

1. $\langle C \rangle \rightarrow \langle V \rangle = \langle D \rangle \mid \langle V \rangle$
 $\langle V \rangle \rightarrow x \mid y \mid z$
 $\langle D \rangle \rightarrow \langle E \rangle ? \langle D \rangle : \langle D \rangle \mid \langle E \rangle$
 $\langle E \rangle \rightarrow \langle E \rangle \mid \mid \langle F \rangle \mid \langle F \rangle$
 $\langle F \rangle \rightarrow \langle F \rangle \ \&\& \ \langle G \rangle \mid \langle G \rangle$
 $\langle G \rangle \rightarrow !\langle G \rangle \mid \langle H \rangle$
 $\langle H \rangle \rightarrow (\langle H \rangle) \mid \langle I \rangle$
 $\langle I \rangle \rightarrow \text{true} \mid \text{false}$

2. Static Semantic Attributes:

type	=	{integer, double}	(synthesized)
typetable($\langle \text{var} \rangle$)	=	{integer, double, error}	(inherited)
inittable($\langle \text{var} \rangle$)	=	{true, false, error}	(inherited)
typebinding	=	($\langle \text{var} \rangle$, {integer, double})	(synthesized)
initialized	=	($\langle \text{var} \rangle$, {true, false})	(synthesized)

Attribute Rules:

```

<start1> → <stmt3> ; <start3>
<start1>.type := N/A
<start1>.typetable(<var>) := <stmt3>.typetable
<start1>.inittable(<var>) := <stmt3>.initvar
<start1>.typebinding := N/A
<start1>.initialized := N/A

<stmt3>.type := N/A
<stmt3>.typetable := <start1>.typetable
<stmt3>.inittable := <start1>.inittable
<stmt3>.typebinding := N/A
<stmt3>.initialized := N/A
<start3>.type := N/A
<start3>.typetable := <stmt3>.typetable ∪ <start1>.typetable
<start3>.inittable := <stmt3>.inittable ∪ <start1>.inittable
<start3>.typebinding := N/A
<start3>.initialized := N/A
<start2> → <stmt4>
<start2>.type := N/A
<start2>.typetable(<var>) := ∅
<start2>.inittable(<var>) := ∅
<start2>.typebinding := N/A
<start2>.initialized := N/A

<stmt4>.type := N/A
<stmt4>.typetable := <start2>.typetable
<stmt4>.inittable := <start2>.inittable

```

```

<stmt4>.typebinding := N/A
<stmt4>.initialized := N/A
<start3>.type := N/A
<stmt1> → <declare2>
<stmt1>.type := N/A
<stmt1>.typetable(<var>) := ∅(inherited from <start>)
<stmt1>.inittable(<var>) := ∅(inherited from <start>)
<stmt1>.typebinding := N/A
<stmt1>.initialized := N/A

<stmt2> → <assign2>
<stmt2>.type := N/A
<stmt2>.typetable(<var>) := ∅(inherited from <start>)
<stmt2>.inittable(<var>) := ∅(inherited from <start>)
<stmt2>.typebinding := N/A
<stmt2>.initialized := N/A

<declare1>.type :=
<declare1>.typetable(<var>) :=
<declare1>.inittable(<var>) :=
<declare1>.typebinding :=
<declare1>.initialized :=

<type1>.type :=
<type1>.typetable(<var>) :=
<type1>.inittable(<var>) :=
<type1>.typebinding :=
<type1>.initialized :=

<type2>.type :=
<type2>.typetable(<var>) :=
<type2>.inittable(<var>) :=
<type2>.typebinding :=
<type2>.initialized :=

<assign1>.type :=
<assign1>.typetable(<var>) :=
<assign1>.inittable(<var>) :=
<assign1>.typebinding :=
<assign1>.initialized :=

<expression1>.type :=
<expression1>.typetable(<var>) :=
<expression1>.inittable(<var>) :=
<expression1>.typebinding :=
<expression1>.initialized :=

<expression2>.type :=
<expression2>.typetable(<var>) :=
<expression2>.inittable(<var>) :=
<expression2>.typebinding :=
<expression2>.initialized :=

<value1>.type :=

```

```

<value1>.typetable(<var>) :=
<value1>.inittable(<var>) :=
<value1>.typebinding :=
<value1>.initialized :=

<value2>.type :=
<value2>.typetable(<var>) :=
<value2>.inittable(<var>) :=
<value2>.typebinding :=
<value2>.initialized :=

<value3>.type :=
<value3>.typetable(<var>) :=
<value3>.inittable(<var>) :=
<value3>.typebinding :=
<value3>.initialized :=

```

Table 1: Attribute Rules

3. asdf

4. Loop Invariants:

Outer (while) loop invariant: The elements $A[\text{bound} \dots n - 1]$ are in non-decreasing order

Inner (for) loop invariant: The elements $A[\text{bound} \dots n - 1]$ are in non-decreasing order and $A[0 \dots i] \leq A[i + 1] \leq A[\text{bound} - 1]$

Precondition: $n \geq 0$ and A contains n elements indexed from 0

```

bound = n;
while (bound > 0) {
  t = 0;
  for (i = 0; i < bound - 1; i++) {
    if (A[i] > A[i+1]) {
      swap = A[i];
      A[i] = A[i+1];
      A[i+1] = swap;
      t = i + 1;
    }
  }
  bound = t;
}

```

Postcondition: $A[0] \leq A[1] \leq \dots \leq A[n - 1]$

5. $M_{state}(\langle \text{var} \rangle = \langle \text{expression} \rangle, S) =$

```

{
  // test that <var> is a legal name in the language
  if  $M_{name}(\langle \text{var} \rangle) = \text{'Error'}$ 
    return 'Error'

  // test that <var> has already been declared
  if  $Lookup(M_{name}(\langle \text{var} \rangle), S) = \text{'Error'}$ 
    return 'Error'

  // calculate the value of <expression> using the old state
   $V = M_{value}(\langle \text{expression} \rangle, S)$ 
}

```

```

    if V = 'Error'
        return 'Error'

    // calculate a new state including any side effects from evaluating <expression>
    S1 = Mstate(<expression>, S)

    // remove <var> from the new state
    Remove(Mname(<var>), S)

    // return the new state with the updated value of <var> added
    return Add(Mname(<var>), V, S1)
}

Mstate(if <condition> then <statement1> else <statement2>, S) =
{
    S1 = Mstate(<condition>, S)

    if Mboolean(<condition>, S1) = true
        return Mstate(<statement1>, S1)
    else if Mboolean(<condition>, S1) = false
        return Mstate(<statement2>, S1)
    else
        return 'Error'
}

Mstate(while <condition> <loop body>, S) =
{
    S1 = Mstate(<condition>, S)

    if Mboolean(<condition>, S1) = true
        return Mstate(while <condition> <loop body>, Mstate(<condition>, S1))
    else if Mboolean(<condition>, S1) = false
        return S1
    else
        return 'Error'
}

```