

1.  $\langle C \rangle \rightarrow \langle V \rangle = \langle D \rangle \mid \langle V \rangle$   
 $\langle V \rangle \rightarrow x \mid y \mid z$   
 $\langle D \rangle \rightarrow \langle E \rangle ? \langle D \rangle : \langle D \rangle \mid \langle E \rangle$   
 $\langle E \rangle \rightarrow \langle E \rangle \mid \mid \langle F \rangle \mid \langle F \rangle$   
 $\langle F \rangle \rightarrow \langle F \rangle \ \&\& \ \langle G \rangle \mid \langle G \rangle$   
 $\langle G \rangle \rightarrow !\langle G \rangle \mid \langle H \rangle$   
 $\langle H \rangle \rightarrow (\langle H \rangle) \mid \langle I \rangle$   
 $\langle I \rangle \rightarrow \text{true} \mid \text{false}$

2. Static Semantic Attributes:

type	=	{integer, double}	(synthesized)
typetable( $\langle \text{var} \rangle$ )	=	{integer, double, error}	(inherited)
inittable( $\langle \text{var} \rangle$ )	=	{true, false, error}	(inherited)
typebinding	=	( $\langle \text{var} \rangle$ , {integer, double})	(synthesized)
initialized	=	( $\langle \text{var} \rangle$ , {true, false})	(synthesized)

Attribute Rules:

```

<start1> → <stmt3> ; <start3>
<start1>.type := N/A
<start1>.typetable(<var>) := <stmt3>.typetable
<start1>.inittable(<var>) := <stmt3>.initvar
<start1>.typebinding := N/A
<start1>.initialized := N/A

<stmt3>.type := N/A
<stmt3>.typetable := <start1>.typetable
<stmt3>.inittable := <start1>.inittable
<stmt3>.typebinding := N/A
<stmt3>.initialized := N/A
<start3>.type := N/A
<start3>.typetable := <stmt3>.typetable ∪ <start1>.typetable
<start3>.inittable := <stmt3>.inittable ∪ <start1>.inittable
<start3>.typebinding := N/A
<start3>.initialized := N/A
<start2> → <stmt4>
<start2>.type := N/A
<start2>.typetable(<var>) := ∅
<start2>.inittable(<var>) := ∅
<start2>.typebinding := N/A
<start2>.initialized := N/A

<stmt4>.type := N/A
<stmt4>.typetable := <start2>.typetable
<stmt4>.inittable := <start2>.inittable

```

```

<stmt4>.typebinding := N/A
<stmt4>.initialized := N/A
<start3>.type := N/A
<stmt1> → <declare2>
<stmt1>.type := N/A
<stmt1>.typetable(<var>) := ∅(inherited from <start>)
<stmt1>.inittable(<var>) := ∅(inherited from <start>)
<stmt1>.typebinding := N/A
<stmt1>.initialized := N/A

<stmt2> → <assign2>
<stmt2>.type := N/A
<stmt2>.typetable(<var>) := ∅(inherited from <start>)
<stmt2>.inittable(<var>) := ∅(inherited from <start>)
<stmt2>.typebinding := N/A
<stmt2>.initialized := N/A

<declare1>.type :=
<declare1>.typetable(<var>) :=
<declare1>.inittable(<var>) :=
<declare1>.typebinding :=
<declare1>.initialized :=

<type1>.type :=
<type1>.typetable(<var>) :=
<type1>.inittable(<var>) :=
<type1>.typebinding :=
<type1>.initialized :=

<type2>.type :=
<type2>.typetable(<var>) :=
<type2>.inittable(<var>) :=
<type2>.typebinding :=
<type2>.initialized :=

<assign1>.type :=
<assign1>.typetable(<var>) :=
<assign1>.inittable(<var>) :=
<assign1>.typebinding :=
<assign1>.initialized :=

<expression1>.type :=
<expression1>.typetable(<var>) :=
<expression1>.inittable(<var>) :=
<expression1>.typebinding :=
<expression1>.initialized :=

<expression2>.type :=
<expression2>.typetable(<var>) :=
<expression2>.inittable(<var>) :=
<expression2>.typebinding :=
<expression2>.initialized :=

<expression3>.type :=

```

```

<expression3>.typetable(<var>) :=
<expression3>.inittable(<var>) :=
<expression3>.typebinding :=
<expression3>.initialized :=

<expression4>.type :=
<expression4>.typetable(<var>) :=
<expression4>.inittable(<var>) :=
<expression4>.typebinding :=
<expression4>.initialized :=

<expression5>.type :=
<expression5>.typetable(<var>) :=
<expression5>.inittable(<var>) :=
<expression5>.typebinding :=
<expression5>.initialized :=

<value1>.type :=
<value1>.typetable(<var>) :=
<value1>.inittable(<var>) :=
<value1>.typebinding :=
<value1>.initialized :=

<value2>.type :=
<value2>.typetable(<var>) :=
<value2>.inittable(<var>) :=
<value2>.typebinding :=
<value2>.initialized :=

<value3>.type :=
<value3>.typetable(<var>) :=
<value3>.inittable(<var>) :=
<value3>.typebinding :=
<value3>.initialized :=
<value4>.type :=
<value4>.typetable(<var>) :=
<value4>.inittable(<var>) :=
<value4>.typebinding :=
<value4>.initialized :=

```

Table 1: Attribute Rules

3. (a) ‘The type of the expression must match the type of the variable in all assignment statements’

1. <assign<sub>1</sub>>: <var>.type = <expression<sub>3</sub>>.type

- (b) ‘A variable must be declared before it is used’

1. <assign<sub>1</sub>>: <var>.typetable != ‘Error’

- (c) ‘A variable must be assigned a value as its first use in the program’

1. <assign<sub>1</sub>>: if <var>.inittable = ‘Error’

4. Loop Invariants:

**Outer (while) Loop Goal:** The elements  $A[0 \dots n - 1]$  are sorted in non-decreasing order

**Outer (while) Loop Invariant:** The elements  $A[bound \dots n - 1]$  are in non-decreasing order  $\wedge$  the elements  $A[t \dots bound - 1]$  have yet to be sorted.

(The last condition may be redundant but I felt it necessary to include  $t$  in the outer loop invariant since it is initialized outside of the inner loop and also interacts with a variable ( $bound$ ) in the outer loop.)

**Inner (for) Loop Goal:** the elements  $A[t \dots n - 1]$  are sorted in non-decreasing order

**Inner (for) Loop Invariant:** The elements  $A[bound \dots n - 1]$  are in non-decreasing order  $\wedge$   
 $A[0 \dots i] \leq A[t] \wedge$   
 $A[t] \leq A[bound]$ .

**Precondition:**  $n \geq 0$  and  $A$  contains  $n$  elements indexed from 0

```
bound = n;
while (bound > 0) {

    // Assume Outer Loop Invariant is true
    t = 0;

    for (i = 0; i < bound - 1; i++) {

        // Assume Inner Loop Invariant is true
        if (A[i] > A[i+1]) {

            // WP (Inner):
            // A[bound...n-1] are in non-decreasing order  $\wedge$ 
            // A[0...i-1]  $\leq$  A[i]  $\wedge$ 
            // A[i]  $\leq$  A[bound]
            swap = A[i];

            // WP (Inner):
            // A[bound...n-1] are in non-decreasing order  $\wedge$ 
            // A[0...i-1]  $\leq$  A[i+1]  $\wedge$ 
            // A[i+1]  $\leq$  A[bound]
            A[i] = A[i+1];

            // WP (Inner):
            // A[bound...n-1] are in non-decreasing order  $\wedge$ 
            // A[0...i]  $\leq$  swap  $\wedge$ 
            // swap  $\leq$  A[bound]
            A[i+1] = swap;

            // WP (Inner):
            // A[bound...n-1] are in non-decreasing order  $\wedge$ 
            // A[0...i]  $\leq$  A[i+1]  $\wedge$ 
            // A[i+1]  $\leq$  A[bound]
            t = i + 1;
        }
        // (loop termination: i=bound-1, t='the last i+1 for which A[i] > A[i+1]')
        // i=bound-1  $\wedge$  A[t]  $\leq$  A[bound]  $\wedge$  A[0...i]  $\leq$  A[t]  $\wedge$ 
        // A[bound...n-1] are in non-decreasing order  $\rightarrow$ 
        // A[t...n-1] are sorted in non-decreasing order
        // i++
    }

    // WP (Outer):
    // A[bound...n-1] are in non-decreasing order  $\wedge$ 
    // A[t...bound-1] have yet to be sorted
    bound = t;
}
```

```

}
// (loop termination: bound=0, t=0)
// bound=0  ^
// A[0...n-1] are sorted in non-decreasing order  ^
//   A[0...-1] have yet to be sorted (trivially true)  →
//       array A is sorted in non-decreasing order

```

**Postcondition:**  $A[0] \leq A[1] \leq \dots \leq A[n-1]$  (i.e., array  $A$  is sorted in non-decreasing order)

5.  $M_{state}(\langle var \rangle = \langle expression \rangle, S) =$

```

{
    // test that <var> is a legal name in the language
    if M_name(<var>) = 'Error'
        return 'Error'

    // test that <var> has already been declared
    if Lookup(M_name(<var>), S) = 'Error'
        return 'Error'

    // calculate the value of <expression> using the old state
    V = M_value(<expression>, S)
    if V = 'Error'
        return 'Error'

    // calculate a new state including any side effects from evaluating <expression>
    S1 = M_state(<expression>, S)

    // remove <var> from the new state
    Remove(M_name(<var>), S)

    // return the new state with the updated value of <var> added
    return Add(M_name(<var>), V, S1)
}

```

$M_{state}(\text{if } \langle condition \rangle \text{ then } \langle statement_1 \rangle \text{ else } \langle statement_2 \rangle, S) =$

```

{
    S1 = M_state(<condition>, S)

    if M_boolean(<condition>, S1) = true
        return M_state(<statement1>, S1)
    else if M_boolean(<condition>, S1) = false
        return M_state(<statement2>, S1)
    else
        return 'Error'
}

```

$M_{state}(\text{while } \langle condition \rangle \langle loop \text{ body} \rangle, S) =$

```

{
    S1 = M_state(<condition>, S)

    if M_boolean(<condition>, S1) = true
        // evaluate the loop body and call the while-loop again
        return M_state(while <condition> <loop body>, M_state(<loop body>, S1))
    else if M_boolean(<condition>, S1) = false
        return S1
    else
        return 'Error'
}

```

}