

1. $\langle C \rangle \rightarrow \langle V \rangle = \langle D \rangle \mid \langle V \rangle$
 $\langle V \rangle \rightarrow x \mid y \mid z$
 $\langle D \rangle \rightarrow \langle E \rangle ? \langle D \rangle : \langle D \rangle \mid \langle E \rangle$
 $\langle E \rangle \rightarrow \langle E \rangle \mid \mid \langle F \rangle \mid \langle F \rangle$
 $\langle F \rangle \rightarrow \langle F \rangle \&\& \langle G \rangle \mid \langle G \rangle$
 $\langle G \rangle \rightarrow !\langle G \rangle \mid \langle H \rangle$
 $\langle H \rangle \rightarrow (\langle H \rangle) \mid \langle I \rangle$
 $\langle I \rangle \rightarrow \text{true} \mid \text{false}$

2. Static Semantic Attributes:

type	=	{integer, double}	(synthesized)
typetable($\langle \text{var} \rangle$)	=	{integer, double, error}	(inherited)
inittable($\langle \text{var} \rangle$)	=	{true, false, error}	(inherited)
typebinding	=	($\langle \text{var} \rangle$, {integer, double})	(synthesized)
initialized	=	($\langle \text{var} \rangle$, {true, false})	(synthesized)

Attribute Rules:

$\langle \text{start}_1 \rangle \rightarrow \langle \text{stmt}_3 \rangle ; \langle \text{start}_3 \rangle$

$\langle \text{start}_1 \rangle.\text{type} := N/A$

$\langle \text{start}_1 \rangle.\text{typetable}(\langle \text{var} \rangle) := \langle \text{stmt}_3 \rangle.\text{typetable}$

$\langle \text{start}_1 \rangle.\text{inittable}(\langle \text{var} \rangle) := \langle \text{stmt}_3 \rangle.\text{initvar}$

$\langle \text{start}_1 \rangle.\text{typebinding} := N/A$

$\langle \text{start}_1 \rangle.\text{initialized} := N/A$

$\langle \text{stmt}_3 \rangle.\text{type} := N/A$

$\langle \text{stmt}_3 \rangle.\text{typetable} := \langle \text{start}_1 \rangle.\text{typetable}$

$\langle \text{stmt}_3 \rangle.\text{inittable} := \langle \text{start}_1 \rangle.\text{inittable}$

$\langle \text{stmt}_3 \rangle.\text{typebinding} := N/A$

$\langle \text{stmt}_3 \rangle.\text{initialized} := N/A$

$\langle \text{start}_3 \rangle.\text{type} := N/A$

$\langle \text{start}_3 \rangle.\text{typetable} := \langle \text{stmt}_3 \rangle.\text{typetable} \cup \langle \text{start}_1 \rangle.\text{typetable}$

$\langle \text{start}_3 \rangle.\text{inittable} := \langle \text{stmt}_3 \rangle.\text{inittable} \cup \langle \text{start}_1 \rangle.\text{inittable}$

$\langle \text{start}_3 \rangle.\text{typebinding} := N/A$

$\langle \text{start}_3 \rangle.\text{initialized} := N/A$

$\langle \text{start}_2 \rangle \rightarrow \langle \text{stmt}_4 \rangle$

$\langle \text{start}_2 \rangle.\text{type} := N/A$

$\langle \text{start}_2 \rangle.\text{typetable}(\langle \text{var} \rangle) := \emptyset$

$\langle \text{start}_2 \rangle.\text{inittable}(\langle \text{var} \rangle) := \emptyset$

$\langle \text{start}_2 \rangle.\text{typebinding} := N/A$

$\langle \text{start}_2 \rangle.\text{initialized} := N/A$

$\langle \text{stmt}_4 \rangle.\text{type} := N/A$

$\langle \text{stmt}_4 \rangle.\text{typetable} := \langle \text{start}_2 \rangle.\text{typetable}$

$\langle \text{stmt}_4 \rangle.\text{inittable} := \langle \text{start}_2 \rangle.\text{inittable}$
 $\langle \text{stmt}_4 \rangle.\text{typebinding} := N/A$
 $\langle \text{stmt}_4 \rangle.\text{initialized} := N/A$

$\langle \text{stmt}_1 \rangle \rightarrow \langle \text{declare}_2 \rangle$
 $\langle \text{stmt}_1 \rangle.\text{type} := N/A$
 $\langle \text{stmt}_1 \rangle.\text{typetable}(\langle \text{var} \rangle) := \langle \text{start} \rangle.\text{typetable}$
 $\langle \text{stmt}_1 \rangle.\text{inittable}(\langle \text{var} \rangle) := \langle \text{start} \rangle.\text{inittable}$
 $\langle \text{stmt}_1 \rangle.\text{typebinding} := N/A$
 $\langle \text{stmt}_1 \rangle.\text{initialized} := N/A$

$\langle \text{declare}_2 \rangle.\text{type} := N/A$
 $\langle \text{declare}_2 \rangle.\text{typetable}(\langle \text{var} \rangle) := \langle \text{stmt}_1 \rangle.\text{typetable}$
 $\langle \text{declare}_2 \rangle.\text{inittable}(\langle \text{var} \rangle) := \langle \text{stmt}_1 \rangle.\text{inittable}$
 $\langle \text{declare}_2 \rangle.\text{typebinding} := N/A$
 $\langle \text{declare}_2 \rangle.\text{initialized} := N/A$

$\langle \text{stmt}_2 \rangle \rightarrow \langle \text{assign}_2 \rangle$
 $\langle \text{stmt}_2 \rangle.\text{type} := N/A$ (I don't think it makes sense to keep propagating that attribute to this non-terminal (the synthesized attribute type should stop propagating upwards after the non-terminal which uses it to add an entry to the typetable and do error-checking))
 $\langle \text{stmt}_2 \rangle.\text{typetable}(\langle \text{var} \rangle) := \langle \text{start} \rangle.\text{typetable}$
 $\langle \text{stmt}_2 \rangle.\text{inittable}(\langle \text{var} \rangle) := \langle \text{start} \rangle.\text{inittable}$
 $\langle \text{stmt}_2 \rangle.\text{typebinding} := N/A$
 $\langle \text{stmt}_2 \rangle.\text{initialized} := N/A$

$\langle \text{assign}_2 \rangle.\text{type} := (\text{see } \langle \text{assign}_1 \rangle)$
 $\langle \text{assign}_2 \rangle.\text{typetable}(\langle \text{var} \rangle) := \langle \text{stmt}_2 \rangle.\text{typetable}$
 $\langle \text{assign}_2 \rangle.\text{inittable}(\langle \text{var} \rangle) := (\langle \text{stmt}_2 \rangle.\text{inittable} - (M_{\text{name}}(\langle \text{var} \rangle), \text{false})) \cup (M_{\text{name}}(\langle \text{var} \rangle), \text{true})$ (checking would need to be done here to ensure that the entry $(M_{\text{name}}(\langle \text{var} \rangle), \text{false})$ actually exists)
 $\langle \text{assign}_2 \rangle.\text{typebinding} := N/A$
 $\langle \text{assign}_2 \rangle.\text{initialized} := N/A$

$\langle \text{declare}_1 \rangle \rightarrow \langle \text{type}_3 \rangle \langle \text{var} \rangle$
 $\langle \text{declare}_1 \rangle.\text{type} := \langle \text{type}_3 \rangle.\text{type}$
 $\langle \text{declare}_1 \rangle.\text{typetable}(\langle \text{var} \rangle) := \langle \text{stmt} \rangle.\text{typetable} \cup (M_{\text{name}}(\langle \text{var} \rangle), \langle \text{type}_3 \rangle.\text{type})$ (checking would need to be done here to ensure that $\langle \text{var} \rangle$ has not already been declared, i.e., the inittable entry for this variable must be $(M_{\text{name}}(\langle \text{var} \rangle), \text{error})$)
 $\langle \text{declare}_1 \rangle.\text{inittable}(\langle \text{var} \rangle) := \langle \text{stmt} \rangle.\text{inittable} \cup (M_{\text{name}}(\langle \text{var} \rangle), \text{false})$
 $\langle \text{declare}_1 \rangle.\text{typebinding} := (M_{\text{name}}(\langle \text{var} \rangle), \langle \text{type}_3 \rangle.\text{type})$
 $\langle \text{declare}_1 \rangle.\text{initialized} := \langle \text{var} \rangle.\text{initialized}$

$\langle \text{type}_3 \rangle.\text{type} := (\text{see } \langle \text{type}_1 \rangle \text{ and } \langle \text{type}_2 \rangle)$
 $\langle \text{type}_3 \rangle.\text{typetable}(\langle \text{var} \rangle) := \langle \text{declare}_1 \rangle.\text{typetable}$
 $\langle \text{type}_3 \rangle.\text{inittable}(\langle \text{var} \rangle) := \langle \text{declare}_1 \rangle.\text{inittable}$
 $\langle \text{type}_3 \rangle.\text{typebinding} := N/A$
 $\langle \text{type}_3 \rangle.\text{initialized} := N/A$

$\langle \text{type}_1 \rangle \rightarrow \text{int}$
 $\langle \text{type}_1 \rangle.\text{type} := \text{int}$
 $\langle \text{type}_1 \rangle.\text{typetable}(\langle \text{var} \rangle) := \langle \text{declare} \rangle.\text{typetable}$
 $\langle \text{type}_1 \rangle.\text{inittable}(\langle \text{var} \rangle) := \langle \text{declare} \rangle.\text{inittable}$

$\langle \text{type}_1 \rangle.\text{typebinding} := N/A$
 $\langle \text{type}_1 \rangle.\text{initialized} := N/A$

$\langle \text{type}_2 \rangle \rightarrow \text{double}$
 $\langle \text{type}_2 \rangle.\text{type} := \text{double}$
 $\langle \text{type}_2 \rangle.\text{typetable}(\langle \text{var} \rangle) := \langle \text{declare} \rangle.\text{typetable}$
 $\langle \text{type}_2 \rangle.\text{inittable}(\langle \text{var} \rangle) := \langle \text{declare} \rangle.\text{inittable}$
 $\langle \text{type}_2 \rangle.\text{typebinding} := N/A$
 $\langle \text{type}_2 \rangle.\text{initialized} := N/A$

$\langle \text{assign}_1 \rangle \rightarrow \langle \text{var} \rangle \langle \text{expression}_3 \rangle$
 $\langle \text{assign}_1 \rangle.\text{type} := \langle \text{var} \rangle.\text{type}$ (checking would need to be done here to ensure that $\langle \text{var} \rangle$ actually has a type (i.e., it has already been declared), and also to ensure that $\langle \text{var} \rangle.\text{type} = \text{jexpression}_i.\text{type}$)
 $\langle \text{assign}_1 \rangle.\text{typetable}(\langle \text{var} \rangle) := \langle \text{stmt} \rangle.\text{typetable}$
 $\langle \text{assign}_1 \rangle.\text{inittable}(\langle \text{var} \rangle) := \langle \text{stmt} \rangle.\text{inittable}$
 $\langle \text{assign}_1 \rangle.\text{typebinding} := \langle \text{var} \rangle.\text{typebinding}$
 $\langle \text{assign}_1 \rangle.\text{initialized} := (M_{\text{name}}(\langle \text{var} \rangle), \text{true})$

$\langle \text{expression}_3 \rangle.\text{type} := (\text{see } \langle \text{expression}_1 \rangle \text{ and } \langle \text{expression}_2 \rangle)$
 $\langle \text{expression}_3 \rangle.\text{typetable}(\langle \text{var} \rangle) := \langle \text{assign}_1 \rangle.\text{typetable} \cup \langle \text{var} \rangle.\text{typetable}$
 $\langle \text{expression}_3 \rangle.\text{inittable}(\langle \text{var} \rangle) := \langle \text{assign}_1 \rangle.\text{inittable} \cup \langle \text{var} \rangle.\text{inittable}$
 $\langle \text{expression}_3 \rangle.\text{typebinding} := \langle \text{value}_4 \rangle.\text{typebinding}$ (jexpression_i eventually has to produce jvalue_i ; however this attribute only makes sense for jexpression_i if jvalue_i produces $\langle \text{var} \rangle$ (rather than jinteger_i or jfloat_i))
 $\langle \text{expression}_3 \rangle.\text{initialized} := \langle \text{value}_4 \rangle.\text{initialized}$

$\langle \text{expression}_1 \rangle \rightarrow \langle \text{expression}_4 \rangle \langle \text{op} \rangle \langle \text{expression}_5 \rangle$
 $\langle \text{expression}_1 \rangle.\text{type} :=$

```

switch ( $\langle \text{op} \rangle$ ):
  case +:
  case -:
    if  $\langle \text{expression}_4 \rangle.\text{type} = \text{float} \parallel \langle \text{expression}_5 \rangle.\text{type} = \text{float}$ 
       $\langle \text{expression}_1 \rangle.\text{type} = \text{float}$ 
    else
       $\langle \text{expression}_1 \rangle.\text{type} = \text{int}$ 
    break;
  case *:
  case /:
     $\langle \text{expression}_1 \rangle.\text{type} = \text{float}$ 

```

$\langle \text{expression}_1 \rangle.\text{typetable}(\langle \text{var} \rangle) := \langle \text{assign} \rangle.\text{typetable}$
 $\langle \text{expression}_1 \rangle.\text{inittable}(\langle \text{var} \rangle) := \langle \text{assign} \rangle.\text{inittable}$
 $\langle \text{expression}_1 \rangle.\text{typebinding} := N/A$
 $\langle \text{expression}_1 \rangle.\text{initialized} := N/A$

$\langle \text{expression}_4 \rangle.\text{type} := (\text{see } \langle \text{expression}_1 \rangle \text{ and } \langle \text{expression}_2 \rangle)$
 $\langle \text{expression}_4 \rangle.\text{typetable}(\langle \text{var} \rangle) := \langle \text{expression}_1 \rangle.\text{typetable}$
 $\langle \text{expression}_4 \rangle.\text{inittable}(\langle \text{var} \rangle) := \langle \text{expression}_1 \rangle.\text{inittable}$
 $\langle \text{expression}_4 \rangle.\text{typebinding} := (\text{see } \langle \text{expression}_2 \rangle \text{ and comment for } \langle \text{expression}_3 \rangle)$
 $\langle \text{expression}_4 \rangle.\text{initialized} := (\text{see } \langle \text{expression}_2 \rangle \text{ and comment for } \langle \text{expression}_3 \rangle)$

$\langle \text{expression}_5 \rangle.\text{type} := (\text{see } \langle \text{expression}_1 \rangle \text{ and } \langle \text{expression}_2 \rangle)$
 $\langle \text{expression}_5 \rangle.\text{typetable}(\langle \text{var} \rangle) := \langle \text{expression}_1 \rangle.\text{typetable} \cup \langle \text{op} \rangle.\text{typetable}$

```

<expression5>.inittable(<var>) := <expression1>.typetable ∪ <op>.inittable
<expression5>.typebinding := (see <expression2> and comment for <expression3>)
<expression5>.initialized := (see <expression2> and comment for <expression3>)

<expression2> → <value4>
<expression2>.type := <value4>.type
<expression2>.typetable :=

if produced by <assign1>:
    <expression2>.typetable := <assign1>.typetable \union \nterm{var}.typetable
elif produced by <expression1>:
    if <expression2> is <expression4>:
        <expression2>.typetable := <expression1>.typetable
    elif <expression2> is <expression5>:
        <expression2>.typetable := <expression1>.typetable \union <op>.typetable

<expression2>.inittable := (same rules as typetable)
<expression2>.typebinding := <value4>.typebinding
<expression2>.initialized := <value4>.initialized

<value4>.type := (see <value1>, <value2>, and <value3>)
<value4>.typetable(<var>) := <expression3>.typetable
<value4>.inittable(<var>) := <expression3>.inittable
<value4>.typebinding := (see <value1>)
<value4>.initialized := (see <value1>)

<value1> → <var>
<value1>.type := <var>.type
<value1>.typetable(<var>) := <expression>.typetable
<value1>.inittable(<var>) := <expression>.inittable
<value1>.typebinding := <var>.typebinding
<value1>.initialized := <var>.initialized

<value2> → <integer>
<value2>.type := <integer>.type
<value2>.typetable(<var>) := <expression>.typetable
<value2>.inittable(<var>) := <expression>.inittable
<value2>.typebinding := N/A
<value2>.initialized := N/A

<value3> → <float>
<value3>.type := <float>.type
<value3>.typetable(<var>) := <expression>.typetable
<value3>.inittable(<var>) := <expression>.inittable
<value3>.typebinding := N/A
<value3>.initialized := N/A

```

3. (a) ‘The type of the expression must match the type of the variable in all assignment statements’

1. <assign₁>: <var>.type = <expression₃>.type
- (b) ‘A variable must be declared before it is used’

1. $\langle \text{assign}_1 \rangle$: $\langle \text{var} \rangle.\text{typetable} \neq \text{'Error'}$
- (c) 'A variable must be assigned a value as its first use in the program'

1. $\langle \text{assign}_1 \rangle$: if $\langle \text{var} \rangle.\text{inittable} = \text{'Error'}$

4. Loop Invariants:

Outer (while) Loop Goal: The elements $A[0 \dots n - 1]$ are sorted in non-decreasing order

Outer (while) Loop Invariant: The elements $A[\text{bound} \dots n - 1]$ are in non-decreasing order \wedge the elements $A[t \dots \text{bound} - 1]$ have yet to be sorted.

(The last condition may be redundant but I felt it necessary to include t in the outer loop invariant since it is initialized outside of the inner loop and also interacts with a variable (bound) in the outer loop.)

Inner (for) Loop Goal: the elements $A[t \dots n - 1]$ are sorted in non-decreasing order

Inner (for) Loop Invariant: The elements $A[\text{bound} \dots n - 1]$ are in non-decreasing order \wedge
 $A[0 \dots i] \leq A[t] \wedge$
 $A[t] \leq A[\text{bound}]$.

Precondition: $n \geq 0$ and A contains n elements indexed from 0

```

bound = n;
while (bound > 0) {

    // Assume Outer Loop Invariant is true
    t = 0;

    for (i = 0; i < bound - 1; i++) {

        // Assume Inner Loop Invariant is true
        if (A[i] > A[i+1]) {

            // WP (Inner):
            // A[bound...n-1] are in non-decreasing order  $\wedge$ 
            // A[0...i-1]  $\leq$  A[i]  $\wedge$ 
            // A[i]  $\leq$  A[bound]
            swap = A[i];

            // WP (Inner):
            // A[bound...n-1] are in non-decreasing order  $\wedge$ 
            // A[0...i-1]  $\leq$  A[i+1]  $\wedge$ 
            // A[i+1]  $\leq$  A[bound]
            A[i] = A[i+1];

            // WP (Inner):
            // A[bound...n-1] are in non-decreasing order  $\wedge$ 
            // A[0...i]  $\leq$  swap  $\wedge$ 
            // swap  $\leq$  A[bound]
            A[i+1] = swap;

            // WP (Inner):
            // A[bound...n-1] are in non-decreasing order  $\wedge$ 
            // A[0...i]  $\leq$  A[i+1]  $\wedge$ 
            // A[i+1]  $\leq$  A[bound]
            t = i + 1;
        }
    }
    // (loop termination: i=bound-1, t='the last i+1 for which A[i] > A[i+1] ')
    // i=bound-1  $\wedge$  A[t]  $\leq$  A[bound]  $\wedge$  A[0...i]  $\leq$  A[t]  $\wedge$ 

```

```

    //   A[bound...n-1] are in non-decreasing order →
    //   A[t...n-1] are sorted in non-decreasing order
    // i++
}
// WP (Outer):
// A[bound...n-1] are in non-decreasing order ∧
// A[t...bound-1] have yet to be sorted
bound = t;
}
// (loop termination: bound=0, t=0)
// bound=0 ∧
// A[0...n-1] are sorted in non-decreasing order ∧
// A[0...-1] have yet to be sorted (trivially true) →
//   array A is sorted in non-decreasing order

```

Postcondition: $A[0] \leq A[1] \leq \dots \leq A[n-1]$ (i.e., array A is sorted in non-decreasing order)

5. $M_{state}(\langle \text{var} \rangle = \langle \text{expression} \rangle, S) =$

```

{
    // test that <var> is a legal name in the language
    if  $M_{name}(\langle \text{var} \rangle) = \text{'Error'}$ 
        return 'Error'

    // test that <var> has already been declared
    if  $\text{Lookup}(M_{name}(\langle \text{var} \rangle), S) = \text{'Error'}$ 
        return 'Error'

    // calculate the value of <expression> using the old state
     $V = M_{value}(\langle \text{expression} \rangle, S)$ 
    if  $V = \text{'Error'}$ 
        return 'Error'

    // calculate a new state including any side effects from evaluating <expression>
     $S_1 = M_{state}(\langle \text{expression} \rangle, S)$ 

    // remove <var> from the new state
     $\text{Remove}(M_{name}(\langle \text{var} \rangle), S)$ 

    // return the new state with the updated value of <var> added
    return  $\text{Add}(M_{name}(\langle \text{var} \rangle), V, S_1)$ 
}

```

$M_{state}(\text{if } \langle \text{condition} \rangle \text{ then } \langle \text{statement}_1 \rangle \text{ else } \langle \text{statement}_2 \rangle, S) =$

```

{
     $S_1 = M_{state}(\langle \text{condition} \rangle, S)$ 

    if  $M_{boolean}(\langle \text{condition} \rangle, S_1) = \text{true}$ 
        return  $M_{state}(\langle \text{statement}_1 \rangle, S_1)$ 
    else if  $M_{boolean}(\langle \text{condition} \rangle, S_1) = \text{false}$ 
        return  $M_{state}(\langle \text{statement}_2 \rangle, S_1)$ 
    else
        return 'Error'
}

```

$M_{state}(\text{while } \langle \text{condition} \rangle \langle \text{loop body} \rangle, S) =$

```

{
     $S_1 = M_{state}(\langle \text{condition} \rangle, S)$ 

```

```

if  $M_{boolean}(\langle condition \rangle, S_1) = true$ 
    // evaluate the loop body and call the while-loop again
    return  $M_{state}(while \langle condition \rangle \langle loop \ body \rangle, M_{state}(\langle loop \ body \rangle, S_1))$ 
else if  $M_{boolean}(\langle condition \rangle, S_1) = false$ 
    return  $S_1$ 
else
    return 'Error'
}

```