

Top 14 patterns that can be used to solve any coding interview question.

1. Sliding Window

The Sliding Window pattern is used to perform a required operation on a specific window size of a given array or linked list, such as finding the longest subarray containing all 1s. Sliding Windows start from the 1st element and keep shifting right by one element and adjust the length of the window according to the problem that you are solving. In some cases, the window size remains constant and in other cases the sizes grows or shrinks.

2. Two Pointers or Iterators

Two Pointers is a pattern where two pointers iterate through the data structure in tandem until one or both of the pointers hit a certain condition. Two Pointers is often useful when searching pairs in a sorted array or linked list; for example, when you have to compare each element of an array to its other elements.

Two pointers are needed because with just pointer, you would have to continually loop back through the array to find the answer. This back and forth with a single iterator is inefficient for time and space complexity — a concept referred to as asymptotic analysis. While the brute force or naive solution with 1 pointer would work, it will produce something along the lines of $O(n^2)$. In many cases, two pointers can help you find a solution with better space or runtime complexity.

3. Fast and Slow pointers

The Fast and Slow pointer approach, also known as the Hare & Tortoise algorithm, is a pointer algorithm that uses two pointers which move through the array (or sequence/linked list) at different speeds. This approach is quite useful when dealing with cyclic linked lists or arrays.

By moving at different speeds (say, in a cyclic linked list), the algorithm proves that the two pointers are bound to meet. The fast pointer should catch the slow pointer once both the pointers are in a cyclic loop.

4. Merge Intervals

The Merge Intervals pattern is an efficient technique to deal with overlapping intervals. In a lot of problems involving intervals, you either need to find overlapping intervals or merge intervals if they overlap. The pattern works like this:

Given two intervals ('a' and 'b'), there will be six different ways the two intervals can relate to each other:

Understanding and recognizing these six cases will help you solve a wide range of problems from inserting intervals to optimizing interval merges.

5. Cyclic sort

This pattern describes an interesting approach to deal with problems involving arrays containing numbers in a given range. The Cyclic Sort pattern iterates over the array one number at a time, and if the current number you are iterating is not at the correct index, you swap it with the number at its correct index. You could try placing the number in its correct index, but this will produce a complexity of $O(n^2)$ which is not optimal, hence the Cyclic Sort pattern.

6. In-place reversal of linked list

In a lot of problems, you may be asked to reverse the links between a set of nodes of a linked list. Often, the constraint is that you need to do this in-place, i.e., using the existing node objects and without using extra memory. This is where the above mentioned pattern is useful.

This pattern reverses one node at a time starting with one variable (current) pointing to the head of the linked list, and one variable (previous) will point to the previous node that you have processed. In a lock-step manner, you will reverse the current node by pointing it to the previous before moving on to the next node. Also, you will update the variable “previous” to always point to the previous node that you have processed.

7. Tree BFS

This pattern is based on the Breadth First Search (BFS) technique to traverse a tree and uses a queue to keep track of all the nodes of a level before jumping onto the next level. Any problem involving the traversal of a tree in a level-by-level order can be efficiently solved using this approach.

The Tree BFS pattern works by pushing the root node to the queue and then continually iterating until the queue is empty. For each iteration, we remove the node at the head of the queue and “visit” that node. After removing each node from the queue, we also insert all of its children into the queue.

8. Tree DFS

Tree DFS is based on the Depth First Search (DFS) technique to traverse a tree.

You can use recursion (or a stack for the iterative approach) to keep track of all the previous (parent) nodes while traversing.

The Tree DFS pattern works by starting at the root of the tree, if the node is not a leaf you need to do three things:

Decide whether to process the current node now (pre-order), or between processing two children (in-order) or after processing both children (post-order).

Make two recursive calls for both the children of the current node to process them.

9. Two heaps

In many problems, we are given a set of elements such that we can divide them into two parts. To solve the problem, we are interested in knowing the smallest element in one part and the biggest element in the other part. This pattern is an efficient approach to solve such problems.

This pattern uses two heaps; A Min Heap to find the smallest element and a Max Heap to find the biggest element. The pattern works by storing the first half of numbers in a Max Heap, this is because you want to find the largest number in the first half. You then store the second half of numbers in a Min Heap, as you want to find the smallest number in the second half. At any time, the median of the current list of numbers can be calculated from the top element of the two heaps.

10. Subsets

A huge number of coding interview problems involve dealing with Permutations and Combinations of a given set of elements. The pattern Subsets describes an efficient Breadth First Search (BFS) approach to handle all these problems.

The pattern looks like this:

Given a set of [1, 5, 3]

Start with an empty set: [[]]

Add the first number (1) to all the existing subsets to create new subsets: [[], [1]];

Add the second number (5) to all the existing subsets: [[], [1], [5], [1,5]];

Add the third number (3) to all the existing subsets: [[], [1], [5], [1,5], [3], [1,3], [5,3], [1,5,3]].

Here is a visual representation of the Subsets pattern:

11. Modified binary search

Whenever you are given a sorted array, linked list, or matrix, and are asked to find a certain element, the best algorithm you can use is the Binary Search. This pattern describes an efficient way to handle all problems involving Binary Search.

The patterns looks like this for an ascending order set:

First, find the middle of start and end. An easy way to find the middle would be: $middle = (start + end) / 2$. But this has a good chance of producing an integer overflow so it's recommended that you represent the middle as: $middle = start + (end - start) / 2$

If the key is equal to the number at index middle then return middle

If 'key' isn't equal to the index middle:

Check if $key < arr[middle]$. If it is reduce your search to $end = middle - 1$

Check if $key > arr[middle]$. If it is reduce your search to $end = middle + 1$

Here is a visual representation of the Modified Binary Search pattern:

12. Top K elements

Any problem that asks us to find the top/smallest/frequent 'K' elements among a given set falls under this pattern.

The best data structure to keep track of 'K' elements is Heap. This pattern will make use of the Heap to solve multiple problems dealing with 'K' elements at a time from a set of given elements. The pattern looks like this:

Insert 'K' elements into the min-heap or max-heap based on the problem.
Iterate through the remaining numbers and if you find one that is larger than what you have in the heap, then remove that number and insert the larger one.

13. K-way Merge

K-way Merge helps you solve problems that involve a set of sorted arrays.

Whenever you're given 'K' sorted arrays, you can use a Heap to efficiently perform a sorted traversal of all the elements of all arrays. You can push the smallest element of each array in a Min Heap to get the overall minimum. After getting the overall minimum, push the next element from the same array to the heap. Then, repeat this process to make a sorted traversal of all elements.

The pattern looks like this:

Insert the first element of each array in a Min Heap.

After this, take out the smallest (top) element from the heap and add it to the merged list.

After removing the smallest element from the heap, insert the next element of the same list into the heap.

Repeat steps 2 and 3 to populate the merged list in sorted order.

14. Topological sort

Topological Sort is used to find a linear ordering of elements that have dependencies on each other. For example, if event 'B' is dependent on event 'A', 'A' comes before 'B' in topological ordering.

This pattern defines an easy way to understand the technique for performing topological sorting of a set of elements.

The pattern works like this:

Initialization

a) Store the graph in adjacency lists by using a HashMap

b) To find all sources, use a HashMap to keep the count of in-degrees
Build the graph and find in-degrees of all vertices

Build the graph from the input and populate the in-degrees HashMap.

Find all sources

a) All vertices with '0' in-degrees will be sources and are stored in a Queue.

Sort

a) For each source, do the following things:

—i) Add it to the sorted list.

— ii)Get all of its children from the graph.

— iii)Decrement the in-degree of each child by 1.

— iv)If a child's in-degree becomes '0', add it to the sources Queue.

b) Repeat (a), until the source Queue is empty.

Common Data Structure Operations

Data Structure	Time Complexity								Space Complexity
	Average				Worst				Worst
	Access	Search	Insertion	Deletion	Access	Search	Insertion	Deletion	
<u>Array</u>	$\Theta(1)$	$\Theta(n)$	$\Theta(n)$	$\Theta(n)$	$\Theta(1)$	$O(n)$	$O(n)$	$O(n)$	$O(n)$
<u>Stack</u>	$\Theta(n)$	$\Theta(n)$	$\Theta(1)$	$\Theta(1)$	$O(n)$	$O(n)$	$O(1)$	$O(1)$	$O(n)$
<u>Queue</u>	$\Theta(n)$	$\Theta(n)$	$\Theta(1)$	$\Theta(1)$	$O(n)$	$O(n)$	$O(1)$	$O(1)$	$O(n)$
<u>Singly-Linked List</u>	$\Theta(n)$	$\Theta(n)$	$\Theta(1)$	$\Theta(1)$	$O(n)$	$O(n)$	$O(1)$	$O(1)$	$O(n)$
<u>Doubly-Linked List</u>	$\Theta(n)$	$\Theta(n)$	$\Theta(1)$	$\Theta(1)$	$O(n)$	$O(n)$	$O(1)$	$O(1)$	$O(n)$
<u>Skip List</u>	$\Theta(\log(n))$	$\Theta(\log(n))$	$\Theta(\log(n))$	$\Theta(\log(n))$	$O(n)$	$O(n)$	$O(n)$	$O(n)$	$O(n \log(n))$
<u>Hash Table</u>	N/A	$\Theta(1)$	$\Theta(1)$	$\Theta(1)$	N/A	$O(n)$	$O(n)$	$O(n)$	$O(n)$
<u>Binary Search Tree</u>	$\Theta(\log(n))$	$\Theta(\log(n))$	$\Theta(\log(n))$	$\Theta(\log(n))$	$O(n)$	$O(n)$	$O(n)$	$O(n)$	$O(n)$
<u>Cartesian Tree</u>	N/A	$\Theta(\log(n))$	$\Theta(\log(n))$	$\Theta(\log(n))$	N/A	$O(n)$	$O(n)$	$O(n)$	$O(n)$
<u>B-Tree</u>	$\Theta(\log(n))$	$\Theta(\log(n))$	$\Theta(\log(n))$	$\Theta(\log(n))$	$O(\log(n))$	$O(\log(n))$	$O(\log(n))$	$O(\log(n))$	$O(n)$
<u>Red-Black Tree</u>	$\Theta(\log(n))$	$\Theta(\log(n))$	$\Theta(\log(n))$	$\Theta(\log(n))$	$O(\log(n))$	$O(\log(n))$	$O(\log(n))$	$O(\log(n))$	$O(n)$
<u>Splay Tree</u>	N/A	$\Theta(\log(n))$	$\Theta(\log(n))$	$\Theta(\log(n))$	N/A	$O(\log(n))$	$O(\log(n))$	$O(\log(n))$	$O(n)$
<u>AVL Tree</u>	$\Theta(\log(n))$	$\Theta(\log(n))$	$\Theta(\log(n))$	$\Theta(\log(n))$	$O(\log(n))$	$O(\log(n))$	$O(\log(n))$	$O(\log(n))$	$O(n)$
<u>KD Tree</u>	$\Theta(\log(n))$	$\Theta(\log(n))$	$\Theta(\log(n))$	$\Theta(\log(n))$	$O(n)$	$O(n)$	$O(n)$	$O(n)$	$O(n)$

Array Sorting Algorithms

Algorithm	Time Complexity			Space Complexity
	Best	Average	Worst	Worst
<u>Quicksort</u>	$\Omega(n \log(n))$	$\Theta(n \log(n))$	$O(n^2)$	$O(\log(n))$
<u>Mergesort</u>	$\Omega(n \log(n))$	$\Theta(n \log(n))$	$O(n \log(n))$	$O(n)$
<u>Timsort</u>	$\Omega(n)$	$\Theta(n \log(n))$	$O(n \log(n))$	$O(n)$
<u>Heapsort</u>	$\Omega(n \log(n))$	$\Theta(n \log(n))$	$O(n \log(n))$	$O(1)$
<u>Bubble Sort</u>	$\Omega(n)$	$\Theta(n^2)$	$O(n^2)$	$O(1)$
<u>Insertion Sort</u>	$\Omega(n)$	$\Theta(n^2)$	$O(n^2)$	$O(1)$
<u>Selection Sort</u>	$\Omega(n^2)$	$\Theta(n^2)$	$O(n^2)$	$O(1)$
<u>Tree Sort</u>	$\Omega(n \log(n))$	$\Theta(n \log(n))$	$O(n^2)$	$O(n)$
<u>Shell Sort</u>	$\Omega(n \log(n))$	$\Theta(n(\log(n))^2)$	$O(n(\log(n))^2)$	$O(1)$
<u>Bucket Sort</u>	$\Omega(n+k)$	$\Theta(n+k)$	$O(n^2)$	$O(n)$
<u>Radix Sort</u>	$\Omega(nk)$	$\Theta(nk)$	$O(nk)$	$O(n+k)$
<u>Counting Sort</u>	$\Omega(n+k)$	$\Theta(n+k)$	$O(n+k)$	$O(k)$
<u>Cubesort</u>	$\Omega(n)$	$\Theta(n \log(n))$	$O(n \log(n))$	$O(n)$