



Data structure and algorithms lab

DEBUGGING WITH GDB & TREE

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Topics of this week

- How to use debugger tool(gdb)
- Tree data structure
 - Binary Tree
 - Binary Search Tree
- Recursive processing on Tree



`gdb` for debugging (1)

- `gdb`: the Gnu DeBugger
- <http://www.cs.caltech.edu/courses/cs11/material/c/mike/misc/gdb.html>
- Use when program core dumps
- or when want to walk through execution of program line-by-line



`gdb` for debugging (2)

- Before using `gdb`:
 - Must compile C code with additional flag:
`-g`
 - This puts all the source code into the binary executable
- Then can execute as: `gdb myprogram`
- Brings up an interpreted environment



gdb for debugging (3)

gdb> run

- Program runs...
- If all is well, program exits successfully, returning you to prompt
- If there is (e.g.) a core dump, **gdb** will tell you and abort the program



gdb – basic commands (1)

- Stack backtrace ("**where**")
 - Your program core dumps
 - Where was the last line in the program that was executed before the core dump?
 - That's what the **where** command tells you

gdb – basic commands (2)

`gdb> where` *last call* *last call in your code*

```
#0 0x4006cb26 in free () from /lib/libc.so.6
#1 0x4006ca0d in free () from /lib/libc.so.6
#2 0x8048951 in board_updater (array=0x8049bd0,
ncells=2) at 1dCA2.c:148
#3 0x80486be in main (argc=3, argv=0xbffff7b4) at
1dCA2.c:44
#4 0x40035a52 in __libc_start_main () from
/lib/libc.so.6
```

stack backtrace

gdb – basic commands (3)

- Look for topmost location in stack backtrace that corresponds to your code
- Watch out for
 - freeing memory you didn't allocate
 - accessing arrays beyond their maximum elements
 - dereferencing pointers that don't point to part of a `malloc()`ed block



gdb – basic commands (4)

- **break**, **continue**, **next**, **step** commands
- **break** causes execution to stop on a given line

```
gdb> break foo.c: 100
```

 (setting a breakpoint)
- **continue** resumes execution from that point
- **next** executes the next line, then stops
- **step** executes the next statement
 - goes into functions if necessary (**next** doesn't)



gdb – basic commands (5)

- **print** and **display** commands
- **print** prints the value of any program expression

```
gdb> print i
```

```
$1 = 100
```
- **display** prints a particular value every time execution stops

```
gdb> display i
```

gdb – printing arrays (1)

- `print` will print arrays as well

```
int arr[] = { 1, 2, 3 };
```

```
gdb> print arr
```

```
$1 = {1, 2, 3}
```

- N.B. the `$1` is just a name for the result

```
print $1
```

```
$2 = {1, 2, 3}
```

gdb – printing arrays (2)

- `print` has problems with dynamically-allocated arrays

```
int *arr;
```

```
arr = (int *)malloc(3 * sizeof(int));
```

```
arr[0] = 1; arr[1] = 2; arr[2] = 3;
```

```
gdb> print arr
```

```
$1 = (int *) 0x8094610
```

- Not very useful...

gdb – printing arrays (3)

- Can print this array by using `@` (gdb special syntax)

```
int *arr;  
arr = (int *)malloc(3 * sizeof(int));  
arr[0] = 1; arr[1] = 2; arr[2] = 3;
```

```
gdb> print *arr@3
```

```
$2 = {1, 2, 3}
```

gdb – abbreviations

- Common gdb commands have abbreviations

`p` (same as `print`)

`c` (same as `continue`)

`n` (same as `next`)

`s` (same as `step`)

- More convenient to use when interactively debugging

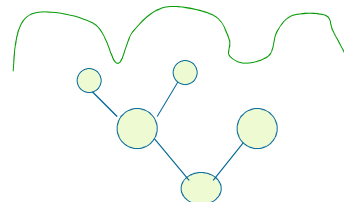
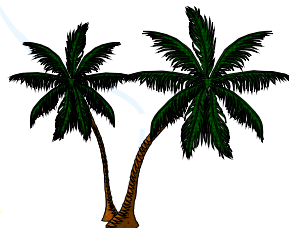
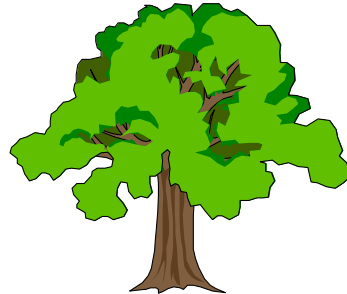
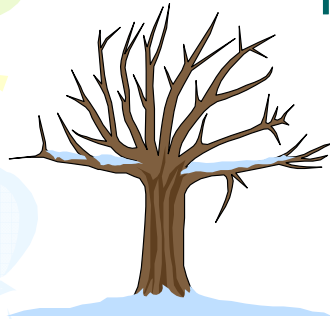
other instruction

- clear : delete break point of current file.
- delete [break position]: delete breakpoint at a specific file and position
- Conditional break

```
gdb> break foo.c: 100 if i== -1
```

- quit
- run: restart from beginning.

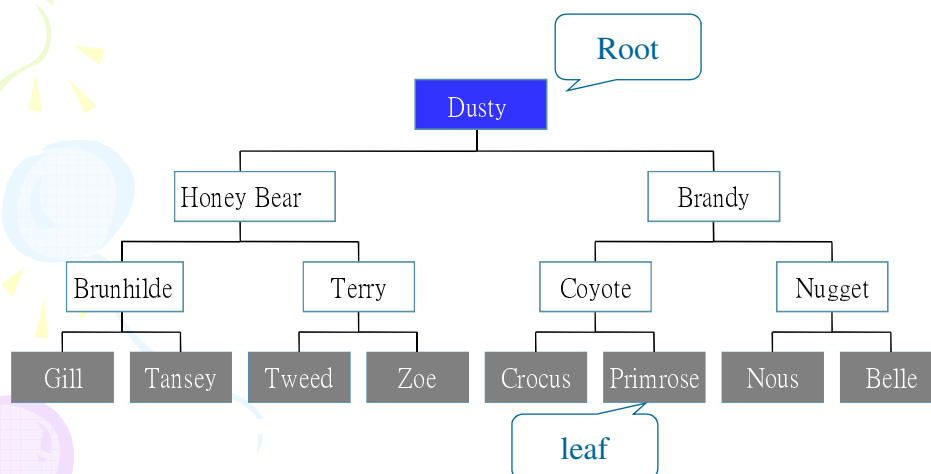
Tree



Trees, Binary Trees, and Binary Search Trees

- Linked lists are **linear structures** and it is difficult to use them to organize a **hierarchical** representation of objects.
- Although stacks and queues reflect some hierarchy, they are limited to only **one dimension**.
- To overcome this limitation, we create a new data type called a **tree** that consists of **nodes** and **arcs**. Unlike natural trees, these trees are **depicted upside down** with the **root** at the top and the **leaves** at the bottom.

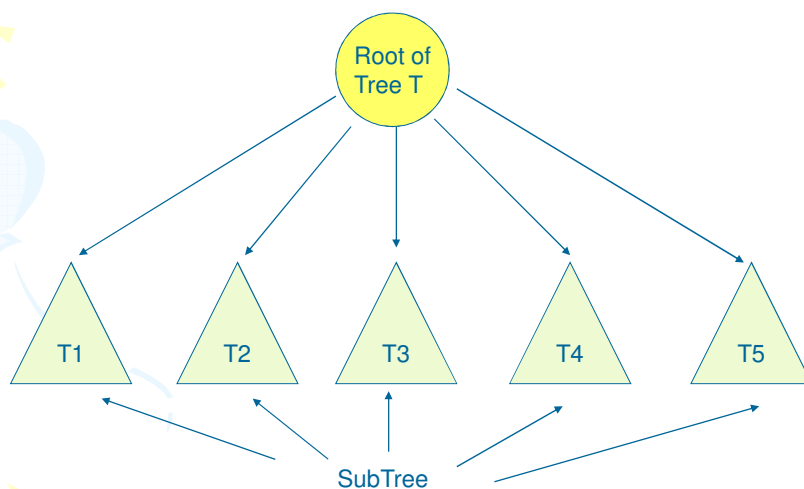
Family Tree



Definition of tree

- A tree is a finite set of one or more nodes such that:
- There is a specially designated node called the root.
- The remaining nodes are partitioned into $n \geq 0$ disjoint sets T_1, \dots, T_n , where each of these sets is a tree.
- We call T_1, \dots, T_n the subtrees of the root.

Recursive definition



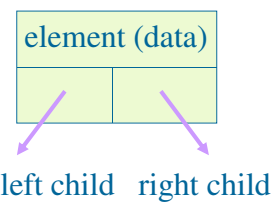
Binary Tree

- A binary tree is a tree in which no node can have more than two children.
- Each node has 0, 1, or 2 children

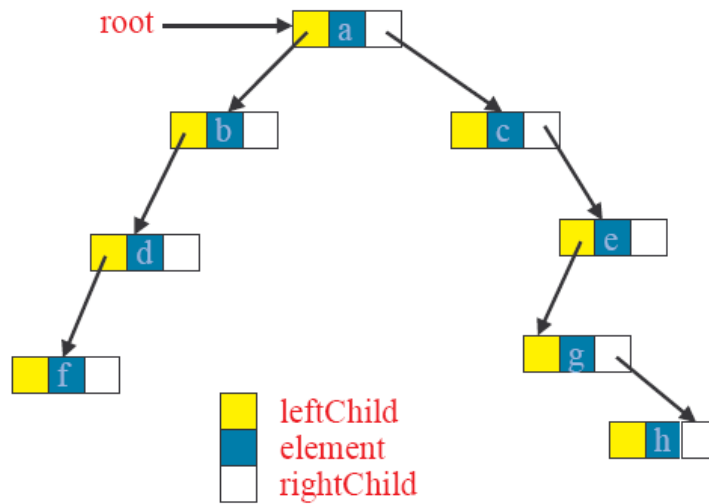
Linked Representation

- Each tree node is represented as an object whose data type is
- The space required by an n node binary tree is $n * (\text{space required by one node})$

```
typedef ... elmType;
//whatever type of element
typedef struct nodeType {
    elmType element;
    struct nodeType *left, *right;
};
typedef struct nodeType *treetype;
```



A linked binary tree



Binary Tree ADT

- `makenullTree(treetype *t)`
- `creatnewNode()`
- `isEmpty()`



Tree initialization and verification

```
typedef ... elmType;
typedef struct nodeType {
    elmType element;
    struct nodeType *left, *right;
} node_Type;

typedef struct nodeType *treetype;

void MakeNullTree(treetype *T){
    (*T)=NULL;
}

int EmptyTree(treetype T){
    return T==NULL;
}
```



Access left and right child

```
treetype LeftChild(treetype n)
{
    if (n!=NULL) return n->left;
    else return NULL;
}

treetype RightChild(treetype n)
{
    if (n!=NULL) return n->right;
    else return NULL;
}
```



create a new node

```
node_type *create_node(elmtype NewData)
{
    node_type *N;
    N=(node_type*)malloc(sizeof(node_type));
    if (N != NULL)
    {
        N->left = NULL;
        N->right = NULL;
        N->element = NewData;
    }
    return N;
}
```



check if a node is a leaf

```
int IsLeaf(treetype n) {
    if (n!=NULL)
        return (LeftChild(n)==NULL) && (Right
            Child(n)==NULL);
    else return -1;
}
```

Recursive processing: Number of nodes

- As tree is a recursive data structure, recursive algorithms are useful when they are applied on tree.

```
int nb_nodes(treetype T){  
    if(EmptyTree(T)) return 0;  
    else return 1+nb_nodes(LeftChild(T))+  
        nb_nodes(RightChild(T));  
}
```

Creat a tree from two sub- trees

```
treetype createfrom2(elmttype v,  
    treetype l, treetype r){  
    treetype N;  
    N=(node_type*)malloc(sizeof(node_type));  
    N->element=v;  
    N->left=l;  
    N->right=r;  
    return N;  
}
```

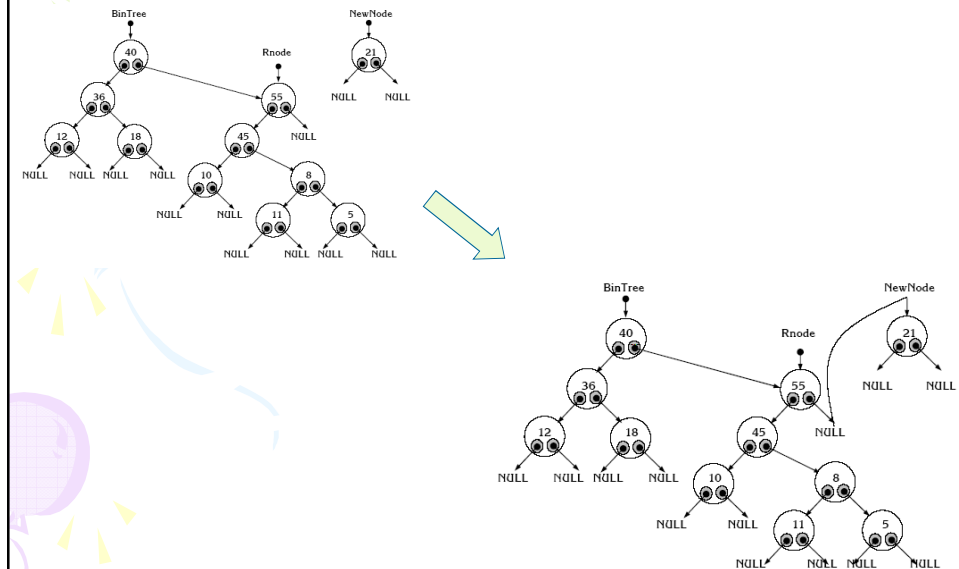
Adding a new node to the left most position

```
treetype Add_Left(treetype *Tree, elmttype NewData){
    node_type *NewNode = Create_Node(NewData);
    if (NewNode == NULL) return (NewNode);
    if (*Tree == NULL)
        *Tree = NewNode;
    else{
        node_type *Lnode = *Tree;
        while (Lnode->left != NULL)
            Lnode = Lnode->left;
        Lnode->left = NewNode;
    }
    return (NewNode);
}
```

Adding a new node to the right most position

```
treetype Add_Right(treetype *Tree, elmttype NewData){
    node_type *NewNode = Create_Node(NewData);
    if (NewNode == NULL) return (NewNode);
    if (*Tree == NULL)
        *Tree = NewNode;
    else{
        node_type *Rnode = *Tree;
        while (Rnode->right != NULL)
            Rnode = Rnode->right;
        Rnode->right = NewNode;
    }
    return (NewNode);
}
```


Illustration

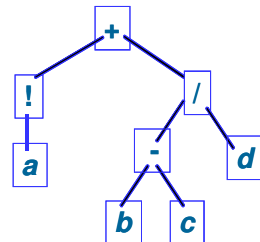


Exercise (BTVN)

- Develop the following helper functions for a tree:
 - return the height of a binary tree.
 - return the number of leafs
 - return the number of internal nodes
 - count the number of right children.

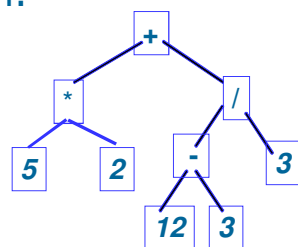
Exercise

- A binary tree can represent an arithmetic expression: The leaves are operands and the other nodes are operators.
- The left and right subtrees of an operator node represent **subexpressions** that must be evaluated **before** applying the operator at the root of the subtree.
- For example
 $a + (b - c)/d$
- Write a program create a tree representing this expression



Exercise at Home

- Write an menu program that take a valid arithmetic expression as input and:
 - Store and represent it in a tree
 - Evaluate the expression.



Homework

- Create a text file AUOpen.txt containing the names of tennis players participating in the Australian Open tournament, one player per line. The number of player is the power of 2 and a minimum of 16 people.
- Build a tree describing the competition results until the final. Initially, the (16) opponents are leaf nodes, the winner will be stored at the parent node of a pair of players. The winner of a match is randomly selected.
- Print the results of the matches (tree content) to the screen and file treegame.txt

Section 1 [edit]

| First Round | Second Round | Third Round | Fourth Round |
|----------------------------|-----------------------------|---------------------------|----------------------|
| 1 R Nadal 6 6 6 | 1 R Nadal 6 7 6 | 1 R Nadal 6 6 6 | 1 R Nadal 6 3 7 7 |
| H Dellien 2 3 0 | F Delbonis 3 6 1 | 27 P Carreño Busta 1 2 4 | 23 N Kyrgios 3 6 6 6 |
| F Delbonis 6 6 7 | Q P Gojowczyk 4 1 6 4 | | |
| J Sousa 3 4 6 | 27 P Carreño Busta 6 6 1 6 | | |
| C Eubanks 6 3 6 0 | | | |
| Q P Gojowczyk 7 6 4 6 | 23 N Kyrgios 6 6 4 7 | | |
| J Kovalik 4 6 1 6 | G Simon 2 4 6 5 | 23 N Kyrgios 6 7 6 6 7 10 | |
| 27 P Carreño Busta 6 3 6 7 | M Ymer 2 6 4 6 6 | 16 K Khachanov 2 6 7 6 6 | |
| 23 N Kyrgios 6 7 7 7 | 16 K Khachanov 6 2 6 3 7 10 | | |
| L Sonogo 2 6 6 | | | |
| P Cuevas 1 3 3 | | | |
| G Simon 6 6 6 | | | |
| Y Uchiyama 4 1 2 | | | |
| M Ymer 6 6 6 | | | |
| M Vilella Martinez 6 4 6 3 | | | |
| 16 K Khachanov 4 6 7 6 | | | |

Section 2 [edit]

| First Round | Second Round | Third Round | Fourth Round |
|------------------------------|-------------------------|--------------------|--------------------|
| 10 G Monfils 6 6 6 | 10 G Monfils 4 7 10 6 7 | 10 G Monfils 7 6 6 | 10 G Monfils 2 4 4 |
| PR Y-h Lu 1 4 2 | I Karlović 6 6 4 6 | Q E Gulbis 6 4 3 | 5 D Thiem 6 6 6 |
| I Karlović 7 6 7 | A Bedene 5 3 2 | | |
| PR V Pospisil 6 4 5 | Q E Gulbis 7 6 6 | | |
| J Duckworth 4 7 7 2 4 | | | |
| A Bedene 6 6 6 6 6 | | | |
| E Gulbis 7 4 7 6 | | | |
| 20 F Auger-Aliassime 5 6 6 4 | | | |
| 20 T Fritz 6 6 6 | 20 T Fritz 4 6 5 7 6 6 | | |
| T Griekspoor 3 3 3 | K Anderson 6 7 6 2 2 | | |
| Q I Ivashka 4 6 6 4 6 | | | |
| K Anderson 6 2 4 6 7 10 | | | |
| A Bolt 7 1 6 6 6 | | | |
| A Ramos Villos 6 6 7 1 4 | WC A Bolt 2 7 7 1 2 | 20 T Fritz 2 4 7 4 | |
| A Mannarino 3 5 2 | 5 D Thiem 6 5 6 6 6 | 5 D Thiem 6 6 6 6 | |
| 5 D Thiem 6 7 6 | | | |