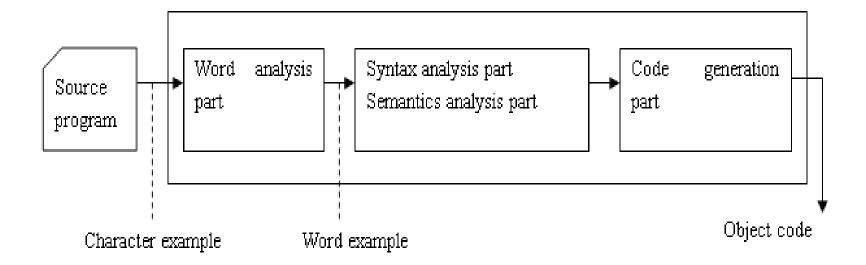


Unit 2. The phases of a Compiler

Main phases of a compiler





Phases of a compiler

Lexical Analysis

Stream of characters making up the source program is read from left to right and grouped into tokens (sequences of characters having a collective meaning)

Syntax Analysis

Group the tokens of the source program into grammatical phrases that are used by the compiler to synthesize output



phases of a compiler

- Semantic Analysis: Check the source program for semantic errors and gather type information for the subsequent code generation part.
- Intermediate Code Generation: Generate an intermediate representation as a program for an abstract machine.

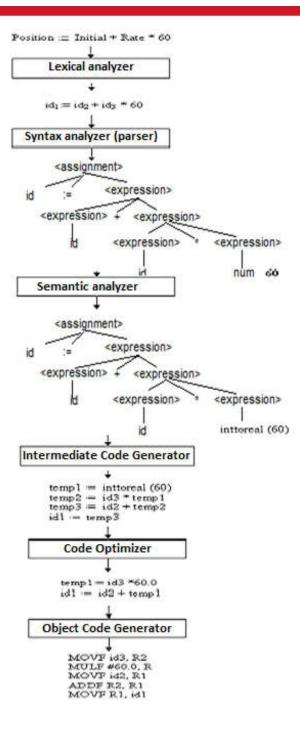


phases of a compiler

- Code optimization: Improve the intermediate code so that faster running code will result
- Code generation: Generation of target code, consisting normally of relocatable machine code or assembly code



Translation of a statement



Details of the phases of a Compiler

phase	Output	Sample
Programmer (source code producer)	Source string	Position:= inition * rate + 60
Scanner (performs lexical analyzer)	Token string	position', \':=', \'inition', \'+', \'60', And symbol table with identifier
Parser (performs syntax analysis based on the grammar of the programming language)	Parse tree or abstract syntax tree	ass_st id = + id := exp id num exp exp exp id exp * exp id num
Semantic analyzer (type checking, etc)	Annotated parse tree or abstract syntax tree	Convert integer (60) to real
Intermediate code generator	Three-address code	temp1 = inttoreal (60) temp2 = id3 * temp1 temp3 = id2 + temp2 id1 = temp3
Optimizer	Three-address code	$temp1 = id3 \approx 60.0$ $id1 := id2 + temp1$
Code generator VIỆN CÔNG NGHỆ THÔNG TIN	Assembly code VÀ TRUYỀN THÔNG	MOVF id3, R2 MULF #60.0, R2 MOVF id2, R1 ADDF R2, R1 MOVF R1, id1

The Grouping of phases

- Compiler *front* and *back ends*:
 - Front end: *analysis* (*machine independent*)
 - Back end: *synthesis* (*machine dependent*)
- Compiler passes:
 - A collection of parts is done only once (*single pass*) or multiple times (*multi pass*)
 - Single pass: usually requires everything to be defined before being used in source program
 - Multi pass: compiler may have to keep entire program representation in memory



Phase 1:Lexical Analysis

- Scanner: Converts the stream of input characters into a stream of tokens that becomes the input to the following phase (parsing)
- Tasks of a scanner

Group characters into tokens

Token: the syntax unit

Categorization of tokens.

• Token types: Identifier, Number, Character constant, operators.....



Phase 2: Parsing

- The process of determining if a string of token can be generated by a grammar
- Is executed by a parser

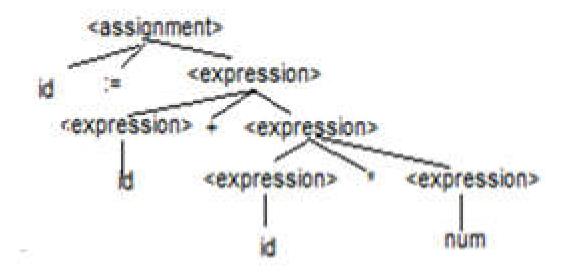


phase 2: Parsing

- Output of a parser:
 - Parse tree (if any)
 - Error Message (otherwise)
- If a parse tree is built successfully, the program is grammatically correct



Parse tree of statement id := id + id *num



Syntax Rules: Grammar (context free)

<assignment> → id := <expression>

<expression> → <expression> + <expression>

<expression> → <expression> * <expression>

 $\langle expression \rangle \rightarrow id$

 $\langle expression \rangle \rightarrow num$



Grammars, languages, BNF, syntax diagrams

- The parser takes the token produced by scanner as input and generates a parse tree (or syntax tree). Token arrangements are checked against the grammar of the source language.
- Notations for grammar:
 - BNF (Backus-Naur Form) is is a meta language used to express grammars of programming languages
 - Syntax Diagrams: A pictorial diagram showing the rules for forming an instruction in a programming language, and how the components of the statement are related. Syntax diagrams are like directed graphs.



BNF

- BNF (and formal grammars) use 2 types of symbol
- Terminals :
 - Tokens of the language
 - Never appear in the left side of any production
- Nonterminals
 - Intermediate symbol to express structures of a language
 - Must be in a left side of at lease one production
 - Enclose in <>
- Start symbol
 - Nonterminal of the first level
 - Appear at the root of parse tree



Parsing: Concept and Techniques

- Continuously apply grammatical rules until a string of terminal is generated.
- If the parser convert first symbol into the input string, it is syntactically correct
- Otherwise, string is not syntactically correct



Parsing: Concept and techniques

- The most important thing of a compiler: grammar
- Grammar includes all structures of a program
- Not includes any other rule



Parsing: Concept and Techniques

Grammar must be unambiguous

• If grammar is ambiguous, more than one parse tree can be created. A program can output 2 different result with the same input.

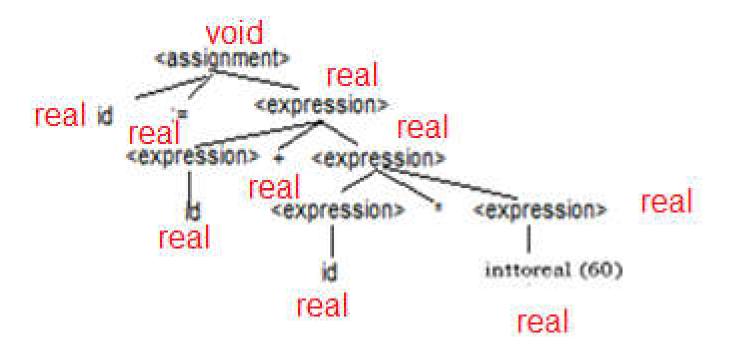


Phase 3: Semantic Analysis

- Certain check are performed to ensure that the components of a program fit together meaningfully
- To generate code, source program must be syntactically and semantically correct

Type checking of position := initial + rate *60

Assume position, initial and rate are declared as real variables





Phase 4: Intermediate code generation

- Source program is transferred to an equivalent program in intermediate code by intermediate code generator
- Intermediate code is close to the target code, which makes it suitable for register and memory allocation, instruction set selection, etc.
- It is good for machine-dependent optimizations.

```
temp1 = inttoreal (60)
temp2 = id3 * temp1
temp3 = id2 + temp2
id1 - temp3
```



Advantages of Intermediate Code

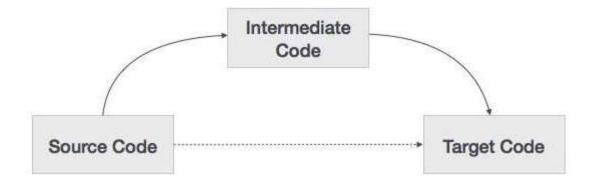
- 1. Easy to translate into object code.
- 2.Code optimizer can be applied before code generation
- 3.Decrease time cost



Phase 5: Code Generator

- Input: Intermediate code of source program
- Output: Object program
 - Assembly code
 - Virtual machine code

MOVF id3, R2 MULF #60.0, R MOVF id2, R1 ADDF R2, R1 MOVF R1, id1





Problems

- Input
- Output
- Object machine
 - Set of instruction
 - Register allocation

