

- Arrays :-

- Majority Element (Moore's voting algorithm) :-

Time complexity  $\rightarrow O(N)$   $\rightarrow$  space  $O(1)$

1	4	3	2	4	4	4	2
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pseudo code  $\longrightarrow$

$\uparrow$   
element = array[0]

count = 1;

if (element != array[i])

count --;

if (count == 0)

count = 1;

element = array[i]

else-if (element == array[i])

count ++

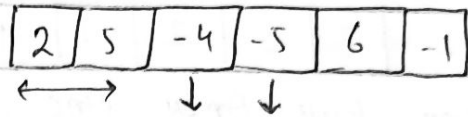
Here we check the next element  
if it is not then we use  
decrement.

$\rightarrow$  This Problem can be solved using  
map.

But its time complexity is  
 $O(N)$   $\rightarrow$  space  $O(N)$ .

## • Largest Sum Contiguous Subarray (Kadane's algorithm) :-

- if a negative number occurs we move to the next element and store sum.



function largestSum(array)

{

let maxSum = 0, curSum = 0;

for (let i = 0; i < array.length; i++)  
{

curSum = curSum + array[i]

if (curSum > maxSum)

maxSum = curSum;

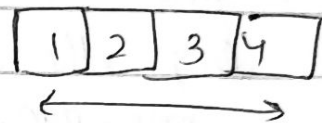
if (curSum <= 0)

cur = 0;

}

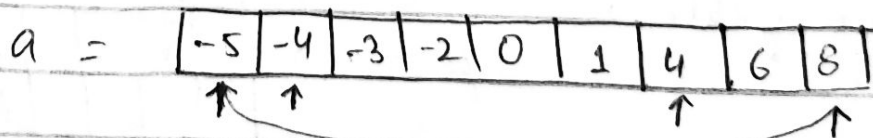
return maxSum;

}



Steps  
→ 3; → maxSum  
→ 1  
→ 6 → maxSum  
→ 12 → maxSum  
→ 9

## • Check sum Zero (array must be sorted) :-



Approach :-

Left = 0; Right = array.length - 1, sum;

sum = a[left] + a[right]

if sum > 0 right--

else left++;

sum == 0; console.log.

- Best time to Buy & sell stocks → LeetCode - 121.

3	5	1	7	4	9	3
---	---	---	---	---	---	---

you can buy only one time ∴

$min = \text{Math.min}(min, a[i]);$

$max = \text{Math.max}(max, a[i]);$

after loop the return

return  $max - min;$

- Best time to Buy & sell stocks → LeetCode 122 :-

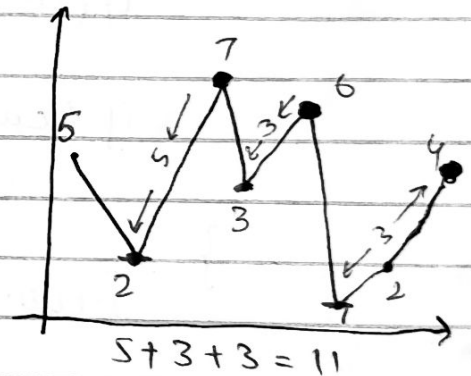
5	2	7	3	6	1	2	4
---	---	---	---	---	---	---	---

Let profit = 0;

if ( $array[i] > array[i-1]$ )

profit +=  $array[i] - array[i-1];$

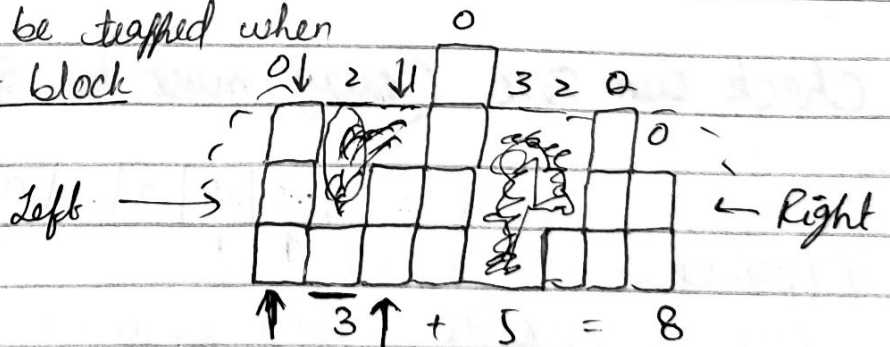
return profit;



- Trapping Rain water - LeetCode 42 :-

3	1	2	4	0	1	3	2
---	---	---	---	---	---	---	---

Logic :- water will only be trapped when  
Left block and right block  
will be greater.



Left-Max = 3, 3, 3, 4, 4, 4, 4, 4

Right-Max = 4, 4, 4, 4, 3, 3, 3, 2

value =  $\text{Min}(\text{Left}[i], \text{right}[i]) - a[i]$

## • Sorting Algorithms:-

### → Bubble Sort :-

Swapping is used and the largest element will move to the end.

8	4	1	5	9	2
---	---	---	---	---	---

in first iteration  $\rightarrow i = n \rightarrow j = 0 ; j \leq i ; j++$   
if (array[j] > array[j+1])

swap (array[j], array[j+1])

4	8	1	5	9	2
4	1	5	8	9	2

4	1	8	5	9	2
4	1	5	8	2	9

→ Largest at end.

### • insertion sort :-

Take array and split it in two parts

8	4	1	5	9	2
---	---	---	---	---	---

sorted

non-sorted.

store element in temp =

for (let i = 1; i < n; i++)

let temp = a[i];

let j = i - 1;

while (j >= 0 && a[j] > temp)

array[j+1] = array[j]

j--;

}

array[j+1] = temp

bcs of while loop we incremented one time.

8	4	1	5	9	2
---	---	---	---	---	---

4	8	1	5	9	2
---	---	---	---	---	---

1	4	8	5	9	2
---	---	---	---	---	---

1	4	5	8	9	2
---	---	---	---	---	---

1	4	5	8	9	2
---	---	---	---	---	---

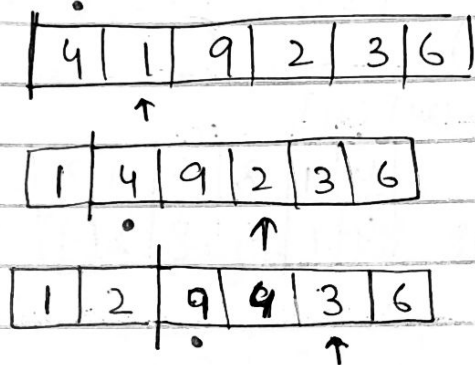
1	2	4	5	8	9
---	---	---	---	---	---

## • Insert Selection Sort :-

Take min value and arrange it / swap it with first value.

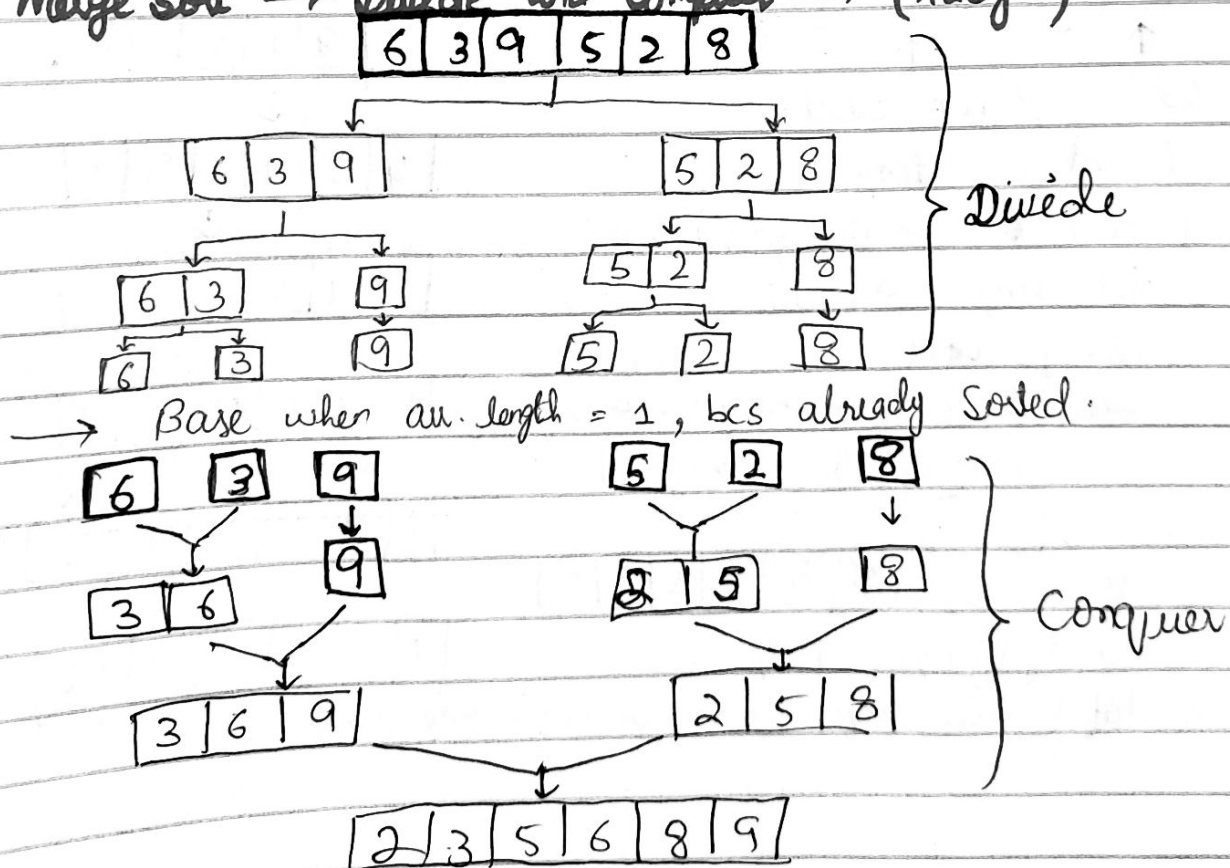
Outer loop runs  $< n-1$  times.

Inner loop runs  $< n$  times.



Dot represents min value and arrow  $j^{\text{th}}$  value.

## • Merge Sort $\rightarrow$ Divide and Conquer $\rightarrow (n \log n)$





## • Quick Sort :-

Pivot and Partition

start = 0

end = n-1

6	3	9	5	2	8
---	---	---	---	---	---

i = start - 1;

↑  
Pivot.

Pivot

- random
- mid
- first
- last

Traversing

for (let j = start; j < end; j++) {

if (array[j] < pivot)

i++

swap;

}

i++

swap (array[end], array[i]);

return i;

}

Loop will end at j = 4

2 8

6	3	5	9	9	9
6	3	9	5	2	8
↑	↑	↑	↑	↑	↑
i=0	i=1	i=1	i=2	i=3	
j=0	j=1	↑	j=3		

bcs condition is false.  
swapped

function quicksort (array)

{ pivotIndex = partition (array, start, end);  
when i reached at the index - 1 times from higher position

quicksort (array, start, pivotIndex - 1);

quicksort (array, pivotIndex + 1, end);

}