CS35L

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Terms

- Operating System: most important software, links hardware with other processes
- multi-user os: allows many users to work on a single system at the same time
- multi-process os: running different tasks/processes at the same time
- command line interface vs. graphical user interface:
 - cli has a steeper learning curve, allows pure control, is faster, uses less resources, shows less graphical changes
- **kernel**: core of any OS, handles allocation of memory / interface of applications with hardware
 - eg. Linux is a kernel, Debian GNU/Linux is a clone of UNIX
 - controls access to system resources: memory, I/O, CPU
 - lowest layer above CPU, provides protections and fair allocations
- shell is the outermost layer of the kernel, interface between kernel and user
 - eg. command prompt in Windows, shell in UNIX, terminal in Mac
- the Linux file system is *hierarchal*
 - usually with single root, but allows for multiple home directories (multiuser)
- sudo is administrator, ie. super-user
- everything in the file system is either a process or a file

Introduction CS35L

- processes are identified by a PID, files are a collection of data
- file permissions:
 - 1s -1 (list) prints out the following attributes, from left to right:
 - * file type, user permissions, group permissions, other permissions, number of hard links, owner name, group name, size, date last modified
 - * d for directory, for file, l for link
 - eg. drwxr-xr-x 1 someUser someGroup
 - different permissions: read, write, executable, denied
 - can also be written in octal:
 - * 3 digits: owner, group, other
 - * values: execute 1, write 2, read 4
 - · eg. 721: full access for user, write only for group, execute only for other
 - special permissions: setuid, set user ID on execution
 - * s flag replaces x flag for the user, users act as the owner
 - * eg. passwd command, chmod u+s file
 - setgid grants permssion of the group owner
 - * s flag replaces x flag for the group, chmod g+s file
- chmod: change mode/permissions
 - classes: u user, g group, o other, a all
 - * adds modes, removes modes, = sets modes
 - eg. chmod u-w file removes write permission from the user

Character Sets and Encodings

- ASCII is both a character set and an encoding
 - set of 128 characters (128-255 not used), fits within a single byte
- *Unicode* is a character set with over 1,000,000 code points
- *UTF-8* is an encoding with a variable length between 1 and 4
 - ASCII is incorporated within UTF-8

Environment Variables

- variables that can be accessed from any process
 - dynamic variables
 - eg. \$HOME to home directory and \$PATH to search for commands
 - change with export VARIABLE=...
- locale is a set of parameters that defines a user's cultural preferences
 - default locale is C

Basic Shell Commands CS35L

- eg. language, country, etc.
- parameteres can be checked with locale command
- different \$LC_* environment variables within the locale
 - * eg. LC_TIME, LC_COLLATE (differs, eg. US sort is case insensitive vs C sort)

Basic Shell Commands

• !! replace with previous command, !str refer to previous command with str

• redirection:

- > file writes stdout to a file
- >> file appends stdout to a file
- < file use contents of file as stdin</p>
- cd: change directory, mkdir: make directory, rm: removes files, use -r flag for directories or rmdir on empty directories
- cp: copy a file (src, dest), mv: moves a file or just renames it (src, dest)
- pwd prints working directory
- ps: processes, du: memory usage
- find: (directory) -type -perm -name -user -maxdepth
 - -mtime for recently modified files
 - -type d for directory, f for file, l for symbolic link
 - eg. find . -name my* (. searches current directory)
 - can use regex and linux wildcards
- · whatis, whereis locates binary, source, and man, which locates binary
- diff: print the line differences between files, -u unified format, -p check C functions, -r recursive
- cmp: print the byte differences between files
- wget and curl: download files
- man: open manual for some command
 - -k searches for relevant manuals, apropros
- ls: print directory:
 - a all files, -l list line-by-line, -t sort by modification time, -r reverse order,
 -i prints inode or memory references to files
 - -d list only directories, -s show size
 - can combine flags together, eg. 1s -altr
- 1n: link files (original, link)
 - when making a hardlink, removing the original file does not affect the hardlinked file
 - * even though they share the *same* inode/contents

Basic Shell Commands CS35L

- * cannot link to a directory
- * permissions AND contents are updated
- * mirror link
- use -s flag for soft or symbolic links that become dangling links when the original file is removed
 - * the inode is *not* shared, can link to directories
 - * permissions are not updated
 - * actual link
- touch: update access and modification time to current time
 - -t set time, eg. 201101311759.30
- echo: prints text
 - can use redirection operators:
 - * echo TEXT > file.txt replaces text in file with TEXT, or creates new file
 - * echo TEXT >> file.txt appends TEXT to file, or creates new file
- head: prints first 10 lines of some file (analagous to tail)
 - can use with pipe operator:
 - * cat file.txt | head -n 2 prints the first two lines of the file
- wc: word count, -l for lines, -m for characters, -c for byte counts, -w for words
- cmp: compare
- cat: output contents of a file
- ssh: connect to server, !ssh calls most recent ssh command
 - scp: secure copy of files between servers
- tar: archive files
 - -x extracts, -v verbose, -f file, -c creates, -transform changes file names
- gzip: compress or expand files
 - -1 fast compression, -9 slowest compression
- du estimates file space usage, ps reports current processes, kill takes PID and kills process

Text Editing Commands

- sort sorts lines of text files, sort order depends on locale
 - f ignores case, -n numeric sort, -M month sort, -r reverse, -u unique, -o output
 - defaults to ASCII sort in C locale
- comm compares two sorted files line by line
 - outputs three columns: words present in first file, second file, and both
 - * supress with -1, -2, -3
 - may have to suppress to check for repeats
- tr translate or delete characters

- usage: tr [OPTION] ... SET1 [SET2]
- d flag deletes characters, -s squeezes repeats, -c uses the complement of SET1
- eg. tr '12' 'ab', tr -d [:digit:]
 - * other special sets: [:lower:], [:upper:]
 - · alpha, blank, space
- sed stream editor command
 - modifies the input as specified
 - sed -n '1p', sed -n '1,10p', sed -n '1~2p' prints specific lines
 - * -n suppresses default print of all lines, -i edits in place instead of stream
 - can also delete text: sed '1~2d', sed '/pattern/d', sed '/pattern/!d',
 - * !d deletes inverse
 - modifying text: sed 's/cat/dog/', sed 's/cat/dog/g, sed 's/<[^>]*>//g'
 - g flag global, number replaces only the numberth match of every line
 - sed -r, -E evaluates ? or +
- grep command to search files or text for the occurence that match a pattern
 - grep -r '*.txt', -c gives the count, -n gives the line number, -r recursively,
 i ignore case, -w whole word boundaries
 - -e specify patterns
 - -l and -L suppresses normal outputs, -v inverts sense of matching
 - E or egrep to use extended regex
- awk is a programing language by itself, utility/language for data extraction
 - views text as fields
 - * follows pattern awk '/pattern/ {actions}' file
 - awk '/text/ {print;}' file.txt prints lines which match text
 - awk '{print \$1,\$2;}' file.txt prints specific fields
 - awk -F 'delim' '{prints \$2}' file.txt prints second column between specified pattern
 - built in variables:
 - * NR count of input records, NF number of fileds, FS field separator (default whitespace), RS record separator

Regular Expressions

- regular expressions allow searching for a specific pattern
- UNIX wildcards: characters that can stand for all members of some class of characters

- the * wildcard matches zero or more characters in a file or directory name
- the ? wildcard will match exactly one character
- the [] represents a single character matching any of the characters enclosed between them, eg. [∅-9]
 - * used in combination with another expression
- regex rules are slightly different...
- quantification:
 - how many times of the previous expression...
 - operators: ? (0 or 1), * (0 or more), + (1 or more)
 - eg. n[0-9]a matches n0a, n1a, n9a, but not na, nxa, or n10a
 - * but n[0-9]*a also matches na and n10a
 - {n} exactly n times
 - {n,m} from n to m times
 - {n,} n or more times
 - quantifiers are greedy by default
 - following a quantifier with a ? makes them lazy
 - * eg. \d+? matches 1 in 12345, A*? matches empty in AAA, \w{2,4}? matches ab in abcd
- alternation:
 - which choices...
 - operators: [] (without hyphens) and | (or operator)
 - eg. Hello | world, [A B C]
- anchors:
 - where...
 - characters: ^ beginning, \$ end, \b word boundary
- [] match any enclosed characters, [^] match complement of enclosed character, . matches a single character of any value
 - inside character classes, only metacharacters are:] \ ^ -
 - * all other special characters do not have to be escaped
 - * if placed in class correctly, these do not have to be escaped either
- use \ to escape character
- regex matches whitespace as well
- parentheses: allow quantifiers to sequences of characters
 - $\text{ eg. } (a[0-9]z)^* \text{ vs. } a[0-9]z^*$
- basic regular expression vs. extended regular expression
 - BRE takes things more literally, characters such as ? or + lose their meaning
 - * BRE only recognizes the metacharacters ^ \$. [] *
 - -E means extended regular expressions in grep command
 - \star ERE adds the other metacharacters () { } ? + |
- character classes:

Shell Scripting CS35L

- − \d digits, \D non-digits
- \w word character, \\ non-word character
- \s whitespace character, \S non-whitespace character
- capturing groups/contents:
 - $\text{ eg. } (\d\d) + (\d\d) = \2 + 1 \text{ matches } 112+65=65+112$
 - non-capturing group: (?:...) vs. (...)
- positive look-ahead and look-behind:
 - eg. d(?=r) matches d if it is followed by r, but r will not be part of match
 - eg. (?<=r)d matches d if it is preceded by r, but r will not be part of match
- negative look-ahead and look-behind:
 - eg. q(?!)u matches q not followed by a u
 - eg. (?<!a)b matches b not preceded by a

Shell Scripting

- in a compiled language, the code must be compiled and is not translated to machine code if there is an error
 - runs faster because it has been compiled
 - eg. C, C++, Java
- in a scripting language, there is no compilation required, it is directly interpreted/translated
 - slower
 - eg. Python, JavaScript, Shell Scripting
 - can be used for automation, less code intensive actions
- shell script: a program designed to be run on a shell
 - all shell commands can be executed inside a script
 - simple, portable, no compilation

Syntax

- shell recognizes built-in commands (eg. echo), shell functions, and external commands
 - eg. echo \$myvar vs. number=`ls | wc -l`
- first line of shell script must be a shebang to indicate what kind of script it is
 - when the shell runs a program, it asks the kernel to start a new process and run the given program in that process
 - #!/bin/sh, need to tell kernel which shell to use (eg. csh, awk, sh)
- variables in shell script can start with letter or underscore

Shell Scripting CS35L

 typical var rules, start with letter or underscore, followed by letters, digits, or underscores

- cannot have spaces between the variable declaration and name
- eg. myvar=text, echo \$text, fullname="\$first \$middle \$last"
 - * multiple assignments in same line allowed
 - * need double quotes when defining whitespace
 - * single quotes take the *literal* value of what is enclosed

```
• eg. myvar='$anothervar'
```

- * usual escape character rules
- can use backticks to set a variable equal to a command

```
- eg. myvar=`ls -l`
```

- special variables: \$ PID, # num arguments, n nth argument, ? exit status of last account
- array_name[index]=value, echo \${array_name[index]}
- can redirect files using
 - redirecting stderr to a file: 2>errorFile
 - redirecting stderr to stdout: 2>&1
 - redirecting both to file: &>file
- conditionals:
 - [[-z STRING]] empty string
 - [[-n STRING]] not empty string
 - [[STRING == STRING]] equality
 - [[STRING != STRING]] non-equality
 - [[STRING =~ STRING]] regex
 - [[NUM -eq NUM]] equality
 - [[NUM -ne NUM]] non-equality
 - [[NUM -1t NUM]] less than
 - [[NUM -le NUM]] less than or equal
 - same for greater than...
- file conditionals: (for current directory)
 - [[-e FILE]] exists
 - [[-r FILE]] readable
 - [[-h FILE]] symlink
 - [[-d FILE]] directory
 - [[-w FILE]] writable
 - [[-s FILE]] size greater than 0
 - [[-f FILE]] file
 - [[-x FILE]] executable
 - [[FILE1 -nt FILE2]] newer than
 - [[FILE1 -ot FILE2]] older than

Shell Scripting CS35L

```
- [[ FILE1 -ef FILE2 ]] equal files
```

• can pass arguments to a script:

```
#!/bin/bash
echo $1 $2 $3

args=("$0") # array of arguments
echo ${args[0]}}

echo $# # number of arguments
echo $0 # base arg

echo $0 # all args
echo $* # all args, starting from first
```

• loops in a script: (similar syntax for while loops)

Software CS35L

done

• conditionals:

```
if [ $a == $b ]
then
  echo "a is equal to b"
elif [ $a -gt $b ]
then
  echo "a is >= b"
elif [ $a -lt $b ]
  echo "a is <= b"
else
  echo "no conditons met"
fi
FRUIT="kiwi"
case $FRUIT in
  "apple") echo "apple pie"
  "kiwi") echo "kiwis"
  ;;
esac
#unconditionals
break
continue
```

Software

- installing software is different for different OS's
- for Linux
 - can consolidate installation with configure, make, make install
 - software usually in the *tarball* format (.tgz or .gz)

Software CS35L

 other sources: rpm (Redhat Package Management), apt-get (Advanced Package Tool)

- to decompress, use tar
 - eg. tar -xzvf filename.tar.gz
 - options:
 - * -x extract, -v verbose, -f file
 - * -j use bzip2, -J use xz, -z use gzip or gunzip, -Z use compress
- compilation process (eg. for C++):
 - 1. **preprocessor** deals with preprocessor directives
 - removes comments
 - includes some header files
 - source code -> expanded source code (temporary file, can be printed on stdout)
 - 2. **compiler** creates the assembly level language
 - expanded source code -> assembler file (.s)
 - 3. assembler translates assembly into machine-readable code
 - assembler file -> object code file (.o)
 - 4. linker links all object and library files (specifically, their headers) together
 - object code file -> executable file
 - eg. g++ -Wall shoppingList.cpp item.cpp shop.cpp -o shop
 - -Wall enables all warnings
- if only a single header or source file is modified out of hundreds in a large project
 - not efficient to completely recompile every file (some files not dependent on the single modified file)
 - instead, have separate object code files for every source file, can link them all together faster
 - c creates object files, -o flag links object files
 - also difficult to keep track of which files to recompile when the project is large
 - * can use make

Make

- utility for managing large software projects
- · helps keeps files compiled and up to date
- is a shell script with:
 - targets, dependencies, and actions
 - will try to create dependencies using other rules
 - * will only recompile files whose timestamps have changed, efficient compilation

Software CS35L

- syntax ex. for a single rule: target : dependencies action
- by default, first target is run if target not specified (eg. make)
- by default, will output executing lines (can suppress by prepending 1 to line)
- have to escape \$ with \$\$
 - special variables:
 - * \$1 current target in rule
 - * \$< first input file in rule
 - * \$^ all input files (no duplicates)
 - * \$? all input files newer than the target
 - * \$* stem part matched from % in rule definition
 - * can grab directory and file: eg. \$(aD) and \$(aF)
- use .PHONY to specify non-file rules

```
- eg. .PHONY : all clean check
```

- can embed shell scripts
- each line is ran in different shell, so must escape newline for embedded scripts

```
all: shop
shop : item.o shoppingList.o shop.o
  g++ -g -Wall -o shop item.o shoppingList.o shop.o
item.o : item.cpp item.h
  g++ -g -Wall -c item.cpp
clean:
  rm -f item.o shoppingList.o shop.o shop
#using variables:
OPTIMIZE = -02
CC = gcc
CFLAGS = $(OPTIMIZE) -g3 -Wall -Wextra -march=native -mtune=native -mrdrnd
randall: randall.c
    $(CC) $(CFLAGS) randall.c -o $0
#using script:
tests = 1-test.ppm 2-test.ppm 4-test.ppm 8-test.ppm
check: baseline.ppm $(tests)
```

Python CS35L

```
for file in $(tests); do \
    diff -u baseline.ppm $$file || exit; \
    done
$(tests): srt
    time ./srt $0 >$0.tmp
    mv $0.tmp $0
```

- build process for installing software:
 - 1. configure checks for machine details / compatibility or dependency issues
 - creates a makefile
 - 2. make uses makefile to compile program code and create executbales in current directory
 - 3. make install copies executables into final, system directories (the shell path)

Patching

- *patch*: piece of software designed to fix problems or update a computer program
- is a diff file that includes the changes made to a file
- diff unified format:

```
- diff -u og_file mod_file
- --- path/to/og_file
- +++ path/to/mod_file
- @@ -l,s +l,s @@ where:
```

- * @@ refers to the beginning of the 'hunk'
- * l is the beginning line number
- * s is the number of lines the change hunk applies
- * lines with sign were deleted
- lines with + sign were added
- · apply patches with patch command
 - eg. patch -pnum <patch_file</pre>
 - pnum strips off leading slashes in order to generate a relative path

Python

- · not just a scripting language, also object oriented
- compiled and interpreted
- not as fast as C

Python CS35L

• fewer syntactical constructions, English keywords instead of symbols

- automatic garbage collection
- interactive and script modes
- no semicolons / braces for blocks of codes
 - * instead uses indents, strictly enforced
 - * blocks must be equally indented
- however, ignores spaces within statements
- python is *loosely* typed
- python variables:
 - start with underscore or letters
 - followed by other letters, underscores, or digits
 - no reserved words
- in python2, print is a statement
 - in python3, print is a function "'py #!/usr/bin/python counter = 100 miles
 = 1000.0 name = "John" print counter print miles print name

if (counter == 100): print "correct" print "good-bye"

#here is a comment """here is a multiline comment"" - python list is like a C array but is dynamically sized - can also hold objects of different types - python dictionary is a hash table - key-value pair storage capability - keys unique, values not - keys and values can be different types - to add to a list: - `my_list.append(obj)` appends an object to the end of the list - `my_list.extend(iterable)` appends each element of the iterable onto the list - for a string, extend will append chars onto the listpy list = [1, 3.14, 'str'] list1 = [1, 2, 3, 4] list2 = [2, 3, 4, 5] list1 += list2 print list[0] list1[0] = 100 print list1

```
dict = {} dict['france'] = " paris" print dict['france']
```

if (dict['france'] == "paris"): print "correct" elif (dict['france'] == "europe"): print "maybe" else: print "wrong"

del dict['france'] del dict

```
dict = { 'name':'test', 'class':3 }
```

for i in list1: print i for i in range(len(list1)): # from 0 up to length of list print i - string operations:py s = `Hello' #s[1:4] is ell #s[1:] is ello #s[1:100] is ello

#can also use negative based indexing #s[-1] is o #s[-4] is e #s[-3] is He #s[-3:] is llo #can split using delimietrs x = "blue, green, red" arr = x.split(",") # arr is ['blue', 'green',

```
'red'] - functions:py def printme(some str): print some str return
def find sum(some list): sum = 0 for element in some list: sum += element return
sum - classes:py class Rectangle: def init(self, x, y): self.l = x; self.w = y;
def getArea(self): return self.l * self.b;
def getPerimeter(self): return 2 * (self.l + self.b);
def main(): rect = Rectangle(3, 4) print("Area:", rect.getArea()) print("Perimeter:",
rect.getPerimeter())
          - can use optparse library to parse through argument, options, and argument
-options when writing scripts - I/O basics: - `raw_input("string prompt")` returns
a string with trailing newline removed - `input("string prompt")` assumes input is
valid Python expression - to check if void user input, `if input_str:` - file I/O:
- `f = open(filename, 'r')` creates a file handle - `lines = f.readlines()` creates
 a list of strings from handle - `f.close()` closes handle - optparse library and
error handling:py parser = OptionParser(version=version msg, usage=usage msg)
parser.add_option('-n', '-numlines', action='store', dest='numlines',) # stores next
argument into destination numlines default=1, help='output NUMLINES lines (default
1)') options, args = parser.parse args(sys.arv[1:]) # options, object with all option
args, args is leftover positional args
try: numlines = int(options.numlines) except: parser.error('invalid NUMLINES:
{0}'.format(options.numlines)) # catching exceptions
if ... raise ValueError('error message') # raising an error
if name == 'main': main() # making Python file standalone "'
```

C Overview

- basic data types:
 - int, float, double, char, void
 - no bool!
 - size_t (unsigned integer)
- different ways to pass arguments / stdio to a script:
 - cat file | ./script
 - ./script <file</pre>
 - ./script file (can access using argv variable)

• *double pointer*:

```
int* p = &x, and then int** q = &pq is a double pointer
```

• pointers can point to arrays, and arrays are pointers:

```
int A[5];
int* p;
p = &A[0];
p = A;
*p == A[0];
p + 1 == &A[1];
*(p + 1) == A[1];
A + 2 == &A[2];
// with 2-D arrays
int A[2][3];
A == &A[0];
A + 1 == &A[1];
**(A + 1) == A[1][0];
A[0] + 1 == &B[0][1]; // A[0] is type int*
// can use function pointers to pass a function to another
int (*fp)(int, int);
fp = &add;
sum = (*fp)(2, 3);
// or
fp = add;
sum = fp(2, 3);
// can use with library functions such as qsort
#include <stdlib.h>
void qsort(void* base, size_t num, size_t size,
           int(*compare)(const void*, const void*));
// can pass in a comparison function to gsort
// < 0 if p1 goes before, 0 if equivalent, > 0 if after
```

```
int compare(const void* a, const void* b) {
  return *(int*)a - *(int*)b;
  // for comparing strings/letters, would have to cast from char** type:
  // char* letter = (char**) a;
}
int values[] = { 40, 30, 60 };
qsort(values, 6, sizeof(int), compare);
```

Pointer Example

```
#include <stdio.h>
char *c[] = { "the", "quick brown fox", "jumped", "over the", "lazy dog" };
char **cp[] = { c+3, c+2, c+1, c, c+4 };
char ***cpp = cp;
int main(void)
  printf("%s ", *(c+1));
  printf("%s ", *(c+2));
  printf("%s ", **(cpp+3));
  printf("%s ", **(cp+4)+4);
  printf("%s ", **(++cpp));
  printf("%s ", **cp);
  printf("%s ", **(++cpp)+5);
  /* the
     _dog
     jumped
     over_the
     _brown_fox */
}
```

Structs

- no classes in C
- in structs:
 - package related data
 - no member functions
 - no access specifiers (private by default)
 - no constructors
- can use typedef as syntactic "sugar"

```
struct Student {
   char name[64];
   int age;
   int year;
};
struct Student s;

typedef struct {
   char name[64];
   int age;
   int year;
} Student;
Student s;
Student* sp = &s;
sp->age = 18;
```

Dynamic Memory

- memory that is allocated at runtime, will be allocated on the *heap*
- must know size of the array at compile time
 - for dynamic arrays, must use allocated memory
- void* malloc(size_t size):
 - allocates size bytes and returns pointer to that memory
 - returns NULL if memory not allocated

```
- eg. ptr = (int*)malloc(n * sizeof(int))
```

- void* calloc(size_t num, size_t size):
 - allocates block of memory for array of num elements, each size bytes

Debugging CS35L

```
    void* realloc(void* ptr, size_t size):

            changes the size of the memory block pointed to by ptr to size bytes
            contents of memory is unchanged

    void free(void* ptr):

            frees the block of dynamic memory pointed to by ptr
            double free has undefined behavior
```

I/O

```
read characters from stdin with getchar()
write characters to stdout with putchar(int char)

– these are unbuffered input and output
fp = fopen("file.txt", "w+")
formatted I/O with fprintf() and fscanf()

– int fprintf(FILE* fp, const char* format, ...)

* fp can be file pointer, stdin, stdout, stderr

– eg. fprintf(stdout, "%s has %d points\n", player, score);
can write to stderr as well with perror()
```

Debugging

- debugger allows programmers to:
 - step through source code line by line
 - inspect program at runtime
- GDB for C
- compile with -g flag to use with debugger
 - gdb <exe> or gdb and (gdb) file <exe>
 (gdb) run or (gdb) run [args]
 (gdb) help <command>
- useful functionalities:
 - set breakpoints
 - * (gdb) break file.c:6
 - * (gdb) break my_function
 - * (gdb) break [position] if expression
 - * (gdb) info break
 - * also delete, disable, enable, ignore breakpoints
 - check variables with print [/format] expression
 - * formats: d, x, o, t (binary)
 - watch changes to variables, rwatch expression

- step (steps into subroutines), next (does not trace into subroutines)
- continue (next breakpoint), finish (until current function returns)
- process memory layout:
 - TEXT segment with instructions
 - global variables (initialized and uninitialized)
 - heap segment grows upward (dynamic memory allocation, malloc, free)
 - stack segment grows downward (push, pop, stores)
 - * stack frame set aside for called functions
 - * holds local variables, parameters, return address, etc.
 - (gdb) backtrace shows the call trace / stack
 - (gdb) info frame shows registers, return address of current frame
 - (gdb) info locals shows local variables and values
 - (gdb) info args shows argument values
 - (gdb) info functions shows functions
 - (gdb) info list lists source code lines around current line

System Call Programming

- processor modes place restrictions on type of operations by processes
 - user space vs. supervisor space
 - * user applications and C library are in the user space
 - * in user space, some memory cannot be accessed, cannot do some I/O
 - kernel/supervisor space has unrestricted access
 - * access all areas of memory, can take over cpu
 - * exception will crash the OS
 - * CPU mode bit changes from 1 to 0
- this allows protection of I/O, memory, and CPU
 - prevent processes from messing with each other and the OS
- system calls are a special type of function:
 - making a type of kernel call from user-level processes
 - provide interface between kernel and user programs
 - * only way user program can perform privileged operations
 - when system call is made:
 - * program is interrupted
 - * control is passed to the kernel and interrupt handler
 - * also handle exceptions and *traps*
 - have an overhead
 - expensive and can hurt performance

- * OS must interrupt and save state of program
- * take control of CPU and verify operation
- * restore saved context, and restore CPU to user process
- in C, system calls are used the same way as calling a procedure or a function
 - but only a system call can enter the kernel
 - also provide C library functions to make system calls
 - * fewer system calls, less switches in control

Types

- 1. process control stopping execution of a running programing
- eg. fork, exit, wait
- 2. file management open, close, read, write, create
- 3. device management
- read, write, ioctl (input-output control)
- 4. information management transfer info between user and OS, eg. date or time
- getpid, alarm, sleep
- 5. communication talk between different processes
- pipe, shmget (allocates shared memory segment), mmap (maps files or devices into memory)

Example System Calls

- file descriptor: integer that identifies open file of process
- file descriptor table: collection of integer array indices that are file descriptors
 - elements are pointers to file table entries
 - one for each process
- int open(const char* Path, int flags)
 - path can use relative or absolute path (starts with /)
 - flags include O_RDONLY, O_WRONLY, O_RDWR, O_CREAT (create file), O_EXCL (prevent creation)
 - * O_APPEND file offset set to end of file before each write
 - * 0_TRUNC truncates existing file to size 0
 - returns file descriptor used (starts at 3 since 0-2 are taken by system)
- int creat(char* filename, mode_t mode)
 - mode indicates permission
 - returns first unused file descriptor, generally 3

- open and close return -1 on an error
- to avoid system call overhead, can use equivalent library functions:
 - these make system calls, but have been optimized to make fewer system calls
 - getchar(), putchar()
 - fopen(), fclose() "'c #include <fcntl.h> #include <unistd.h>

// file descriptors: 0 stdin, 1 stdout, 2 stderr // note: ssize_t can take values from -1 to Tmax

ssize_t read(int fildes, void* buf, size_t nbyte); // file descriptor, buffer to write to, num bytes to read // returns actual number of bytes read, -1 is an error, 0 is EOF

// read is unbuffered, direct system call reading one byte at a time // vs. getchar is buffered, library implemented system call, // reads multiple bytes at one time into a block (buffered usually faster)

ssize_t write(int fildes, void* buf, size_t nbyte); // file descriptor, buffer to write from, num bytes to write // returns actual number of bytes written, -1 is an error, 0 is EOF

// to read from a file: fildes = open("foo.txt", O_RDONLY); if (fildes < 0) { perror("error");
exit(1); } close(fildes);</pre>

// lseek() repositions file descriptor offset off_t lseek(int fd, off_t offset, int whence) // SEEK_SET set offset // SEEK_CUR current location + offset // SEEK_END size of file + offset

int fstat(int fildes, struct stat* buf); // returns information about a file from filedes in a structure // negative if error struct stat buf; buf.st_size buf.st_nlink // number of links buf.st_ino // inode S_ISDIR(buf.st_mode) // directory? S_ISLNK(buf.st_mode) // symbolic link? buf.st_mode & S_IRUSR // u+r? buf.st_mode & S_IWUSR // u+w? buf.st_mode & S_IXUSR // u+x? ...S_IRGRP ...S_IROTH

int ret = fstat(STDIN_FILENO, &buf); // STDIN_FILENO is 0 if (S_ISREG(buf.st_mode))
printf("Regular File");

cat file.txt | ./fstat // non regular file, 0 bytes in size ./fstat < file // regular file, n bytes in size ./fstat < /proc/self/status // regular file, 0 bytes in size! (growing file) // can use sleep or GDB breakpoint to debug "' ## Parallelism ***

- multiprocessing: the use of multiple CPU cores to run multiple tasks simultaneously
 - uni-processing system vs. a multiprocessing one
 - on uni-processing system, uses time-sharing, switches between different threads

- * creates an illusion of parallelism
- vs. multiple cores running the threads at the same time
- parallelism: executing several computations simultaneously to gain performance
 - multitasking: several processes scheduled alternatively or simultaneously
 - different processes set different address spaces (bad for sharing, good for protecting)
 - * expensive creation/destruction
 - * insulated from errors in other processes
 - multithreading: same job is broken logically into pieces or threads which may be executed simultaneously
 - * all threads in same process share the same memory (good for sharing, bad for protecting)
 - · can access each others' memory
 - · light-weight threads creation, easy communication
 - * creating and destroying threads is easier to do, as opposed to creating an entirely new process
 - * one error can crash all the threads
 - *embarassingly parallel*: an application where threads do not have to synchronize with each other
 - * good candidates for synchronization
- thread: a flow of instructions or a path of execution within a process
 - smallest unit of scheduling by OS
 - process *consists* of up to one thread
 - on a uni-processor, processor switches between different threads, parallelism is an illusion
 - on a multi-processor, multiple processors run threads at the same time, true parallelism
- if there are global variables or pointers being used, multiple threads can access the same variable
 - sharing the same resources
 - can lead to issues with synchronization, a race condition
 - * eg. incrementing a value twice on two separate threads
 - · may only result in a single increment if they execute at the same time
 - can solve by "locking" and "unlocking" critical sections
- critical sections:
 - mutex: object that allows only one thread into a critical section
 - owned by a thread, forces other threads to wait
 - each resources has a mutex

```
* pthread_mutex_t lock;
* pthread_mutex_init(&lock, NULL);
* pthread_mutex_lock(&lock);
* pthread_mutex_unlock(&lock);
* pthread_mutex_destroy(&lock);
```

- *semaphore*: value in a designated place in the OS that each process can check or change
 - * signaling mechanism
 - * restricts number of simultaneous threads of a shared resource up to a maximum number
 - * requesting access to a resource decrements semaphore
 - · signal on completion increments the semaphore
- C pthread functions found in pthread.h, use -pthread when compiling
- thread id held in pthread_t structure
- can create a thread in C with a function using pthread_create:

```
- int pthread_create(pthread_t* thread, attr, void* (*someFunct) (void*),
    void* arg)
```

- parameters:
 - * address of pthread_t to create thread in
 - * thread attributes (priority, stack size), usually NULL
 - * function to run in thread, must return void* and take void* argument
 - void* argument to pass (only one argument taken, may have to use a struct)
- on success, returns zero, otherwise an error number
- use -lpthread option when compiling
- int pthread_join(pthread_t thread, void** retval) waits for another thread to complete
 - can also catch return values from the thread
 - takes thread ID and exit status
 - error if nonzero return
- pthread_equal() compares thread ID's to see if they are equal
- pthread_exit() can return a value, completely exit thread

```
/* Passing struct to a thread */
struct Point {
  int x, y;
  struct Res* r;
};
struct Res {
```

Linked Libraries CS35L

```
int sum;
};
void* add(void* pointPtr) {
  struct Point* myPoint = (struct Point*)pointPtr;
  for (int i = 0; i < 2; i++) {
    struct Res* myRes = myPoint->r;
    myRes->sum = myPoint->x + myPoint->y;
    myPoint++;
  }
}
int main()
  struct Point arr[4];
  struct Res arr1[4]:
  struct Point* p1 = &arr[0];
  struct Point* p2 = &arr[2];
  pthread_t thread1, thread2;
  pthread_create(&thread1, NULL, add, (void*)p1);
  pthread_create(&thread1, NULL, add, (void*)p2);
  pthread_join(thread1, NULL);
  pthread_join(thread1, NULL);
}
```

Linked Libraries

• an object code library is a previously compiled collection of standard program functions

- static linking links an entire file into a program
 - static libraries have .a extension
 - gcc -static hello.c -o hello-static (treats dynamic libraries as static)
- *dynamic linking* allows a process to add, remove, replace, or relocate object modules during its execution

Linked Libraries CS35L

- not copying the entire library, just **referencing** the path to that library
- thus, the linking is only **completed** at run time
 - * the size *increases* while executing after incorporating libraries
- dynamic/shared libraries have a .so extension
 - * or .dll on Windows
- some advantages:
 - * more space efficient than copying the entire library
 - * multiple programs can use the library
 - * also will *not* have to **recompile** the libraries when they are updated
- some disadvantages:
 - * slightly slower execution
 - * no control over specific version of libraries
 - unusable if the library does not exist
- gcc hello.c -o hello-dynamic
- in *dynamic loading*, only functions that are used from libraries are copied at runtime
 - instead of the entire library, as with *dynamic linking*
 - even slower execution

Creating and Using Libraries

- to create a **static** library:
 - compile without linking: gcc -c lib_mylib.c lib -o lib_mylib.o
 - create a static library: ar rcs lib_mylib.a lib_mylib.o (archiver)
 - * c creates, s writes index, r replaces, q quick append without replacement, v verbose
 - compile main code: gcc -c main.c -o myMain.o
 - link to main code: gcc -o mainCode myMain.o -L. -l_mylib (syntax for library path)
- to create a **shared** library:
 - create C and header file for library: lib_mylib.c, lib_mylib.h
 - compile without linking: gcc -fPIC -c lib_mylib.c -o lib_mylib.o
 - create dynamic library: gcc -shared -o lib_mydynamiclib.so lib_mylib.o (no archiver)
 - link to main code: gcc -W1,-rpath=\$PWD -o mainCode main.c -L. l_mydynamiclib
 - * -Wl specifies where the library comes from during runtime, specifying linker options
 - * two parts:
 - · compiling and outputting errors with the library (-L)

- · the relative path reference to use when linking(-Wl)
- GCC flags:
 - fPIC outputs position independent code, required for shared libraries
 - - *llibrary* links with lib*library*.a
 - static libraries, or object file archives
 - -L, at compile time, find the library from this path
 - by default, checks /usr/lib
 - Wl passes options to linker
 - rpath at runtime finds .so from this path
 - c generate object code from C but do not link
 - shared produces a shared object which can then be linked with other objects to form an executable
- · dynamic loading has a unique API
- to use shared library with **dynamic loading**:
 - void* dlopen(const char* file, int mode) makes an object file accessible to a program
 - * mode flags include:
 - RTLD_LAZY resolve undefined symbols as code is executed
 - · RTLD_NOW resolve all undefined symbols now (for debugging)
 - RTLD_GLOBAL can be or'ed with the other two, allows symbols to be accessible to other loaded libraries
 - * returns a handle to be used by other DL library routines, can check for NULL
 - void* dlhandle = dlopen(...)
 - void* dlsym(void* handle, const char* name) obtains the address of a symbol within an opened object file
 - * use with a function pointer, eg. int (*myfunc)(int*) = dlysym(dl_handle
 , "powerOfNine");
 - char* dlerror() returns a string error of the last error that occurred
 - * should save into a char*, returns NULL after error first returned
 - char* dlclose(void* handle) closes an library file
 - compilation: gcc -ldl -W1,-rpath=\$PWD -o myCode main.c (-ldl similar to -lpthread)

Communication

- over the Internet, want certain guarantees when communicating:
 - confidentiality with encryption

- data integrity
- authentication of host and receiver identification
- authorization, rights to resources
- encryption can be **symmetric**:
 - sender and receiver have the same key to decrypt
 - both parties need to know the keys
 - * key may have to be delivered between parties
 - * could be compromised
 - could use Key Exchange Algorithm, independently calculate same secret key

• or asymmetric:

- each party has their *own* public-private key pair
 - * if a message is encrypted with one key, it has to be decrypted with the other
- data decrypted with private key, encrypted with public key, or vice versa
- public key is published and known to everyone
 - * anyone can encrypt, but cannot decrypt
- more expensive to set up, for prolonged communication:
 - * symmetric encryption is set up after asymmetric communication is established
 - 1. server sends copy of assymetric public key
 - 2. browser creates symmetric session key and encrpyts it with serber's public key
 - 3. server decrypts session key using asymettric private key to get symmetric session key
 - 4. can now use symmetric session key
 - * called a session key
- can *authorize* a party using **asymmetric** encryption: (host validation)
 - encrypt random data with party's public key
 - that party can decrypt that data with their *private* key
 - reverse the data (for example), and re-encrypt the data with their *private* key
 - can decrypt the data with party's *public* key again
 - should be in reverse order if party is indeed authorized
- *public* keys must be *verified* and signed by a certificate authority
 - web browser will check certificates with this authority
- signatures provide for *authentication* of identities
 - an electronic stamp or seal
 - also ensures data integrity
 - can be appended to a document or separate (detached signature)

- * detached is stored separately
- steps:
 - generate a message digest
 - created using hashing algorithms
 - * acts a summary of the message
 - * a slight change in message can create a different digest
 - digest is encrypted using sender's private key
 - * resulting encryption is digital signature
 - attach to message and send
- receiver can decrypt signature into digest with senders' public key
 - only the sender's public key can decrypt the signature
 - generate digest with same algorithm
 - compare digests
 - * if not exactly the same, message has been tampered with
- GPG Handbook
 - gpg [options]
 - gen-key generates new keys
 - armor ASCII format
 - -export, -import
 - detach-sign creates detached signature
 - -verify verifies signature with public key
 - encrypt, –decrypt
 - -list-keys, -search-keys

SSH

- secure socket shell, used to remotely access shell
- encrypted and more authenticated than predecessors
- client ssh's to remote server
 - first time, asks for host validation
 - * checks server's public key against saved public key
 - no central authority for servers, user must manually trust host
- man-in-the-middle attack is where an attacker assumes the keys of another server
 - eg. a server's keys were compromised and updated without the user knowing
 - check host public key with save public key
 - must manually remove key from the trusted host
- password-based authentication vs. key-based authentication
 - key-based:

- generate key pair
- private key is used to to decrpyt challenge message from server
 - * private key still secured by passphrase, user still must enter passphrase
- to avoid reentering passphrase, use ssh-agent, which stores a copy of the private key in memory
 - this copy is used to answer challenge messages
 - still must enter passphrase once to load into memory
 - kernel still protects memory from being read by other processes
- ssh-keygen generates key-pair
 - sudo useradd <user>, sudo passwd <user>
 - sudo mkdir .ssh
 - sudo chown -R <user> .ssh, sudo chmod 700 .ssh
 - ssh-copy-id -i uses local key to authorize logins on remote machine (-i identify file only)
 - ssh-add uses ssh-agent
- can use ifconfig or hostname -I to get hostname and IP addressses

Sample Questions

- 1. a. If you don't trust your server, you cannot use OpenSSH to connect to it, as it can easily corrupt your client.
 - b. If you don't trust your client, you can still use OpenSSH to connect to a trusted server.
 - c. If you don't know the name or IP address of your server, you can use OpenSSH to discover this info in a secure way.
 - d. If you don't know the name or IP address of your client, you can still use OpenSSH to connect to a server.
 - e. If you don't trust your network, you an still use OpenSSH to discover whether your server is running or not.
- 2. a. ssh-agent improves security by making a copy of private keys.
 - b. ssh-agent acts on your behalf by running on the server and executing commands there, under your direction.
 - c. Even if the attacker replaces the ssh-agent program with a modified version your communications will still be secure.
 - d. ssh-agent eliminates all need for password authentication when communicating to SEASnet hosts.
 - e. If you successfully use ssh-agent and then log out from the client and then

- log back in again, you can then connect to the same SSH server again without any additional passwords or passphrases.
- 3. a. OpenSSH typically uses public-key encryption for authentication, because private-key encryption is less secure.
 - b. OpenSSH typically uses private-key encryption for data communication, because public-key encryption is less efficient.
 - c. When you run ssh, it chooses its authentication key randomly from a large key space.
 - d. The OpenSSH client and server are essentially symmetric, so that it's easy and common to use the same program as either a client or a server.
 - e. Once your private keys are 1024 bits long, there's no reason to make them any longer, as they're impossible to break.
- 4. a. OpenSSH is not limited to just one client-server connection. A team can use OpenSSH to communicate information to each other.
 - b. For security, OpenSSH refuses to connect to programs written by other people. A client running OpenSSH code will only connect to a server running OpenSSH code.
 - c. Although port forwarding can be used to display from OpenSSH server to an OpenSSH client, the reverse is not possible.
 - d. For security, port forwarding cannot be chained.
 - e. When using port forwarding to connect to SEASnet, one should take care not to create a forwarding loop, leading to a cycle of packets endlessly circulating on the Internet.
- 5. a. It's not a good idea to connect to a SEASnet server and use GPG on the server to sign a file, because then an attacker on the network can snoop the GPG passphrase.
 - b. A detached signature file must be protected as securely as the private key. Otherwise, an attacker will be able to forge your signature more easily.
 - c. When generating a key pair it's important to use a private entropy pool. Otherwise, an attacker might be able to guess your key by inspecting the public entropy pool.
 - d. Exporting a GPG public key to ASCII format neither improves nor reduces its security.
 - e. Because a deached cleartext signature isn't encrypted, it is easily forged by an attacker with access to the file being signed.

Version Change Management

- software development process involves a lot of changes
 - new features, bug fixes, performance enhancements
 - many different versions and large software teams
- with multiple users across a server, must save file history
 - track changes to code and other file
 - version control software, eg. Git, Subversion, Perforce
- distributed vs. centralized version histories
- in *local VCS*, organizing different version as local folders
 - no server, hard to distribute to other users
- in *centralized VCS*, version history sits on a central server
 - changes must be committed
 - all users can get the changes
 - if central server fails, there are no *local* copies
- in a distributed VCS, version history is replicated at every user's machine
 - users have version control all the time
 - changes can be communicated between users
 - eg. Git

Git

- everything in Git is checksummed
 - hash values for every object, SHA-1 hashing
- terms:
 - head: currently active head, commit object
 - branch: refers to a head and its entire set of ancestor commits
 - master: default branch
- local operations (on local database):
 - git add pushes to the staging area (git rm put into staging for removal from repo)
 - selectively stage commits
 - * can use .gitignore file to ignore filenames/patterns to ignore while staging
 - git commit pushes from staging area to the Git directory (repository)
 - * if files in staging area are modified further, have to re-add
 - git checkout checks out a previous version
 - * eg. git checkout <hash>
- steps to make a Git repository:

- git config, -list, -global
- git init
- git add somefile
- git commit -m "Check-in ..."
 - * do not need to specify files, commit pushes everything in staging area to local repository
- git commit -a -m "Check-in ..." automatically stages modified/deleted files
 - * git commit --amend redoes a commit with a new message
- git reset unstages files
 - * eg. git reset HEAD <file> unstages file to current commit
- git status gives status of files, -s for short
 - files untracked if not added to repository yet
 - files not staged for commit if they have been changed
 - files to be committed if in staging area
- git log gives commit history, -p successive difference in patching
- git diff gives diff of files
 - diff of changes in the working directory
 - staged compares with last commit

Branches

- a *branch* is a pointer to one of the commits in the repo (head) and all ancestor commits
 - pointer to one of commits (HEAD), and all ancestor commits
 - * default master branch points to last commit made, moves forward with every commit
 - branches can create commits without changing the master branch
 - * expirement code without affecting main branch
 - * separate projects that once had a common code base
 - git branch newBranch
 - git checkout someBranch switches branches (on file level, detaches head to look at commit history)
 - * git checkout -b newBranch makes a new branch
 - * git checkout -b newBranch existingBranch bases new branch off of a previous one
 - git merge someBranch tries to merge multiple branches together
 - * from someBranch to current branch
 - * creates a new merged commit
 - * can lead to merge conflicts which must be manually resolved
 - git rebase creates a linear structure of branches

Appendix CS35L

* cleaner working directory with linear commits

Remote Repositories

- remote repos are hosted on a network somewhere
 - git remote show origin
 - git clone implicitly adds origin remote
- communication with remote repository:
 - git clone creates a copy of an existing repository
 - git push can push multiple commits to the server
 - git pull takes changes from the central repository
 - also merges with the working directory
 - * essentially git fetch followed by git merge
 - git fetch takes changes into the local repository
 - * does NOT merge into the working directory
 - * can see what changes have been made on central repository
- Git repository is stored in a linked list structure
 - commits point to previous commits
 - HEAD points to the front of the repo, indicates where new commits will go in repo

More Git Commands

- git checkout -- <file> discards unstaged local changes, does not change history
- git revert reverts commit (this creates new commits)
- git reset can rewrite history without making a new commit
- git ignore ignores extra files, such as .o
- git clean cleans up untracked files
- git show <hash> shows object in repo
- git am applies patches from mailbox
- git tag manages more human-readable tags
 - eg. git tag -a v1.0 -m 'Version 1.0'
 - can checkout branches from their tags as well
- git help

Appendix

Appendix CS35L

Other Commands

- generating a random file:
 - head --bytes=# /dev/urandom > output.txt
- time command checks how long a process took
 - outputs real elapsed time (wall clock time)
 - CPU time used by the process
 - CPU time used by the system on behalf of the process
 - * in the actual kernel
 - * for example, memory allocation
 - CPU time is **cumulative** across threads
- strace intercepts and prints out system calls
 - -c gives a summary
 - o gives an output
- od command formats input
 - -W specifies width
 - c character format

Attributes of Functions

- used to declare certain things about functions in your program
 - helps compiler optimize and check code
 - C is not object oriented, but can be used when calling libraries
- placed before return type of function
- __attribute__ ((__constructor__)) called whenever dlopen is called
- __attribute__ ((destructor)) called whenever dlclose is called