Rebecca Frederick EECS 345 Written Exercise 1 February 4, 2016

```
<C>
        \rightarrow <V> = <D> | <V>
<V>
        \rightarrow x | y | z
        \rightarrow <E> ? <D> : <D> | <E>
<D>
<E>

ightarrow <E> || <F> | <F>
        \rightarrow
             <F> && <G> <G>
<F>
             !<G> | <H>
<G>
        \rightarrow
              (<H>) | <I>
<H>
              true | false
<I>
```

## 2. Static Semantic Attributes:

## Attribute Rules:

```
\langle start_1 \rangle \rightarrow \langle stmt_3 \rangle; \langle start_3 \rangle
<start<sub>1</sub>>.type := N/A
<start1>.typetable(<var>) := <stmt3>.typetable
<start1>.inittable(<var>) := <stmt3>.initvar
{\rm start_1} > . {\rm typebinding} := N/A
{\rm start_1}>.initialized := N/A
<stmt<sub>3</sub>>.type := N/A
<stmt3>.typetable := <start1>.typetable
{\rm <stmt_3>.inittable} := {\rm <start_1>.inittable}
{\rm <stmt_3>.typebinding} := N/A
{\rm <stmt_3>.initialized} := N/A
<start<sub>3</sub>>.type := N/A
{\sf start_3>}.typetable := {\sf stmt_3>}.typetable \cup {\sf start_1>}.typetable
{\sf start_3>.inittable} := {\sf stmt_3>.inittable} \cup {\sf start_1>.inittable}
{\rm start_3}.typebinding := N/A
{\rm start_3}.initialized := N/A
\langle start_2 \rangle \rightarrow \langle stmt_4 \rangle
<start<sub>2</sub>>.type := N/A
<start2>.typetable(<var>) := 0
<start2>.inittable(<var>) := Ø
{\rm start_2}.typebinding := N/A
{\rm <start_2>.initialized} := N/A
<stmt<sub>4</sub>>.type := N/A
<stmt<sub>4</sub>>.typetable := <start<sub>2</sub>>.typetable
{\rm <stmt_4>.inittable} := {\rm <start_2>.inittable}
```

```
<stmt<sub>4</sub>>.typebinding := N/A
{\rm <stmt_4>.initialized} := N/A
{\rm start_3}.type := N/A
\langle stmt_1 
angle \; 	o \; \langle declare_2 
angle
<stmt<sub>1</sub>>.type := N/A
<stmt<sub>1</sub>>.typetable(<var>) := Ø(inherited from <start>)
<stmt₁>.inittable(<var>) := ∅(inherited from <start>)
{\rm <stmt_1>.typebinding} := N/A
{\rm <stmt_1>.initialized} := N/A
\langle stmt_2 
angle \ 	o \ \langle assign_2 
angle
<stmt_2>.type := N/A
<stmt<sub>2</sub>>.typetable(<var>) := \emptyset(inherited from <start>)
<stmt<sub>2</sub>>.inittable(<var>) := \emptyset(inherited from <start>)
{\rm <stmt_2>.typebinding} := N/A
{\rm <stmt_2>.initialized} := N/A
<declare1>.type :=
<declare<sub>1</sub>>.typetable(<var>) :=
<declare<sub>1</sub>>.inittable(<var>) :=
<declare<sub>1</sub>>.typebinding :=
<declare<sub>1</sub>>.initialized :=
<type1>.type :=
<type<sub>1</sub>>.typetable(<var>) :=
<type<sub>1</sub>>.inittable(<var>) :=
<type1>.typebinding :=
<type1>.initialized :=
<type<sub>2</sub>>.type :=
<type<sub>2</sub>>.typetable(<var>) :=
<type<sub>2</sub>>.inittable(<var>) :=
<type2>.typebinding :=
<type2>.initialized :=
\langle assign_1 \rangle .type :=
<assign<sub>1</sub>>.typetable(<var>) :=
<assign<sub>1</sub>>.inittable(<var>) :=
\langle assign_1 \rangle.typebinding :=
\langle assign_1 \rangle.initialized :=
<expression_1>.type :=
<expression1>.typetable(<var>) :=
<expression<sub>1</sub>>.inittable(<var>) :=
<expression<sub>1</sub>>.typebinding :=
<expression<sub>1</sub>>.initialized :=
<expression_2>.type :=
<expression<sub>2</sub>>.typetable(<var>) :=
<expression2>.inittable(<var>) :=
<expression_2>.typebinding :=
<expression_2>.initialized :=
<expression<sub>3</sub>>.type :=
```

```
<expression3>.typetable(<var>) :=
<expression<sub>3</sub>>.inittable(<var>) :=
<expression<sub>3</sub>>.typebinding :=
<expression<sub>3</sub>>.initialized :=
<expression_4>.type :=
<expression<sub>4</sub>>.typetable(<var>) :=
<expression<sub>4</sub>>.inittable(<var>) :=
<expression<sub>4</sub>>.typebinding :=
<expression<sub>4</sub>>.initialized :=
<expression<sub>5</sub>>.type :=
<expression<sub>5</sub>>.typetable(<var>) :=
<expression<sub>5</sub>>.inittable(<var>) :=
<expression_5>.typebinding :=
<expression5>.initialized :=
<value<sub>1</sub>>.type :=
<value<sub>1</sub>>.typetable(<var>) :=
<value<sub>1</sub>>.inittable(<var>) :=
<value<sub>1</sub>>.typebinding :=
<value<sub>1</sub>>.initialized :=
<value2>.type :=
<value<sub>2</sub>>.typetable(<var>) :=
<value<sub>2</sub>>.inittable(<var>) :=
<value<sub>2</sub>>.typebinding :=
<value<sub>2</sub>>.initialized :=
<value<sub>3</sub>>.type :=
<value<sub>3</sub>>.typetable(<var>) :=
<value3>.inittable(<var>) :=
<value3>.typebinding :=
<value3>.initialized :=
<value<sub>4</sub>>.type :=
<value<sub>4</sub>>.typetable(<var>) :=
<value<sub>4</sub>>.inittable(<var>) :=
<value<sub>4</sub>>.typebinding :=
<value<sub>4</sub>>.initialized :=
```

Table 1: Attribute Rules

- 3. (a) 'The type of the expression must match the type of the variable in all assignment statements'
  - 1. <assign<sub>1</sub>>: <var>.type = <expression<sub>3</sub>>.type
  - (b) 'A variable must be declared before it is used'
    - 1. <assign<sub>1</sub>>: <var>.typetable != 'Error'
  - (c) 'A variable must be assigned a value as its first use in the program'
    - 1. <assign<sub>1</sub>>: if <var>.inittable = 'Error'
- 4. Loop Invariants:

Outer (while) Loop Goal: The elements A[0...n-1] are sorted in non-decreasing order

Outer (while) Loop Invariant: The elements A[bound ... n-1] are in non-decreasing order  $\land$  the elements A[t...bound-1] have yet to be sorted.

(The last condition may be redundant but I felt it necessary to include t in the outer loop invariant since it is initiallized outside of the inner loop and also interacts with a variable (bound) in the outer loop.)

Inner (for) Loop Goal: the elements  $A[t \dots n-1]$  are sorted in non-decreasing order

**Inner (for) Loop Invariant:** The elements A[bound...n-1] are in non-decreasing order  $\land$   $A[0...i] \le A[t] \land$   $A[t] \le A[bound].$ 

**Precondition:**  $n \geq 0$  and A contains n elements indexed from 0

```
bound = n;
while (bound > 0) {
          // Assume Outer Loop Invariant is true
          t = 0;
          for (i = 0; i < bound - 1; i++) {
                    // Assume Inner Loop Invariant is true
                   if (A[i] > A[i+1]) {
                              // WP (Inner):
                              // A[bound...n-1] are in non-decreasing order \land
                              // A[0 \dots i-1] \leq A[i] \wedge
                              // A[i] \leq A[bound]
                              swap = A[i];
                              // WP (Inner):
                              // A[bound...n-1] are in non-decreasing order \land
                              // A[0 \dots i-1] \le A[i+1] \land
                              // A[i+1] \leq A[bound]
                              A[i] = A[i+1];
                              // WP (Inner):
                              // A[bound \dots n-1] are in non-decreasing order \wedge
                              //A[0...i] \leq swap \wedge
                              // swap \leq A[bound]
                              A[i+1] = swap;
                              // WP (Inner):
                              // A[bound...n-1] are in non-decreasing order \land
                              // A[0 \dots i] \leq A[i+1] \wedge
                              //\ A[i+1] \leq A[bound]
                             t = i + 1;
                   // (loop termination: i=bound-1, t='the last i+1 for which A[i]>A[i+1]')
                    \textit{//} \quad \textit{i=bound-1} \quad \land \quad A[t] \leq A[bound] \quad \land \quad A[0 \ldots i] \leq A[t] \quad \land \quad A[t] \leq A[t] \quad A[t] \quad \land \quad A[t] \leq A[t] \quad \land \quad A[t] \quad
                                                A[bound...n-1] are in non-decreasing order
                   //
                                                                     A[t \dots n-1] are sorted in non-decreasing order
                   // i++
          }
          // WP (Outer):
          // A[bound...n-1] are in non-decreasing order \land
          // A[t...bound-1] have yet to be sorted
          bound = t;
```

```
// (loop termination: bound=0, t=0)
   // bound=0 \wedge
   // A[0\ldots n-1] are sorted in non-decreasing order \wedge
           A[0\cdots -1] have yet to be sorted (trivially true) 
ightarrow
   //
                array A is sorted in non-decreasing order
   Postcondition: A[0] \le A[1] \le \cdots \le A[n-1] (i.e., array A is sorted in non-decreasing order)
5. M_{state}(\langle var \rangle = \langle expression \rangle, S) =
        // test that <var> is a legal name in the language
        if M_{name}(\langle var \rangle) = 'Error'
              return 'Error'
        // test that <var> has already been declared
        if Lookup(M_{name}(\langle var \rangle), S) = 'Error'
              return 'Error'
        // calculate the value of <expression> using the old state
        V = M_{value} (\langle expression \rangle, S)
        if V = 'Error'
             return 'Error'
        // calculate a new state including any side effects from evaluating <expression>
        S_1 = M_{state} (\langle expression \rangle, S)
        // remove <var> from the new state
        Remove(M_{name}(\langle var \rangle), S)
        // return the new state with the updated value of <var> added
        return Add(M_{name}(\langle var \rangle), V, S_1)
   }
   M_{state} (if <condition> then <statement<sub>1</sub>> else <statement<sub>2</sub>>, S) =
        S_1 = M_{state} (<condition>, S)
        if M_{boolean} (<condition>, S_1) = true
             return M_{state} (<statement<sub>1</sub>>, S_1)
        else if M_{boolean} (<condition>, S_1) = false
             return M_{state} (<statement<sub>2</sub>>, S_1)
        else
             return 'Error'
   }
   M_{state} (while <condition> <loop body>, S) =
        S_1 = M_{state} (\langle condition \rangle, S)
        if M_{boolean} (<condition>, S_1) = true
              // evaluate the loop body and call the while-loop again
             return M_{state} (while <condition> <loop body>, M_{state}(<loop body>, S_1))
        else if M_{boolean} (<condition>, S_1) = false
             return S_1
        else
             return 'Error'
```

}