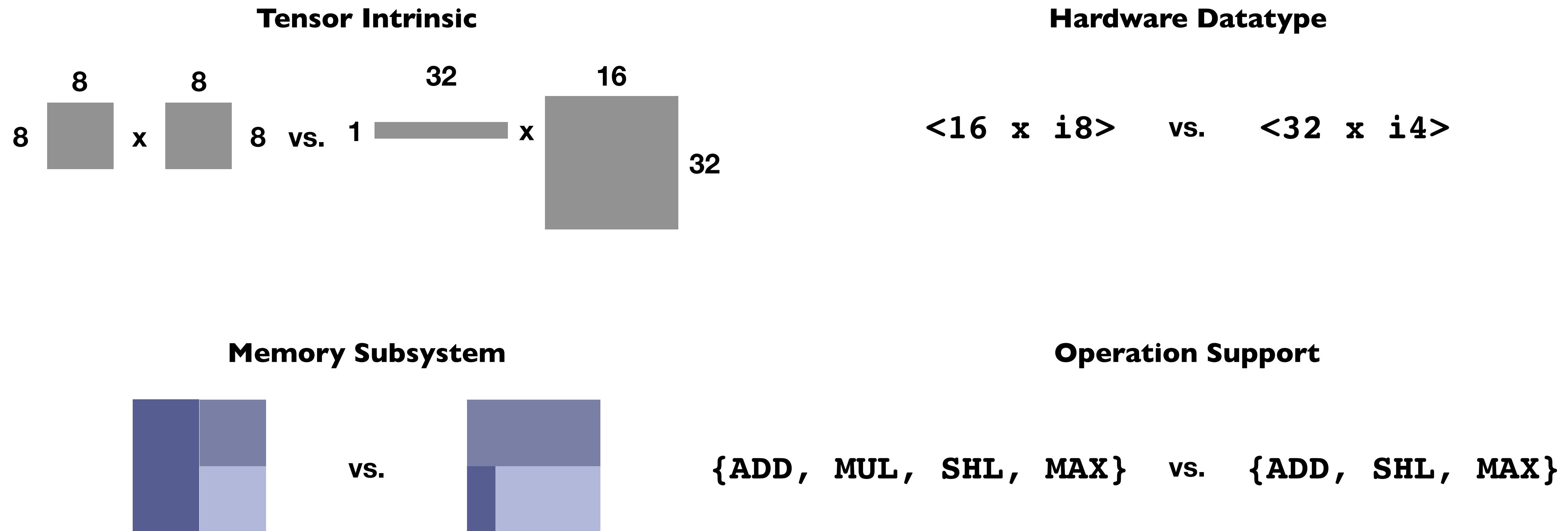
VTA: General DL Architecture



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VTA: General DL Architecture

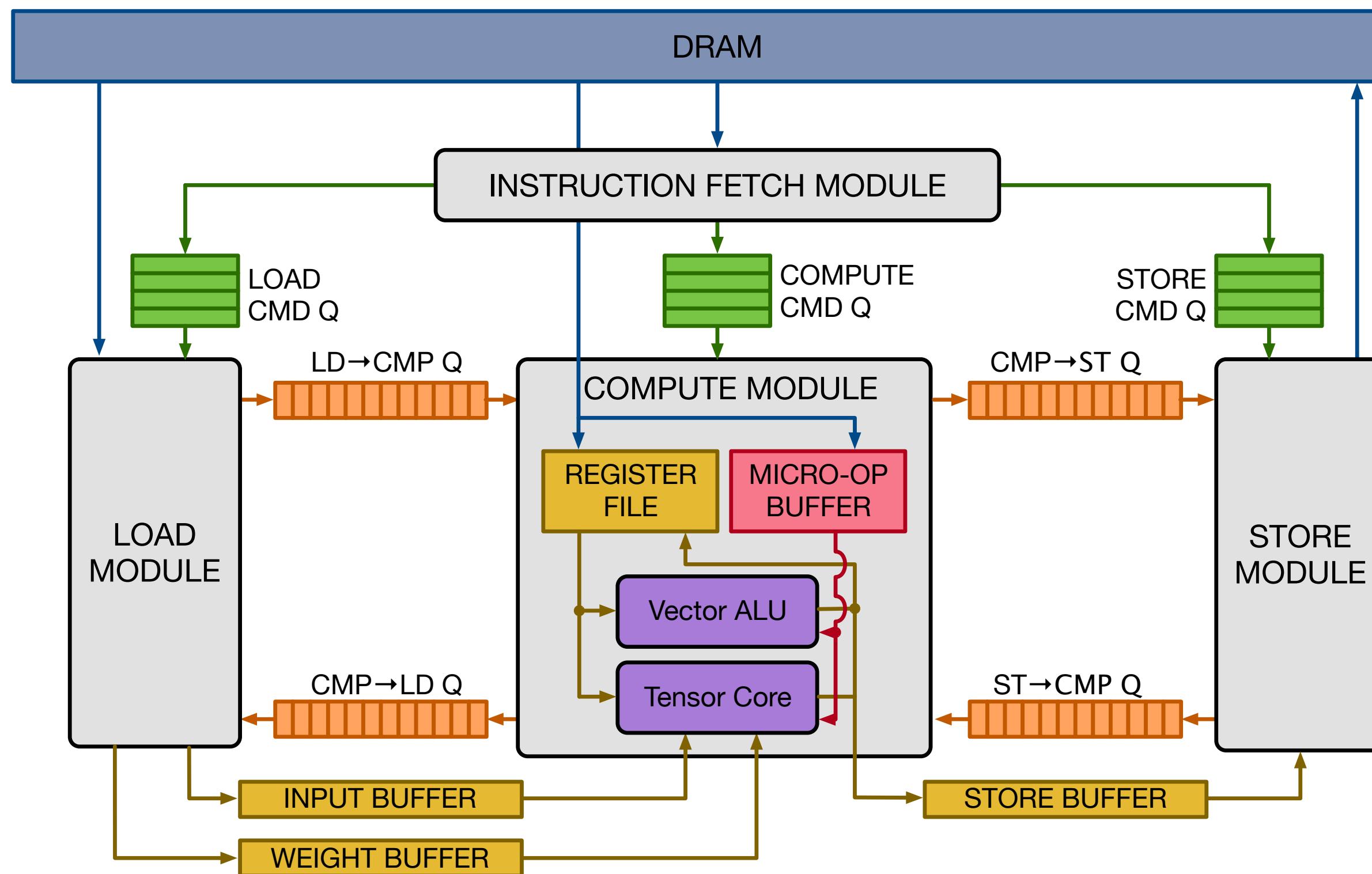


VTA Hardware Architecture

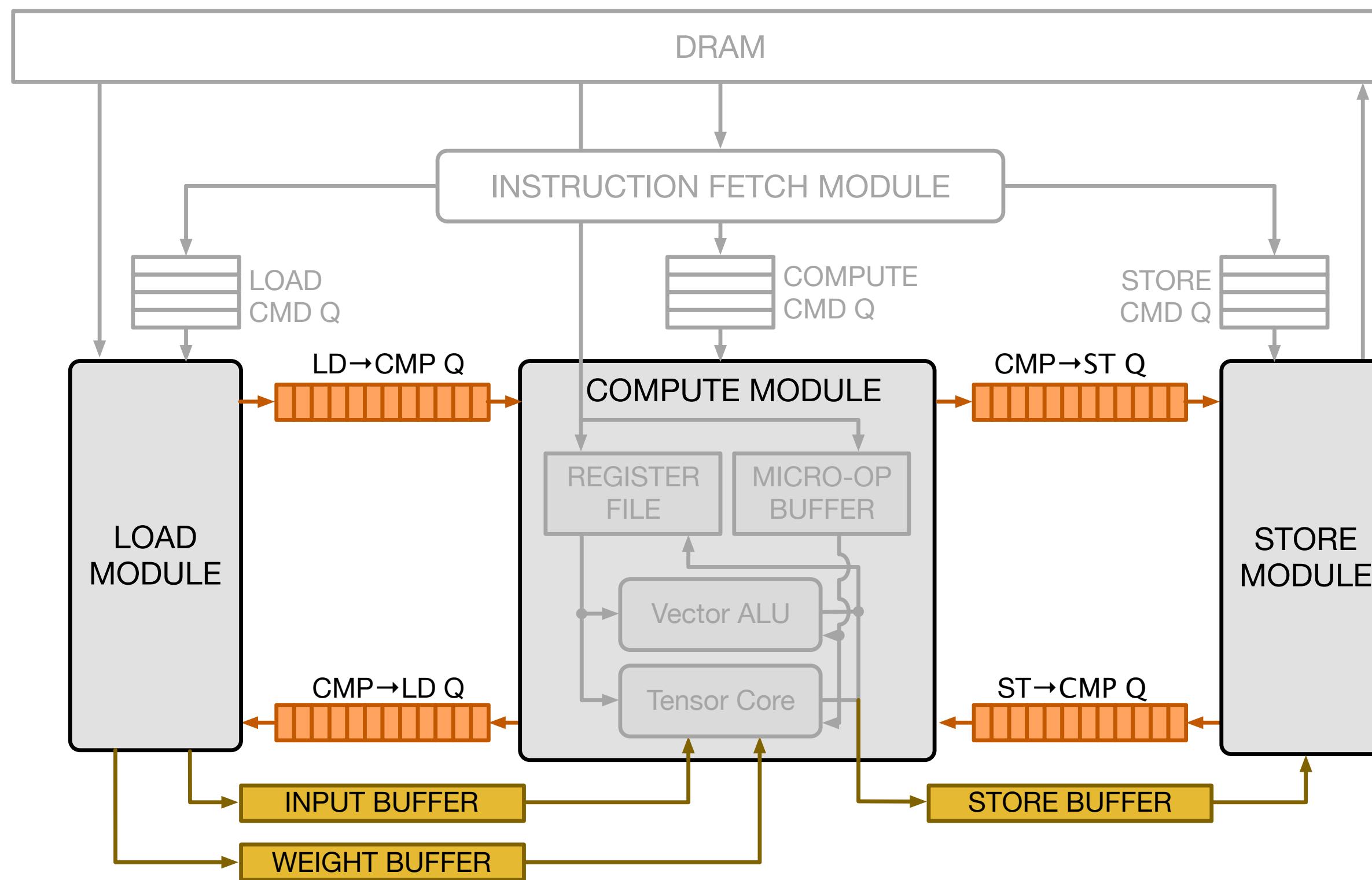
Philosophy: simple hardware, provide software-defined flexibility

VTA Hardware Architecture

Philosophy: simple hardware, provide software-defined flexibility



VTA Hardware Architecture



Pipelining Tasks to Hide Memory Latency

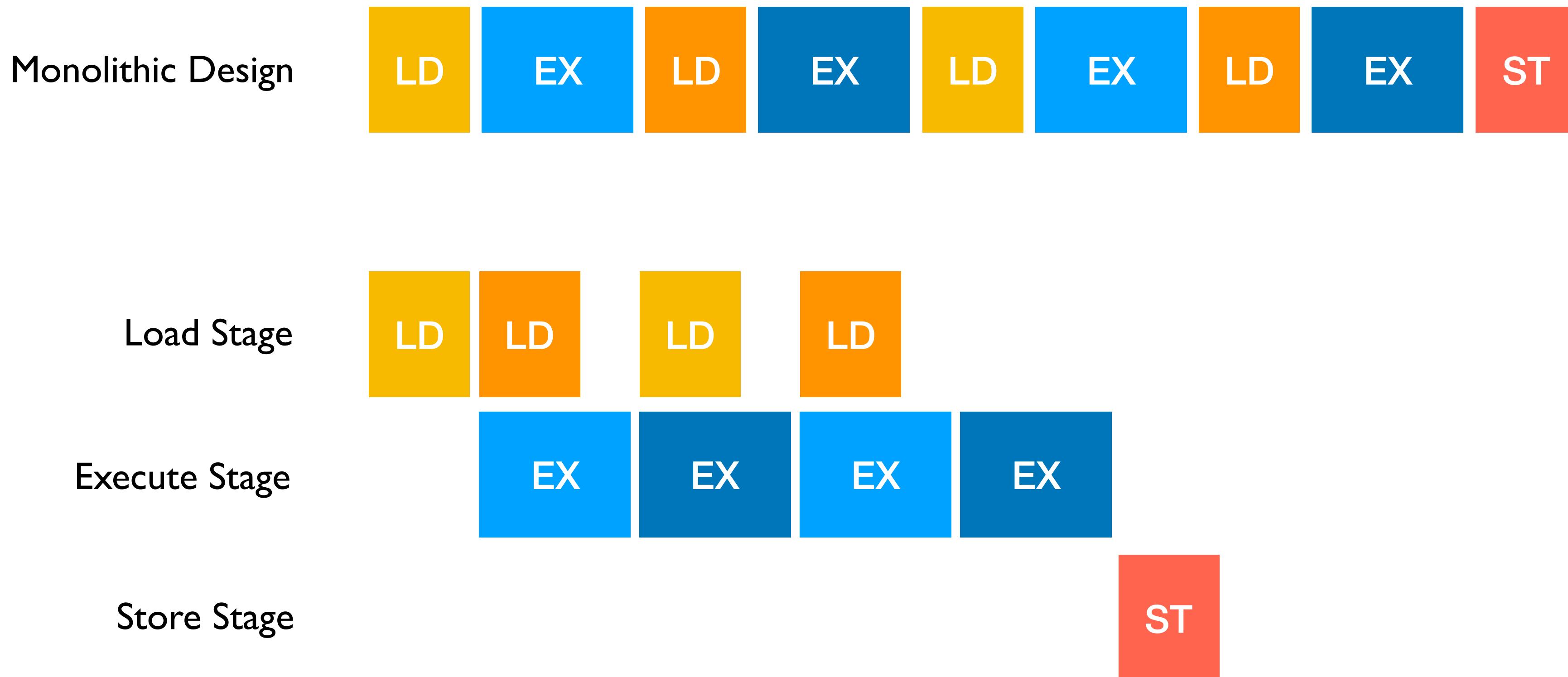


LD: load

EX: compute

ST: store

Pipelining Tasks to Hide Memory Latency

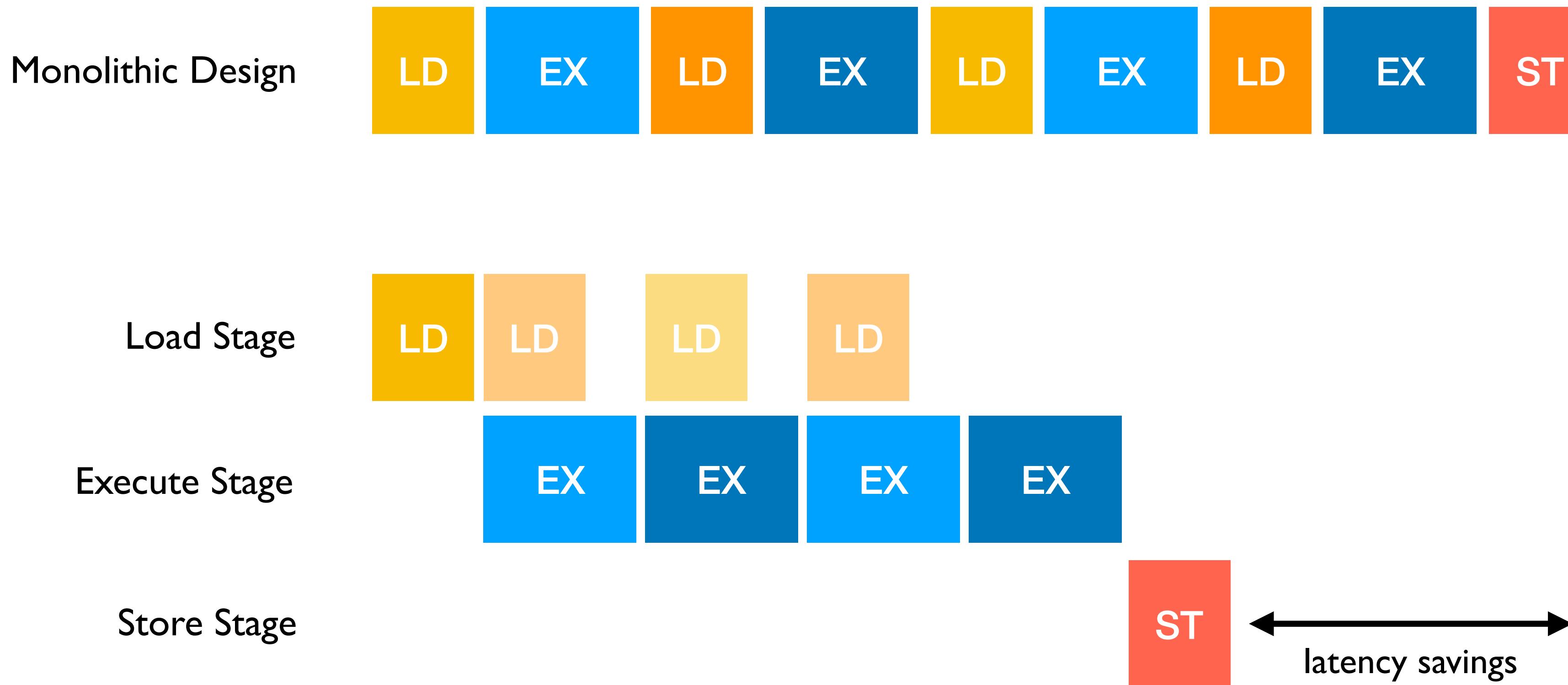


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Pipelining Tasks to Hide Memory Latency

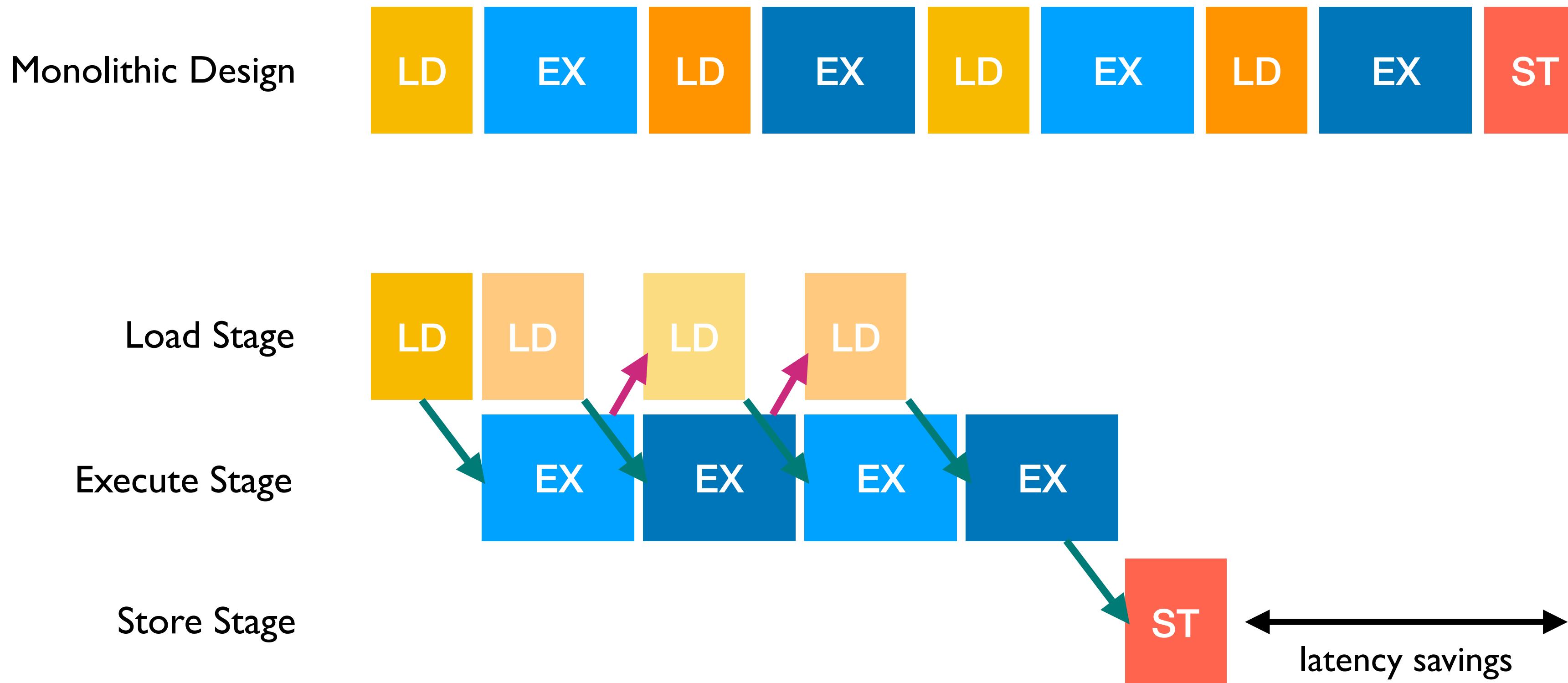


LD: load

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ST: store

Pipelining Tasks to Hide Memory Latency



low-level synchronization between tasks is explicitly managed by the software

LD: load

EX: compute

ST: store

Two-Level ISA Overview

Provides the right tradeoff between expressiveness and code compactness

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- Use CISC instructions to perform multi-cycle tasks



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R0 : R0 + GEMM (A8 , W3)

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- Use CISC instructions to perform multi-cycle tasks



- Use RISC micro-ops to perform single-cycle tensor operations

R0 : R0 + GEMM (A8 , W3)

R2 : MAX (R0 , ZERO)

VTA RISC Micro-Kernels

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multiple RISC instructions define a **micro-kernel**,
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```

```
CONV2D: layout=NCHW, chan=256, kernel=(1,1), padding=(0,0), strides=(2,2)
```

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```
CONV2D: layout=NCHW, chan=256, kernel=(1,1), padding=(0,0), strides=(2,2)
```

```
CONV2D_TRANSPOSE: ...
```

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multiple RISC instructions define a **micro-kernel**,
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CONV2D: layout=NCHW, chan=128, kernel=(3,3), padding=(1,1), strides=(1,1)
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```

```
CONV2D_TRANSPOSE: ...
```

```
GROUP_CONV2D: ...
```

VTA RISC Micro-Kernels

micro-kernel programming gives us
software-defined flexibility



DCGAN

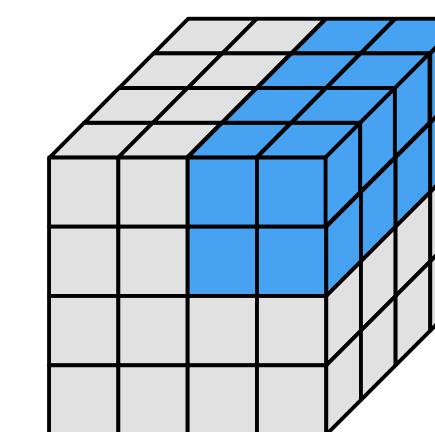
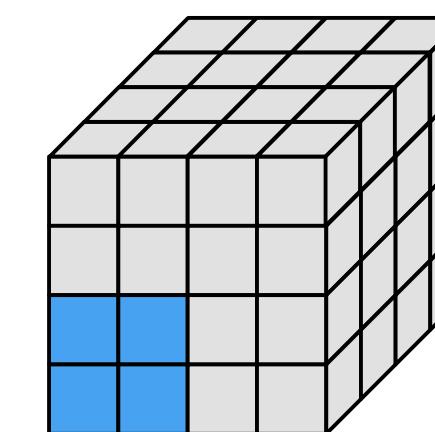
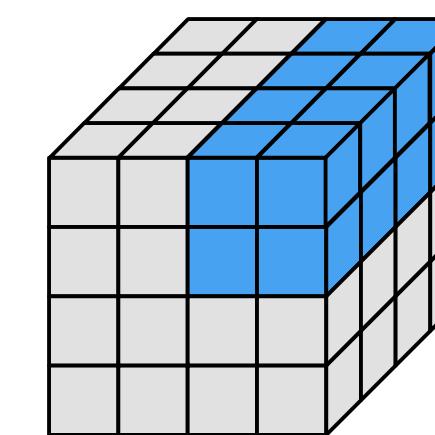
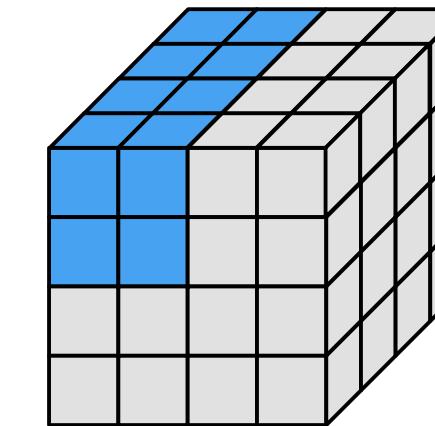


ResNet50

How is VTA Programmed?

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```
// Pseudo-code for convolution program for the VIA accelerator
// Virtual Thread 0
0x00: LOAD(PARAM[ 0-71]) // LD@TID0
0x01: LOAD(ACTIV[ 0-24]) // LD@TID0
0x02: LOAD(LDBUF[ 0-31]) // LD@TID0
0x03: PUSH(LD->EX) // LD@TID0
0x04: POP (LD->EX) // EX@TID0
0x05: EXE (ACTIV[ 0-24],PARAM[ 0-71],LDBUF[ 0-31],STBUF[ 0- 7]) // EX@TID0
0x06: PUSH(EX->LD) // EX@TID0
0x07: PUSH(EX->ST) // EX@TID0
0x08: POP (EX->ST) // ST@TID0
0x09: STOR(STBUF[ 0- 7]) // ST@TID0
0x0A: PUSH(ST->EX) // ST@TID0
// Virtual Thread 1
0x0B: LOAD(ACTIV[25-50]) // LD@TID1
0x0C: LOAD(LDBUF[32-63]) // LD@TID1
0x0D: PUSH(LD->EX) // LD@TID1
0x0E: POP (LD->EX) // EX@TID1
0x0F: EXE (ACTIV[25-50],PARAM[ 0-71],LDBUF[32-63],STBUF[32-39]) // EX@TID1
0x10: PUSH(EX->LD) // EX@TID1
0x11: PUSH(EX->ST) // EX@TID1
0x12: POP (EX->ST) // ST@TID1
0x13: STOR(STBUF[32-39]) // ST@TID1
0x14: PUSH(ST->EX) // ST@TID1
// Virtual Thread 2
0x15: POP (EX->LD) // LD@TID2
0x16: LOAD(PARAM[ 0-71]) // LD@TID2
0x17: LOAD(ACTIV[ 0-24]) // LD@TID2
0x18: LOAD(LDBUF[ 0-31]) // LD@TID2
0x19: PUSH(LD->EX) // LD@TID2
0x1A: POP (LD->EX) // EX@TID2
0x1B: POP (ST->EX) // EX@TID2
0x1C: EXE (ACTIV[ 0-24],PARAM[ 0-71],LDBUF[ 0-31],STBUF[ 0- 7]) // EX@TID2
0x1D: PUSH(EX->ST) // EX@TID2
0x1E: POP (EX->ST) // ST@TID2
0x1F: STOR(STBUF[ 0- 7]) // ST@TID2
// Virtual Thread 3
0x20: POP (EX->LD) // LD@TID3
0x21: LOAD(ACTIV[25-50]) // LD@TID3
0x22: LOAD(LDBUF[32-63]) // LD@TID3
0x23: PUSH(LD->EX) // LD@TID3
0x24: POP (LD->EX) // EX@TID3
0x25: POP (ST->EX) // EX@TID3
0x26: EXE (ACTIV[25-50],PARAM[ 0-71],LDBUF[32-63],STBUF[32-39]) // EX@TID3
0x27: PUSH(EX->ST) // EX@TID3
0x28: POP (EX->ST) // ST@TID3
0x29: STOR(STBUF[32-39]) // ST@TID3
```



(a) Blocked convolution program with multiple thread contexts

```
// Convolution access pattern dictated by micro-coded program.
// Each register index is derived as a 2-D affine function.
// e.g.  $idx_{rf} = a_{rf}y + b_{rf}x + c_{rf}^0$ , where  $c_{rf}^0$  is specified by
// micro op 0 fields.
for y in [0...i)
    for x in [0...j)
        rf[idx_rf^0] += GEVM(act[idx_act^0], par[idx_par^0])
        rf[idx_rf^1] += GEVM(act[idx_act^1], par[idx_par^1])
        ...
        rf[idx_rf^n] += GEVM(act[idx_act^n], par[idx_par^n])
```

(b) Convolution micro-coded program

```
// Max-pool, batch normalization and activation function
// access pattern dictated by micro-coded program.
// Each register index is derived as a 2D affine function.
// e.g.  $idx_{dst} = a_{dst}y + b_{dst}x + c_{dst}^0$ , where  $c_{dst}^0$  is specified by
// micro op 0 fields.
for y in [0...i)
    for x in [0...j)
        // max pooling
        rf[idx_dst^0] = MAX(rf[idx_dst^0], rf[idx_src^0])
        rf[idx_dst^1] = MAX(rf[idx_dst^1], rf[idx_src^1])
        ...
        // batch norm
        rf[idx_dst^m] = MUL(rf[idx_dst^m], rf[idx_src^m])
        rf[idx_dst^{m+1}] = ADD(rf[idx_dst^{m+1}], rf[idx_src^{m+1}])
        rf[idx_dst^{m+2}] = MUL(rf[idx_dst^{m+2}], rf[idx_src^{m+2}])
        rf[idx_dst^{m+3}] = ADD(rf[idx_dst^{m+3}], rf[idx_src^{m+3}])
        ...
        // activation
        rf[idx_dst^{n-1}] = RELU(rf[idx_dst^{n-1}], rf[idx_src^{n-1}])
        rf[idx_dst^n] = RELU(rf[idx_dst^n], rf[idx_src^n])
```

(c) Max pool, batch norm and activation micro-coded program

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0x10: PUSH(EX->LD)
0x11: PUSH(EX->ST)
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// Virtual Thread 3
0x20: POP (EX->LD)
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0x23: PUSH(LD->EX)
0x24: POP (LD->EX)
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0x27: PUSH(EX->ST)
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```



Programming
accelerators is
hard!!!

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```

```
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        ...
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```

(b) Convolution micro-coded program

```
old batch normalization and activation function
pattern dictated by micro-coded program.
register index is derived as a 2D affine function.
 $idx_{dst} = a_{dst}y + b_{dst}x + c_{dst}^0$ , where  $c_{dst}^0$  is specified by
micro op 0 fields.
```

```
[0...n-1)
    for j in [0...n-1)
        x_d = MAX(rf[idx_dst^0], rf[idx_src^0])
        x_d = MAX(rf[idx_dst^1], rf[idx_src^1])
```

max pooling

```
rf[idx_dst^m] = MUL(rf[idx_dst^m], rf[idx_src^m])
rf[idx_dst^{m+1}] = ADD(rf[idx_dst^{m+1}], rf[idx_src^{m+1}])
rf[idx_dst^{m+2}] = MUL(rf[idx_dst^{m+2}], rf[idx_src^{m+2}])
rf[idx_dst^{m+3}] = ADD(rf[idx_dst^{m+3}], rf[idx_src^{m+3}])
```

batch norm

```
rf[idx_dst^{n-1}] = RELU(rf[idx_dst^{n-1}], rf[idx_src^{n-1}])
rf[idx_dst^n] = RELU(rf[idx_dst^n], rf[idx_src^n])
```

(c) Max pool, batch norm and activation
micro-coded program

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Facilitates HW-SW Co-Design

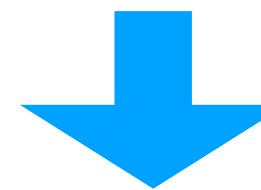
Latency Hiding: An Example of Cross-Stack Design

programmer friendly construct

```
// Virtual Threading
tx, co = s[OUT_L].split(co, factor=2)
s[OUT_L].bind(tx, thread_axis("cthread"))
```

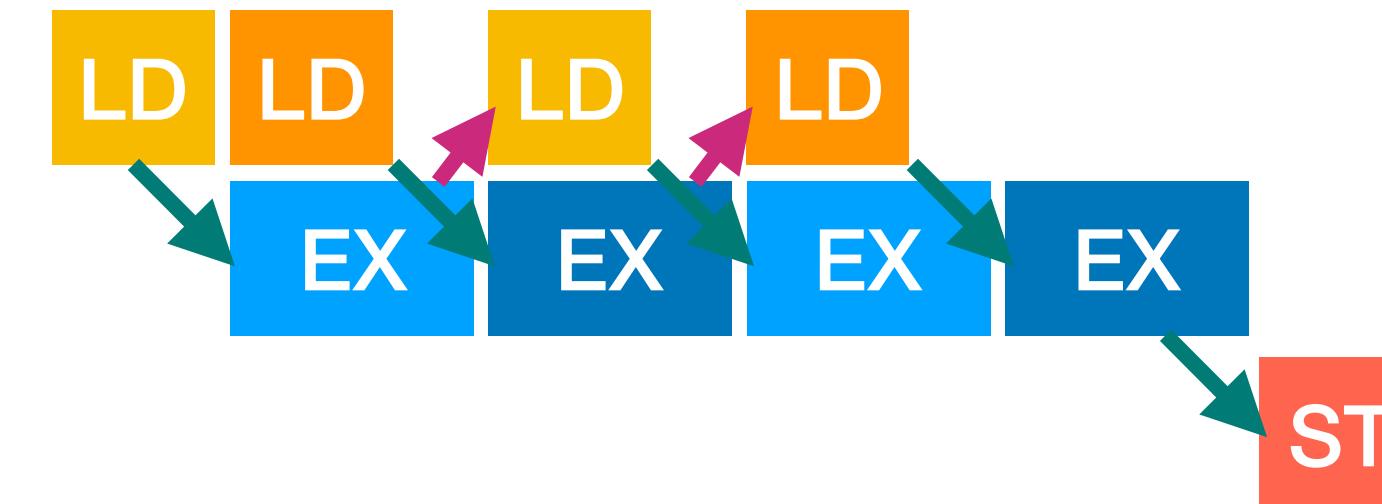
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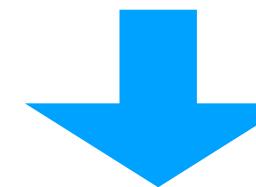
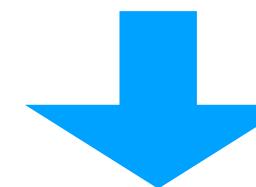
```
// Virtual Threading  
tx, co = s[OUT_L].split(co, factor=2)  
s[OUT_L].bind(tx, thread_axis("cthread"))
```

low-level pipelined execution



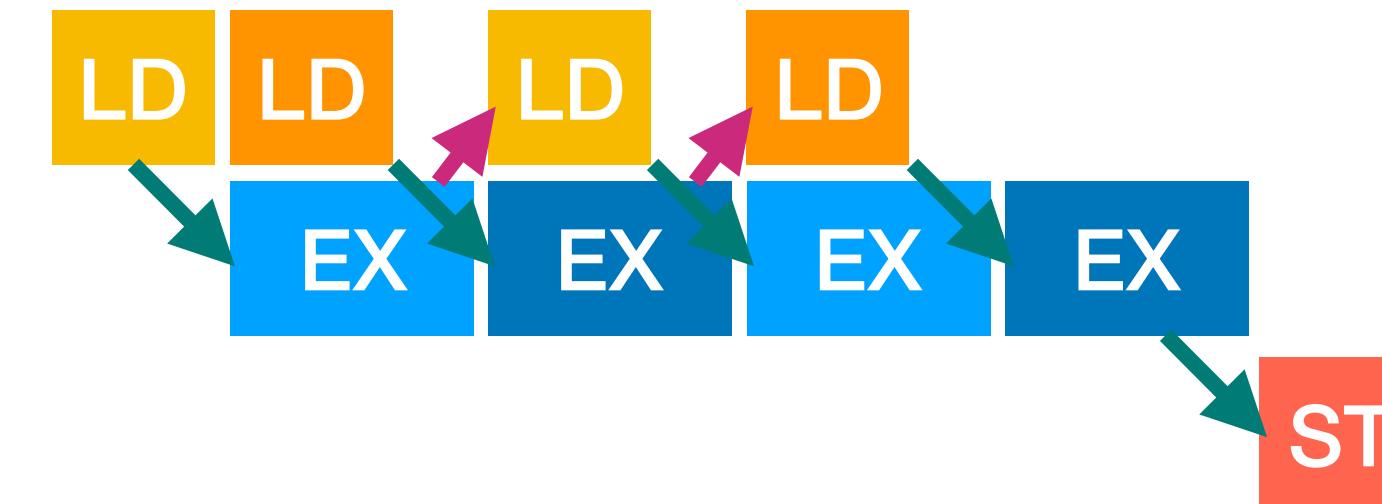
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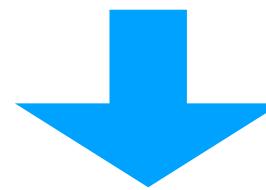
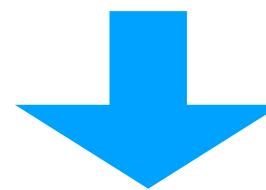
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```



Latency Hiding: An Example of Cross-Stack Design

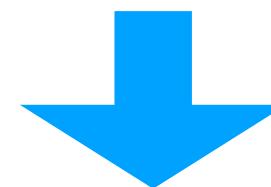
programmer friendly construct



low-level pipelined execution

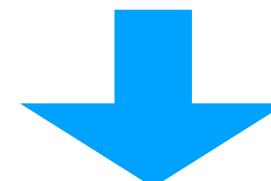
Latency Hiding: An Example of Cross-Stack Design

programmer friendly construct



Tensor Expression Optimizer (TVM)

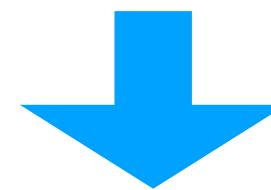
inserts dependence ops based on thread scope



low-level pipelined execution

Latency Hiding: An Example of Cross-Stack Design

programmer friendly construct

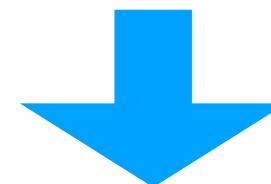


Tensor Expression Optimizer (TVM)

inserts dependence ops based on thread scope

VTA Runtime & JIT Compiler

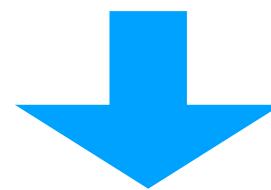
generates instruction stream



low-level pipelined execution

Latency Hiding: An Example of Cross-Stack Design

programmer friendly construct



Tensor Expression Optimizer (TVM)

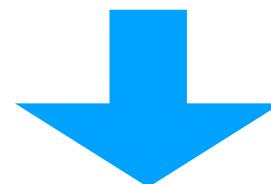
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VTA Runtime & JIT Compiler

generates instruction stream

VTA Hardware/Software Interface (ISA)

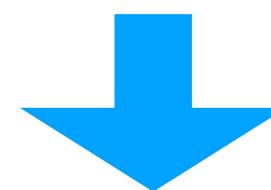
exposes explicit dependences



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VTA Runtime & JIT Compiler

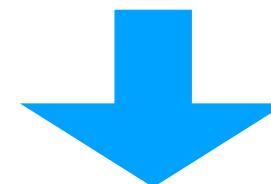
generates instruction stream

VTA Hardware/Software Interface (ISA)

exposes explicit dependences

VTA MicroArchitecture

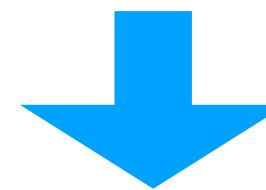
execution predicated on dependences



low-level pipelined execution

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VTA Runtime & JIT Compiler

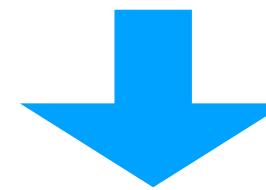
generates instruction stream

VTA Hardware/Software Interface (ISA)

exposes explicit dependences

VTA MicroArchitecture

execution predicated on dependences



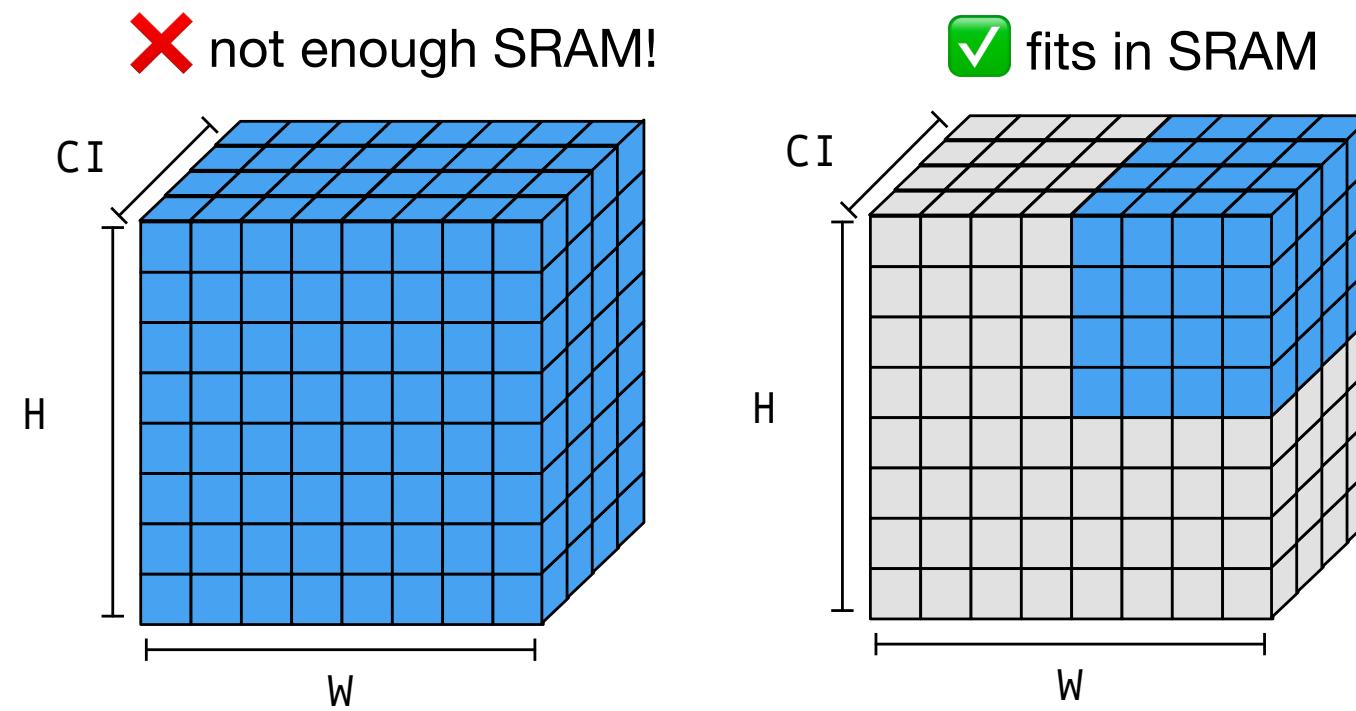
low-level pipelined execution

9-60% better compute utilization

VTA Helped inform ASIC Support in TVM

VTA Helped inform ASIC Support in TVM

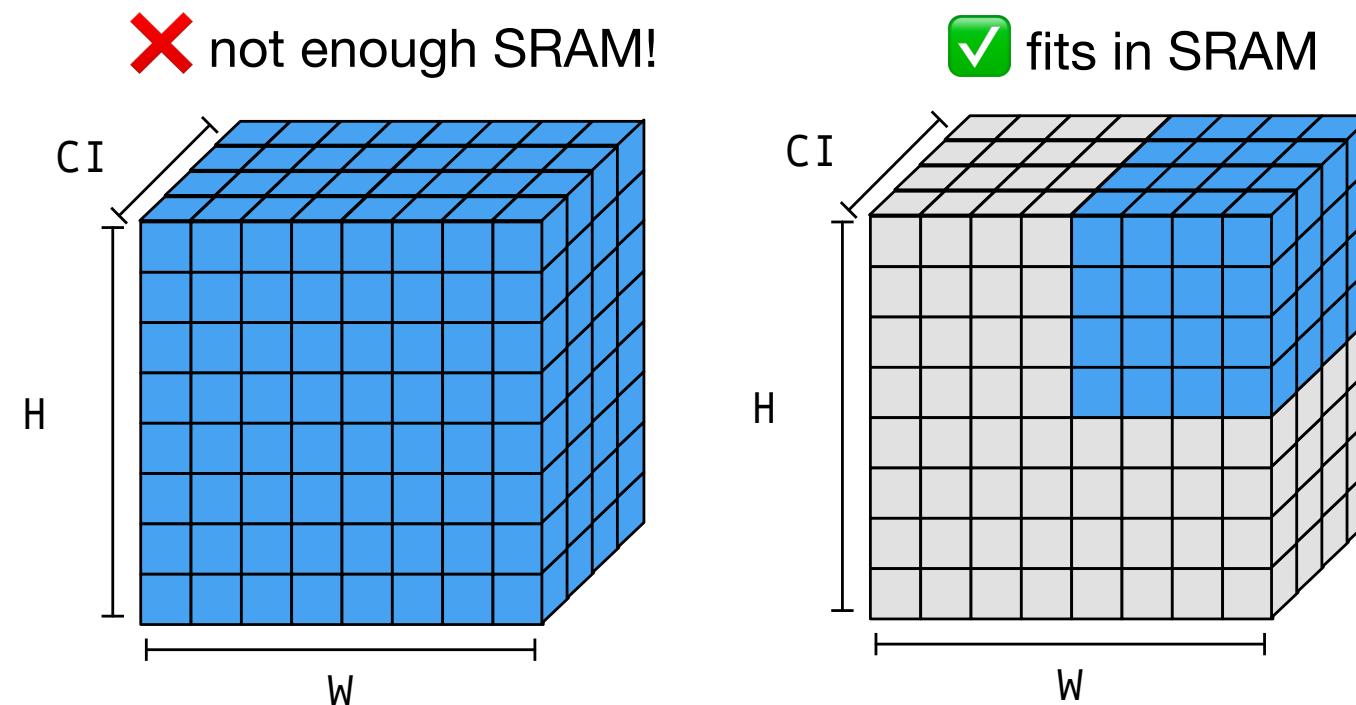
I. How do we partition work and explicitly manage on-chip memories?



```
// Tile  
yo, xo, yi, xi = s[OUT].tile(y, x, 4, 4)  
// Scoped cache read  
INP_L = s.cache_read(INP, vta.inp, [OUT])  
s[INP_L].compute_at(s[OUT], xo)
```

VTA Helped inform ASIC Support in TVM

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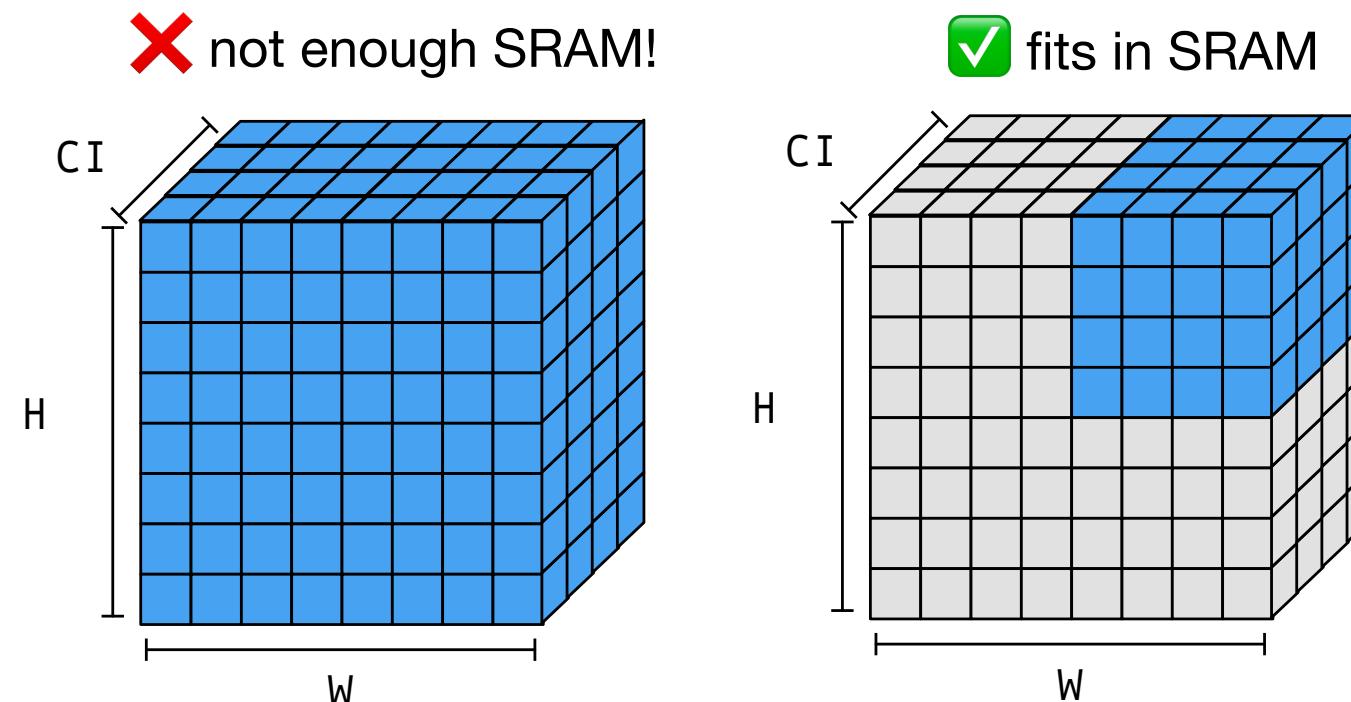
2. How do we take advantage of tensor computation intrinsics?

The diagram shows a large tensor being decomposed into two smaller tensors. On the left is a large tensor with four columns and four rows, colored in a repeating pattern of yellow, orange, pink, and cyan. An equals sign follows it. To the right are two smaller tensors: one with four columns and two rows, colored yellow and orange, and another with two columns and four rows, colored grey.

```
// Tensorize  
s[OUT_L].tensorize(ni)
```

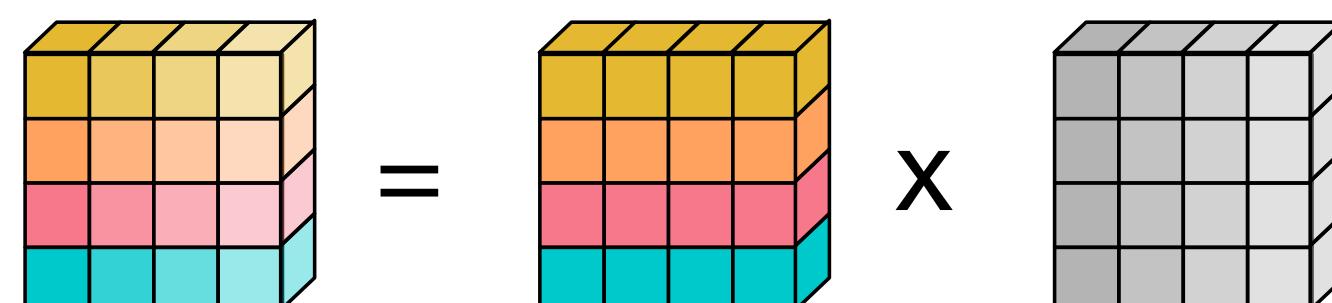
VTA Helped inform ASIC Support in TVM

I. How do we partition work and explicitly manage on-chip memories?



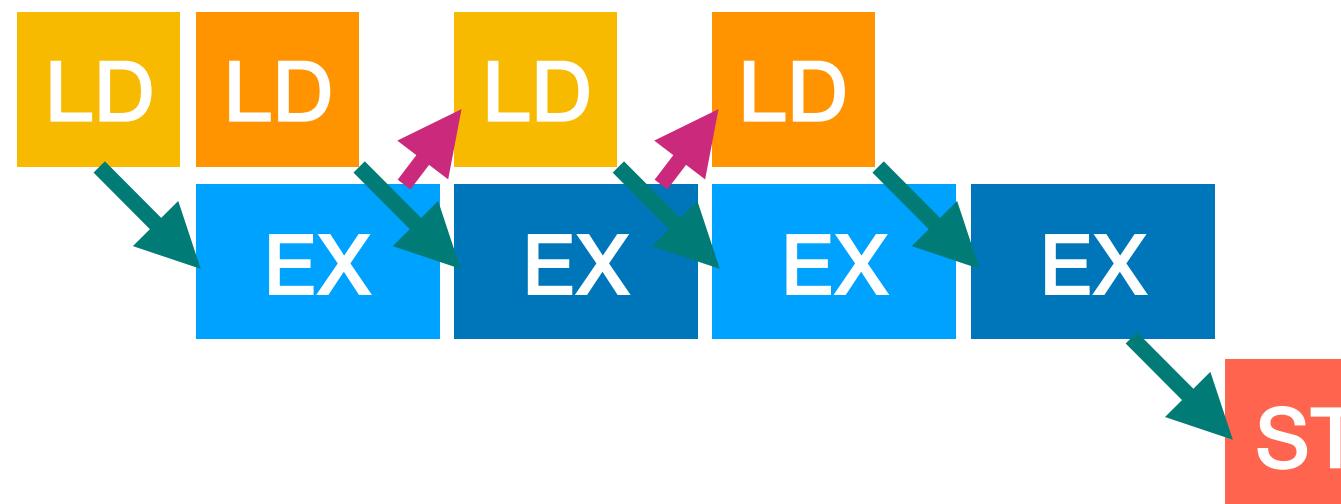
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```

2. How do we take advantage of tensor computation intrinsics?



```
// Tensorize  
s[OUT_L].tensorize(ni)
```

3. How do we hide memory access latency?



```
// Virtual Threading  
tx, co = s[OUT_L].split(co, factor=2)  
s[OUT_L].bind(tx, thread_axis("cthread"))
```

VTA Overview

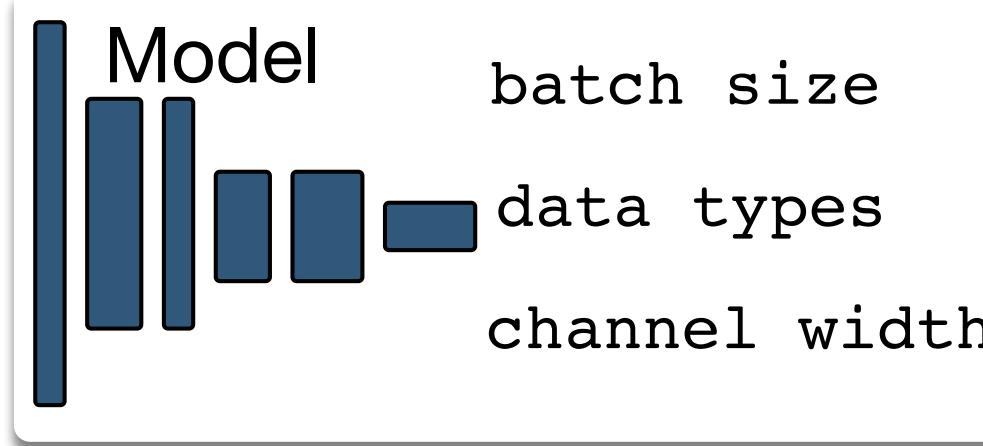
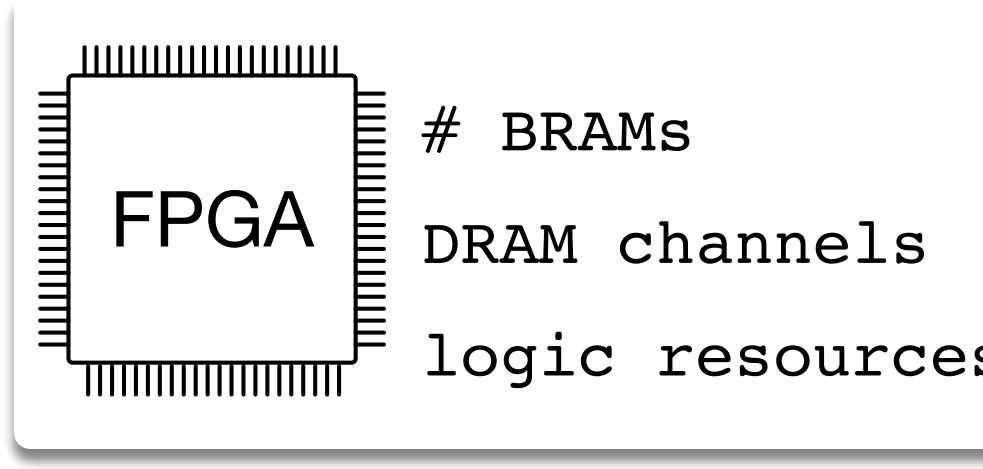
Extensible Hardware Architecture

Programmability Across the Stack

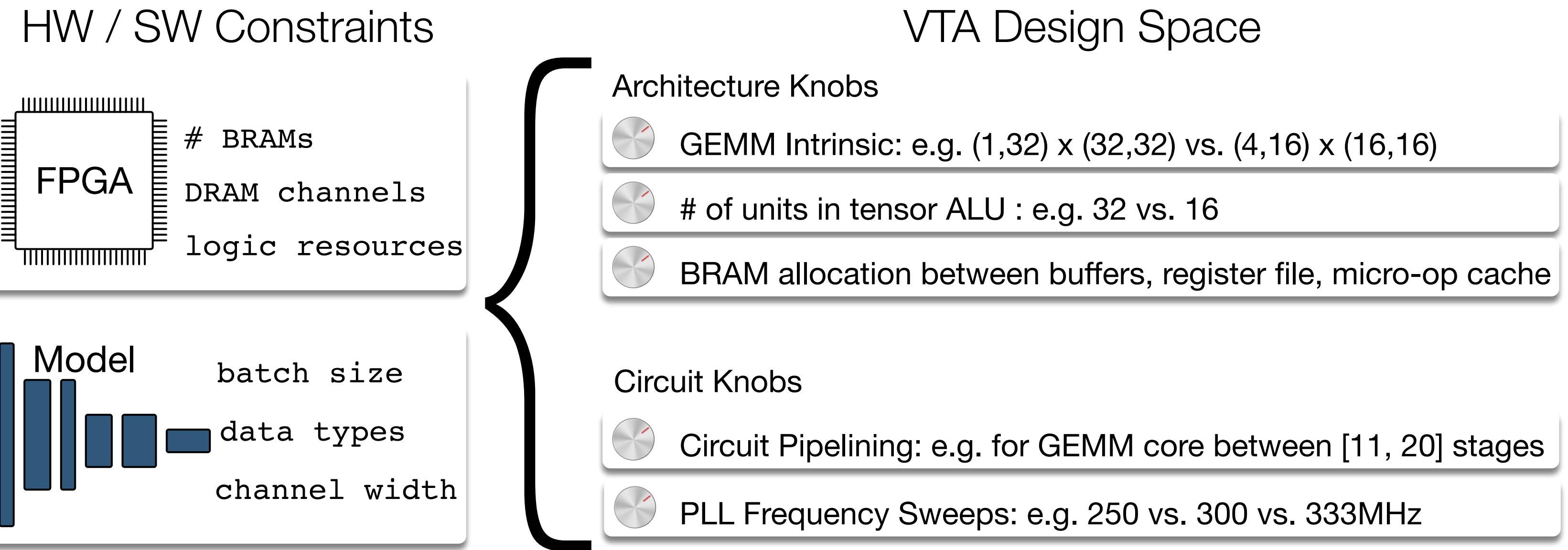
Facilitates HW-SW Co-Design

Hardware Exploration with VTA

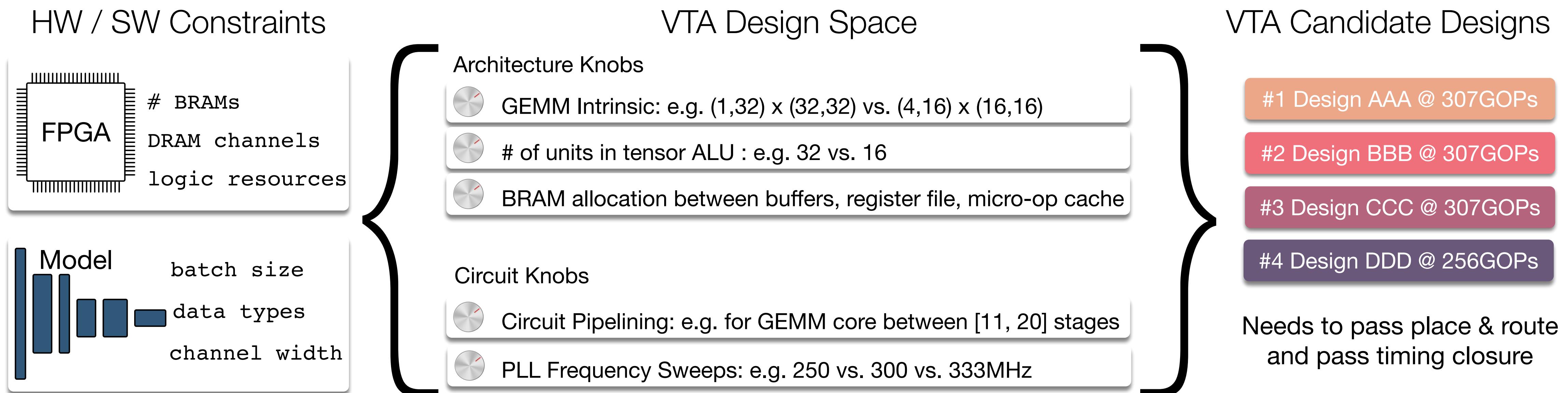
HW / SW Constraints



Hardware Exploration with VTA

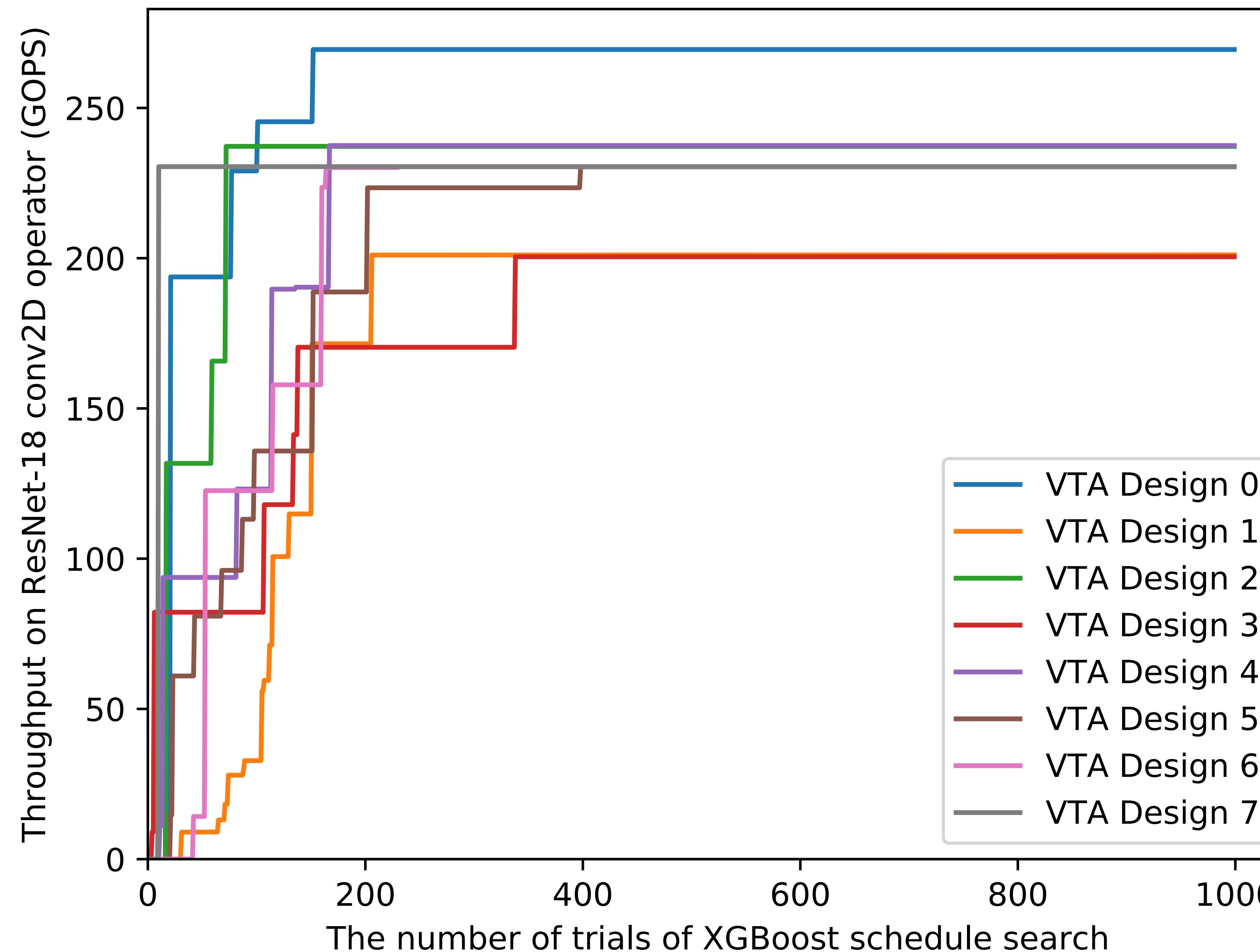


Hardware Exploration with VTA



AutoTVM for Conv2D on Hardware Candidates

AutoTVM for Conv2D on Hardware Candidates



Schedule Exploration with VTA

VTA Candidate Designs

#1 Design AAA @ 307GOPs

#2 Design BBB @ 307GOPs

#3 Design CCC @ 307GOPs

#4 Design DDD @ 256GOPs

Needs to pass place & route
and pass timing closure

Schedule Exploration with VTA

VTA Candidate Designs

#1 Design AAA @ 307GOPs

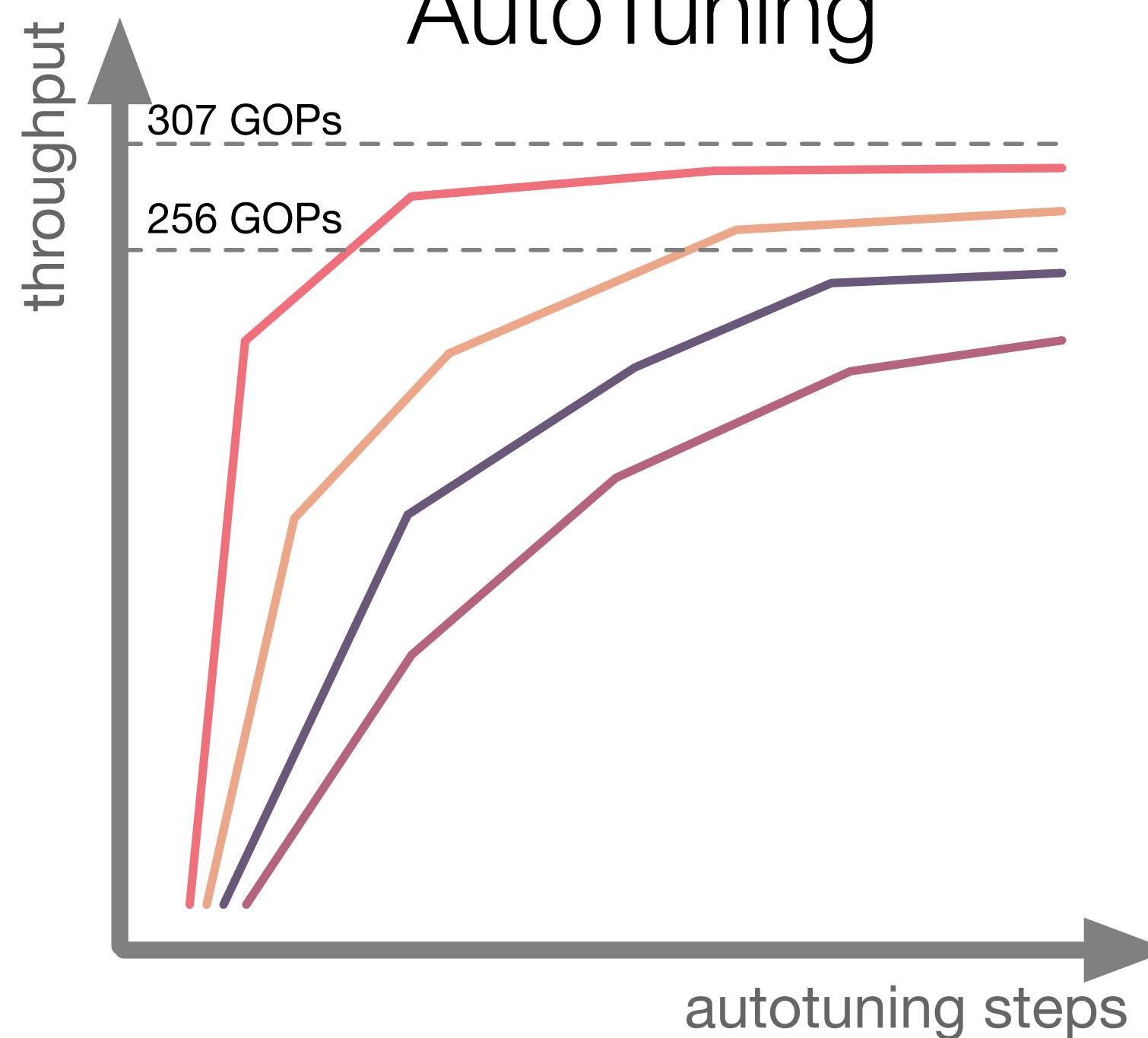
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Operator Performance AutoTuning



Schedule Exploration with VTA

VTA Candidate Designs

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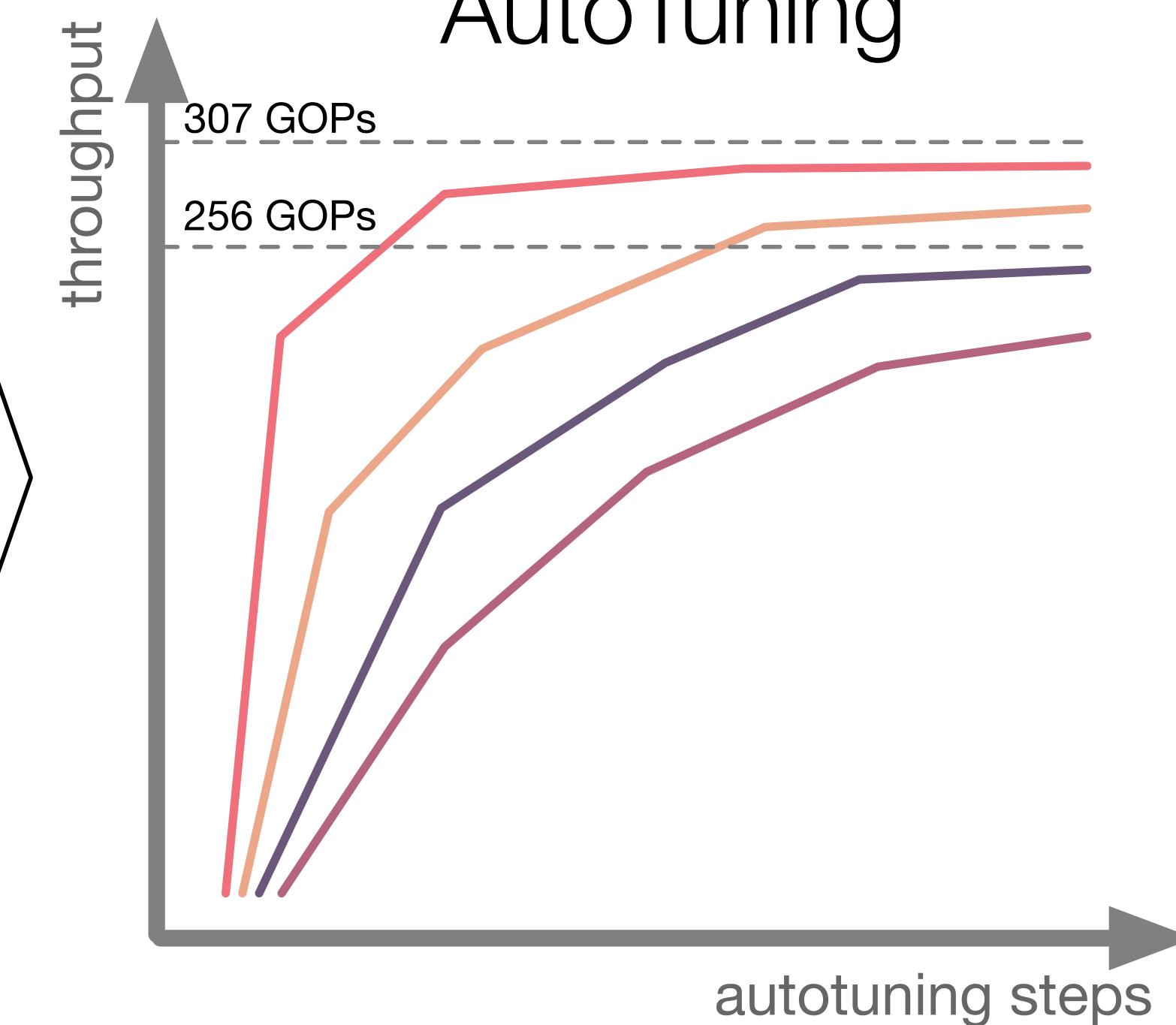
#2 Design BBB @ 307GOPs

#3 Design CCC @ 307GOPs

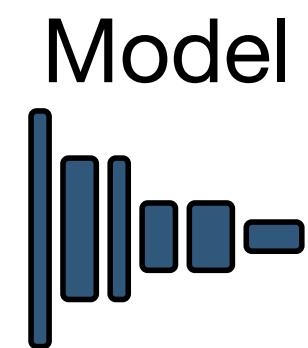
#4 Design DDD @ 256GOPs

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and pass timing closure

Operator Performance AutoTuning



Deliverable



Model



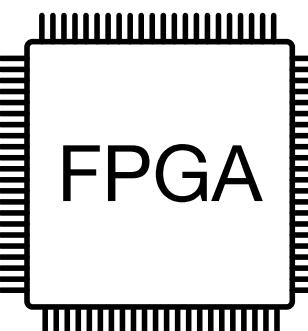
Graph Optimizer



Tuned Operator Lib



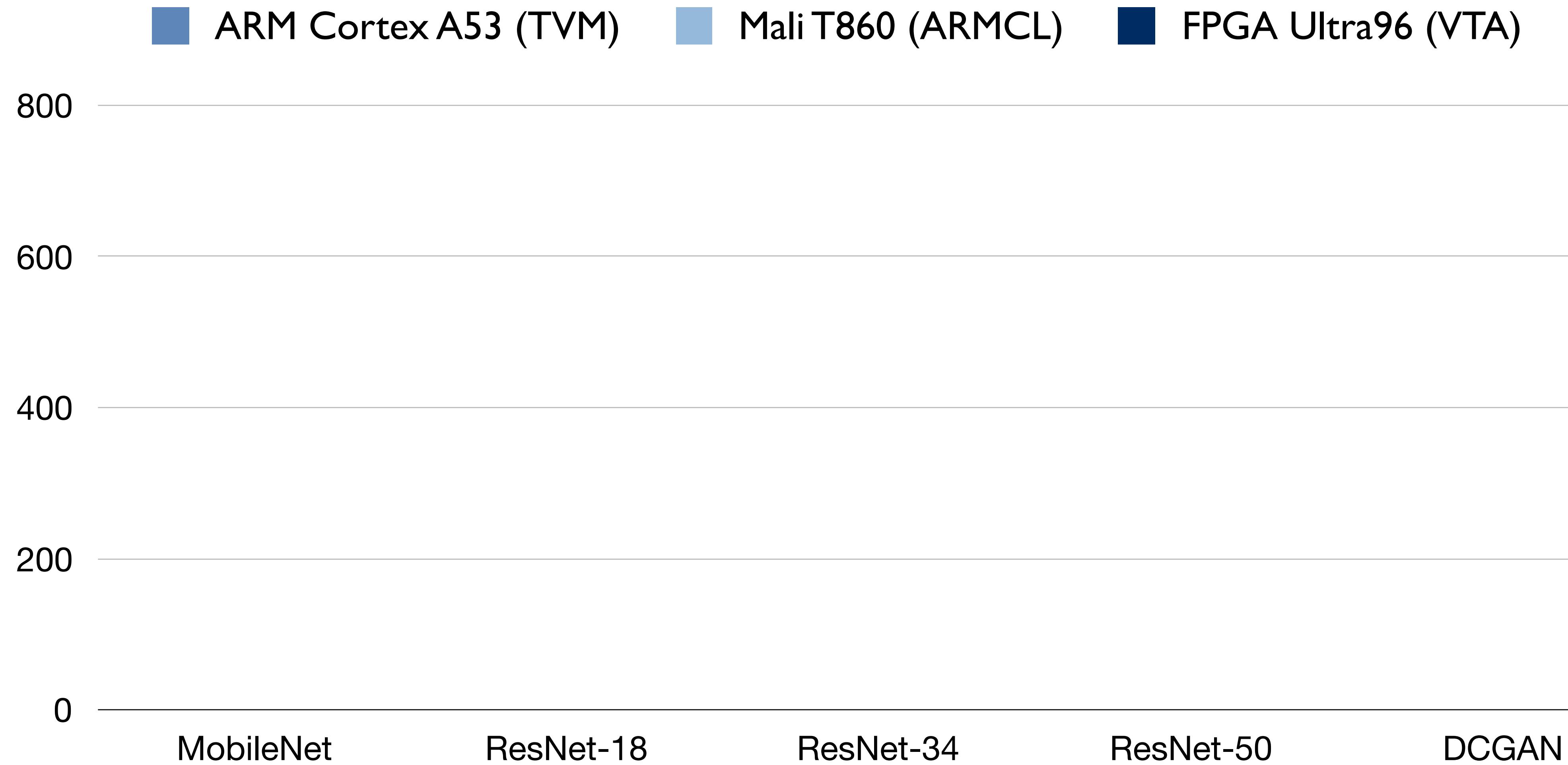
VTA Design BBB



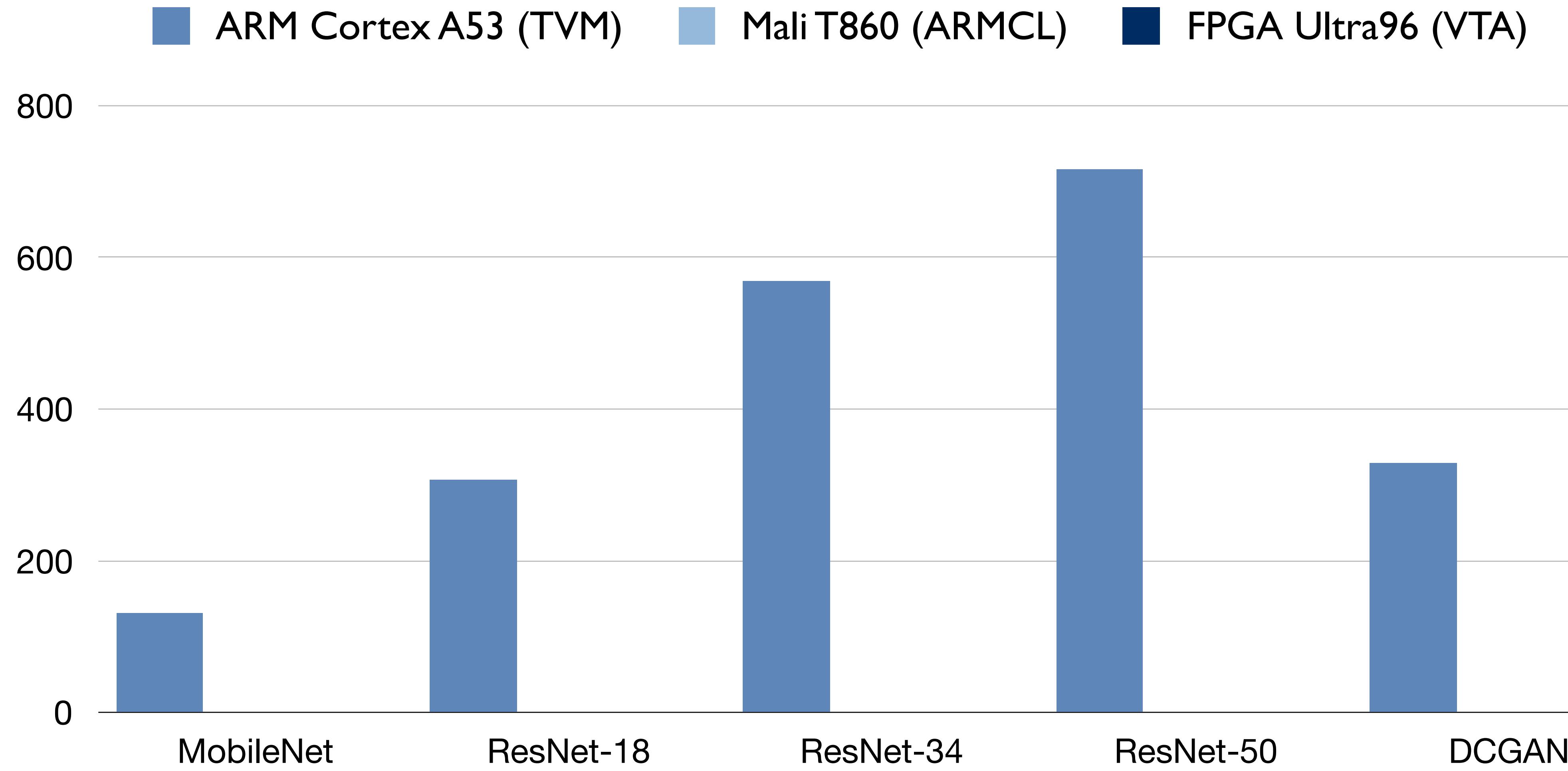
FPGA

End-to-end Performance

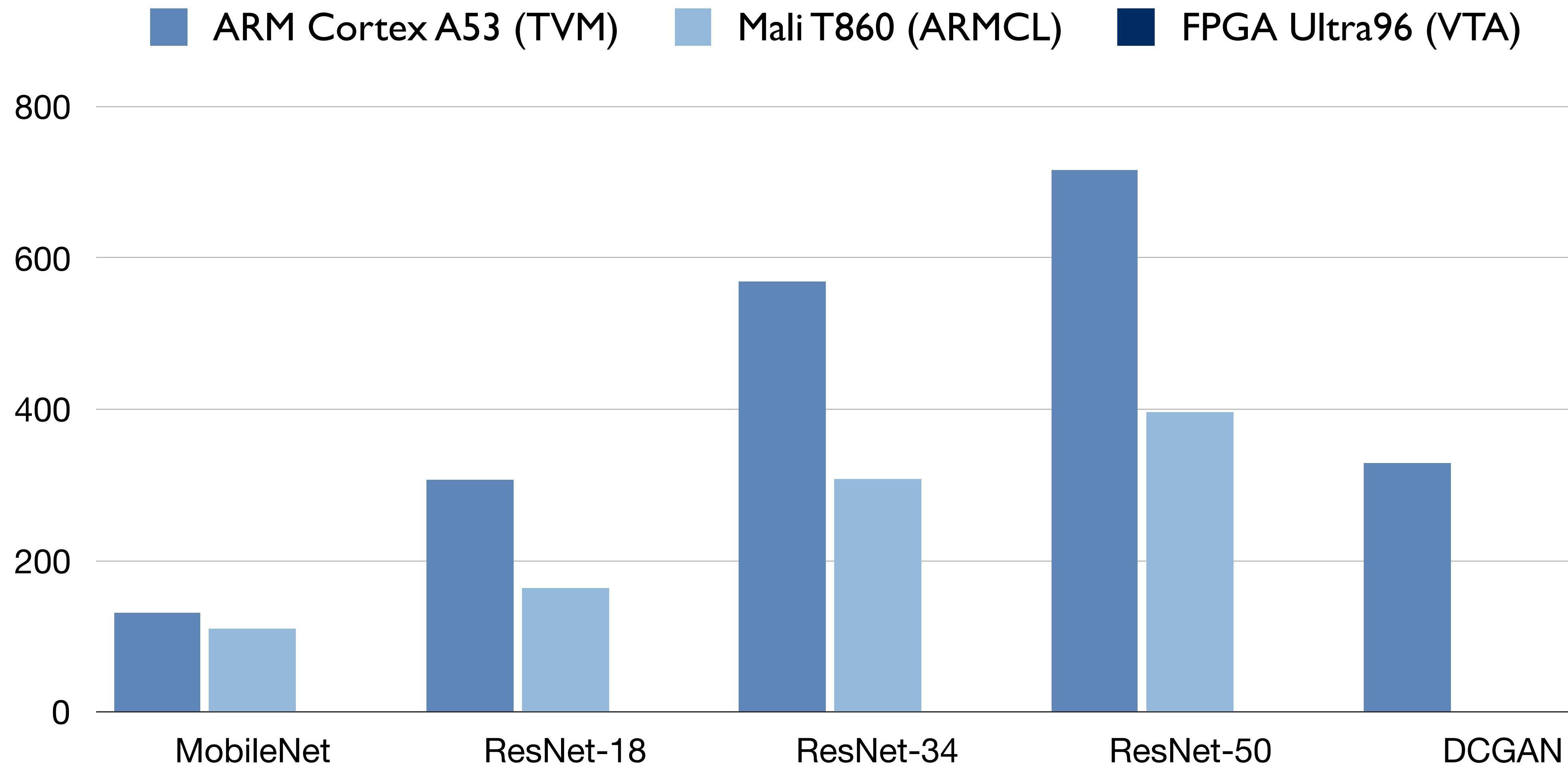
End-to-end Performance



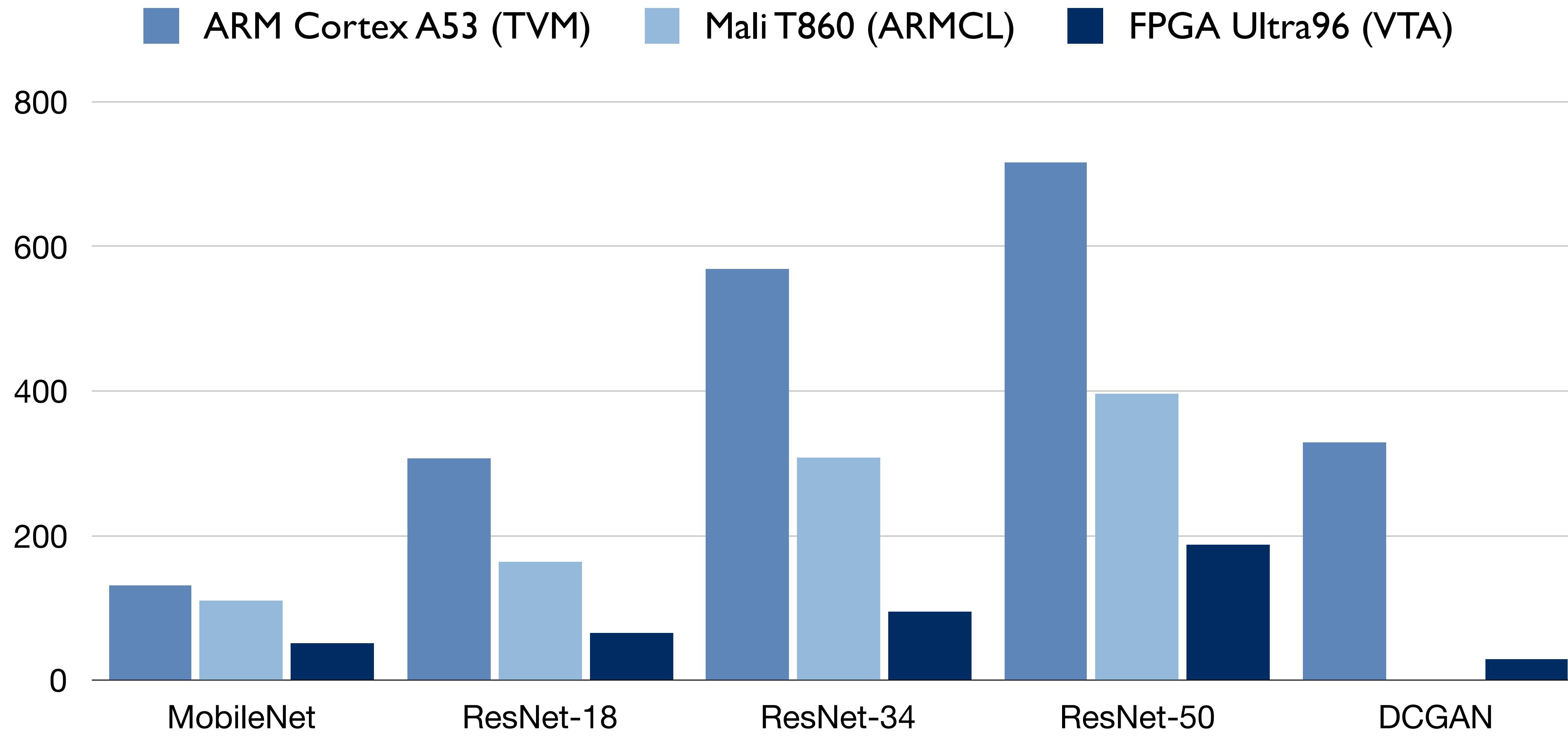
End-to-end Performance



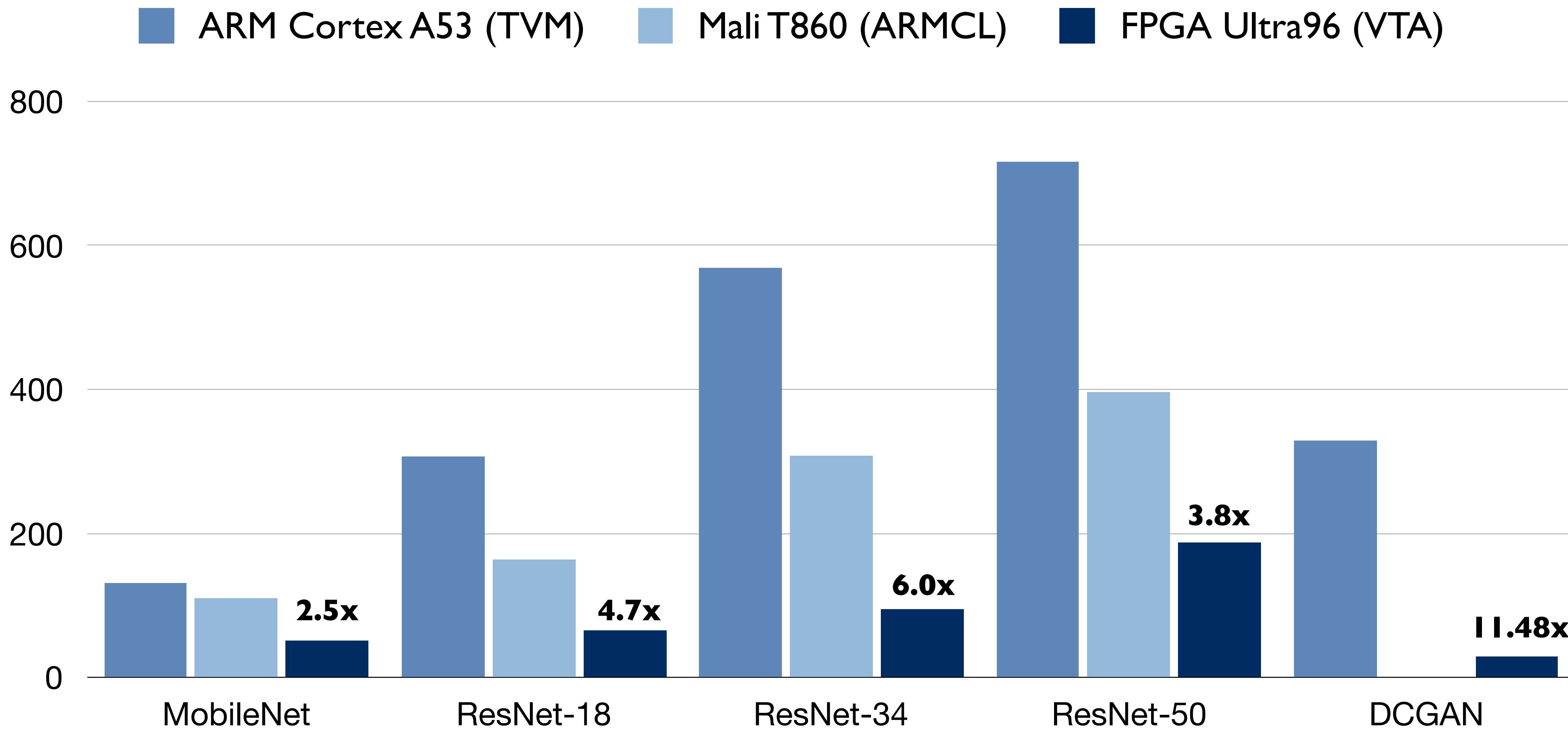
End-to-end Performance



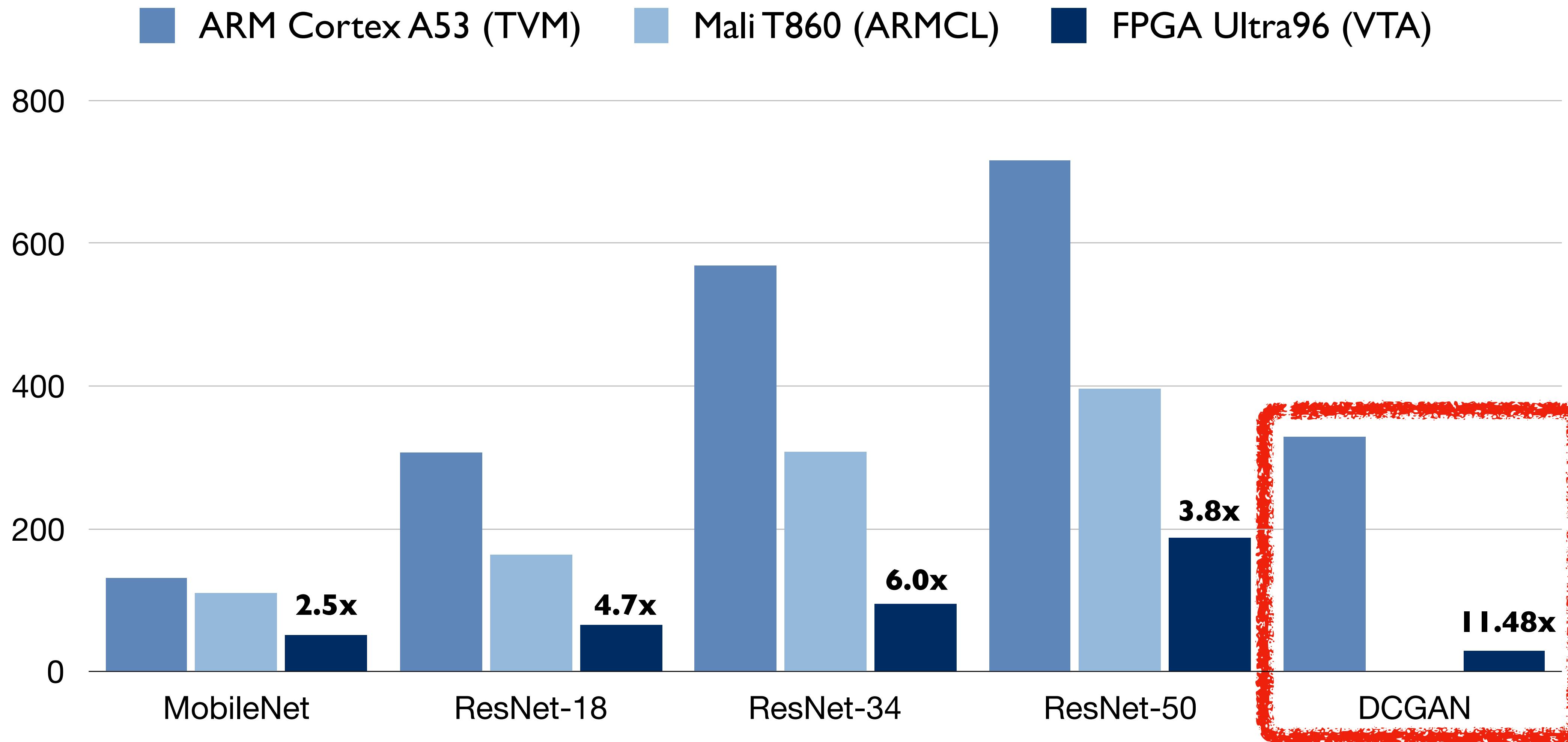
End-to-end Performance



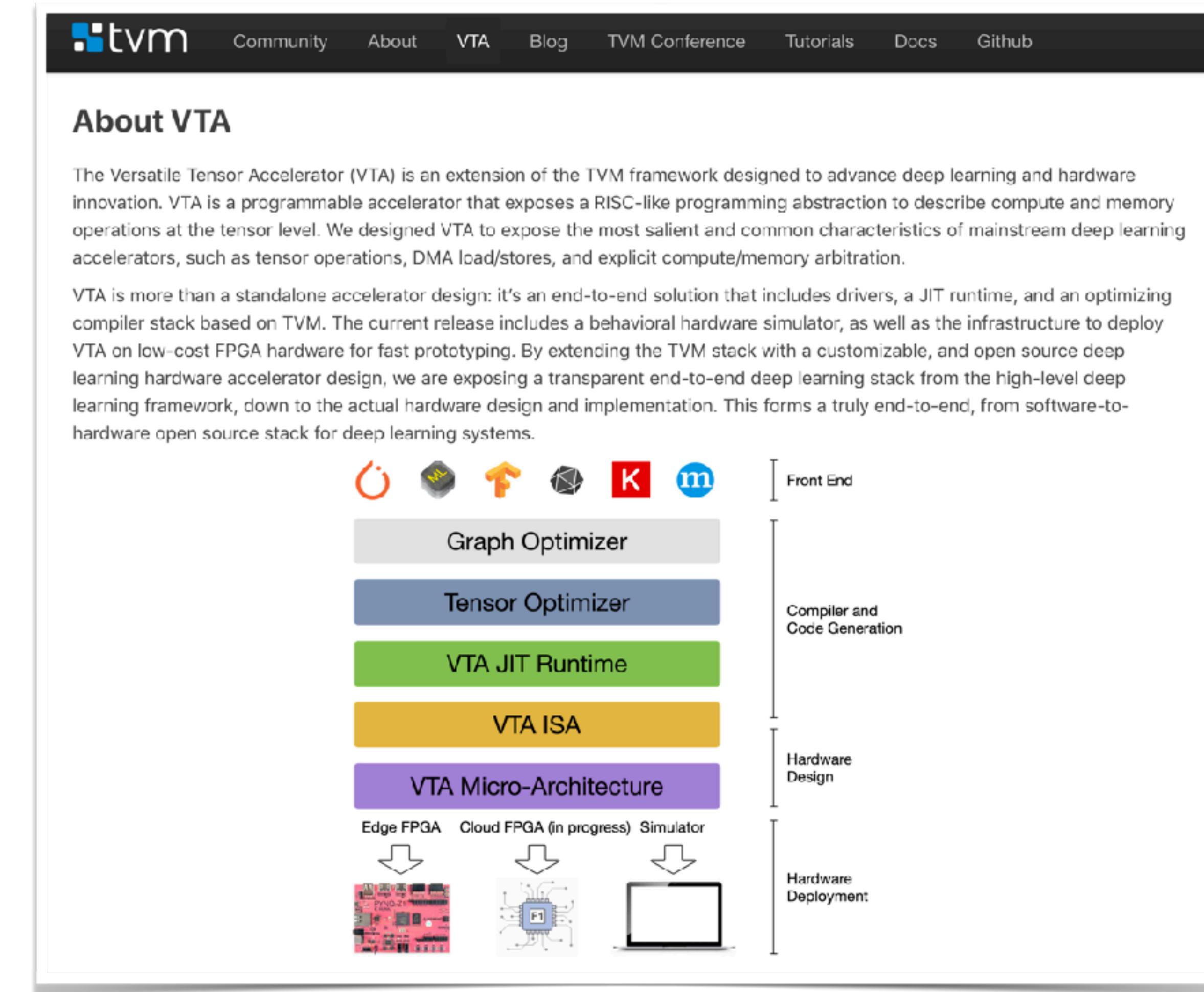
End-to-end Performance



End-to-end Performance

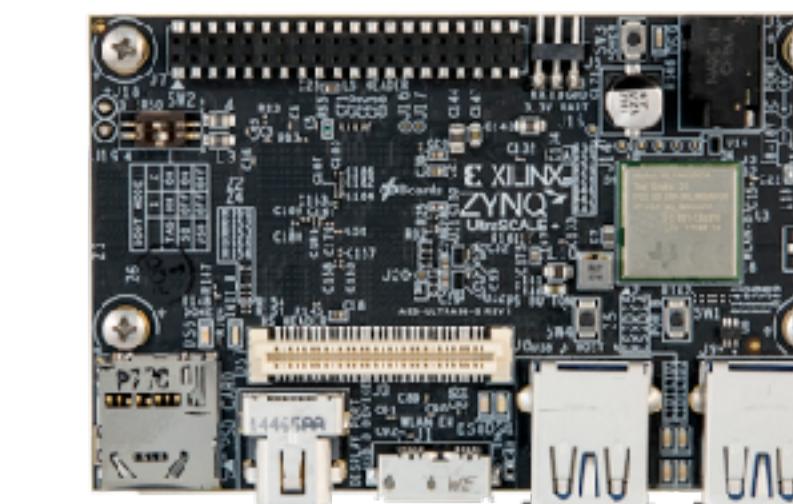
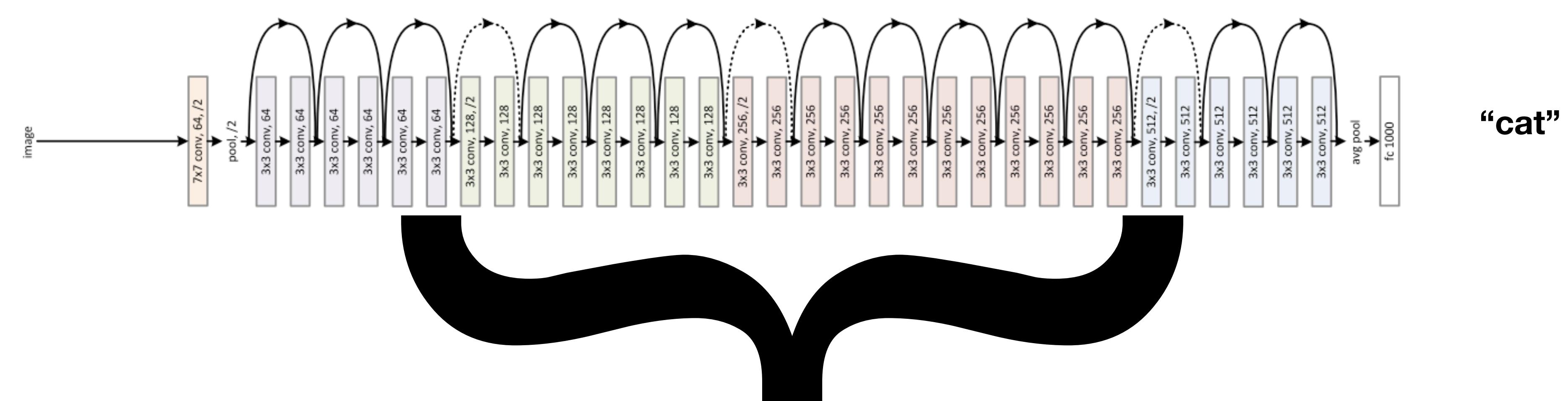


VTA Released in the Summer



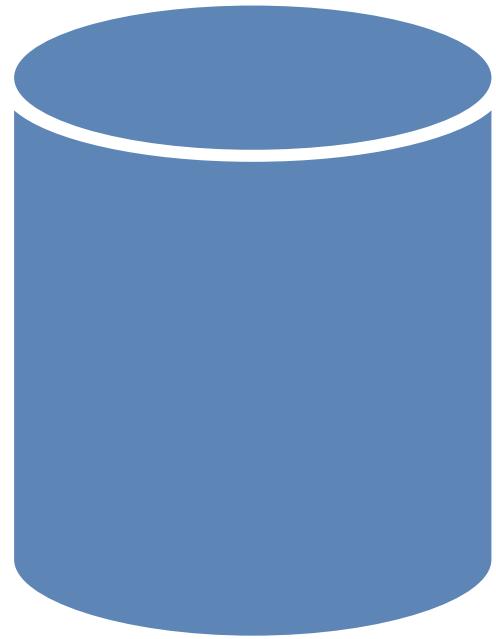
VTA Demonstration

Based on of the box FPGA demo & tutorials that you can try on your own!



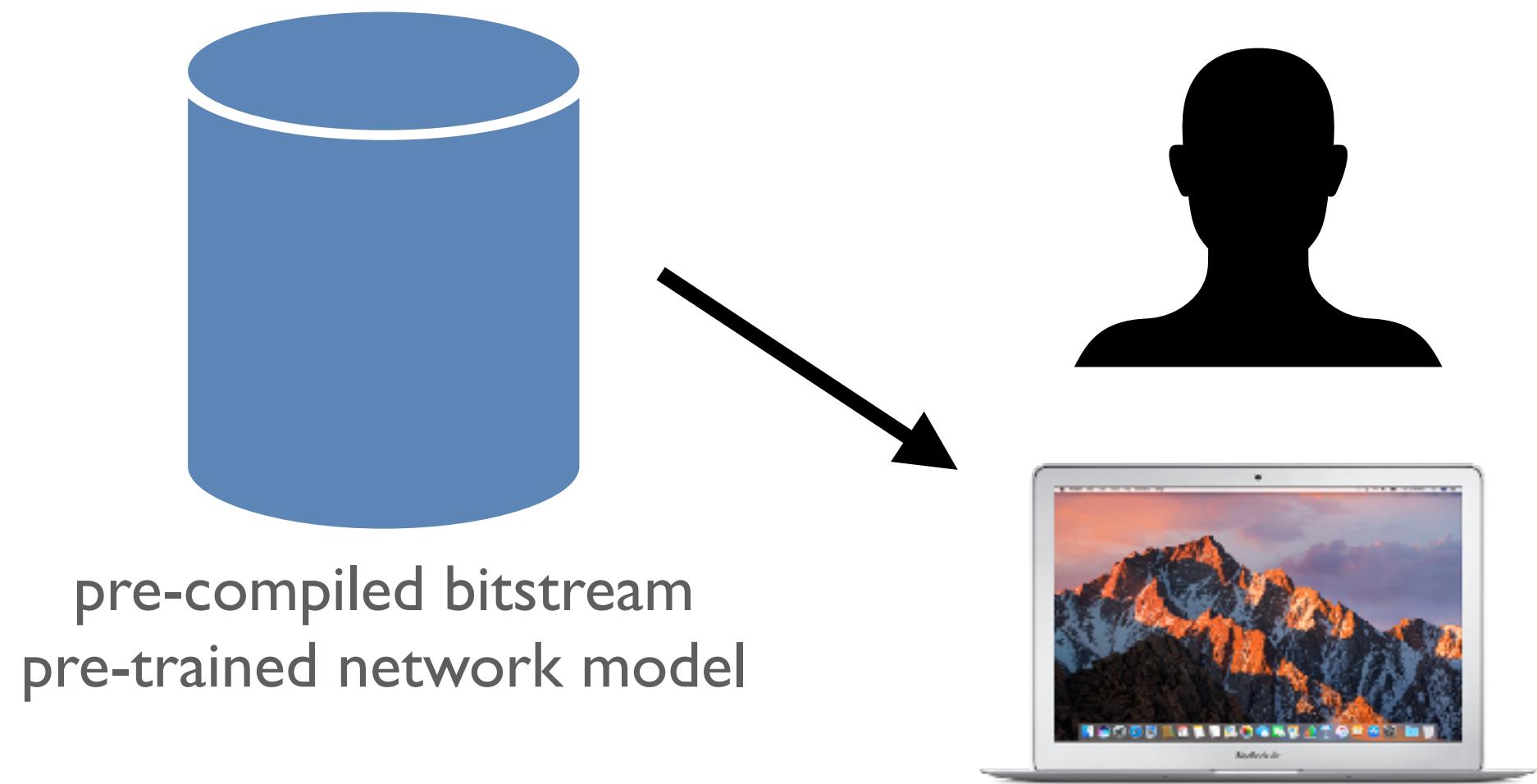
VTA Demonstration

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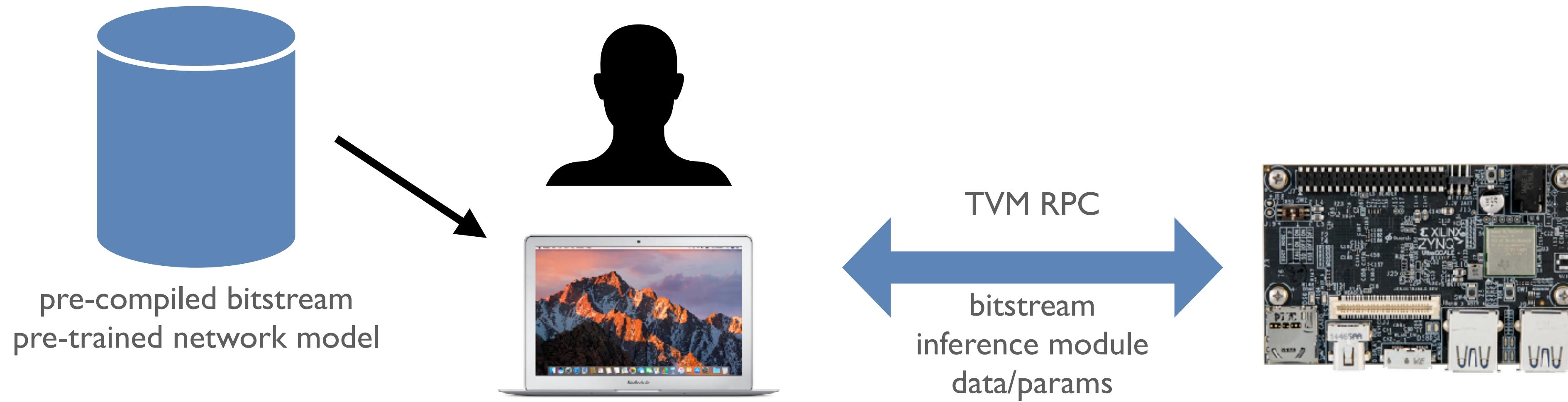


pre-compiled bitstream
pre-trained network model

VTA Demonstration



VTA Demonstration



VTA Demonstration

- I. CPU Only Inference (ResNet34,W8)
- 2.VTA Inference (ResNet34,W8)
3. Fast VTA Inference (ResNet18,W4)

VTA Demonstration

1. CPU Only Inference (ResNet34,W8): 2.6 FPS
2. VTA Inference (ResNet34,W8): 10 FPS
3. Fast VTA Inference (ResNet18,W4): 19 FPS

TVM 0.5 VTA Release Features

TVM 0.5 VTA Release Features

- FPGA Support: Ultra96, ZCU102, Intel DE10Nano
- TOPI Operator Library & AutoTVM support
- Relay graph conversion front end, push-button 8bit quantization

2019 VTA Timeline

2019 VTA Timeline

- Q1:
 - Chisel Generator for ASIC backends
 - Initial Datacenter FPGA Prototype
- Q2:
 - Novel Numerical Representation Support (Posit)
 - Initial Training Prototype

More at tvm.ai/vta

