Rebecca Frederick EECS 345 Written Exercise 1 February 2, 2016

```
<C>
        \rightarrow <V> = <D> | <V>
<V>
        \rightarrow x | y | z
        \rightarrow <E> ? <D> : <D> | <E>
<D>
<E>

ightarrow <E> || <F> | <F>
        \rightarrow
             <F> && <G> <G>
<F>
             !<G> | <H>
<G>
        \rightarrow
              (<H>) | <I>
<H>
              true | false
<I>
```

2. Static Semantic Attributes:

Attribute Rules:

```
\langle start_1 \rangle \rightarrow \langle stmt_3 \rangle; \langle start_3 \rangle
<start<sub>1</sub>>.type := N/A
<start1>.typetable(<var>) := <stmt3>.typetable
<start1>.inittable(<var>) := <stmt3>.initvar
{\rm start_1} > . {\rm typebinding} := N/A
{\rm start_1}>.initialized := N/A
<stmt_3>.type := N/A
<stmt3>.typetable := <start1>.typetable
{\rm <stmt_3>.inittable} := {\rm <start_1>.inittable}
{\rm <stmt_3>.typebinding} := N/A
{\rm <stmt_3>.initialized} := N/A
<start<sub>3</sub>>.type := N/A
{\sf start_3>}.typetable := {\sf stmt_3>}.typetable \cup {\sf start_1>}.typetable
{\sf start_3>.inittable} := {\sf stmt_3>.inittable} \cup {\sf start_1>.inittable}
{\rm start_3}.typebinding := N/A
{\rm start_3}.initialized := N/A
\langle start_2 \rangle \rightarrow \langle stmt_4 \rangle
<start<sub>2</sub>>.type := N/A
<start2>.typetable(<var>) := 0
<start2>.inittable(<var>) := Ø
{\rm start_2}.typebinding := N/A
{\rm <start_2>.initialized} := N/A
<stmt<sub>4</sub>>.type := N/A
<stmt<sub>4</sub>>.typetable := <start<sub>2</sub>>.typetable
{\rm <stmt_4>.inittable} := {\rm <start_2>.inittable}
```

```
<stmt<sub>4</sub>>.typebinding := N/A
{\rm <stmt_4>.initialized} := N/A
{\rm start_3}.type := N/A
\langle stmt_1 
angle \; 	o \; \langle declare_2 
angle
<stmt<sub>1</sub>>.type := N/A
<stmt<sub>1</sub>>.typetable(<var>) := Ø(inherited from <start>)
<stmt₁>.inittable(<var>) := ∅(inherited from <start>)
{\rm <stmt_1>.typebinding} := N/A
{\rm <stmt_1>.initialized} := N/A
\langle stmt_2 \rangle \rightarrow \langle assign_2 \rangle
<stmt_2>.type := N/A
<stmt<sub>2</sub>>.typetable(<var>) := \emptyset(inherited from <start>)
<stmt<sub>2</sub>>.inittable(<var>) := \emptyset(inherited from <start>)
{\rm <stmt_2>.typebinding} := N/A
{\rm <stmt_2>.initialized} := N/A
<declare1>.type :=
<declare<sub>1</sub>>.typetable(<var>) :=
<declare<sub>1</sub>>.inittable(<var>) :=
<declare<sub>1</sub>>.typebinding :=
<declare<sub>1</sub>>.initialized :=
<type1>.type :=
<type<sub>1</sub>>.typetable(<var>) :=
<type<sub>1</sub>>.inittable(<var>) :=
<type1>.typebinding :=
<type1>.initialized :=
<type<sub>2</sub>>.type :=
<type<sub>2</sub>>.typetable(<var>) :=
<type<sub>2</sub>>.inittable(<var>) :=
<type2>.typebinding :=
<type2>.initialized :=
\langle assign_1 \rangle .type :=
<assign<sub>1</sub>>.typetable(<var>) :=
<assign<sub>1</sub>>.inittable(<var>) :=
\langle assign_1 \rangle.typebinding :=
\langle assign_1 \rangle.initialized :=
<expression_1>.type :=
<expression1>.typetable(<var>) :=
<expression<sub>1</sub>>.inittable(<var>) :=
<expression<sub>1</sub>>.typebinding :=
<expression<sub>1</sub>>.initialized :=
<expression_2>.type :=
<expression<sub>2</sub>>.typetable(<var>) :=
<expression2>.inittable(<var>) :=
<expression_2>.typebinding :=
<expression_2>.initialized :=
<value<sub>1</sub>>.type :=
```

```
<value1>.typetable(<var>) :=
<value1>.inittable(<var>) :=
<value1>.typebinding :=
<value1>.initialized :=

<value2>.type :=
<value2>.typetable(<var>) :=
<value2>.typetable(<var>) :=
<value2>.typebinding :=
<value2>.initialized :=

<value3>.type :=
<value3>.type :=
<value3>.typetable(<var>) :=
<value3>.initiable(<var>) :=
<value3>.initiable(<var>) :=
<value3>.initiable(<var>) :=
<value3>.initiable(<var>) :=
<value3>.initiable(<var>) :=
<value3>.initiable(<var>) :=
<value3>.initialized :=
```

Table 1: Attribute Rules

- 3. asdf
- 4. Loop Invariants:

bound = n:

Outer (while) loop invariant: The elements A[bound ...n-1] are in non-decreasing order Inner (for) loop invariant: The elements A[bound ...n-1] are in non-decreasing order and $A[0...i] \le A[i+1] \le A[bound-1]$

Precondition: $n \ge 0$ and A contains n elements indexed from 0

```
while (bound > 0) {
     t = 0;
     for (i = 0; i < bound - 1; i++) {
        if (A[i] > A[i+1]) {
          swap = A[i];
          A[i] = A[i+1];
          A[i+1] = swap;
          t = i + 1;
     }
     bound = t;
  Postcondition: A[0] \le A[1] \le \cdots \le A[n-1]
5. M_{state}(\langle var \rangle = \langle expression \rangle, S) =
        // test that <var> is a legal name in the language
        if M_{name} (<var>) = 'Error'
             return 'Error'
        // test that <var> has already been declared
        if Lookup(M_{name}(\langle var \rangle), S) = 'Error'
             return 'Error'
        // calculate the value of \langle expression \rangle using the old state
        V = M_{value} (\langle expression \rangle, S)
```

```
if V = 'Error'
           return 'Error'
     // calculate a new state including any side effects from evaluating <expression>
     S_1 = M_{state} (\langle expression \rangle, S)
     // remove <var> from the new state
     Remove(M_{name}(\langle var \rangle), S)
     // return the new state with the updated value of <var> added
     return Add(M_{name}(\langle var \rangle), V, S_1)
}
M_{state} (if <condition> then <statement<sub>1</sub>> else <statement<sub>2</sub>>, S) =
     S_1 = M_{state} (< condition>, S)
     if M_{boolean} (<condition>, S_1) = true
          return M_{state} (<statement<sub>1</sub>>, S_1)
     else if M_{boolean} (<condition>, S_1) = false
          return M_{state} (<statement<sub>2</sub>>, S_1)
     else
          return 'Error'
}
M_{state} (while <condition> <loop body>, S) =
     S_1 = M_{state} (<condition>, S)
     if M_{boolean} (<condition>, S_1) = true
          return M_{state} (while <condition> <loop body>, M_{state} (<condition>, S_1))
     else if M_{boolean} (<condition>, S_1) = false
          return S_1
     else
          return 'Error'
}
```