CS 3101 Practice Quiz 1

Na	ıme:	Time allotted: Pages: 4 printed sides 45 minutes + 5 min scan/upload Maximum score: 36 points
		ules dictate strict penalties for any form of academic dishonesty. Looking sideways will be Look at only your own exam at all times.
		questions, some with subparts. Read them carefully to avoid throwing away points!! Write in the space provided. Closed book, closed notes. Calculators are allowed.
Pa	rtial cred	it rule: Must show your intermediate steps clearly for partial credit!
1.	Fill in th	ne blanks with at most a few words: (1 point * 6 = 6 points)
	(a)	The is a part of the Operating System, which when a program is about to run, allocates the memory segments needed for it, and performs dynamic
		program is about to run, allocates the memory segments needed for it, and performs dynamic linking of libraries.
	(b)	In a compiler infrastructure, the part of the compiler that is rewritten to support each new instruction set is called the
	(c)	In lexical analysis, according to the rule, the longest initial substring of the input that can match any regular expression is taken as the next token.
	(d)	Often control and data flow analysis is done on instead of assembly (low level) code.
	(e)	Statements that define non-terminals in terms of terminals and non-terminals are called
	(f)	In Context Free Grammars, the act of finding out if a string is a member of presented grammar by the following a sequence of expansion rules is called a
2.	than on	a subpart (i) to (ii) below, circle all correct answers from among the four given - note that more e , one , or none of the answers may be correct! (<i>Grading note</i> : each part will be graded as four dent questions, each worth ½ point.) (2 points * 2 = 4 points)
	(i)	Regarding parsing,
		an ambiguous grammar is one in which there is more than one derivation for a given input being accepted by the grammar.
	(b)	CFGs are in general implemented by DFAs.
		unambiguous grammars are often more complex than ambiguous grammars for the same input language.
	(d)	CFGs can implement a greater set of languages than regular expressions can.

	(11) In lexical analysis,
	(a) regular expressions are used to specify the set of strings accepted for each token.
	(b) compiler writers in modern compilers usually write the code for the finite automaton(s) required for the language.
	(c) There is only one unique regular expression possible for each language that can be accepted by regular expressions.
	(d) rule priority ensures that "if8" is recognized as a single identifier, instead of a keyword followed by a number.
3.	Write a regular expression for the language of strings containing only a's and b's that contain an odd number of a's.
	(4 points)
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4.	Draw a DFA for the language of strings containing only a's and b's that contain an even number of a's and odd number of b's. (4 points)
	(4 points)

5.	Below is a small subset of the CFG rules for statements and the if-then construct in a programming language:	
	STMTLIST \rightarrow STMTLIST STMT SEMICOLON STMTLIST \rightarrow ϵ	
	IFTHEN \rightarrow IF LPAREN COND RPAREN LCURLY STMTLIST RCURLY	
	Write a suitable abstract CFG for the above grammar subset. (4 point	s)
6.	Consider the following grammar. $(3+3 = 6 \text{ points})$	
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above.

7. The C programming language includes switch statements. Here is an example:

(8 points)

```
switch(i) {
    case 1 : c = '1';
        break;

    case 2 : c = '2';
        break;

    default: c = 'e';
        break;
}
```

Write a context-free grammar for accepting switch case statements like the one above. Your grammar must be unambiguous. <u>IMPORTANT</u>: You should use STATEMENT as a non-terminal without attempting to expand it any further. (In other words, STATEMENT is expanded by other parts of the grammar which you are not responsible for. Please note that STATEMENT non-terminal does not include the SEMICOLON token at its end.)

Here are the assumptions:

- A token ID that recognizes identifiers and variables such as i is available.
- A token NUM that recognizes integer numbers such as 1 or 2 is available.
- Tokens switch, case, default, and break are available to recognize those keywords.
- Switch statements can only accept variables like i as arguments. The only variable type that is passed to a switch statement is an integer.
- Each switch statement must have at least one (non-default) case.
- Each case statement should be followed by a value for the identifier in the switch.
- The body of each case statement is zero or more STATEMENTs which <u>must</u> be followed by a BREAKSTATEMENT.
- The presence of "default" is optional.
- Use descriptive non-terminal and terminal names like SWITCH (non-terminal) and SEMICOLON (terminal). (Caution: The TA will not grade your answer if you use non-descriptive names!!)