# Systems in One Lesson

@tlberglund

## What are Distributed Systems?

Any system too large to fit on one computer.

## What are Distributed Systems?

- Storage
- Transactions
- Computation
- Coordination

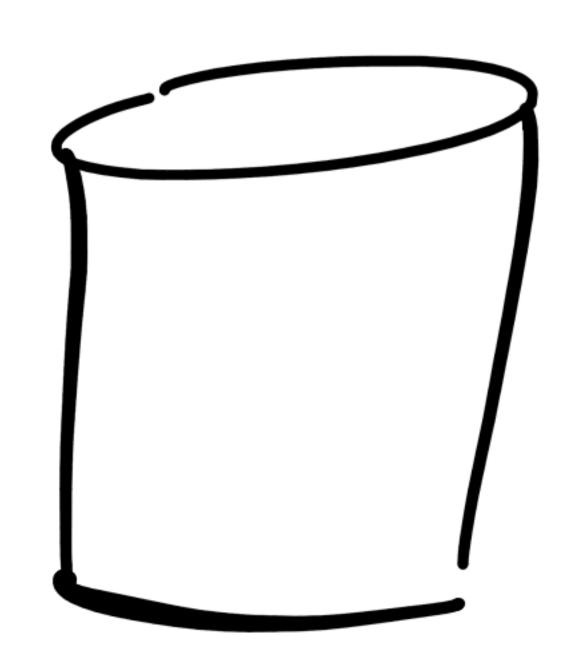
# Distributed Storage

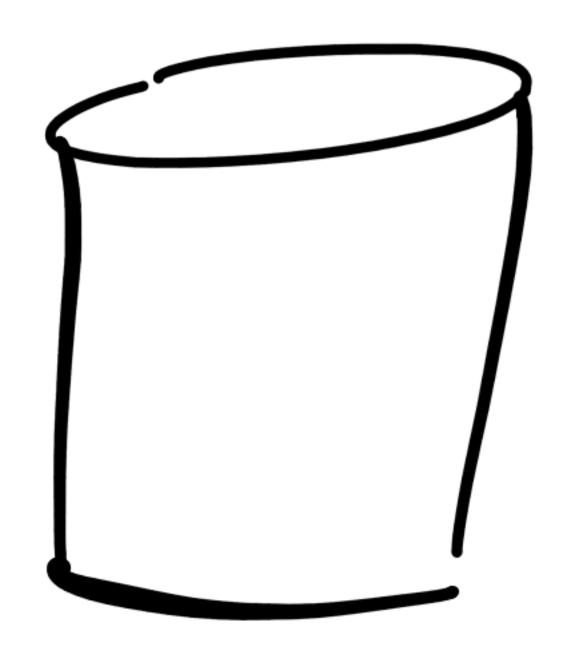
## Distributed Storage

- Read replication
- Sharding
- Consistent hashing
- Distributed filesystems

## Distributed Storage

- MongoDB
- Cassandra
- HDFS

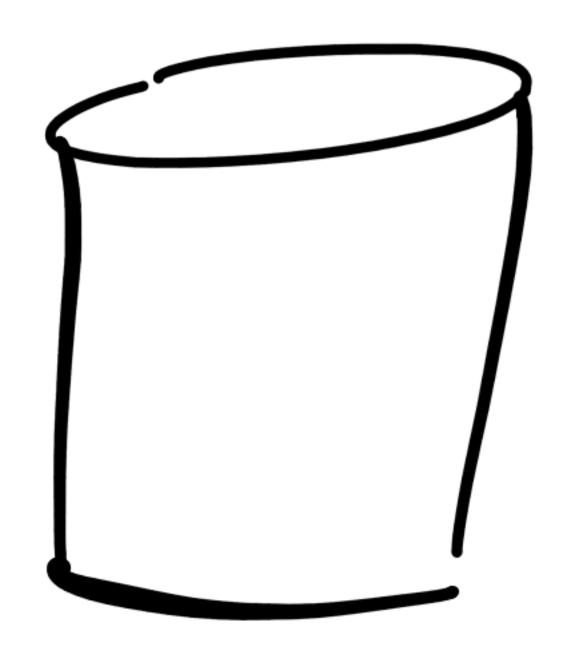


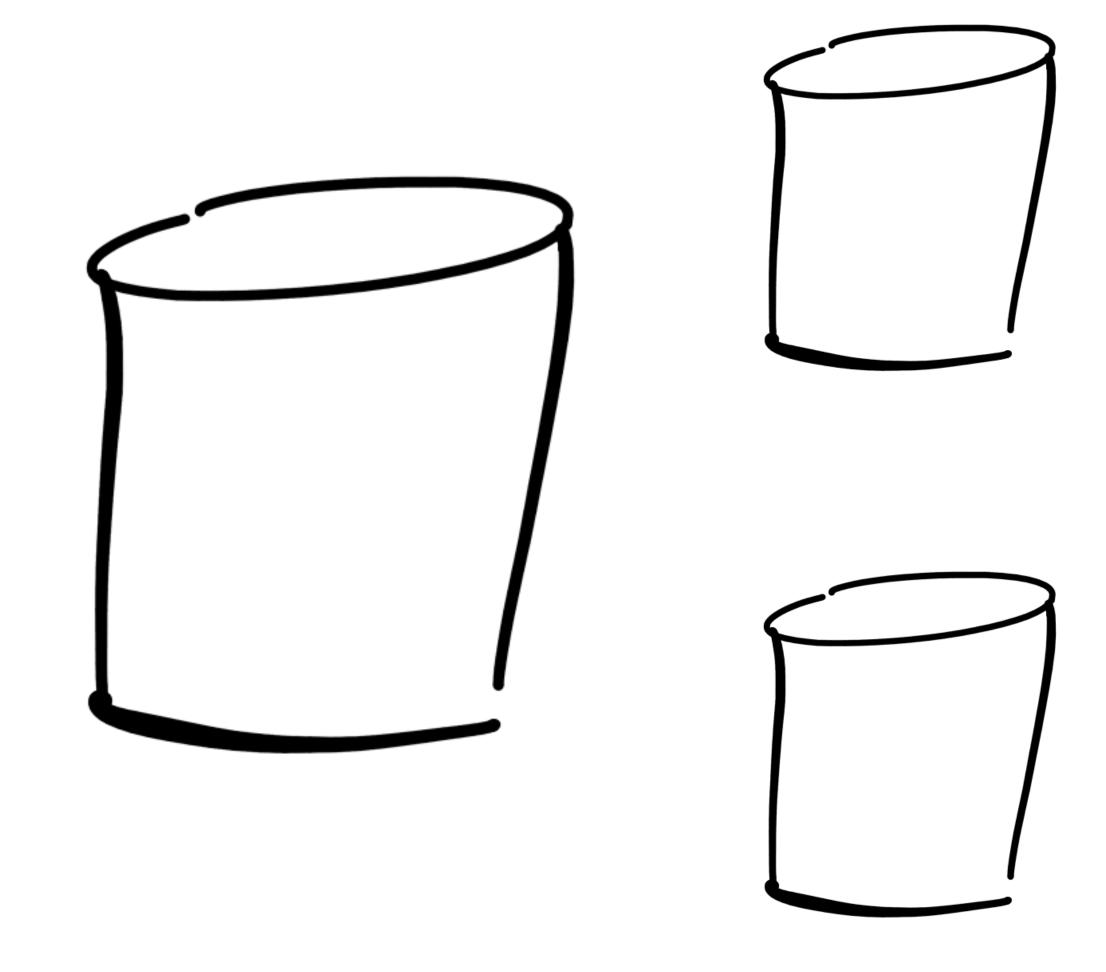


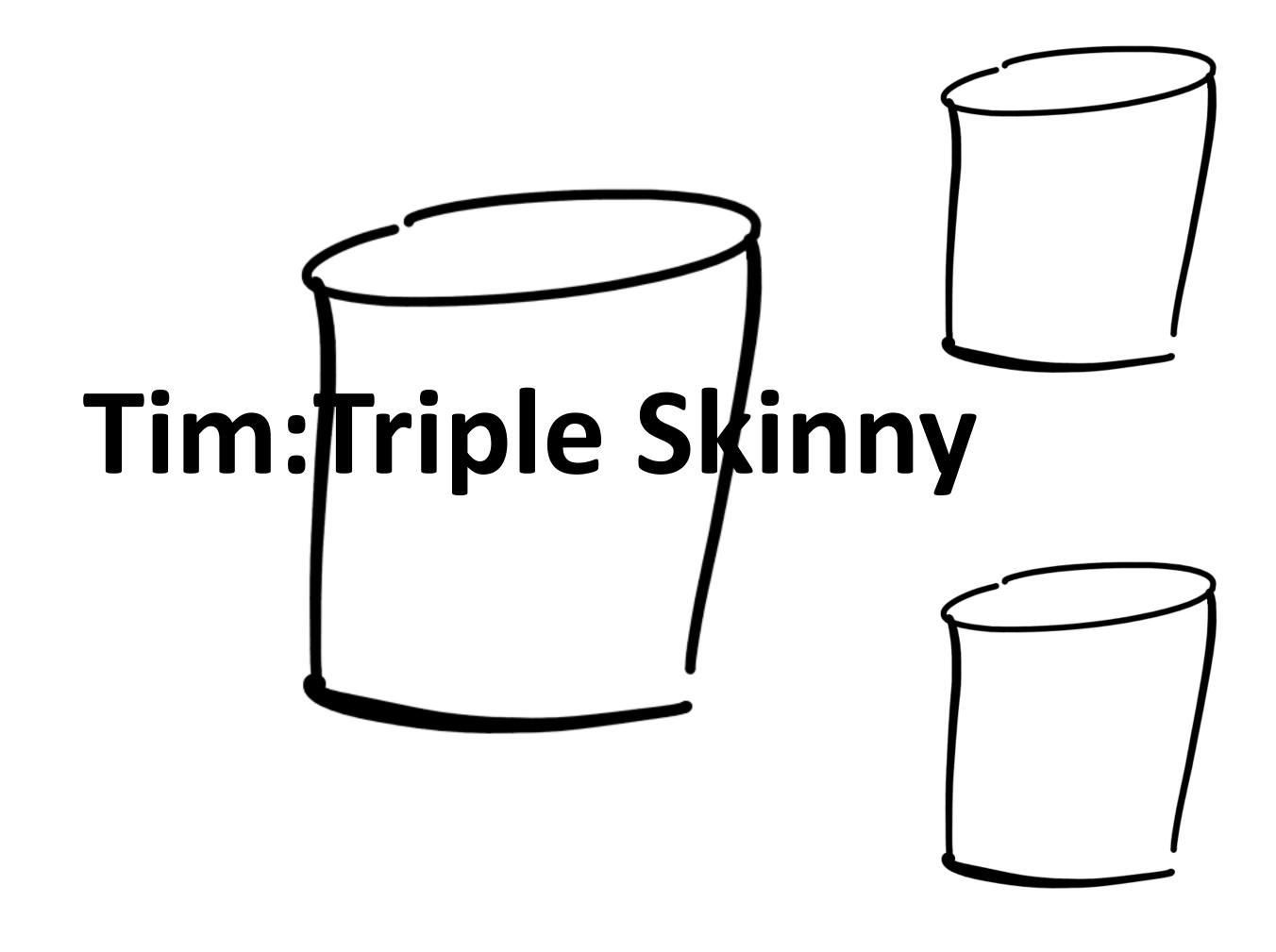


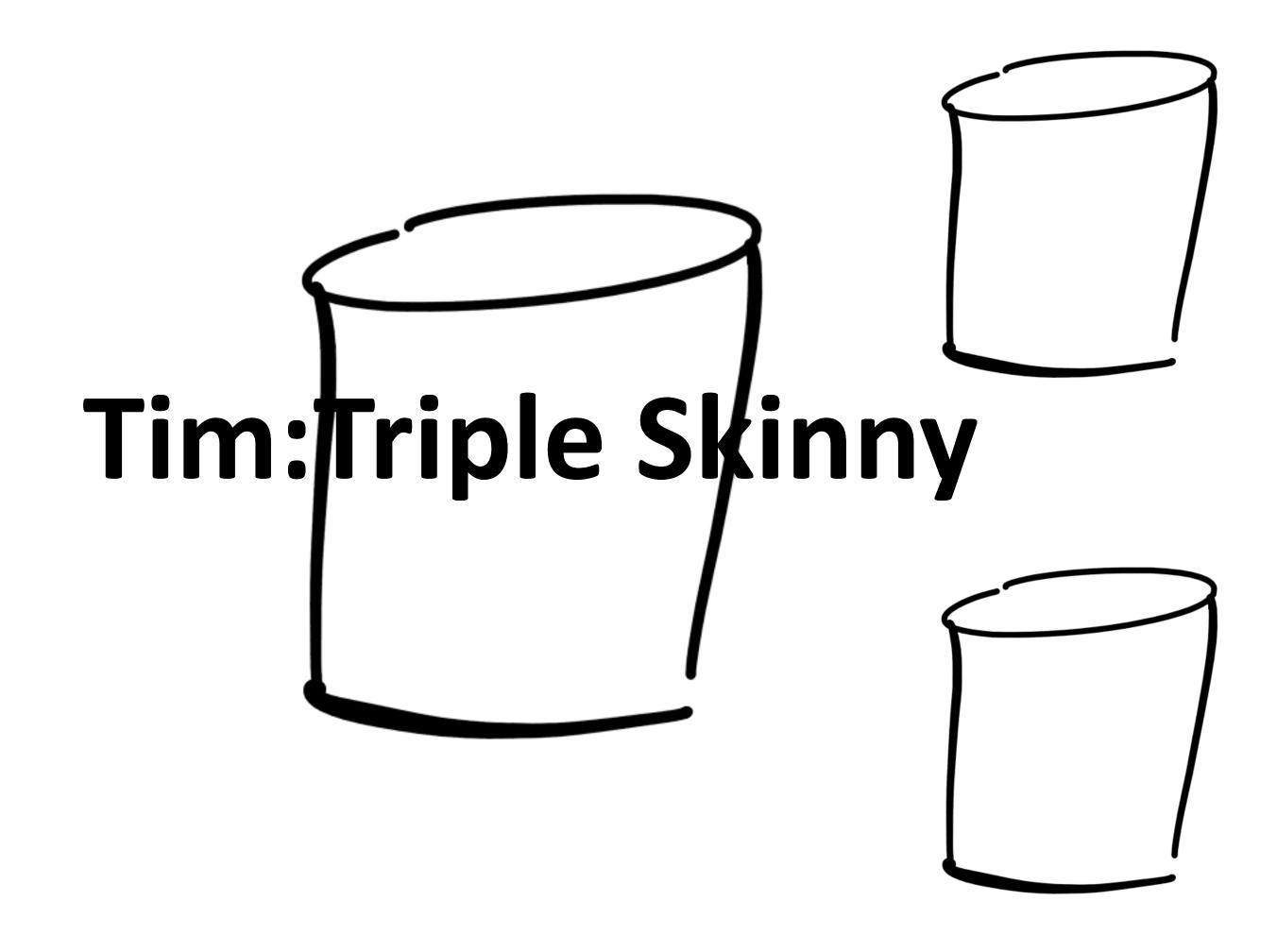
Tim?

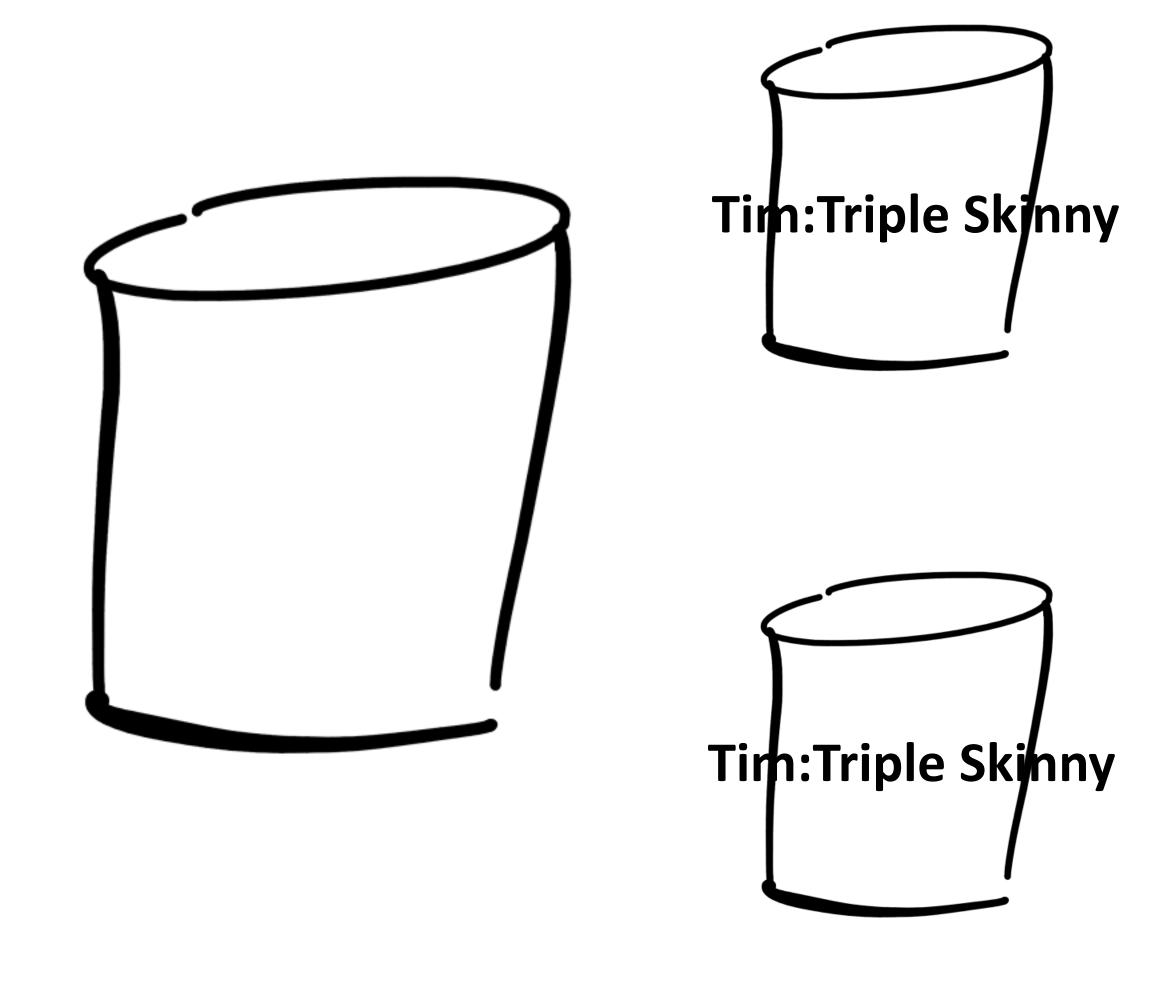




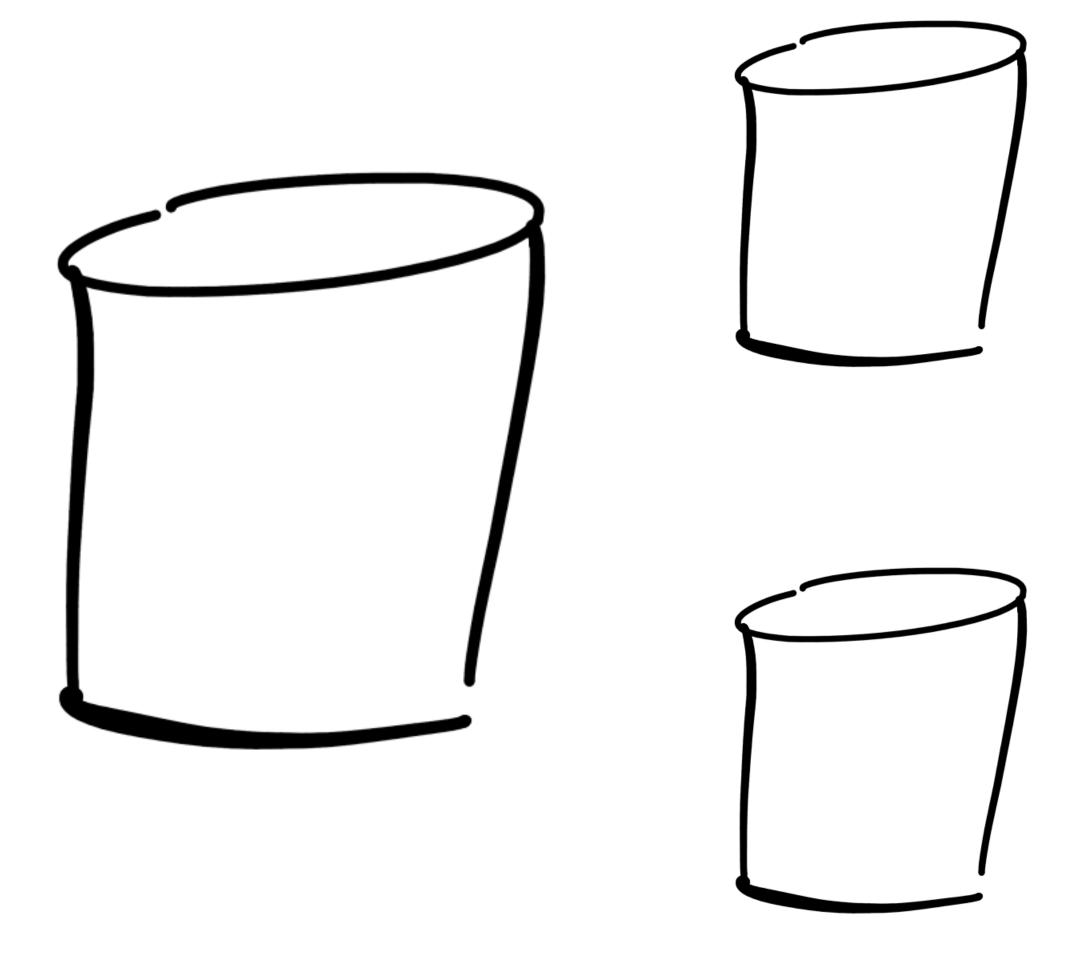








Tim?

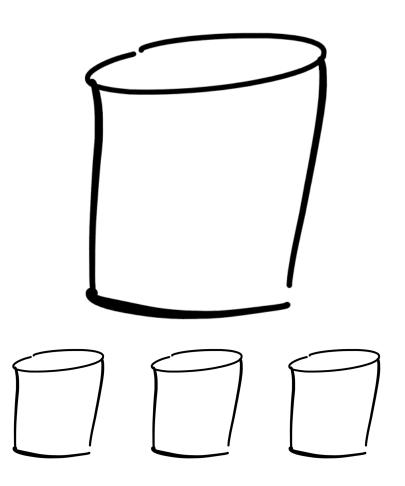


## Replication Problems

- Complexity
- Consistency

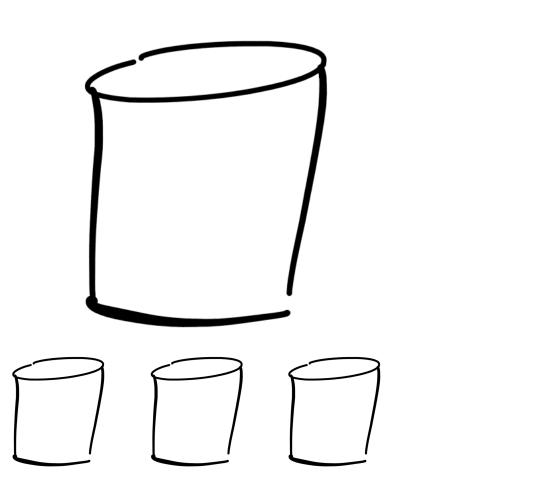
## When Life is REALLY busy

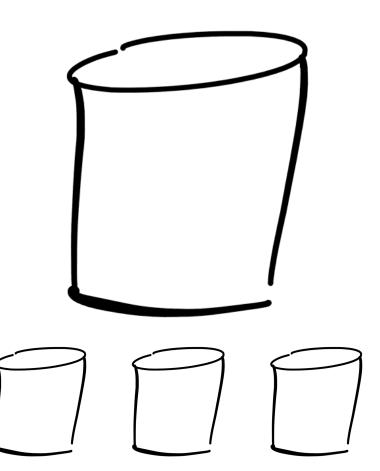
#### Aaron-Faye



## When Life is REALLY busy

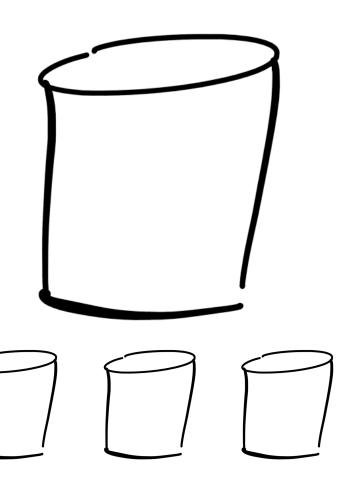
Aaron-Faye Faye-Nancy



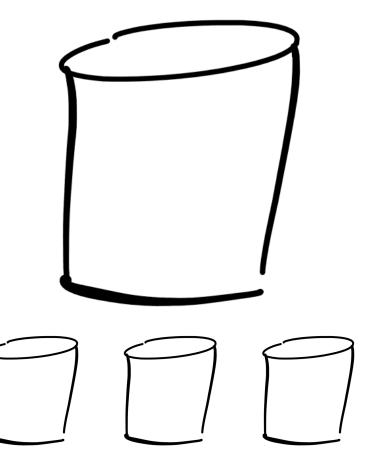


## When Life is REALLY busy

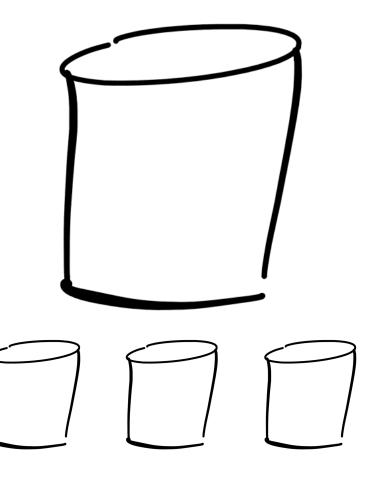
Aaron-Faye



Faye-Nancy

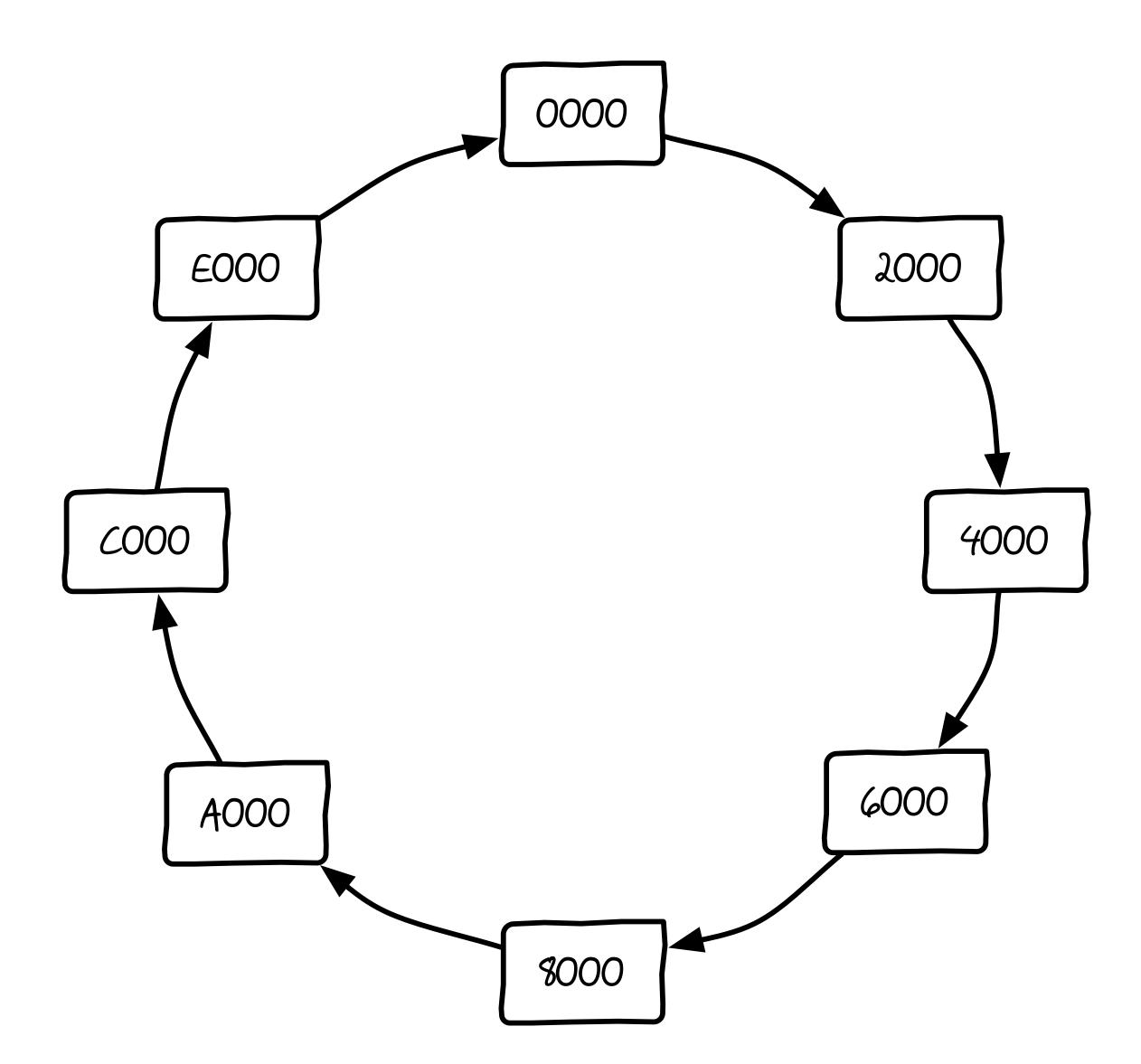


Nancy-Zed

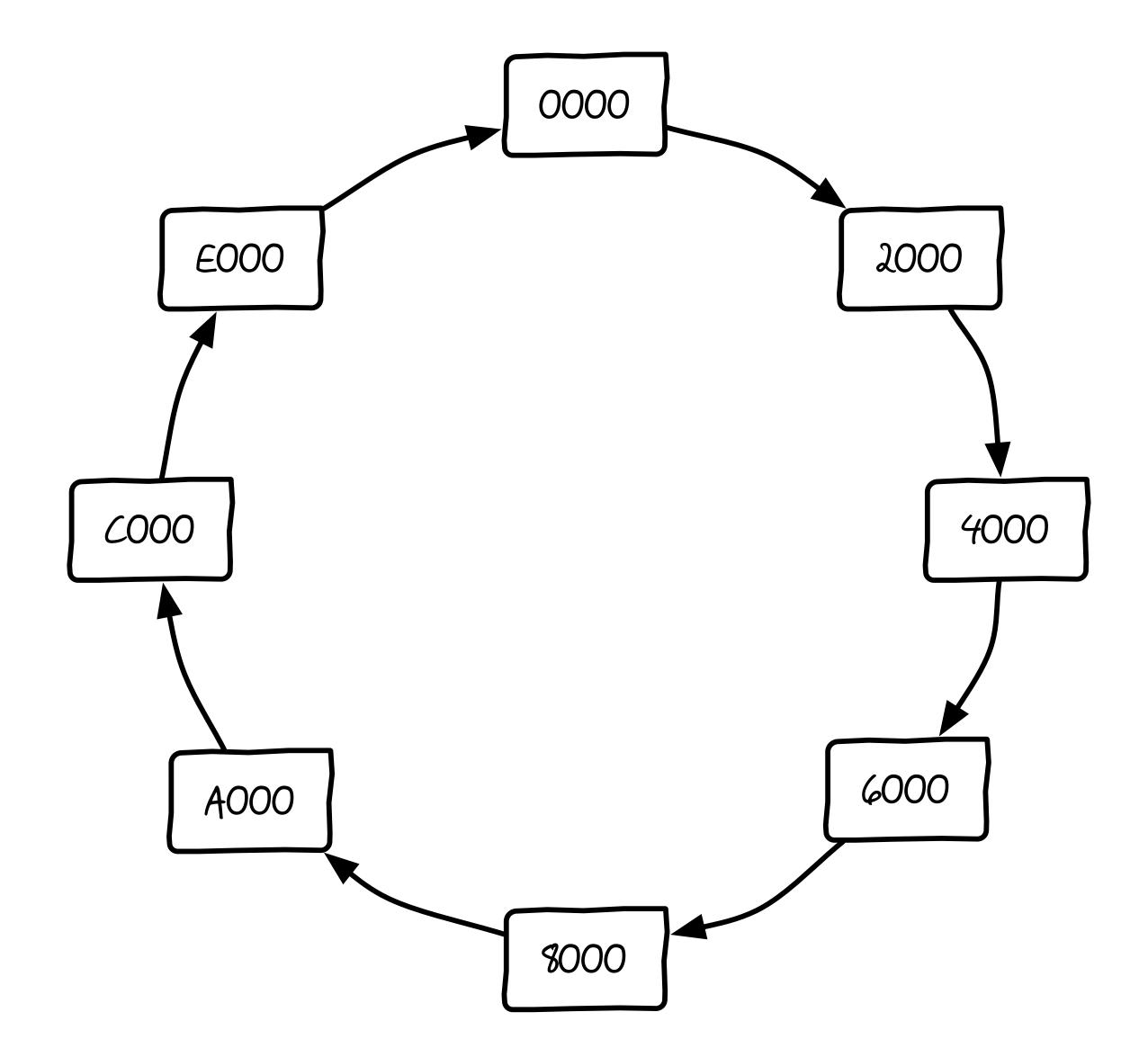


## Sharding Problems

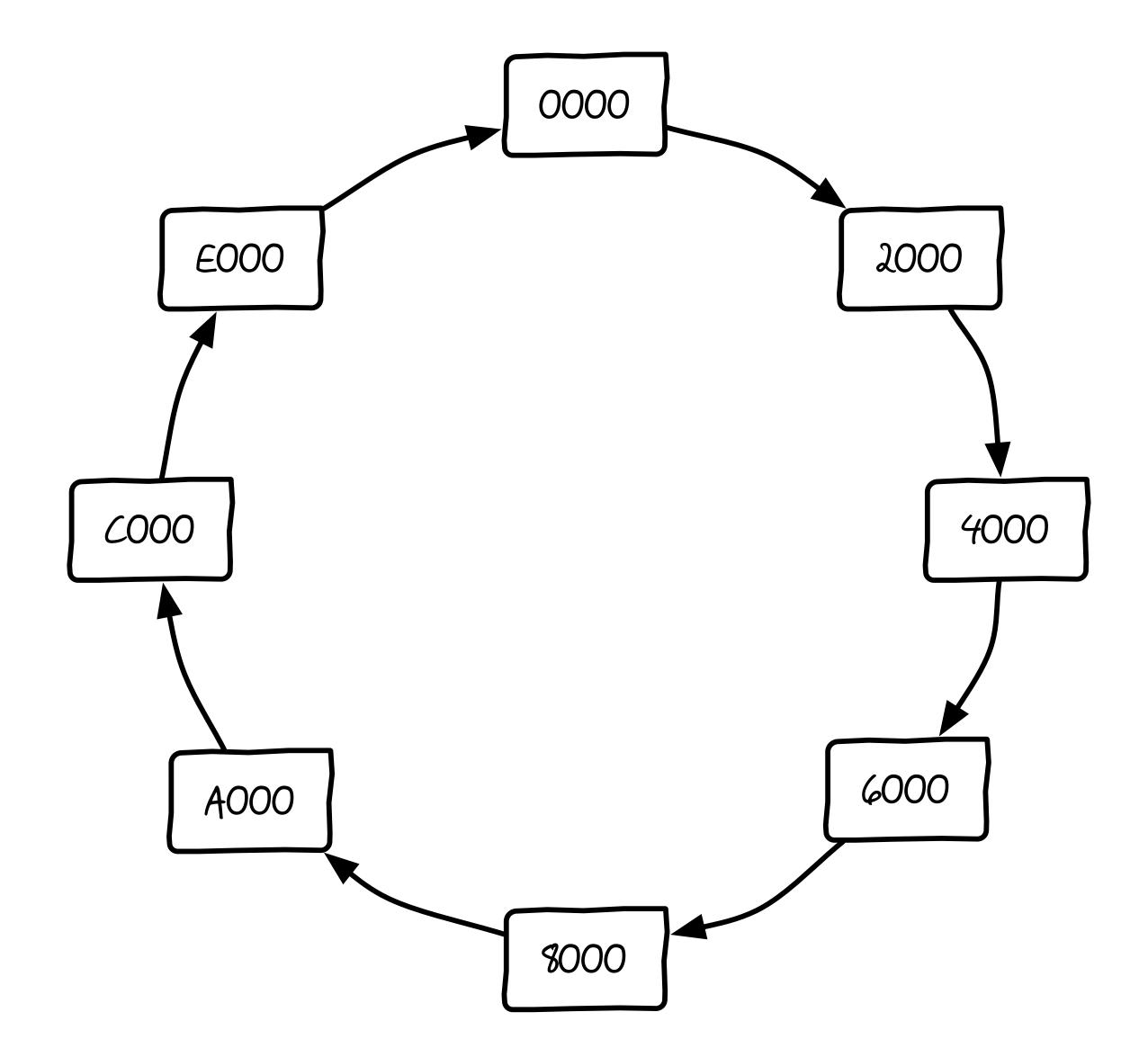
- More Complexity
- Broken data model
- Limited data access patterns



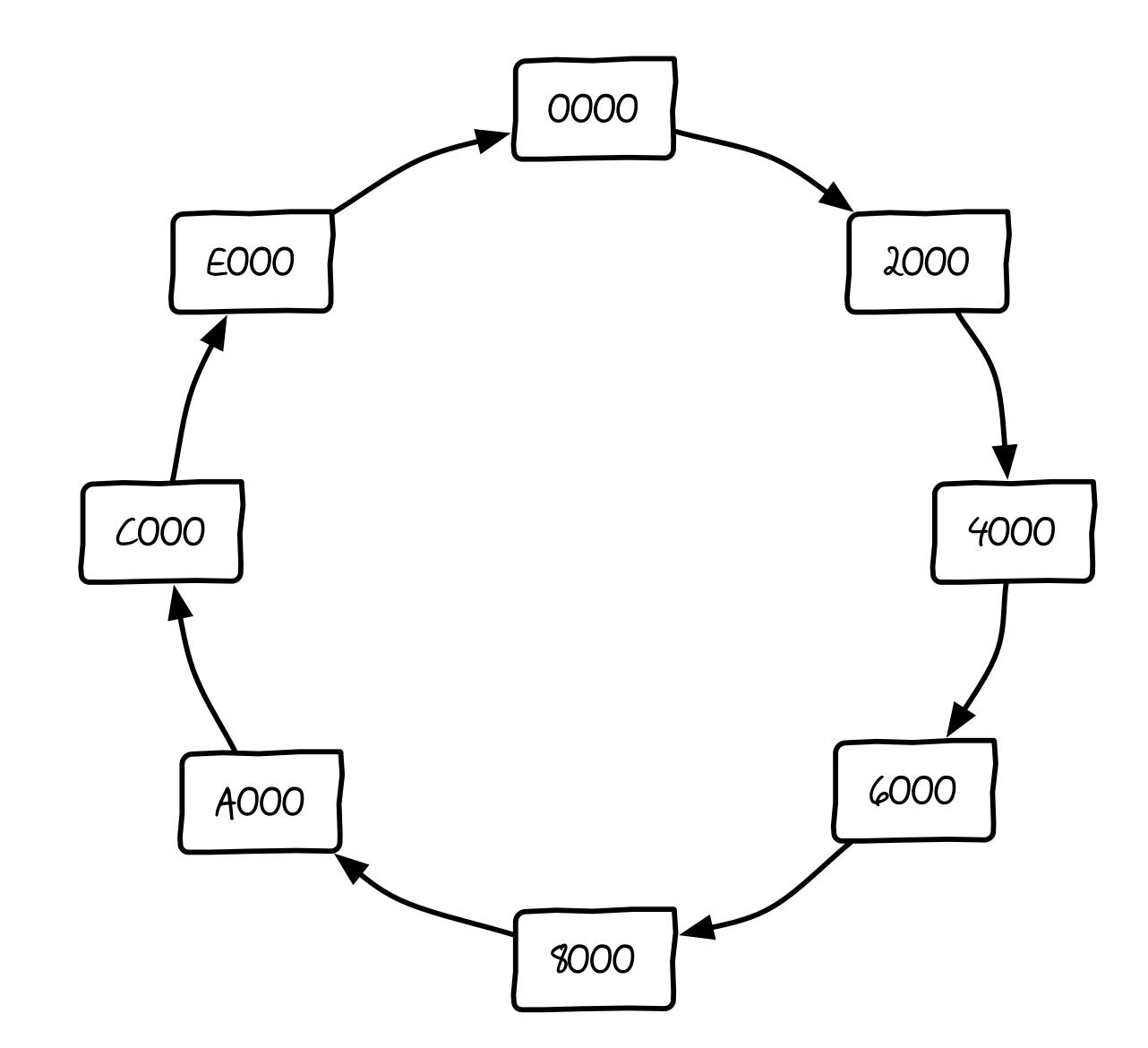
Tim:Americano



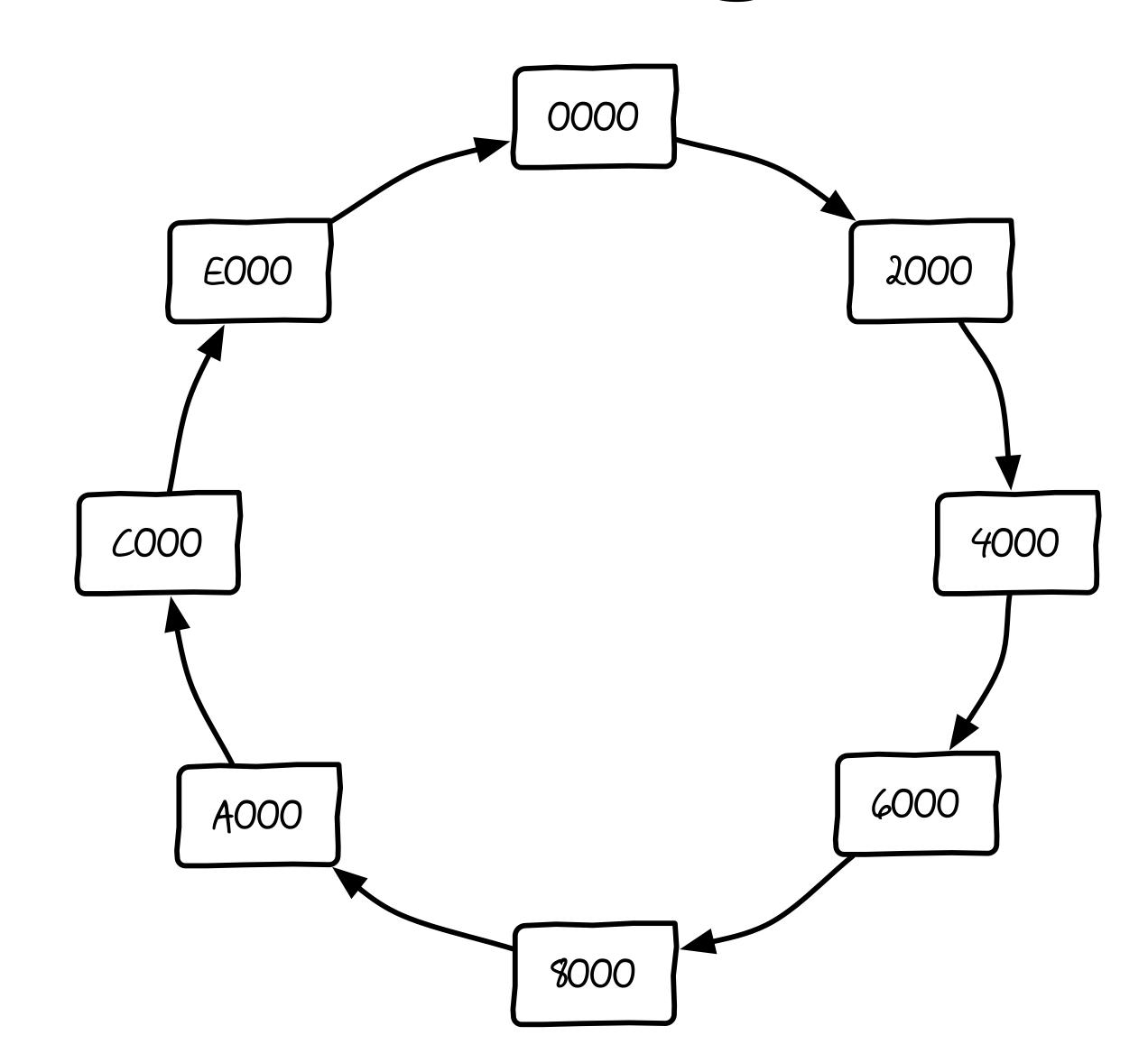
9F72:Americano

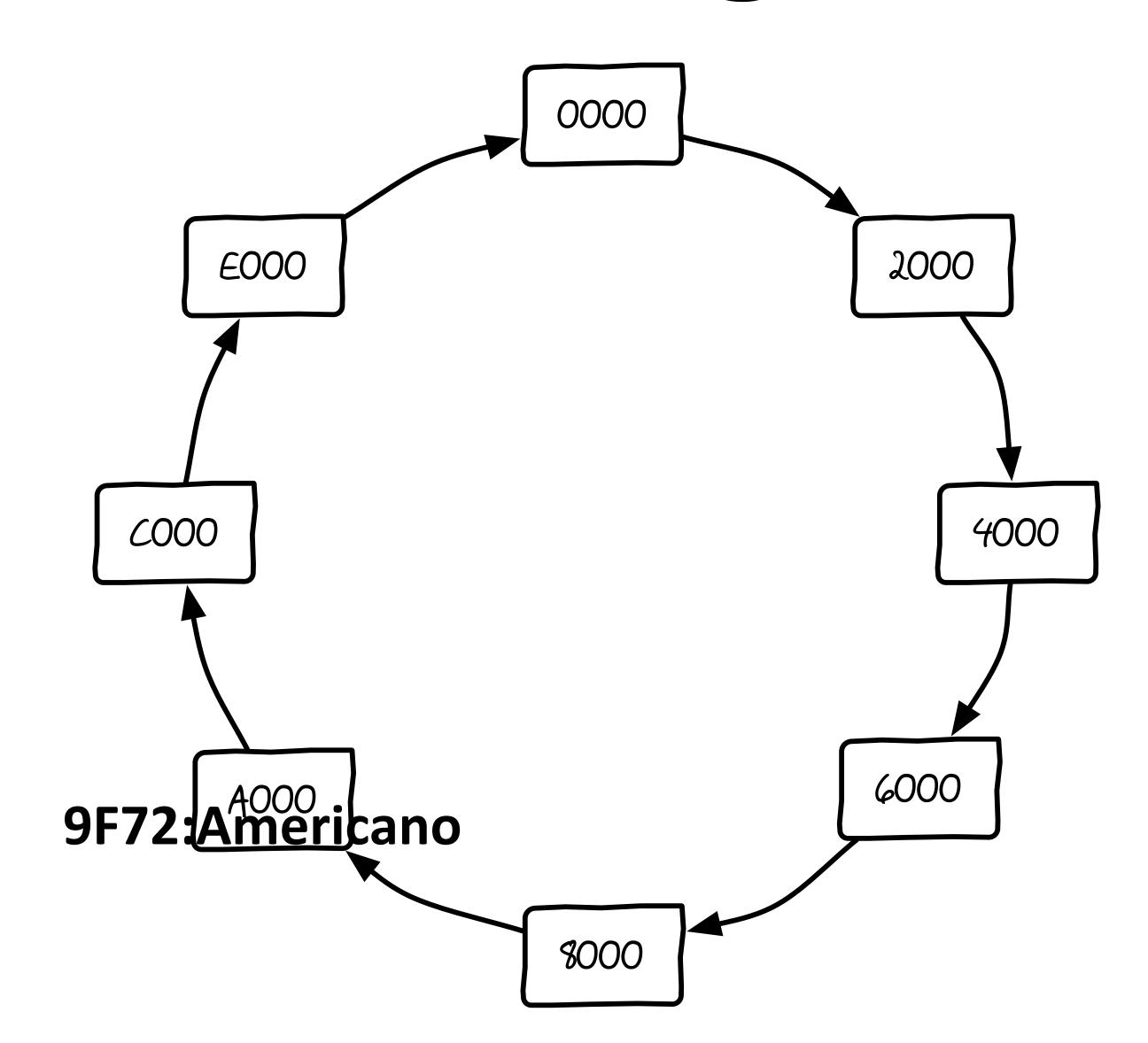


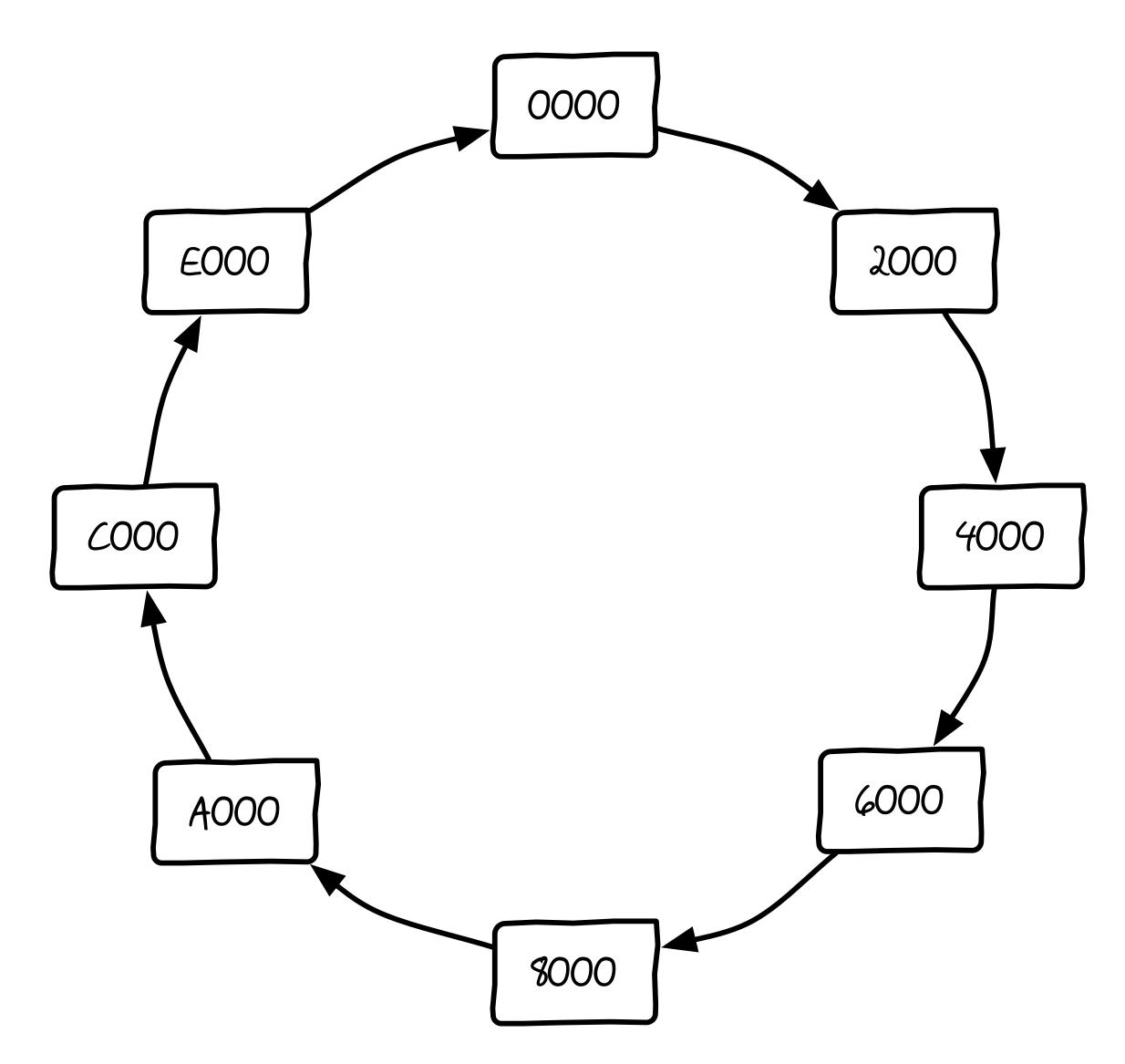
Tim?

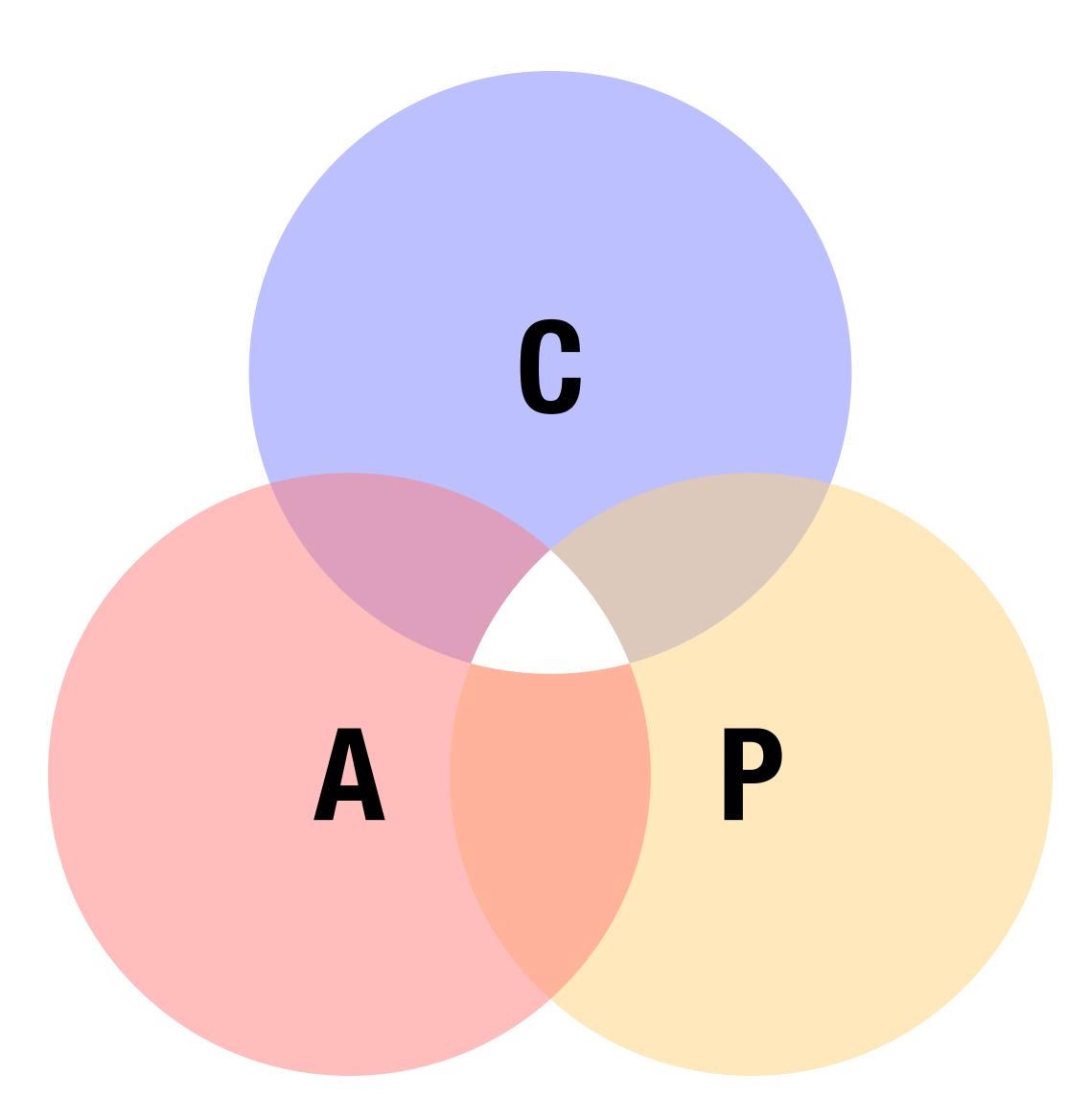


9F72?



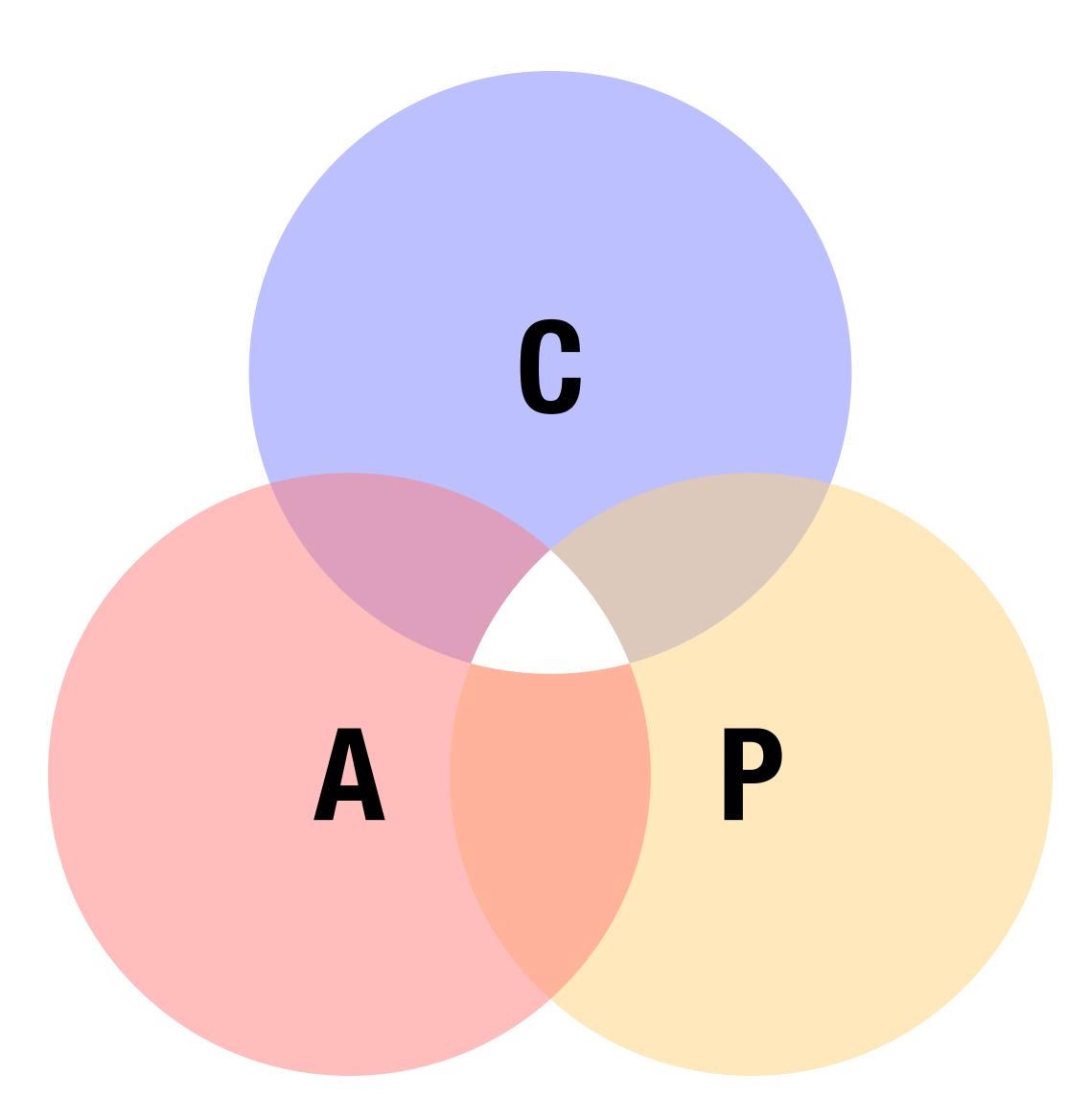






- Consistency
- Availability
- Partition Tolerance

- Shared writing project
- Coffee shop closes
- Synchronizing over the phone
- Battery dies
- Status report!



## Distributed Transactions

## Distributed Transactions



#### Starbucks Does Not Use Two-Phase Commit



November 19, 2004 Hotto Cocoa o Kudasai

I just returned from a 2 week trip to Japan. One of the more familiar sights was the ridiculous number of Starbucks  $(\mathcal{A}\mathcal{P}-\mathcal{N}\mathcal{P}\mathcal{A})$  coffee shops, especially around Shinjuku and Roppongi. While waiting for my "Hotto Cocoa" I started to think about how Starbucks processes drink orders. Starbucks, like most other businesses is primarily interested in maximizing throughput of orders. More orders equals more revenue. As a result they use asynchronous processing. When you place your order the cashier marks a coffee cup with your order and places it into the queue. The queue is quite literally a queue of coffee cups lined up on top of the espresso machine. This queue decouples cashier and barista and allows the cashier to keep taking orders even if the barista is backed up for a moment. It allows them to deploy multiple baristas in a Competing Consumer scenario if the store gets busy.

#### Correlation

By taking advantage of an asynchronous approach Starbucks also has to deal with the same challenges that asynchrony inherently brings. Take for example, correlation. Drink orders are not necessarily completed in the order they were placed. This can happen for two reasons. First, multiple baristas may be processing orders using different equipment. Blended drinks may take longer than a drip coffee. Second, baristas may make multiple drinks in one batch to optimize processing time. As a result, Starbucks has a correlation problem. Drinks are delivered out of sequence and need to be matched up to the correct customer. Starbucks solves the problem with the same "pattern" we use in messaging architectures -- they use a Correlation Identifier. In the US, most Starbucks use an explicit correlation identifier by writing your name on the cup and calling it out when the drink is complete. In other countries, you have to

#### ABOUT ME



Hi, I am Gregor Hohpe, coauthor of the book Enterprise Integration Patterns. I like to work on and write about asynchronous messaging systems, service-oriented architectures, and all sorts of enterprise computing and

#### ACID Transactions

- Atomic
- Consistent
- Isolated
- Durable
- One barista
- Multiple baristas

## Ordering Coffee

- Receive order
- Process payment
- Enqueue order
- Make coffee
- Deliver drink

# Ordering Coffee





- Process paymentEnqueue order

  - Make coffee
  - Deliver drink

## Ordering Coffee

- Why split up order processing?
- What can fail?
- What are the consequences of failure?
- How do we repair failure?

## Ordering Coffee

- How can we design a coffee shop with atomic transactions?
- How does that limit the business?
- Why give up atomicity?

# Distributed Computation

# Distributed Computation

- Scatter/Gather
- MapReduce
- Hadoop
- Spark

#### MapReduce

- All computation in two functions: Map and Reduce
- Keep data (mostly) where it is
- Move compute to data

#### MapReduce

#### MapReduce

#### poems/raven.txt:

Once upon a midnight dreary, while I pondered, weak and weary, Over many a quaint and curious volume of forgotten lore, While I nodded, nearly napping, suddenly there came a tapping, As of some one gently rapping, rapping at my chamber door.

"'Tis some visitor," I muttered, "tapping at my chamber door-Only this, and nothing more."

Ah, distinctly I remember it was in the bleak December, And each separate dying ember wrought its ghost upon the floor. Eagerly I wished the morrow;- vainly I had sought to borrow Clasp a rare and radiant maiden whom the angels name Lenore."

Clasp a rare and radiant maiden whom the angels name Lenore."

"Be that word our sign in parting, bird or fiend," I shrieked, upstarting—
"Get thee back into the tempest and the Night's Plutonian shore!

Leave no blad puberas as kerantial in thy Xou hath spoken!

Leave my loneliness unbroken!- quit the bust above my door!

Take thy beak from out my heart, and take thy form from off my door!"

Quoth the Raven, "Nevermore."

And the Raven, never flitting, still is sitting, still is sitting
On the pallid bust of Pallas just above my chamber door;
And his eyes have all the seeming of a demon's that is dreaming,
And the lamp-light o'er him streaming throws his shadow on the floor;
And my soul from out that shadow that lies floating on the floor
Shall be lifted- nevermore!

# Map Reduce Map

midnight suddenly while gently dreary upon pondered rapping tapping door came my a chamber rapping there tapping Once at a

# MapReduce Map

```
gently:1
                                                  midnight:1
                      upon:1
suddenly:1
                                         dreary:1
          while:1
                           pondered:1
                                         rapping:1
       came:1 a:1
                                                   tapping:1
                    my:1
door:1
       rapping:1 there:1 Once:1 chamber:1
                                             tapping:1
                                                         at:1
 a:1
```

# MapReduce Shuffle

suddenly:1

while:1

midnight:1

door:1

a:1

a:1

chamber:1

Once:1

came:1

rapping:1

rapping:1

gently:1

pondered:1

dreary:1

tapping:1

tapping:1

there:1

upon:1

my:1

at:1

# MapReduce Reduce

```
a:[1,1]
suddenly:[1]
                                                                  there:[1]
                               came:[1]
                                               pondered:[1]
                             rapping:[1,1]
                                                 dreary:[1]
  while:[1]
                                                                  upon:[1]
                                               tapping:[1,1]
midnight:[1]
               chamber:[1]
                                                                   my:[1]
                               gently:[1]
  door:[1]
                                                                   at:[1]
                 Once:[1]
```

# MapReduce Reduce

suddenly:1 a:2 came:1 pondered:1 there:1

while:1 rapping:2 dreary:1 upon:1

midnight:1 chamber:1 tapping:2 my:1

door:1 Once:1 gently:1

#### Real-World MapReduce

- Hadoop
- Cloudera, Hortonworks, MapR
- Nobody write map, reduce functions
- Hive (SQL-like interface)
- Integration with BI front-ends

#### Hadoop

- MapReduce API
- Job Tracker, Task Tracker
- Distributed Filesystem (HDFS)
- Enormous ecosystem

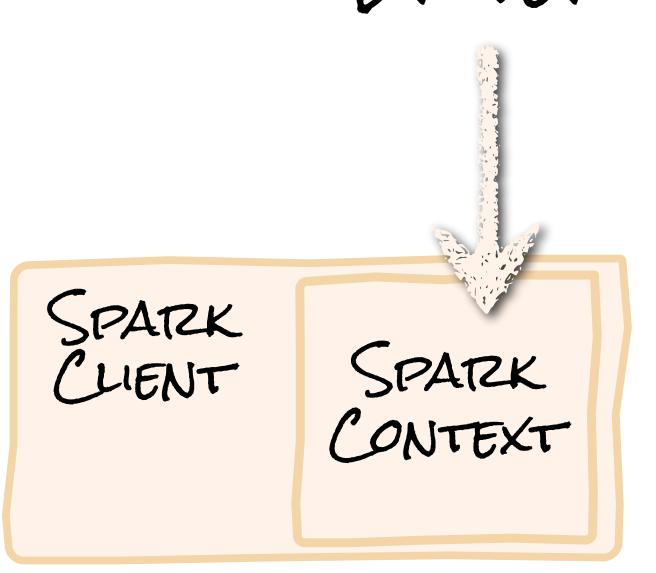
# Spark

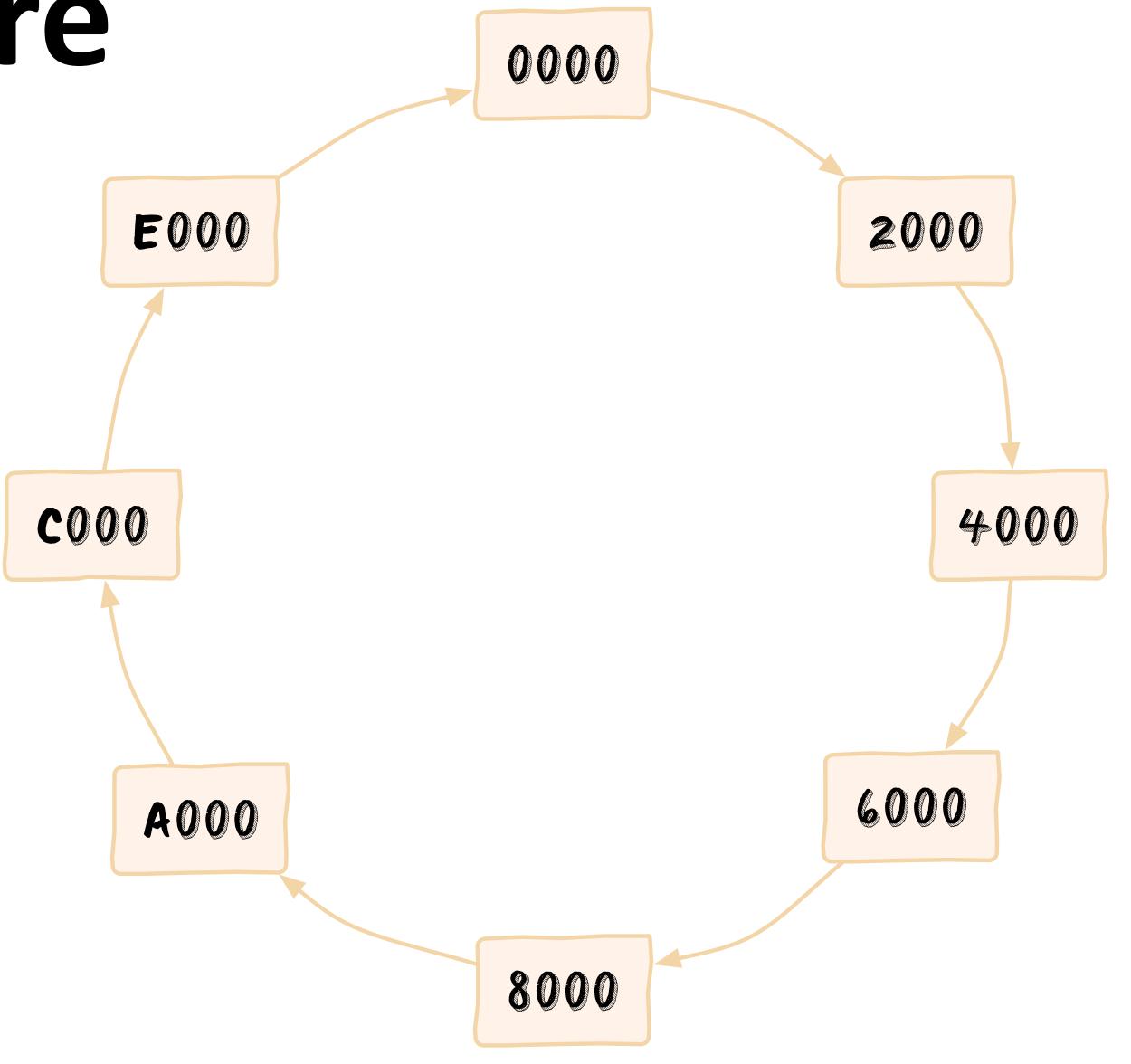
#### Spark

- Scatter/gather paradigm (similar to MapReduce)
- More general data model (RDDs)
- More general programming model (transform/action)
- Storage agnostic

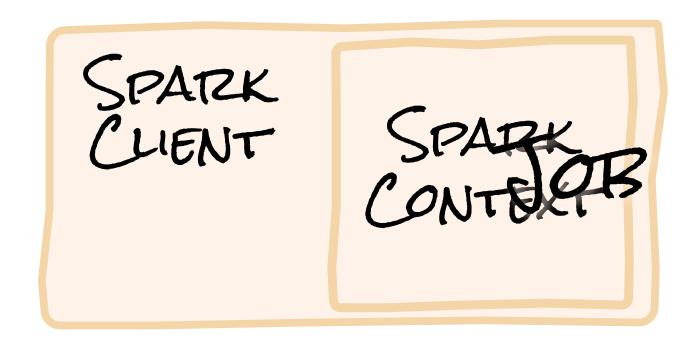
Spark Architecture

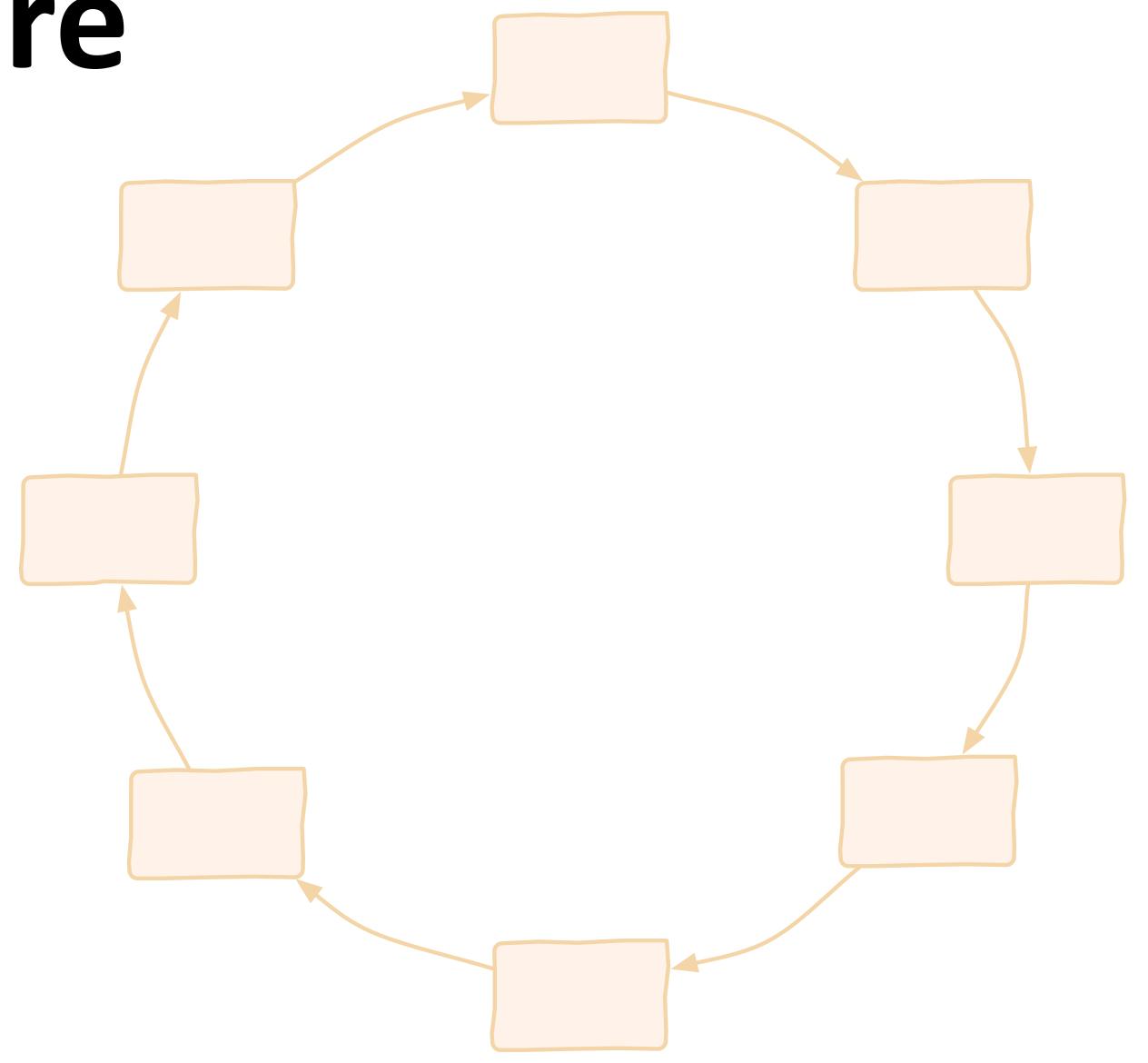
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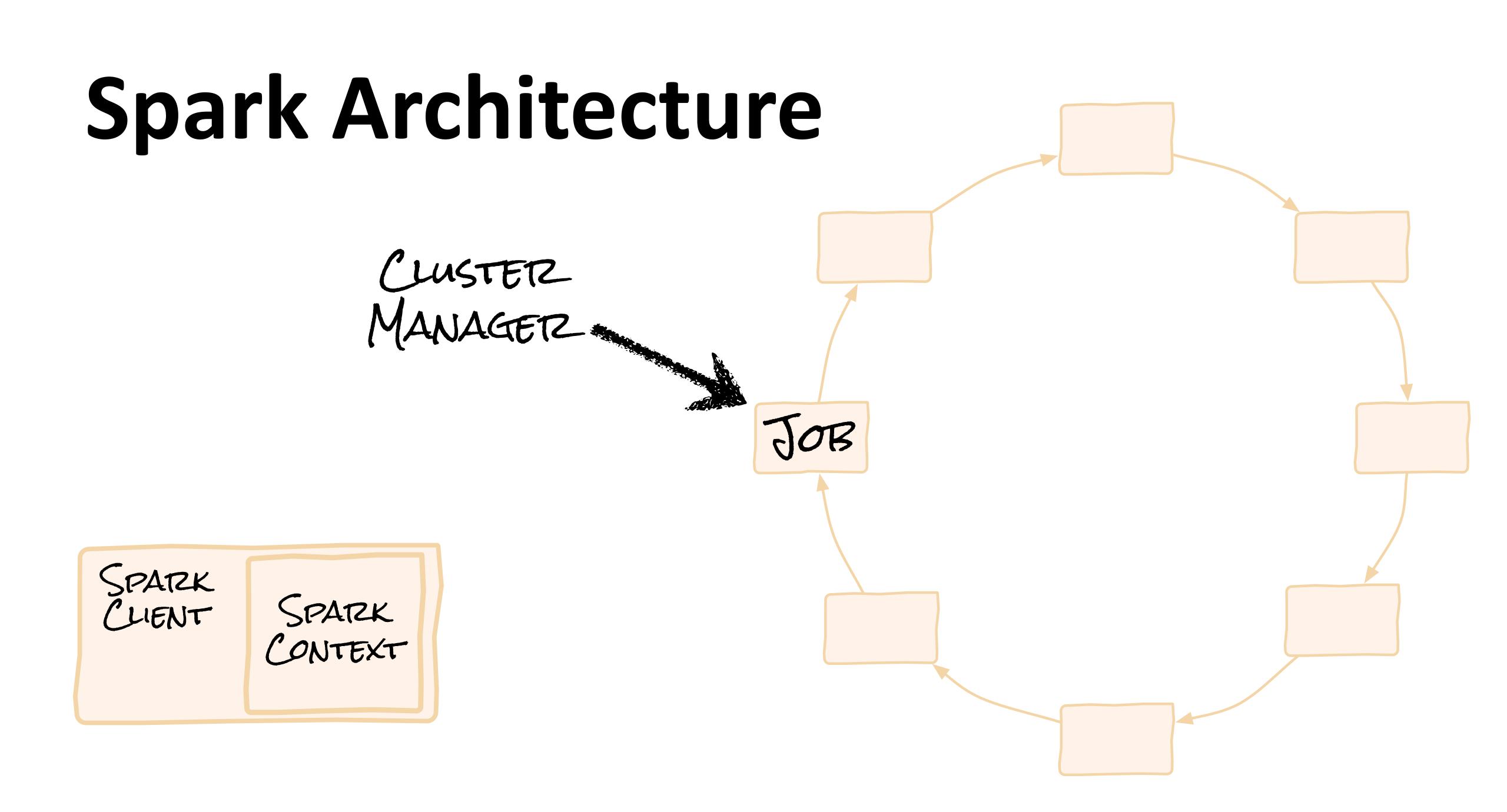


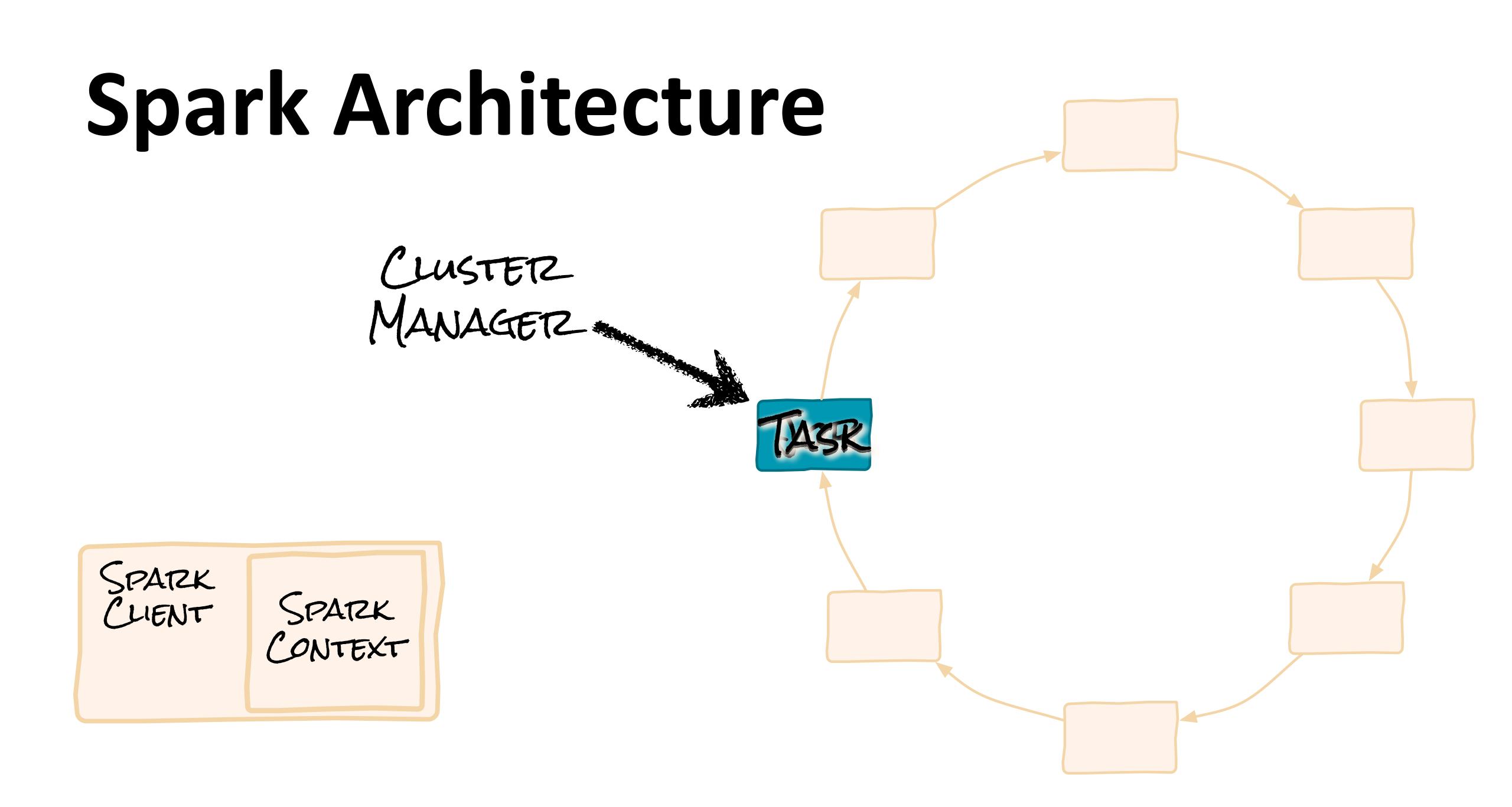


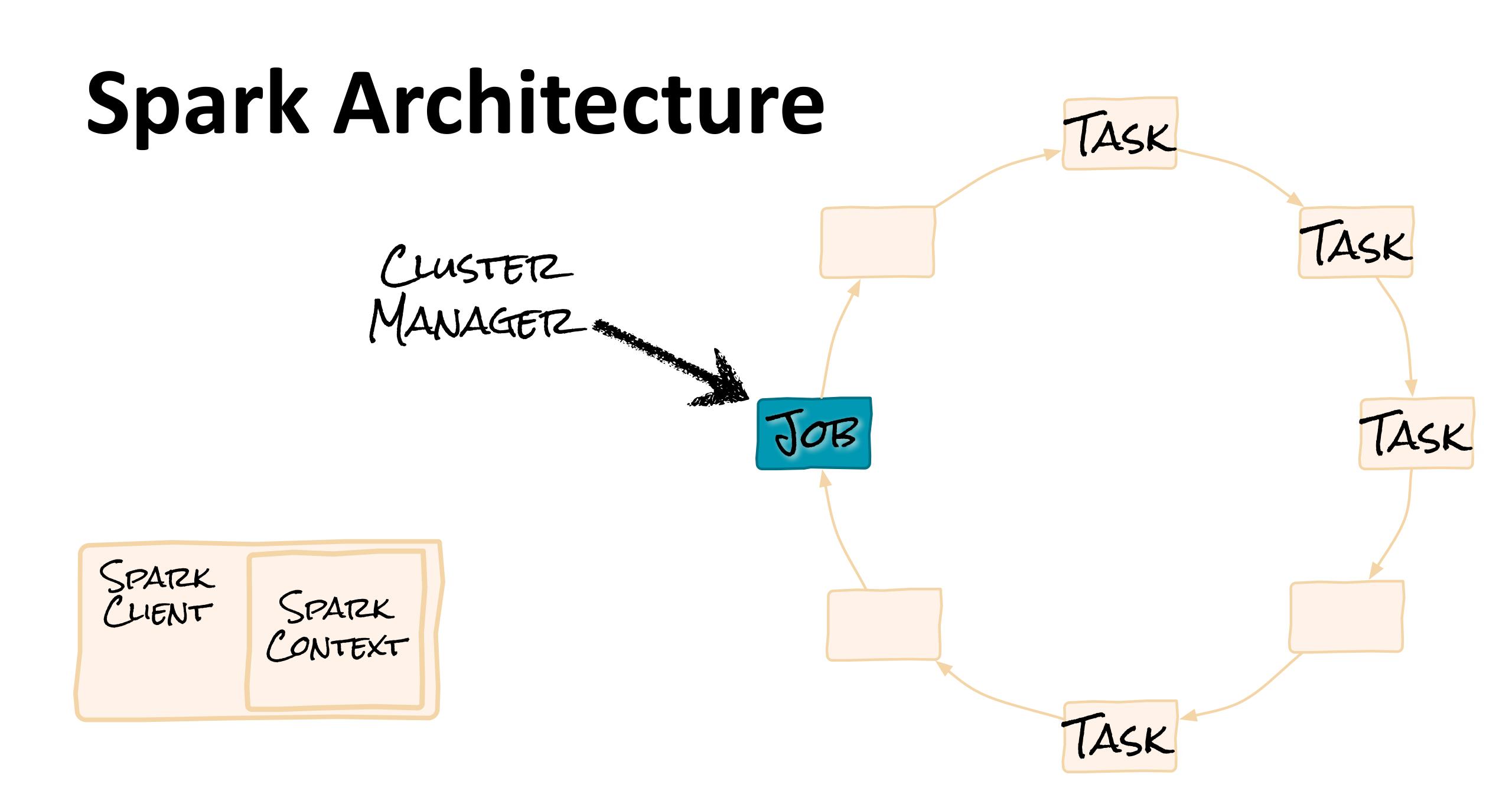
Spark Architecture



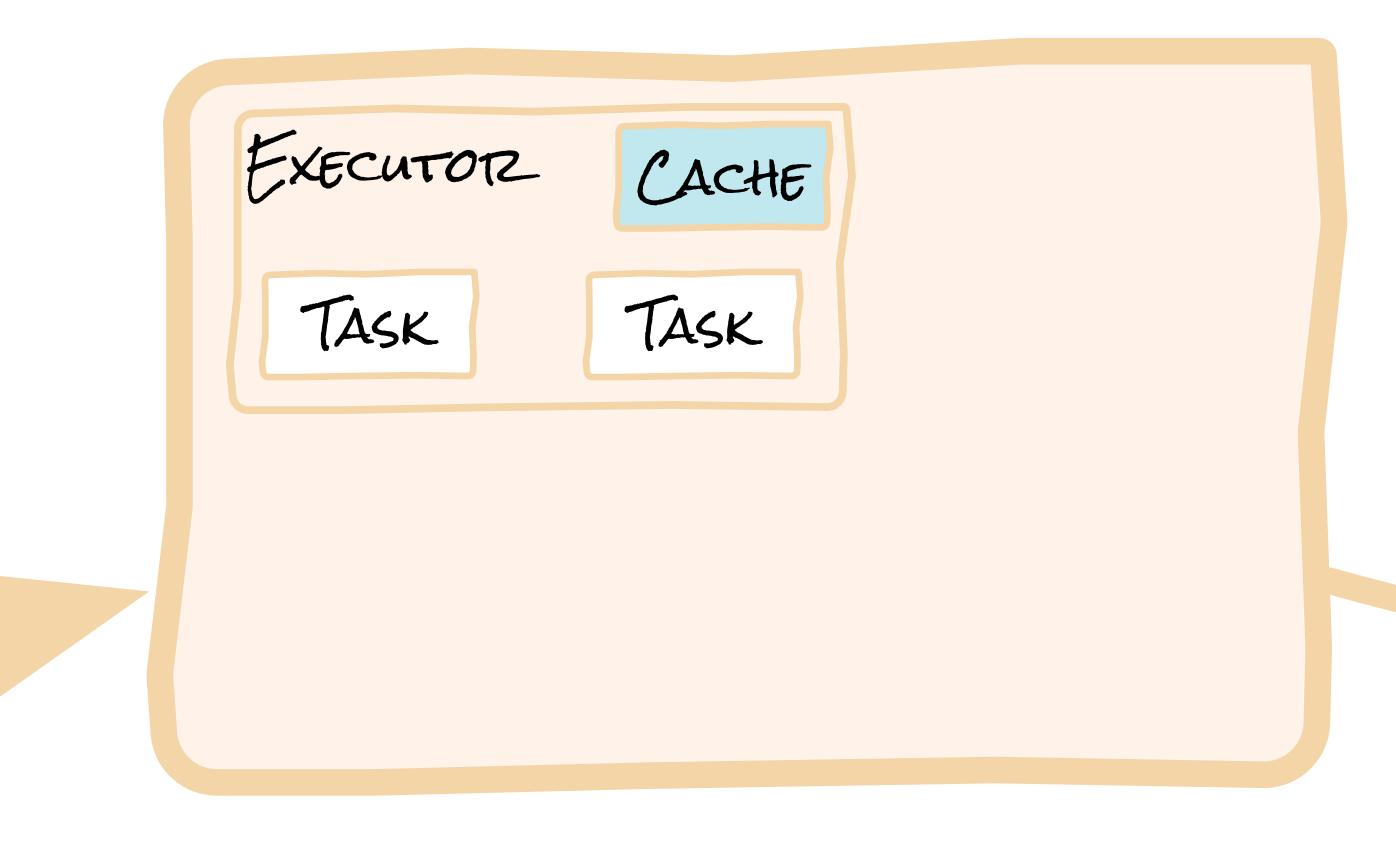




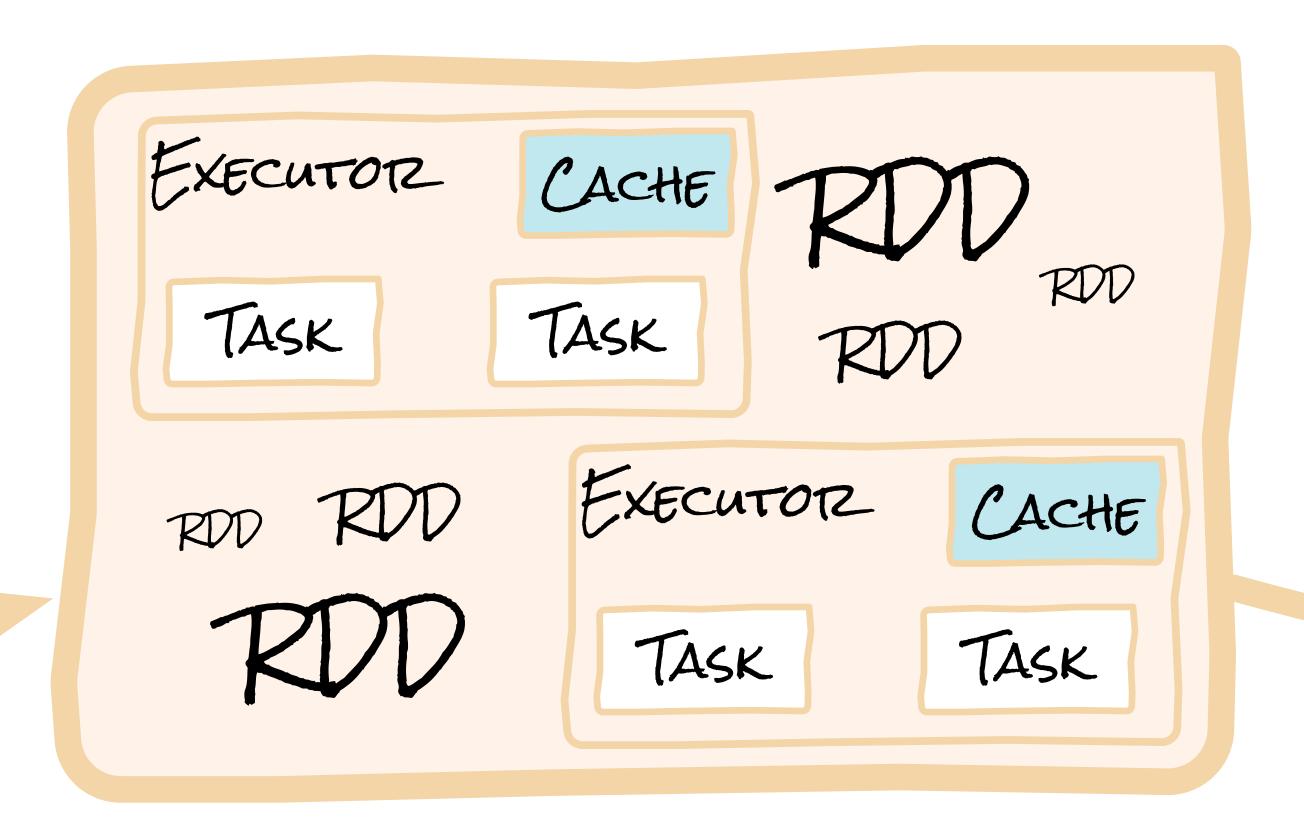




#### Spark Worker Node



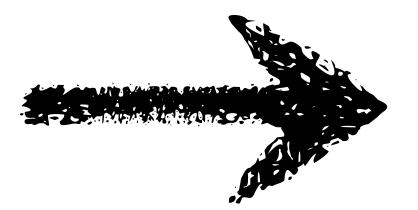
#### Spark Worker Node

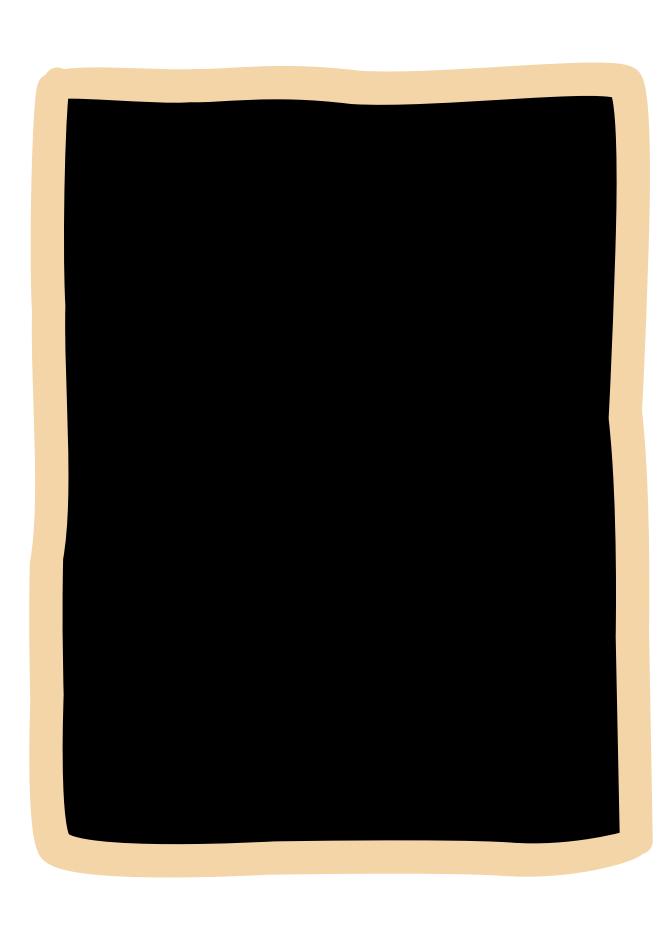


# What's an RDD?

#### What's an RDD?

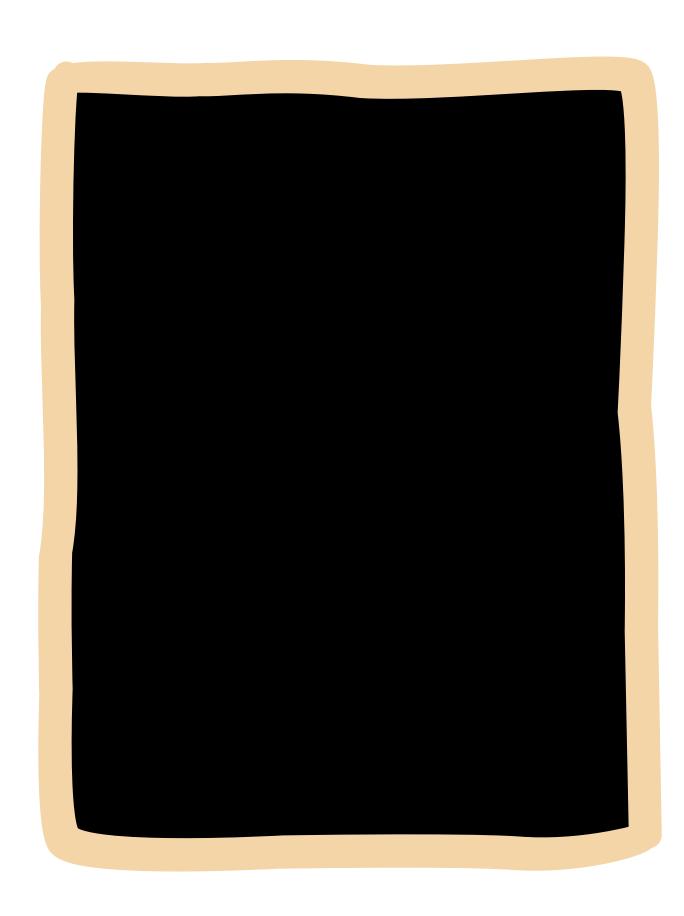
THIS IS



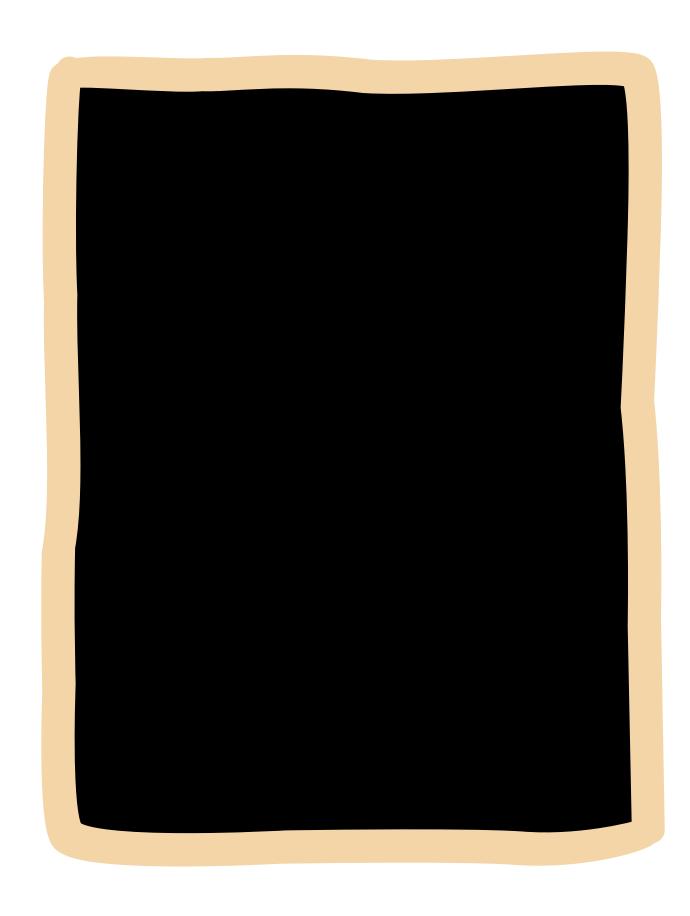


#### What's an RDD?

- Bigger than a computer
- Read from an input source
- Output of a pure function
- Immutable
- Typed
- Ordered
- Lazily evaluated
- Partitioned
- Collection of things



### Distributed, how?



#### Distributed, how?

```
HASH
THESE
```

```
5444 5676 0686 8389:
  List(xn-l, xn-2, xn-3)
4532 4569 7030 1191:
  List(xn-4, xn-5, xn-6)
    5607 6517 6027:
  List(xn-7, xn-凸, xn-月)
4532 4577 0122 2189:
  lict(vn-1.0-vn-1.1-vn-1.2)
```

#### Distributed Coordination

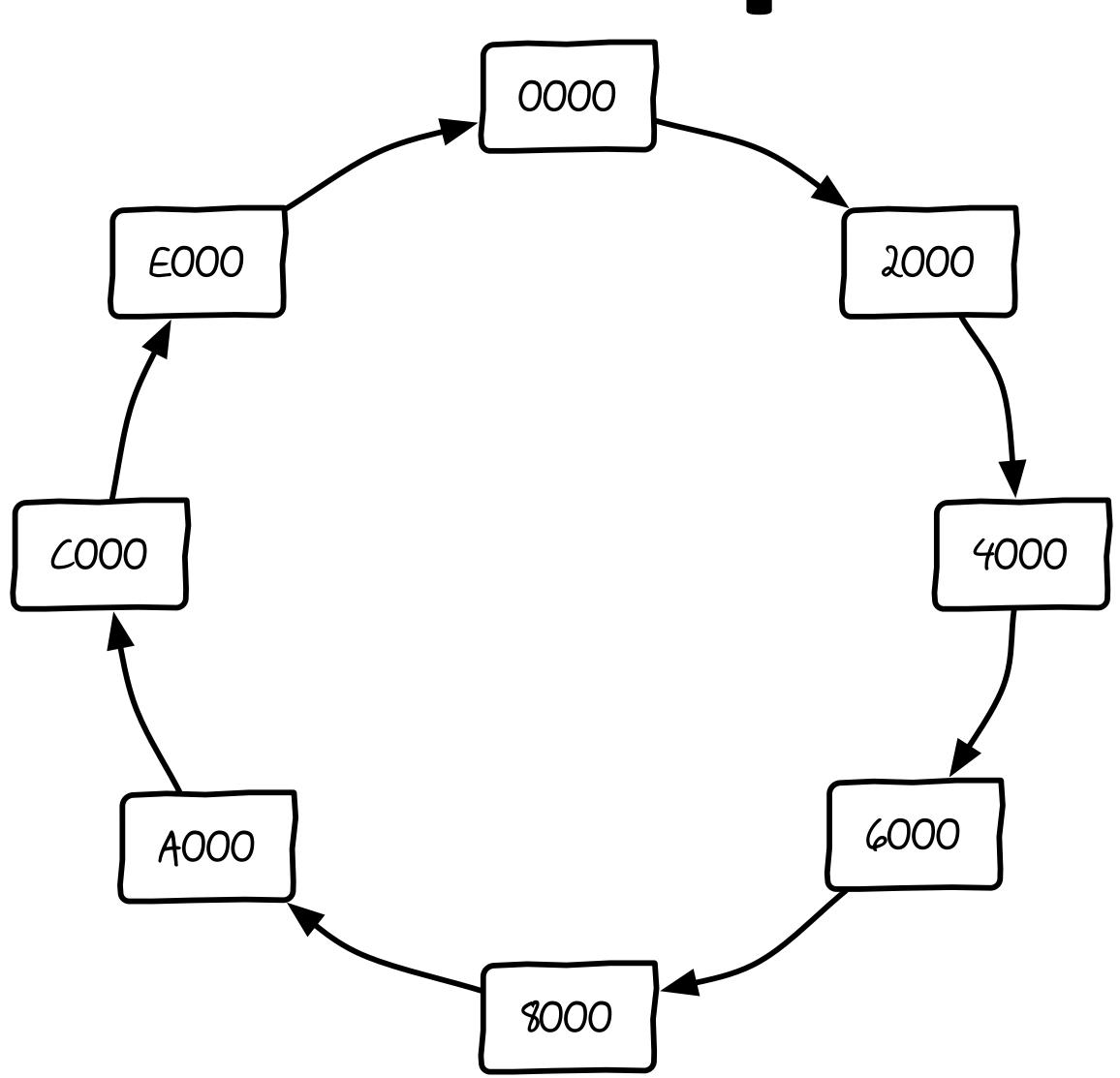
#### Distributed Coordination

- NTP
- Epidemic Protocols
- Paxos

#### NTP

- Distributed time synchronization
- A network of tiered clocks
- Synchronization algorithm
- Often accurate to ±10ms

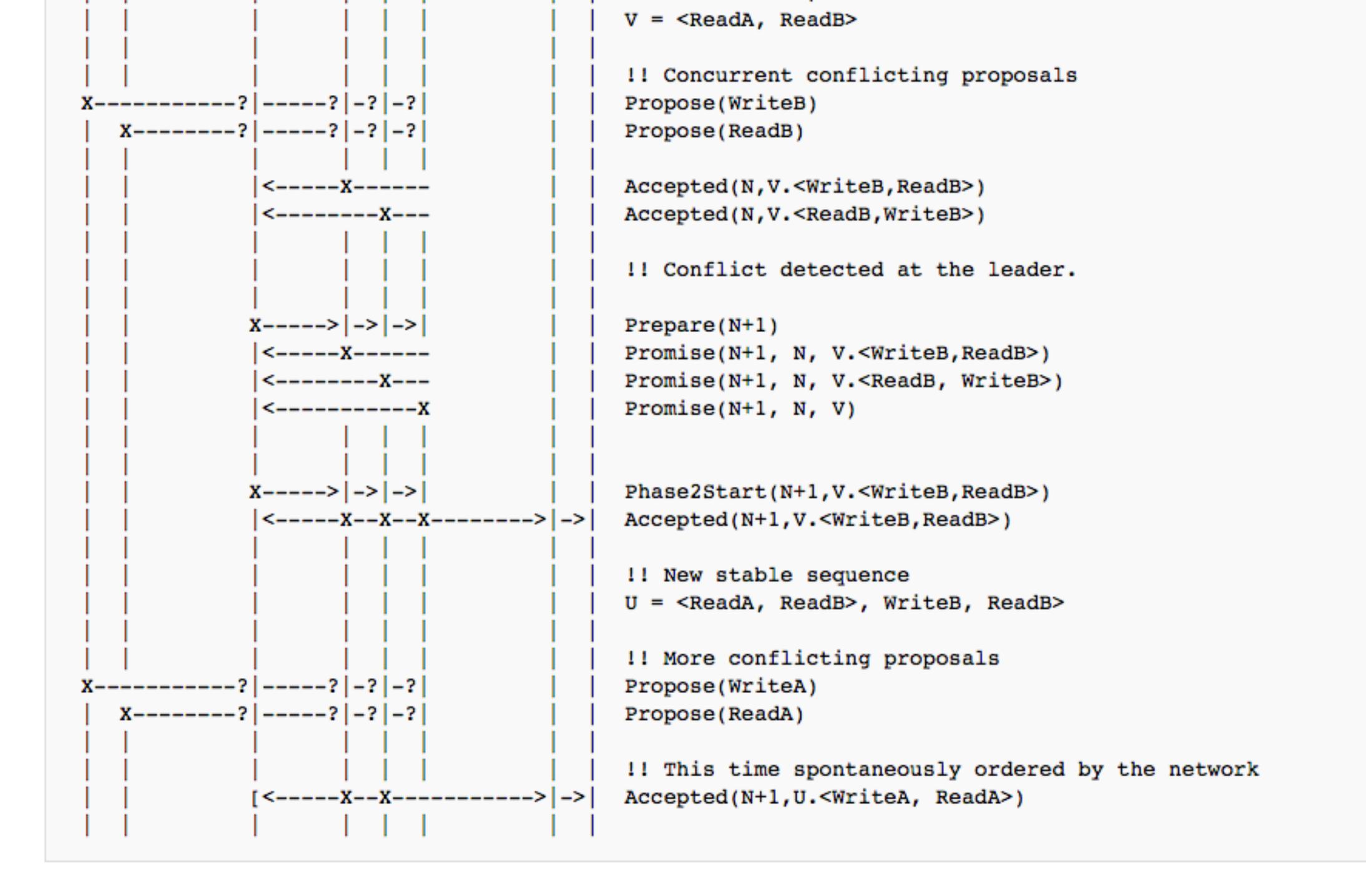
# GOSSID



#### Paxos

- Distributed log protocol
- Assumes distributed, unreliable nodes
- Performs leader election

```
Acceptors
Client
           Leader Acceptor
                                  Learner
                                          !! New Leader Begins Round
             X----> | -> | ->
                                          Prepare(N)
              <----X--X
                                          Promise(N, null)
             X---->|->|->|
                                          Phase2Start(N, null)
                                          !! Concurrent commuting proposals
   X----? | ----? | -? | -? |
                                          Propose(ReadA)
 X----? | ----? | -? | -? |
                                         Propose(ReadB)
              <----X-->-->|->| Accepted(N, <ReadA, ReadB>)
              <----> ->
                                         Accepted(N, < ReadB, ReadA>)
                                          !! No Conflict, both stable
                                          V = <ReadA, ReadB>
                                          !! Concurrent conflicting proposals
 X----? | ----? | -? | -? |
                                          Propose(WriteB)
   X----? | ----? | -? | -?
                                          Propose(ReadB)
                                          Accepted(N, V. < WriteB, ReadB>)
                                          Accepted(N, V. < ReadB, WriteB>)
                                          !! Conflict detected at the leader.
             X----> | -> | -> |
                                          Prepare(N+1)
                                          Promise(N+1, N, V.<WriteB,ReadB>)
                                          Promise(N+1, N, V.<ReadB, WriteB>)
              <----X
                                          Promise(N+1, N, V)
```



#### Review

- Storage
- Transactions
- Computation
- Coordination

