



Coping with failures in datacenters

Lies, Damned lies, and timeouts







If only Bradley's arm was longer. Best photo ever. #oscars



7:06 PM - 2 Mar 2014

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Who am !?

Cache @ Twitter

Now working on performance in general



Why do I want to give this talk?

I've been telling my coworkers the same things for years...

What do I know?

- My views are heavily influenced by two things:
 - In-memory, datacenter-scale caching
 - Twitter's environment

Cache in datacenters

- Distributed in-memory KV store
- Serving many, many requests
- with very tight latency expectation

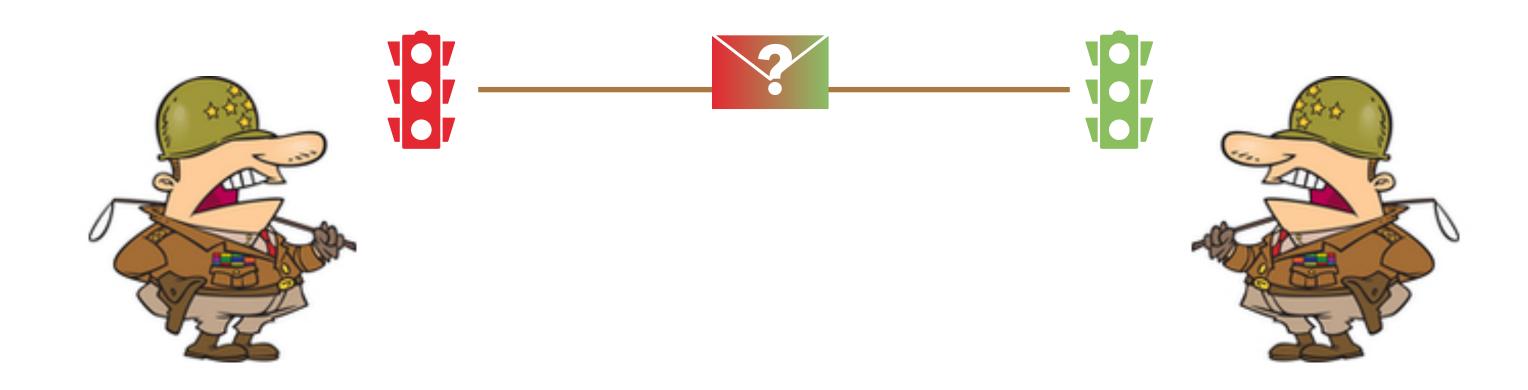
Twitter's runtime

Architecture: monolithic -> SOA/Microservices

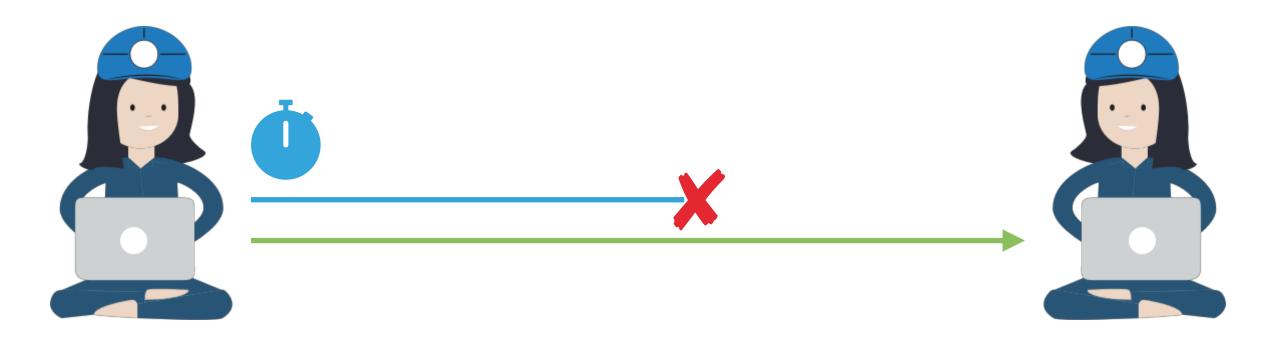
- Owns datacenters
- Services mostly on JVMs
- Jobs deployed with containers
- > Scale: up to many thousands of nodes per service

Living in an imperfect world

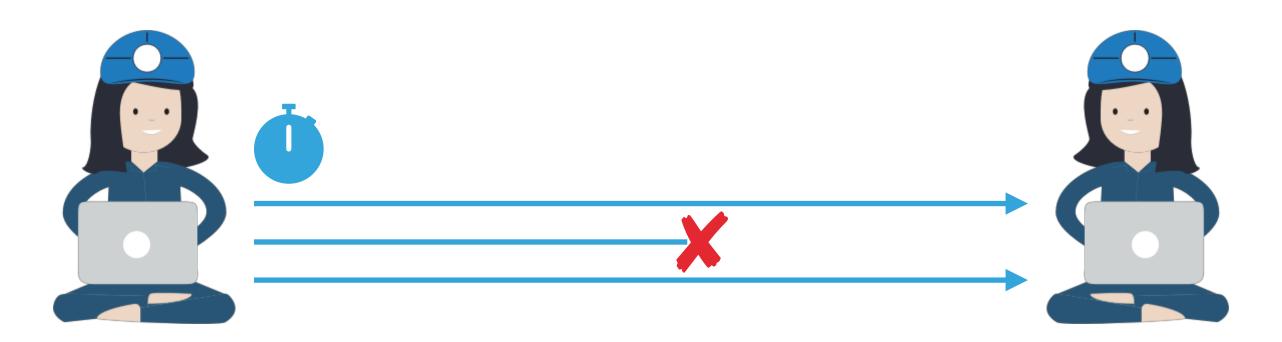
Two unhappy generals

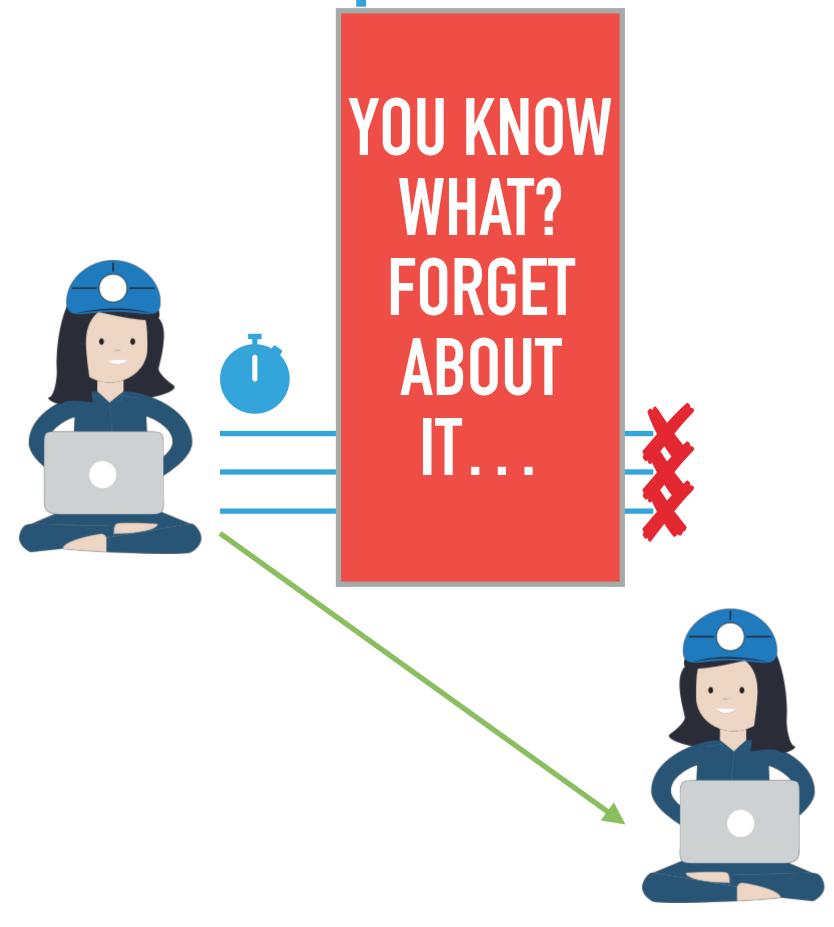


Two General Problem / Paradox







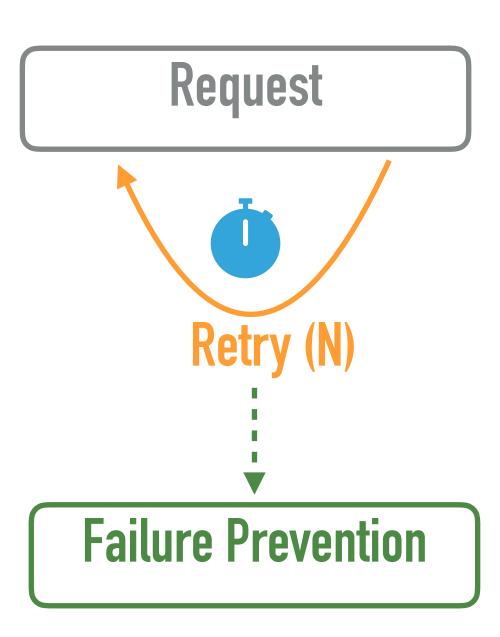




Timeouts, Retries, and preventions

Coping with failure

Coping with failure



Timeout and retry

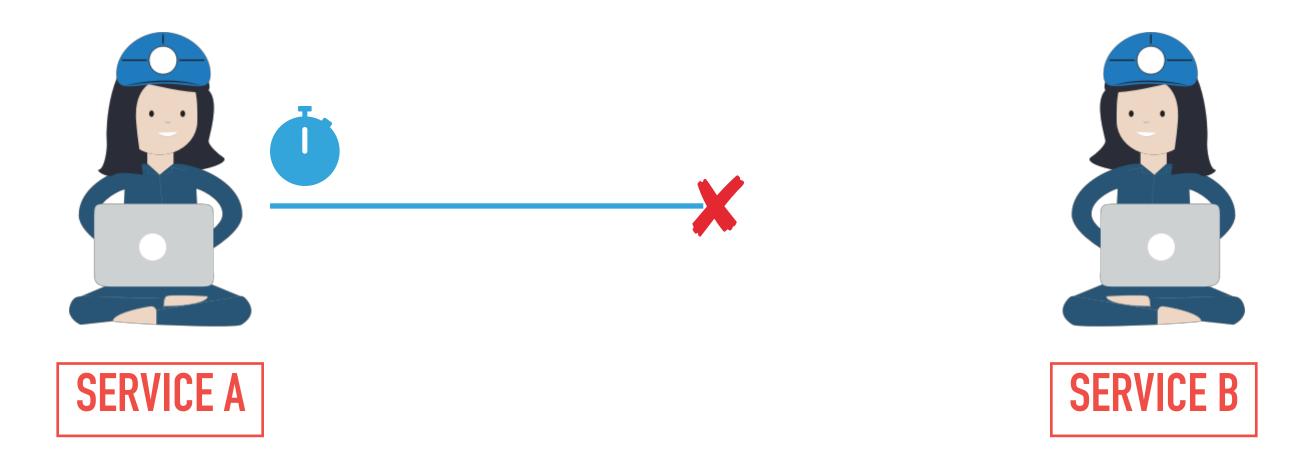
- Timeout: an approximation of failure
 - False positive is possible
- Retry: an effort to reduce failure
 - Has a cost
 - Has an effect on system

Timeout and retry get close to the heart of the difficulty of distributed systems, and yet we treat them casually because they're often presented as configuration.

What works intuitively may lead to catastrophes.

TIMEOUT

Timeout could be misleading





One service

SERVICE

LIBRARY+VM

KERNEL

HARDWARE

NETWORK

One service

Another service

SERVICE A

LIBRARY+VM

KERNEL

HARDWARE

SERVICE B

LIBRARY+VM

KERNEL

HARDWARE

NETWORK

One service **Another service SERVICE A SERVICE B** LIBRARY+VM **KERNEL CARE HARDWARE NETWORK**

One service

SERVICE

resource contention, head-of-line blocking, locking...

LIBRARY+VM

garbage collection, function calls with indeterministic timing...

KERNEL

system calls with indeterministic timing, background tasks, unfair scheduling...

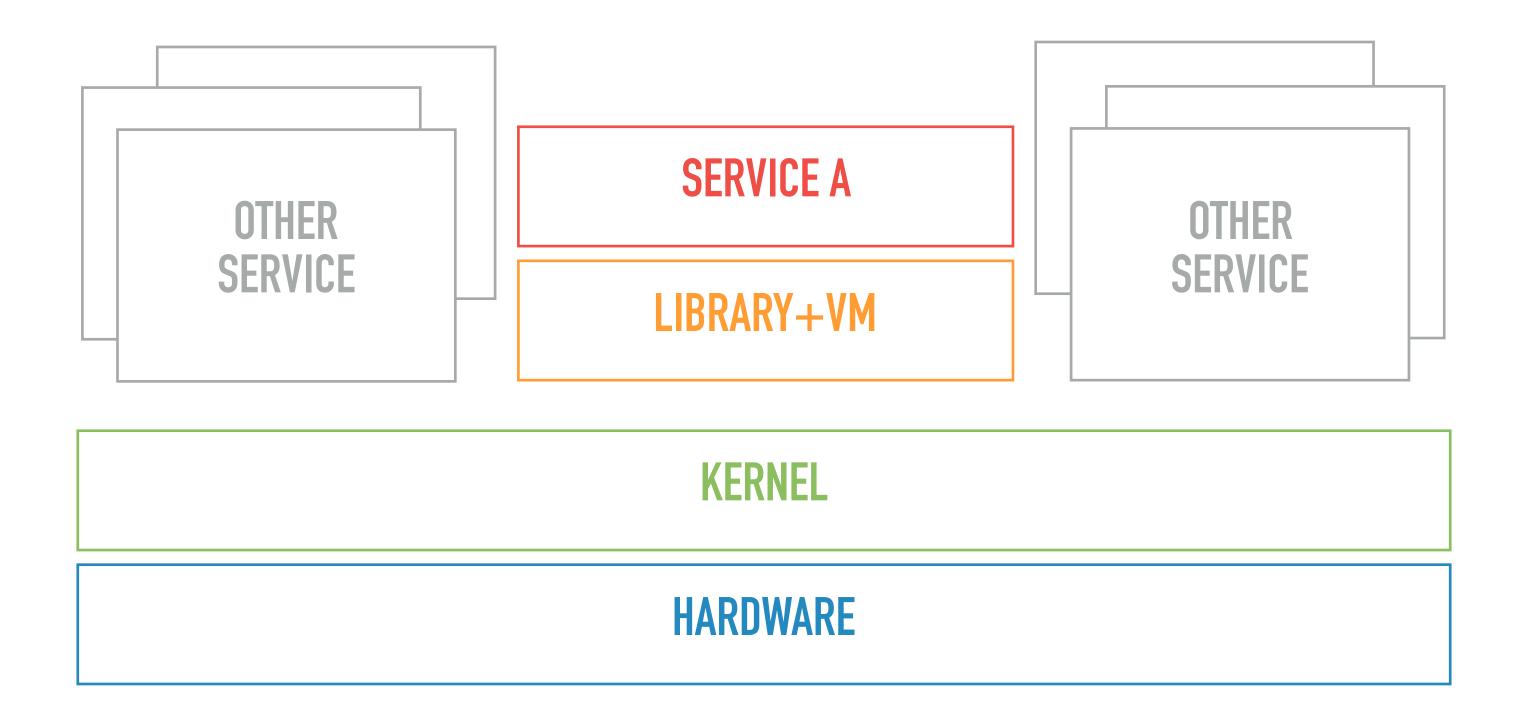
HARDWARE

blocking IO, congestion, cycle stealing...

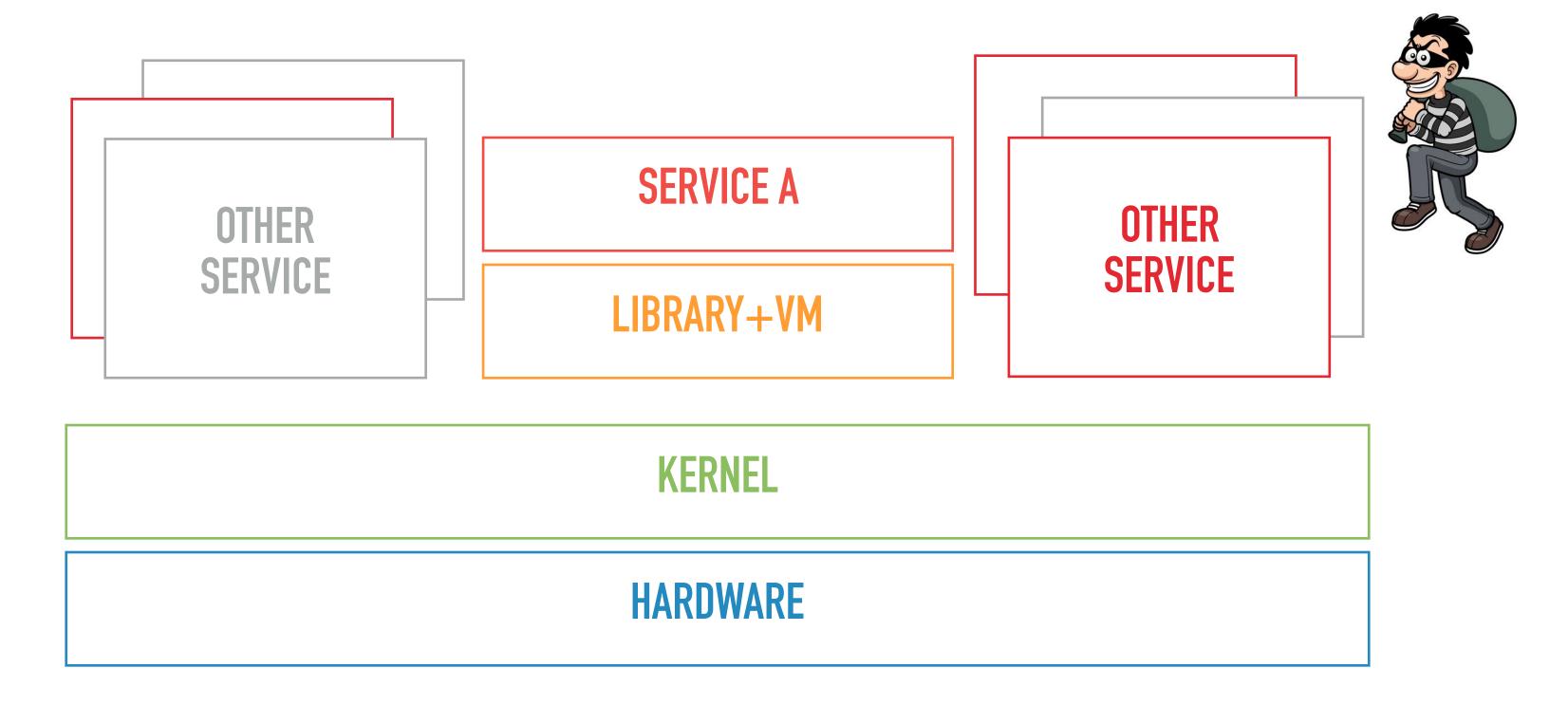
NETWORK

packet drop, queue backup...

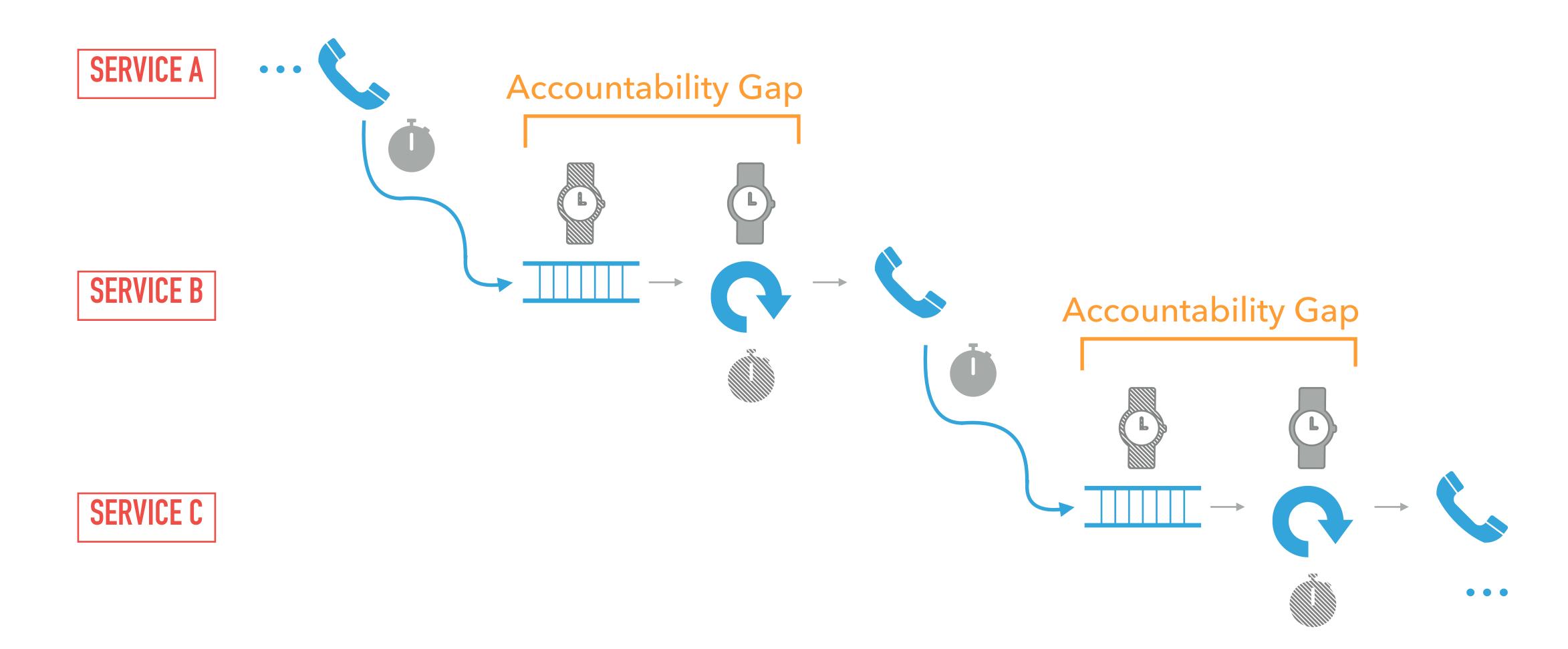
Multitenant runtime



Multi-tenancy with Noisy Neighbors



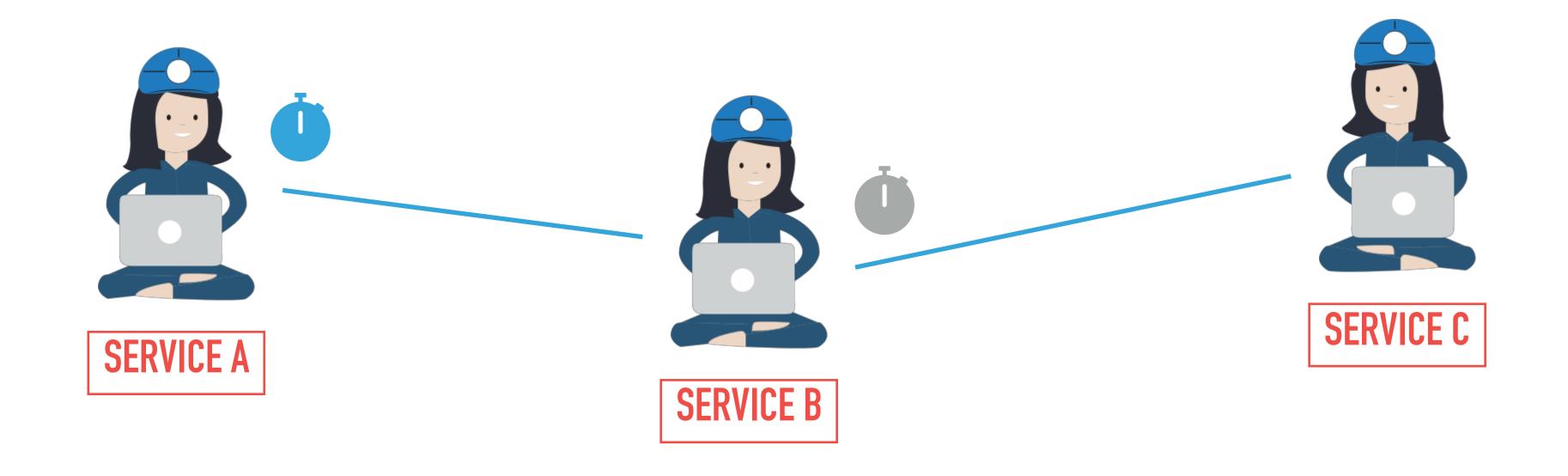
Timeout cascade



Inconvenient truths about timeout

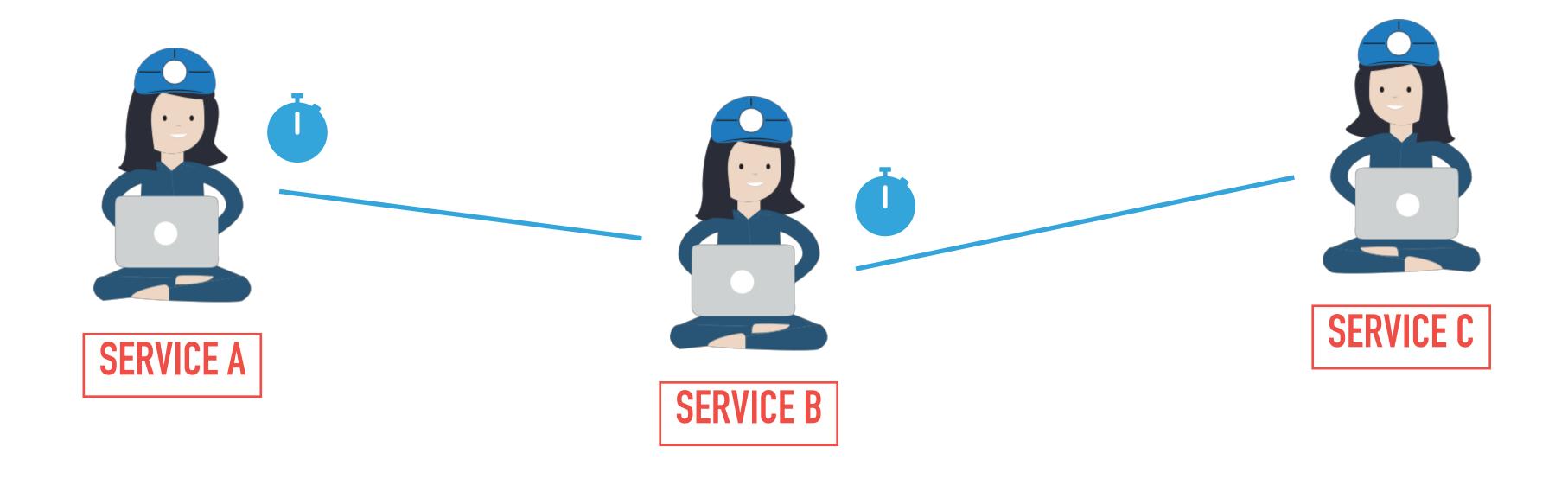
- Timeouts often do not indicate remote service health
- The optimal timeout is a moving target
- Often less predictable in shared environment
- Have gaps in overall timeline

Chained dependencies



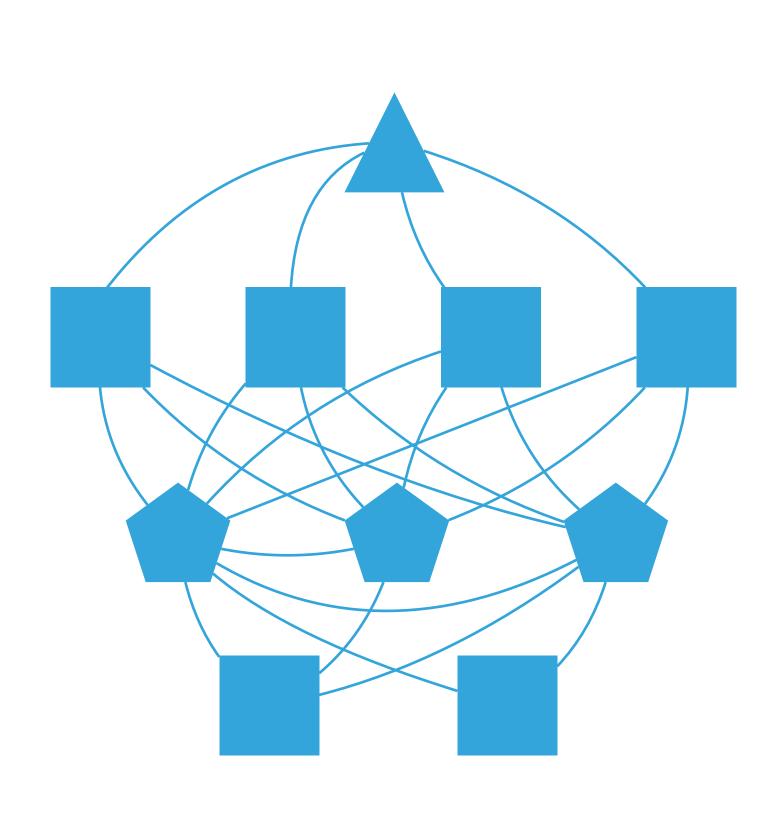
Q: A's requests time out, but not B's, which on-call engineer(s) should be paged?

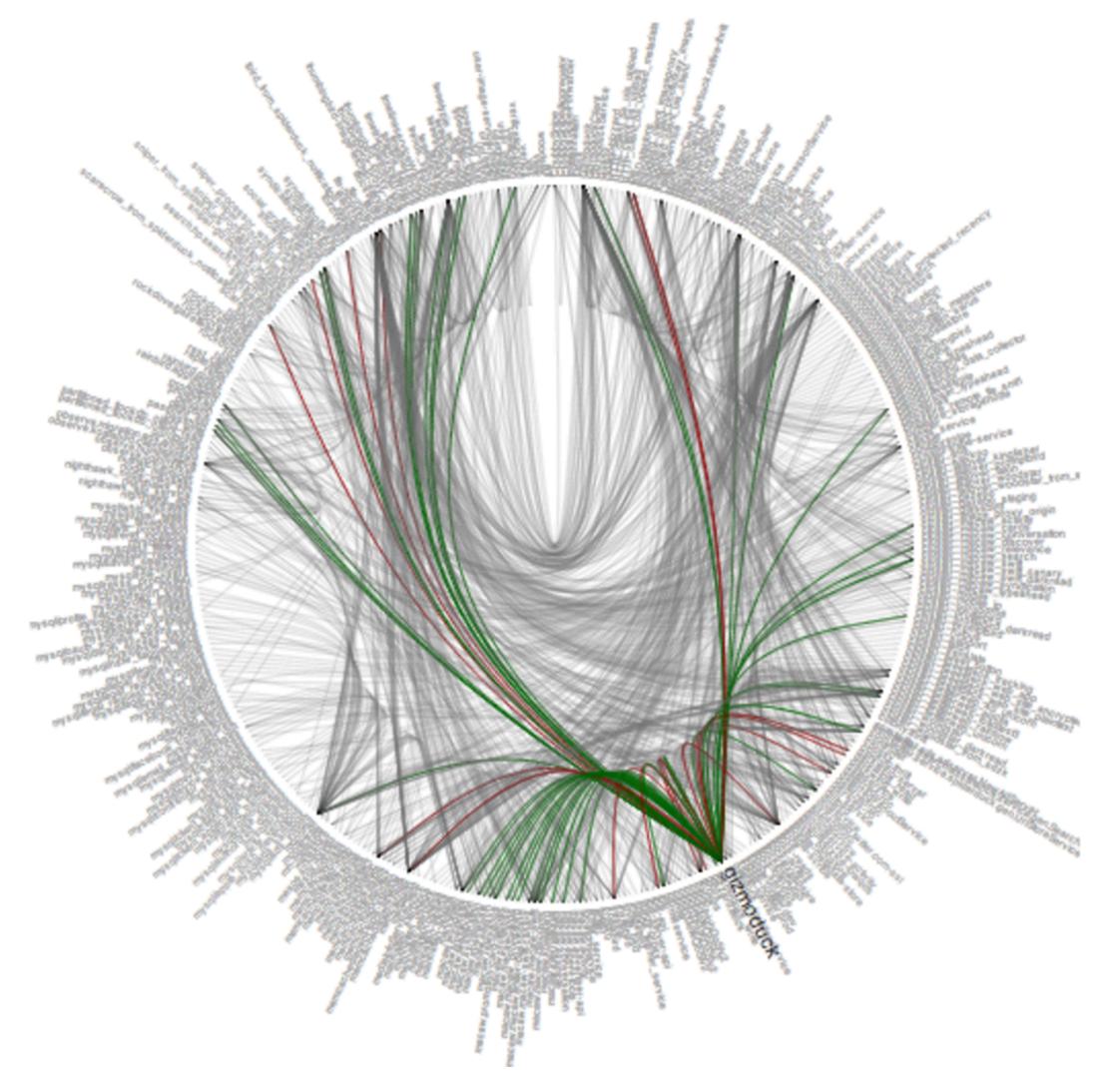
Chained dependencies



Q: A's requests time out, so do B's, which on-call engineer(s) should be paged?

Oftentimes modern application architecture is a mesh

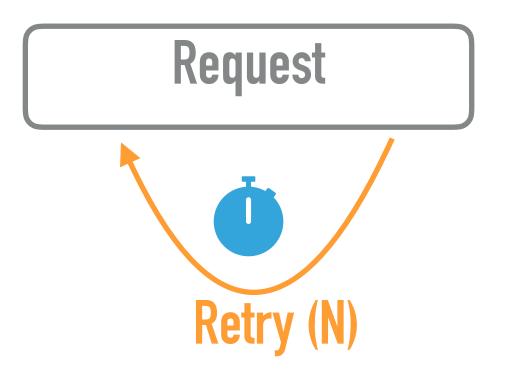




correlation! = causality

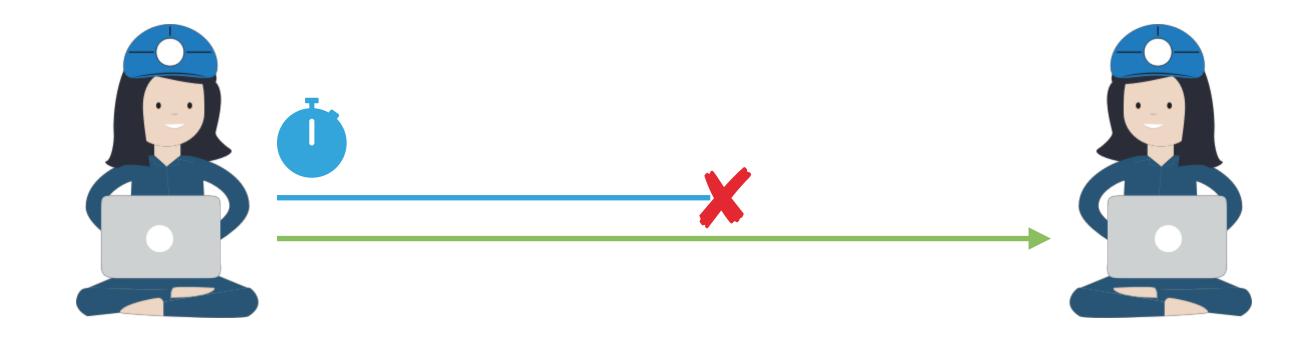
timeouts!= causality

Retries carry risks



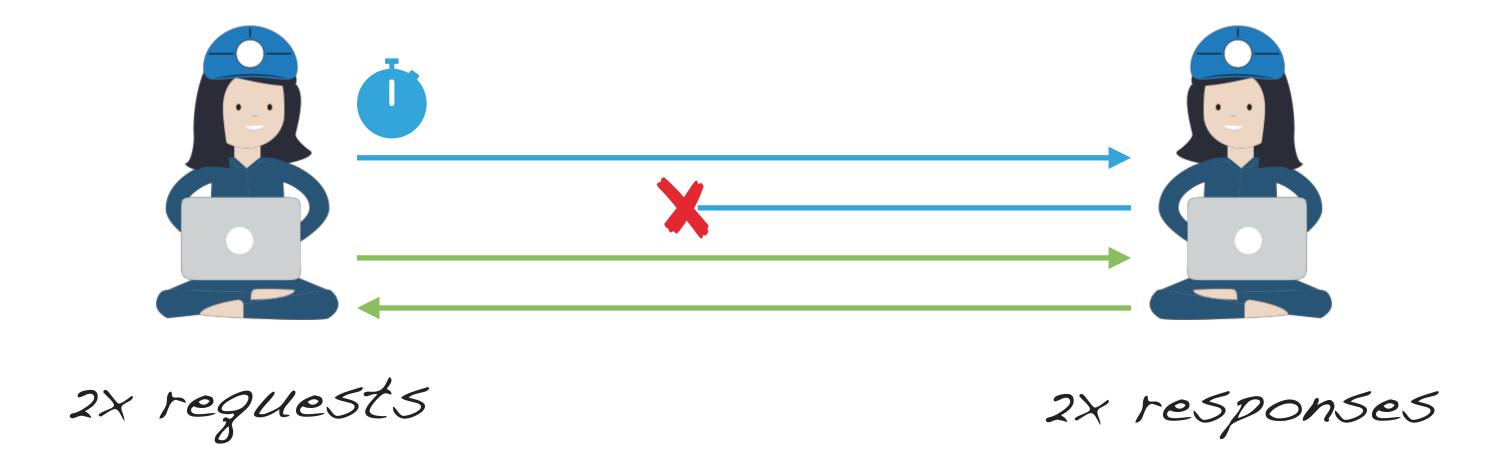


Retry is extra work



2x requests

Retry is extra work



Retry!= Replay

- Potential behavioral change
- State change
- Hidden retries in lower stack
 - e.g. TCP

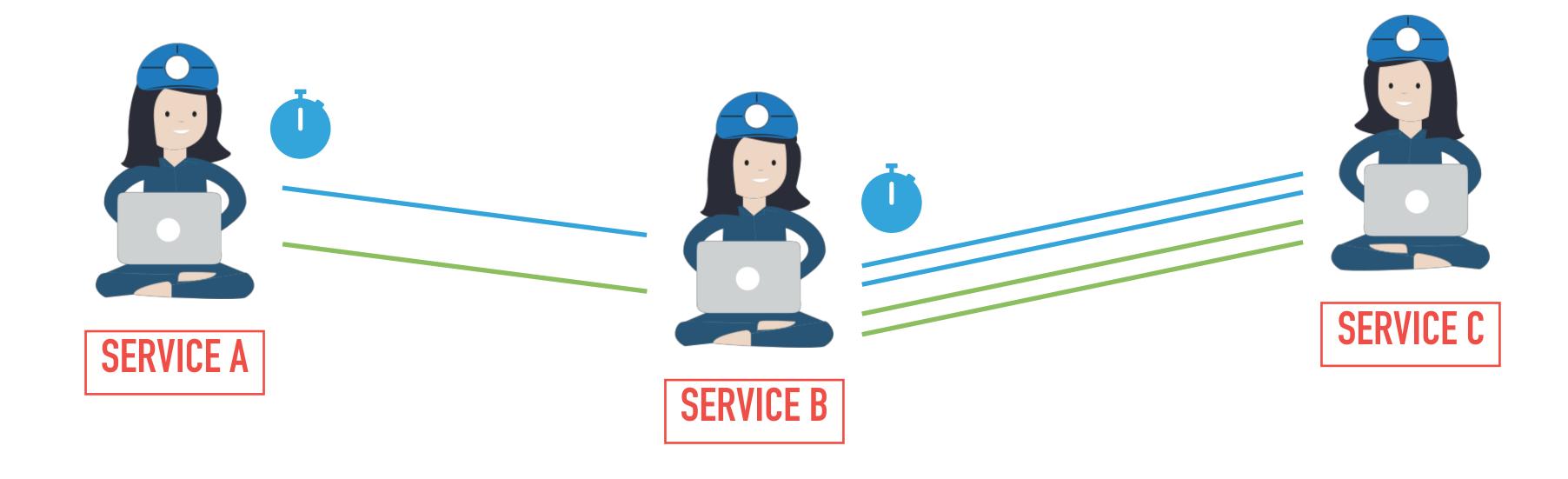
Other retry key decisions

- How many times?
- How to set timeout for each retry?
- How much delay between retries?
- Where to send it?

Inconvenient truths about retry

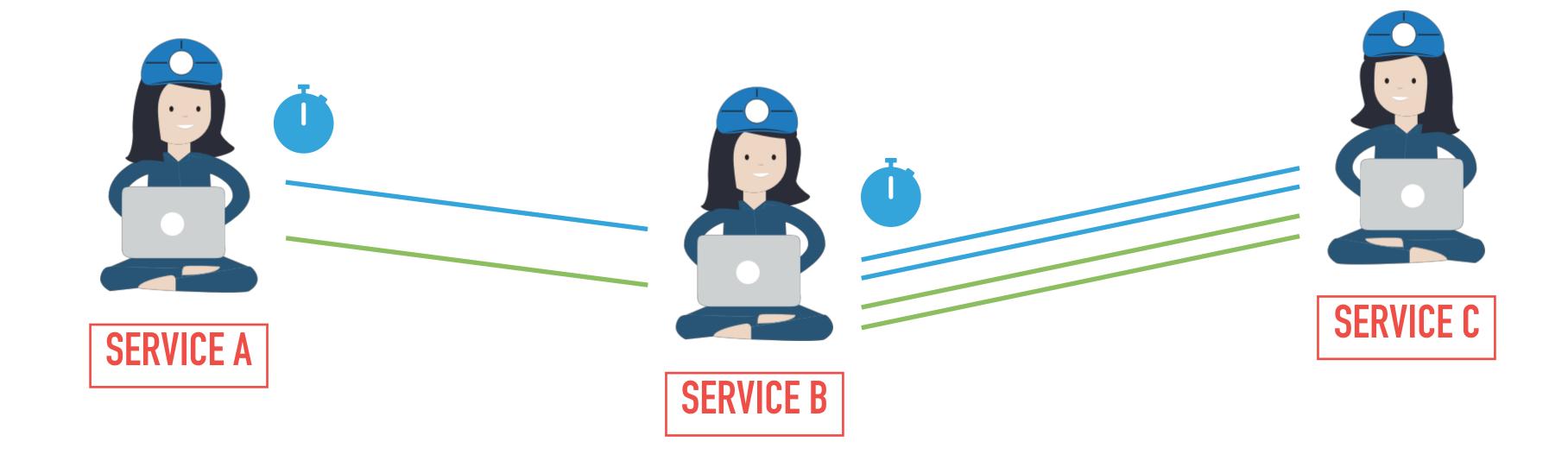
- Retries can create positive feedback loop
- Retries can change system state and behavior
- Sophisticated retry configuration has many knobs

Chained dependencies



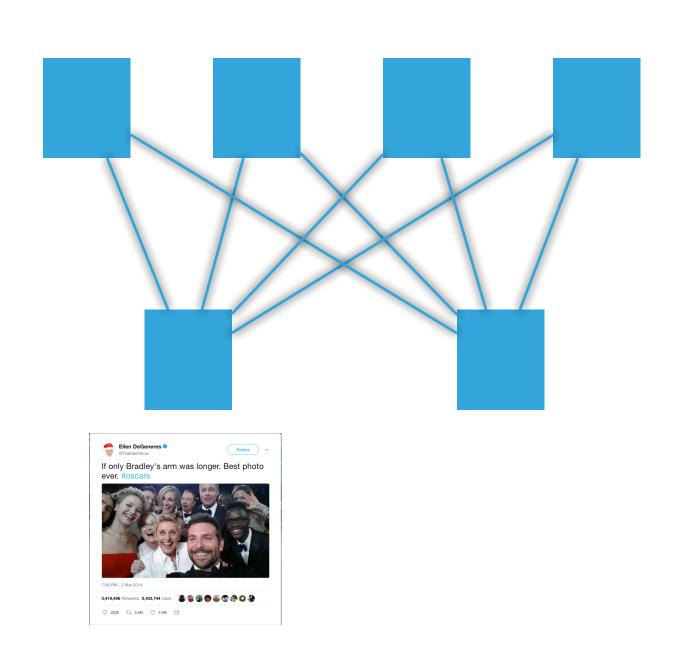
If C is slow...

Chained dependencies



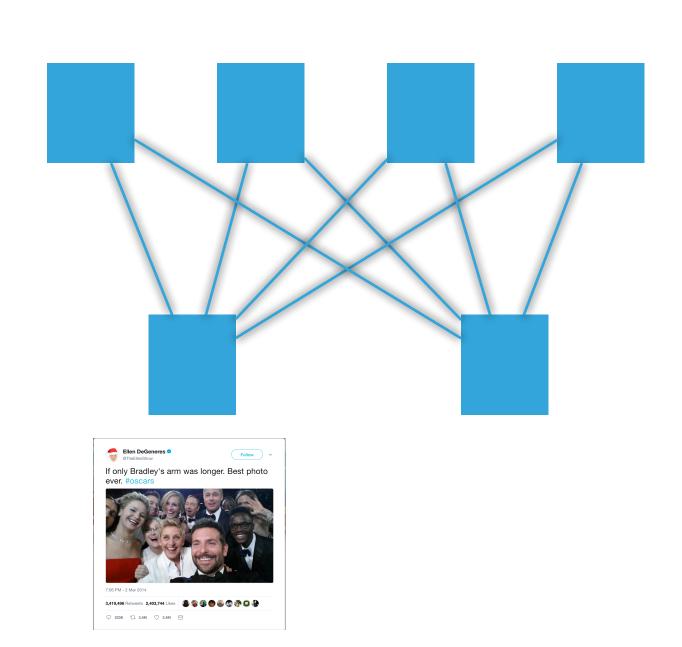
If B is slow...

Retry at scale- "How to DDoS A cache?"



- ► Timeout ⇒ connection teardown
- Large number of clients
- Stateful backend, fixed route
- Full connectivity mesh
- Variable connections per route
- Container-based deploy

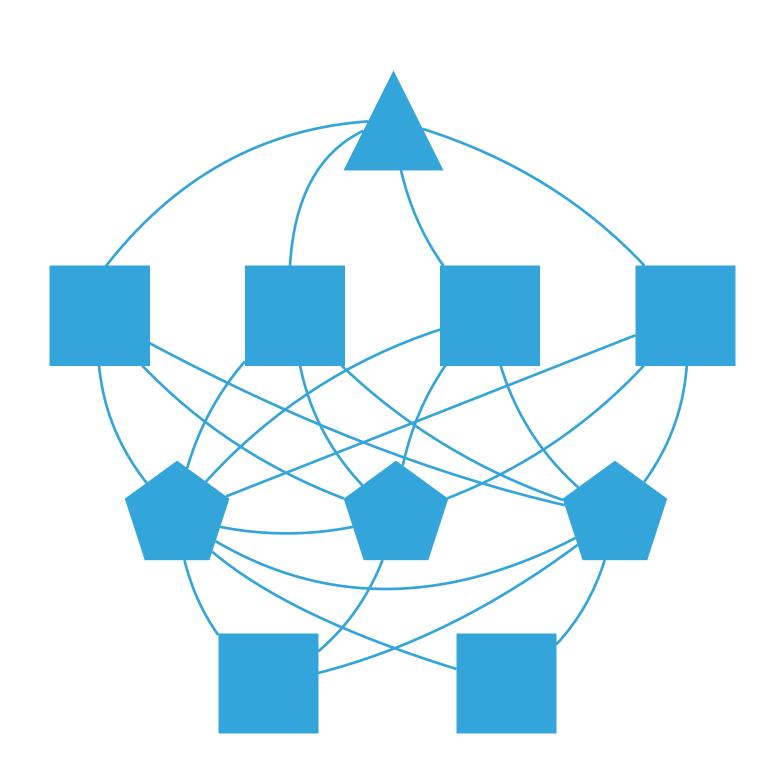
Retry at scale- "How to DDoS A cache?"



Transient hotspot

- ⇒ initial timeouts (closing connections)
- ⇒ connection storm (fixed route)
- ⇒ massive timeouts (full mesh)
- ⇒ bigger storm (multiple, variable connections)
- ⇒ NIC saturation / backend offline
- ⇒ frontend failing, site down

Service mesh can make retries significantly worse



- Top-level retries affect whole system
- Bigger multipliers toward the bottom
- Hard to predict bottleneck

Problems with naive failure coping mechanisms

Wrong implicit assumption, single input, single feedback loop

How to improve?

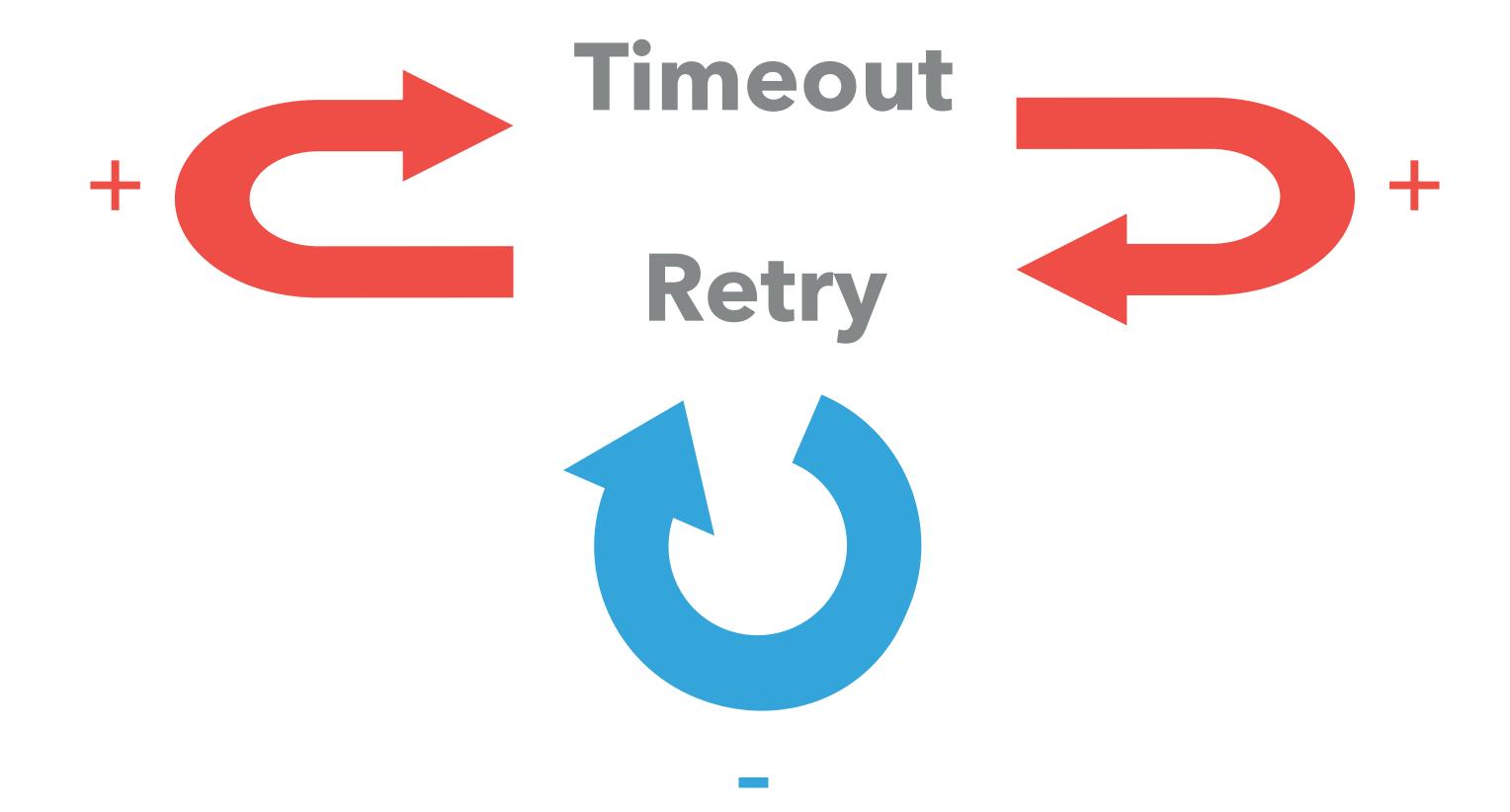
Address the positive feedback loop



Slow down, retry less



Introduce other feedback loop



Good timeouts

- Minimize false positives caused by local stack
- Attempt to root cause with other signals

Good retries

- Customize for purpose of request
- Be as conservative as possible

- Circuit breakers for timeout induced retries
 - State-based rules
 - Back-pressure

Manage impact across services

- Enforce top-down budget
- Apply and act on back pressure

Using micro services? Your task just got much harder...

Tracing helps

drill, baby, drill

- Test common failure scenarios
- Test failures at different parts
- Test partial failures
- ▶ Test relaxed timeout/retry/escalation config

Example: a typical cache config

(Guide w/ examples, 2K+ words excl. code)

- Overall timeout: 500ms*
- Request timeout: 150ms*
- Connect timeout: 200ms
- Pipelining: request timeout does not reset connection
- ▶ Read retry: 2 tries*, no write-back, no backoff
- Write retry: 3 tries*, random backoff (5-200ms)
- Overall retry budget: 20% of requests, minimum 10 retries per second (10-sec credit window)
- ▶ Blackhole: 5 consecutive failures*, revive after 30 seconds
- Centralized topology manager, changes dampened >1min

CACHE SLO: P999 <5MS

Making timeout/retry easier

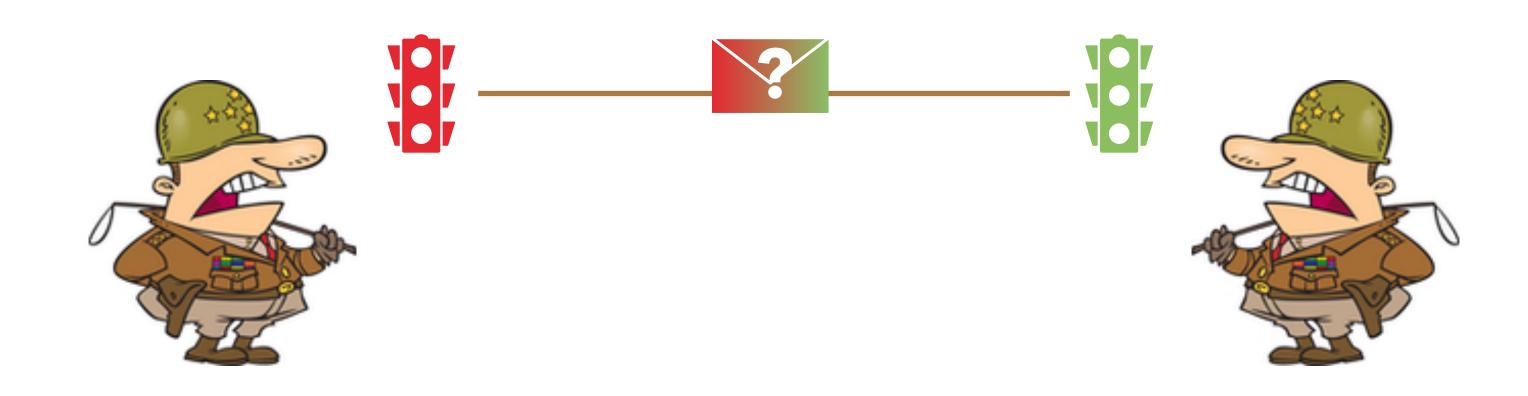
Reduce variance

Example: Making cache more predictable

- Staying off shared hosts
- Rate limiting new connections with iptables
- Setting CPU affinity for sirq, application

- Removing shared locks
- All blocking IO on background threads

There's no magic-Important to know when to give up



Understand common anomalies, Break the loop, Form a bigger picture,

and you'll be happy most of the time.

Questions?