



Basics of Time Complexity

⇒ → String
Agenda → Python Memory
→ Time Complexity

⇒ String →

String ≠ list of characters

"this is a string"

['t', 'h', 'i', 's', ...]

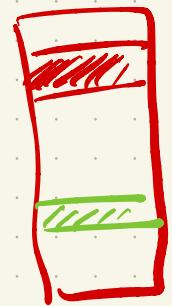
Python Memory →

$b = [[] \times 3] \times 3$

$b \rightarrow [[" ", " ", " "]]$
 $\rightarrow []$
 $\rightarrow []$

$a = " "$

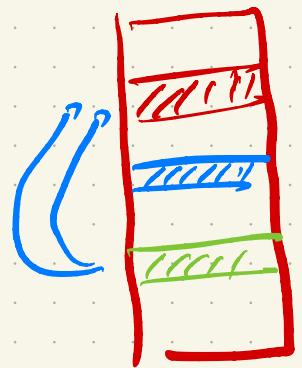
$c = [a, a, a]$



$a = " "$

$c = [list]$

$a = 10$



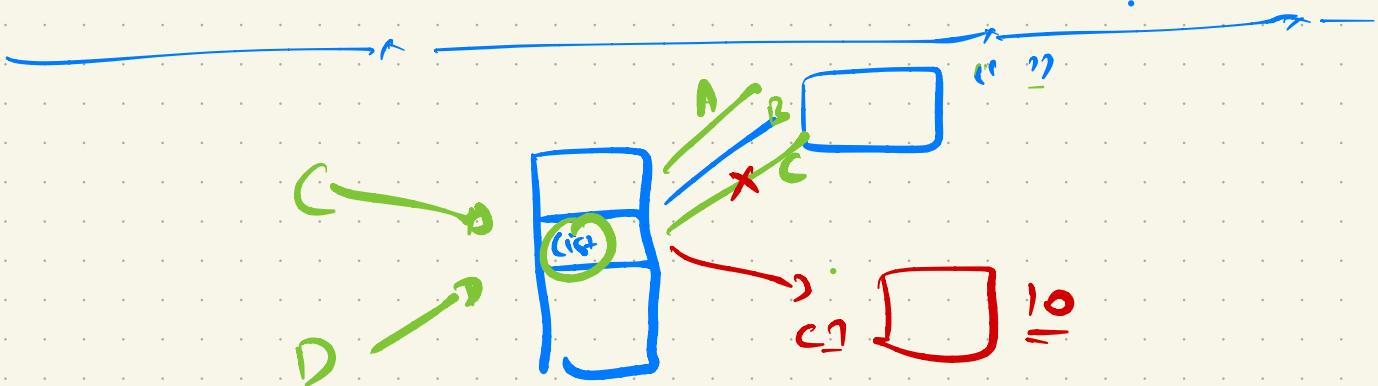
" "

$10 = "a"$

$c[0]$ refers to a String
(immutable)

if we change $c[0]$ only $c[0]$ get changed

$c[1], c[2]$ are not effected.



⇒ All mutable
Objects

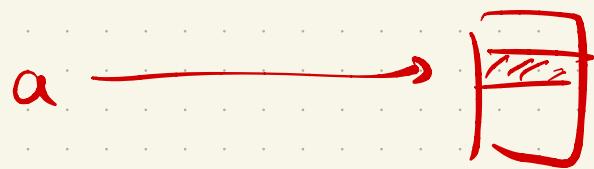
→ They keep
reference
of memory

Further
refer.
at
element

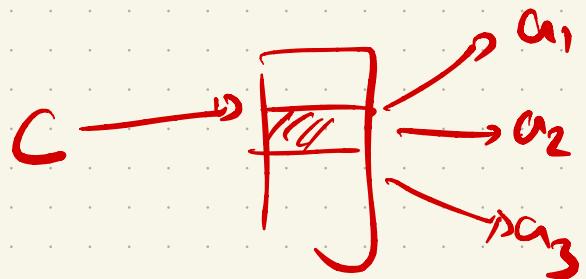
immutable
obj

→ final
value

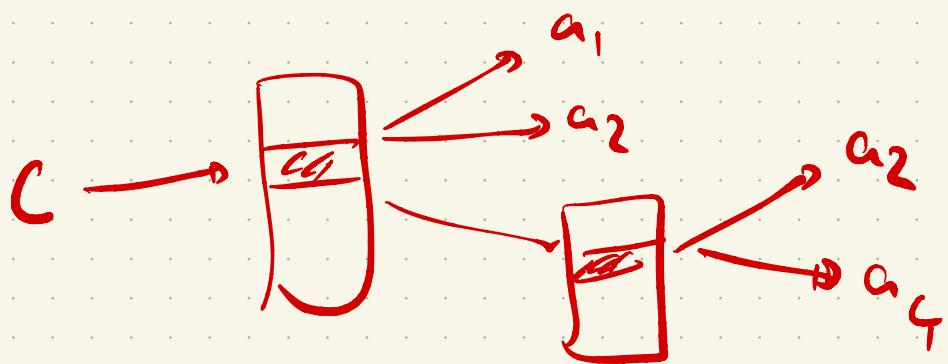
a = 5
b = "String"



c = [a₁, a₂, a₃]



c = [a₁, a₂, [a₂, a₄]]



$d = (a_1, a_2)$ ^{tuple}



d

Time Complexity → Related to algorithm analysis.

Algorithm Analysis

Space
(Memory)
[Space Complexity]

Processing
(CPU)
[Time Complexity]

⇒ 2.5 GHz

The system can perform

2.5×10^9 operations/sec
(Clock Speed)



Add 4 + 5
=

Time Complexity

xx How much time the algo will take?

↳ How many operation a algo will do.

$$1.5 \text{ GHz} \quad \frac{3}{1.5 \times 10^9} \quad 4+5 \rightarrow \underline{\underline{3 \text{ operations}}}$$

$$4.5 \text{ GHz} \quad \frac{3}{4.5 \times 10^9} \quad \underline{\underline{\text{sec}}}$$

\Rightarrow Actual Query time | Execution is dependent on machine.

Time Complexity \rightarrow number of operations

Time Complexity →

Define a Cost function } fact that calculate
no. of steps.

Whatever is your relationship bloom & no. of
steps for that function.

$$C(n) = \underline{k_1 n^2} + k_2 n + k_3 (1)$$

Highest power is considered to be the complexity → Big O notation

$$n^3 + \underline{n^2} + n + 1$$

only biggest power dictates no. of steps to
large n's.

$$O(n^3 + n^2 + n+1) \approx O(n^3)$$

$$O(2n^3) \approx O(n^3)$$

$$\underline{2n^3} > \underline{n^3} \gg \cancel{n^2}$$

$$2n^3 \approx n^3$$

$$\underline{\text{loop } n^3} \quad \text{vs} \quad \underline{n^3}$$

lower the complexity the better is the performance

$$\underline{n^3} \gg \cancel{n^2}$$

$$\underline{n^3 \approx n^3 + n^2}$$

\Rightarrow Integer / Number \rightarrow Boolean

0 \rightarrow False

~~-1
-2
1
2
1~~ \rightarrow True

\Rightarrow Boolean \rightarrow Integer

False \rightarrow 0

True \rightarrow 1

Only complexities are important.

① Compute Resource are limited

\rightarrow Search

\rightarrow Recommendation

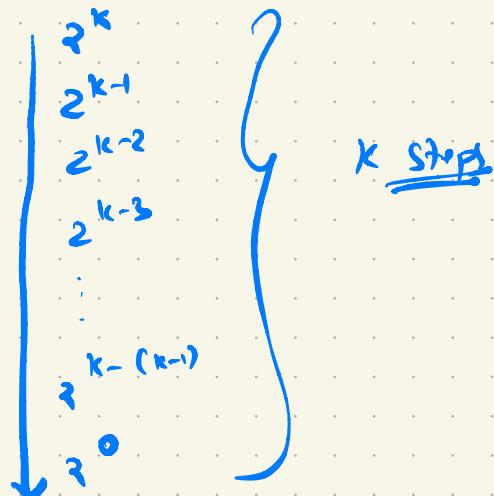
{ for want
fast}

\rightarrow understand how much operational cost your algo will take

Machines \rightarrow Repetitive Task \rightarrow

$$\text{If } \underline{n = 2^k}$$

It takes k steps
to reach 1



for a while loop which reduce n to half everytime

$$\underline{n = 2^k} \rightarrow k \text{ steps}$$

Complexity

$$\underline{n = 2^k} \rightarrow \boxed{k = \log_2 N}$$

(Complexity $\Rightarrow \underline{\log N}$)

for $\rightarrow n \rightarrow n^2 \rightarrow n^3 \rightarrow \dots$

\Rightarrow while $\underline{n/2}$

$$\begin{array}{l} \underline{n/3} \\ \underline{n/5} \\ \underline{n/10} \end{array}$$

$$\begin{array}{l} n \\ n/2 \\ n/3 \\ n/5 \\ n/10 \end{array}$$

$$\begin{array}{ll} n = 2^k & n = 3 \\ n = 5^m & \\ \log_2(n) & \log_3(n) \\ & \log_5(n) \end{array}$$

```

for i in range(n):
    j=i
    while j//2>=1:
        print(j)
    
```

$\log j$ time $\Rightarrow \boxed{\log(n!)}$

$$\log_1 + \log_2 + \log_3 + \dots + \log_n = \log(n!) = \underline{\underline{\log(n!)}}$$

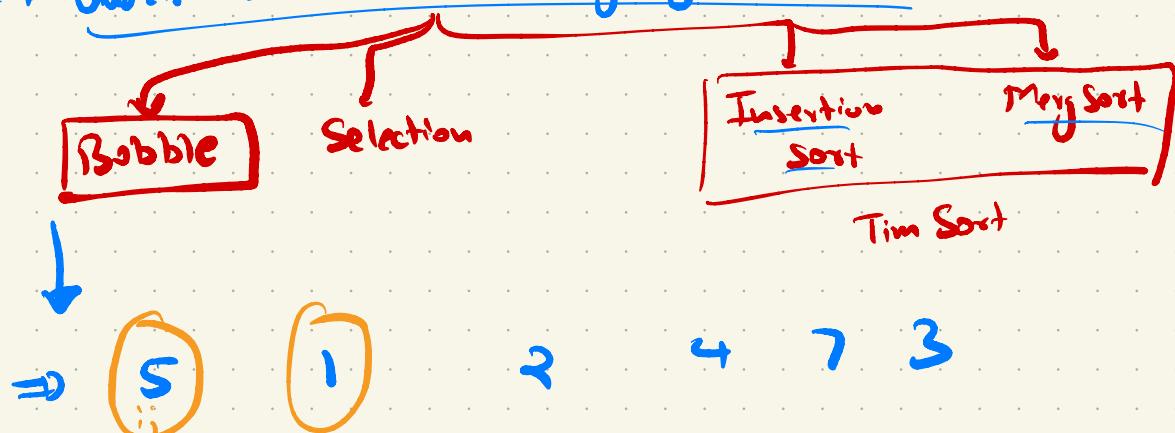
$$\Rightarrow \log \approx \log n + \dots + \log n = \underline{\underline{n \log n}}$$

\Rightarrow which is a better algorithm
 $n \log n$ vs $\log(n!)$

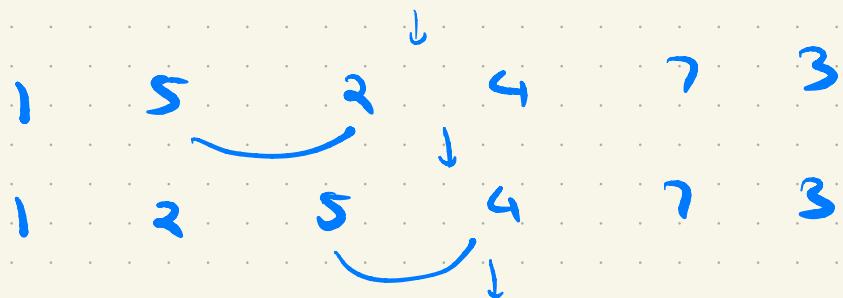
$\log(n!)$ is a better algo

as $\log(n!) < n \log n$

\Rightarrow Understand TC for Sorting Algorithms



Compare Swap $a[i] > a[i+1] \Rightarrow \underline{\text{Swap}}$



1 2 4 5 7 3
 ↓ ↗
 1 2 4 5 7 3
 | ↓
 1 2 4 5 3 7
 | ↓
 1 2 4 3 5 7
 | ↓
 1 2 3 4 5 7

$[1, 2, 3, 4, 5, 7] \rightarrow$
 How many swaps = ?? $\underline{0}$

$[1, 2, 3, 4, 7, 5] \rightarrow$
swaps $\rightarrow 1$

$[2, 1, 3, 4, 7, 5] \rightarrow$
swaps $\rightarrow 2$
 1

$[7, 5, 4, 3, 2, 1] \rightarrow$
swaps \rightarrow

<u>5 times</u>	5 4 3 2 1 7	$5+4+3+2+1$
<u>4 times</u>	4 3 2 1 5 7	<u>15 swaps</u>
<u>3 times</u>	3 2 1 4 5 7	
<u>2 times</u>	2 1 3 4 5 7	
<u>1 time</u>	1 2 3 4 5 7	

\Rightarrow Bubble Sort \rightarrow

Best Case \rightarrow O Swaps

Worst Case \rightarrow

$n-1, n-2, \dots, 1 \rightarrow 1+2+3+\dots+n-1$
swaps

$1+2+\dots+(n-1)$

$$\frac{(n-1)n}{2}$$

~~(n-1)n~~ $\frac{n^2-n}{2}$ swaps

$$0 \leftrightarrow \frac{(n)(n-1)}{2}$$

\Rightarrow "Worst Case Complexity" $\frac{(n)(n-1)}{2} \leftarrow O(n^2)$

If data is provided \rightarrow find no. of steps \rightarrow then give complexity

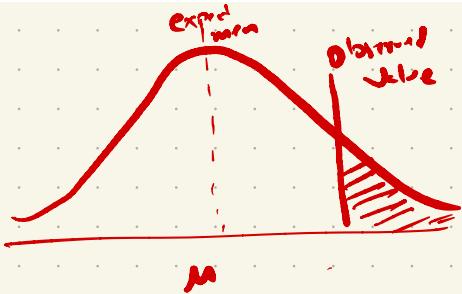
Bubble Sort \rightarrow Best Case \rightarrow O Swaps $\boxed{O(n)}$

at least 1 complete iteration of check would be needed.

1, 2, 3, 4, 5, 6 Comparison \rightarrow $\frac{n-1}{2}$ con

A fitness App claims that its users walk an average of 8,000 steps per day. A random sample of 30 users showed an average of 7,600 steps per day with a standard deviation of 1,200 steps. Conduct a right-tailed Z-test at a 5% significance level to determine if the App's claim is supported. What is the p-value?

Right tail test \rightarrow Null esp on right side of median.
 \hookrightarrow



How much chance / probability

that data will lie on
right side of
observed value

$$\left\{ \begin{array}{l} 30 \text{ samples} \\ \rightarrow 7600 \\ \rightarrow 8000 \\ \rightarrow 8500 \\ \rightarrow 8900 \end{array} \right\} \text{ deviation of } T_{ij}$$

As a data analyst at Global Mariotech, one of your prime duties is to understand customer sentiments and provide a better experience to the customers.

Based on a recent survey, it was found that there were 30% negative responses, and 70% positive. Assuming that we pick one response after another randomly, estimate the number of feedbacks to expect, before a negative feedback appears....solve using geometric distribution..

$$0.7 \rightarrow P \quad 0.3 \rightarrow N$$

$$\rightarrow N N N N N N N P \rightarrow (0.3)(0.7)(0.7) \dots (0.7) \rightarrow (0.3)^7 (0.7) \rightarrow 8$$

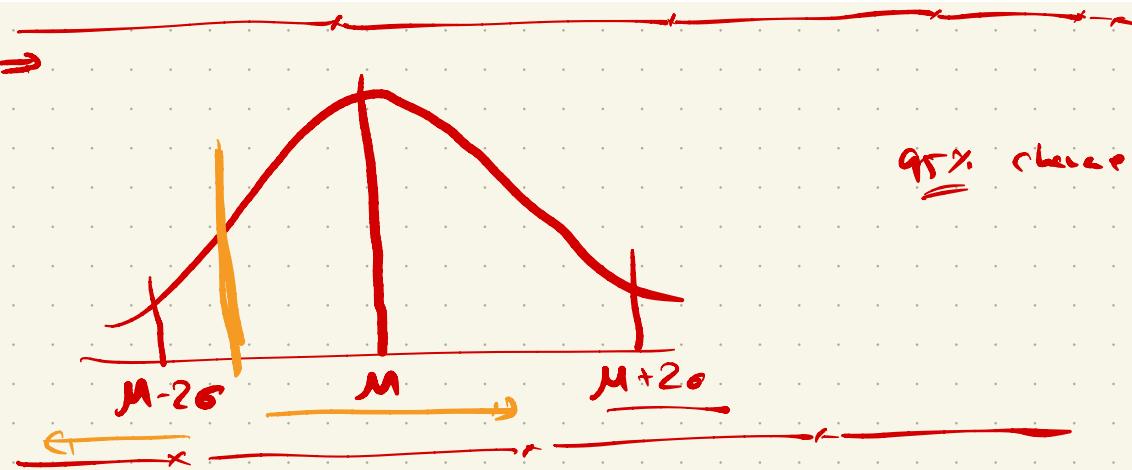
$$\rightarrow N N N P \rightarrow (0.3)^3 (0.7) \rightarrow 4$$

$$\rightarrow P \rightarrow (0.7) \rightarrow 1$$

$\frac{1}{M}$	(0.7)	$\rightarrow 1$	$\frac{(0.7)+(2)(0.7)(0.7)}{(0.7)(0.7)(0.7)}$
$\frac{0}{M}$	$(0.3)(0.7)$	$\rightarrow 2$	$+ \dots$
$\frac{0.2}{M}$	$(0.3)^2 (0.7)$	$\rightarrow 3$	$E(x) =$

$$(0.7) + (0.21) \times 2 + \dots$$

$\overbrace{\hspace{10em}}$



\Rightarrow I just randomly get a sample

$$\frac{M_{\text{sample}}}{\sigma_{\text{sample}}} \rightarrow \text{30 size}$$

$$\frac{\text{Sample mean}}{\text{std dev}}$$

$$\frac{\sigma_{\text{pop}}}{\sqrt{n}}$$

$$95\% \rightarrow \text{any random sample will lies b/w } M_{\text{pop}} \pm \frac{2\sigma_{\text{pop}}}{\sqrt{n}}$$

$$M_{\text{SM}} \in$$

Collab : https://colab.research.google.com/drive/1tWbICOWJ_CmTpsZtSafT3MyTLbWchAp4?usp=sharing