



Arrays: Quicksort





- Explore the Quick sort algorithm
- Understand the following aspects
 - Algorithm mechanism and pseudocode
 - Algorithm iterations on varying input
 - Algorithm time and space complexity





- A "randomized" sort algorithm:
 - Addresses performance issues with Merge sort.
 - Performs well in most scenarios, badly very rarely.
- Algorithm scheme:
 - Randomly shuffle the given array A[0...N-1].
 - ii. Partition the array into 2 pieces, as follows:
 - For some index j, A[j] is in-place.
 - No element to the left of j is larger than A[j].
 - No element to the right of j is smaller than A[j].
 - iii. Sort each piece recursively.





- Characteristics:
 - Quick sort is a comparison-based sort.
 - Quick sort can sort an array in-place.
 - Auxiliary array not required.
 - Starts by randomly shuffling the order of elements.
- Comparison operation:
 - Defined as required for the data type
 - Numbers
 - Strings
 - Objects: by attributes





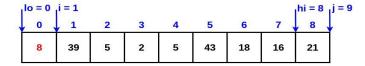
- 1. Shuffle the array.
- 2. Partition this array, so that, for some j:
 - a. Array[j] is in place.
 - b. All entries to the left of j are smaller.
 - c. All entries to the right of j are larger.
- 3. Recursively sort each piece.



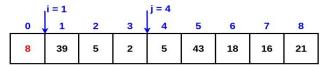


How Quicksort Partitioning Works

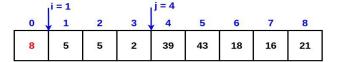
Call to Partition



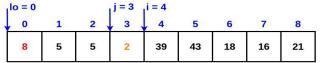
End Of Initial Iterations



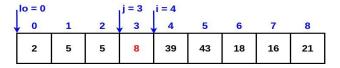
Swap Array[i] And Array[j]



Final Position of j



Swap Array[lo] And Array[j]



Partitioning Complete!

Quick sort pseudocode



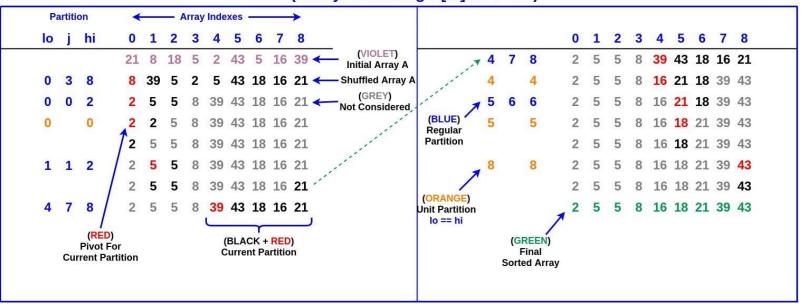
Code For Quick Sort

```
import random
   def swap(array, i, j):
       temp = array[i]
       array[i] = array[j]
       array[j] = temp
def partition(array, lo, hi):
           i = lo
           j = hi+1
                                                           def quick sort(array, lo, hi):
                                                                    if hi <= lo:
         while True:
                                                                        return
            while array[i] < array[lo]:</pre>
                                                              j = partition(array, lo, hi)
               if i == hi:
                  break
                                                              quick sort(array, lo, j-1)
                                                              quick sort(array, i+1, hi)
               i += 1
                                                              def wrapper sort(array):
           i -= 1
                                                                  random.shuffle(array)
           while array[lo] < array[j]:</pre>
                                                                  quick_sort(array, 0, len(array)-1)
              if j == lo:
                  break
              j -= 1
         if i \ge j:
              break
         swap(array, i, j)
      swap(array, lo, j)
          return j
```





Quick Sort Iterations On Array A (Always Selecting A[lo] As Pivot)







Quick Sort Iterations On Pre-Sorted Array A (Always Selecting A[lo] As Pivot)

Partition		Array Indexes -								→	
lo	j	hi	0	1	2	3	4	5	6	7	8
			2	5	5	8	16	18	21	39	43 (VIOLET) Initial Array A
0	6	8	21	18	5	5	39	2	43	8	16 - Shuffled Array A
0	3	5	8	18	5	5	16	2	21	43	39
0	1	2	5	2	5	8	16	18	21	43	39
0	1	1	2	5	5	8	16	18	21	43	39
			2	5	5	8	16	18	21	43	39
4	4	5	2	5	5	8	16	18	21	43	39
			2	5	5	8	16	18	21	43	39
7	8	8	2	5	5	8	16	18	21	43	39
(00	FEAT	v.	2	5	5	8	16	18	21	39	43
	EEN nal d Arr	9-	→ 2	5	5	8	16	18	21	39	43





Quick Sort Iterations On Reverse-Sorted Array A (Always Selecting A[Io] As Pivot)

la		hi	0	1	2	3	Λ	5	6	7	8	
lo	J	hi	U	1		3	4	5	0	-	0	A/IOLET
			43	39	21	18	16	8	5	5	2	(VIOLET) Initial Array A
0	4	8	16	18	2	8	39	21	43	5	5	Shuffled Array A
0	3	5	5	5	2	8	16	21	43	39	18	
0	0	3	2	5	5	8	16	21	43	39	18	
0		0	2	5	5	8	16	21	43	39	18	
			2	5	5	8	16	21	43	39	18	
2	2	3	2	5	5	8	16	21	43	39	18	
			2	5	5	8	16	21	43	39	18	
5	6	8	2	5	5	8	16	21	43	39	18	
5		5	2	5	5	8	16	18	21	39	43	
			2	5	5	8	16	18	21	39	43	
7	7	8	2	5	5	8	16	18	21	39	43	
			2	5	5	8	16	18	21	39	43	
	EEN nal d Arı	10-	→ 2	5	5	8	16	18	21	39	43	





- Analyzing time complexity:
 - Note that the input is always randomly shuffled.
 - This does not eliminate worst behaviour of Quicksort, but greatly reduces its probability.
- Worst case: The total number of compares and entry swaps (during partitioning)
 is:

$$C(N) = N + (N-1) + (N-2) + ... + 1$$
: proportional to N^2 steps

Average case: Proportional to Nlog₂N steps





- Analyzing space complexity:
 - Quicksort is an in-place sorting algorithm.
 - Extra space is not used during the sorting.
 - Best case: Constant space
 - Worst case: Constant space
 - Average case: Constant space





We explored the Quick sort algorithm as follows:

- Algorithm mechanism and pseudocode
- Algorithm iterations on varying input
- Time and space complexity





Thank You