### Quantum information project - reading notes

**Hugo Thomas** 

#### Abstract

This documents gathers reading notes and solutions found during the project.

# Self-testing of quantum systems: a review [ŠB20]

### The self-testing scenario

 $\mathcal{L}(\mathcal{H})$  denotes the set of linear operators acting on Hilbert space  $\mathcal{H}$ . We know there exist measurement operators  $M_{a|x} \in \mathcal{L}(\mathcal{H})$  acting on Alice's Hilbert space and satisfying

$$M_{a|x} \succcurlyeq 0; \forall a, x \sum_{a} M_{a|x} = \mathbb{1}_A$$
 (1)

Similarly, there exist measurement operators  $N_{b|y} \in \mathcal{L}(\mathcal{H})$  acting on Bob's Hilbert space. The measurement operators are therefore projective:

$$\forall a, a' : M_{a|x} M_{a'|x} = \delta_{a,a'} M_{a|x} 
\forall b, b' : N_{b|y} N_{b'|y} = \delta_{b,b'} N_{b|y}$$
(2)

Now, from the Born rule, there must exist some quantum state  $\rho_{AB} \in \mathcal{L}(\mathcal{H}_A \otimes \mathcal{H}_B) \geq 0$ , and  $\operatorname{tr} \rho_{AB} = 1$  such that

$$p(a, b|x, y) = \operatorname{tr}\left[\rho_{AB} M_{a|x} \otimes N_{b|y}\right]$$
 (3)

In self-testing, one aims to infer the form of the state and the measurement in the trace from knowledge of the correlation p(a,b|x,y) alone, i.e. in device-independent scenario.

Born rule: A key postulate of quantum mechanics which gives the probability that a measurement of a quantum system will yield a given result. More formally, for a state  $|\psi\rangle$  and an  $F_i$  POVM element (associated with the measurement outcome i), then the probability of obtaining i when measuring  $|\psi\rangle$  is given by

$$p(i) = \langle \psi | F_i | \psi \rangle \tag{4}$$

### Physical assumptions

- 1. The experiment admits a quantum description (state and measurement)
- 2. The laboratories of Alice and Bob are located in separate location in space and there is no communication between the two laboratories.
- 3. The setting x and y are chosen freely and independently of all other systems in the experiment.
- 4. Each round of the experiment is independent of all other rounds a physically equivalent to all others (i.e. there exists a single density matrix and measurement operators that are valid in every round).

## Impossibility to infer exactly the references

- 1. Unitary invariance of the trace: one can reproduce the statistics of any state  $|\psi\rangle$  and measurement  $\{M_{a|x}\}, \{N_{b|y}\}$  by instead using the rotated state  $U \otimes V |\psi\rangle$  and measurement  $\{UM_{a|x}U^{\dagger}\}, \{VN_{b|y}V^{\dagger}\},$  where U, V are unitary transformations. Hence, one can never conclude that the state was  $|\psi\rangle$  or  $U \otimes V |\psi\rangle$ .
  - On the other hand, considering real reference states  $(|\psi\rangle = |\psi\rangle^*)$ , one can only self-test measurements that are invariant under the complex conjugate \*, since, assuming a real state  $|\psi\rangle$ ,  $p(ab|xy) = \text{tr}[|\psi\rangle\langle\psi|M_{a|x}\otimes N_{b|y}] = \text{tr}[|\psi\rangle\langle\psi|M_{a|x}^*\otimes N_{b|y}^*]$ . Thus, any correlation obtained using  $\{|\psi\rangle, M_{a|x}, N_{b|y}\}$  can also be obtained using  $\{|\psi\rangle, M_{a|x}^*, N_{b|y}^*\}$ ; but the second is not related to the first one via a local isometry (It's an open problem to list the set of state and measurement transformation that do not affect the probabilities).
- 2. Additional degrees of freedom: a state  $|\psi\rangle \otimes |\xi\rangle$  and measurements  $\{M_{a|x} \otimes \mathbb{1}_{\xi}\}, \{N_{b|y} \otimes \mathbb{1}_{\xi}\}$  gives the same correlation as  $|\psi\rangle$  and  $\{M_{a|x}\}, \{N_{b|y}\}$ .

### Extractability relative to a Bell Inequality

The extractability  $\Xi$  is defined as the maximum fidelity of  $\Lambda_A \otimes \Lambda_B[\rho]$  and  $|\psi'\rangle$  over all CPTP (Completely positive and trace preserving) maps:

$$\Xi(\rho \to |\psi'\rangle) = \max_{\Lambda_A, \Lambda_B} F(\Lambda_A \otimes \Lambda_B, |\psi'\rangle) \quad (5)$$

where  $\rho \to |\psi'\rangle$  defines a kind of mapping of the test state  $\rho$  to the target state  $|\psi'\rangle$ . The maximum is taken over all quantum channels (why are the  $\Lambda_{A,B}$  called quantum channels?). This implies that  $\Xi$  return the  $\Lambda_{A,B}$  such that the fidelity to the reference state is maximal.

In order to test the entanglement characteristics of  $\rho$ ,  $|\psi'\rangle$  is assumed to be a state which achieves the maximal quantum violation. Hence, when the maximal quantum violation is observed in a self-testing scenario, the shared unknown state  $(\rho)$  can be mapped to  $|\psi'\rangle$ , and the resulting extractability is 1.

To get the optimal (robustness-wise) self-testing statement, one can minimize the possible extractability (over all states) when a violation of at least  $\beta$  is observed on a Bell inequality B. This quantity can be captured by the function  $\mathcal{Q}$  defined as

$$Q_{\psi,\mathcal{B}_{\mathcal{I}}} = \min_{\rho \in S_{\mathcal{B}}(\beta)} \quad \Xi(\rho \to |\psi'\rangle) \tag{6}$$

where  $S_{\mathcal{B}}(\beta)$  is the set of states  $\rho$  which violate Bell inequality  $\mathcal{B}$  as defined in [ZCP<sup>+</sup>18] with value at least  $\beta$ . One needs to note that the optimal CPTP map generally depends on the observed violation.

## The device-independent outlook on quantum physics [Sca15]

### Formal characterization of local variables

The local variable paradigm is related to the idea of pre-established agreement between the two parties, i.e. the output of each run is fully determined by a variable  $\lambda$ :

$$\mathbb{P}(a, b|x, y, \lambda) = \mathbb{P}(a|x, \lambda)\mathbb{P}(b|y, \lambda) \tag{7}$$

In other words, one can write the probability of having (a, b) given the two inputs (x, y) as:

$$\mathbb{P}(a, b|x, y) = \int d\lambda \mathbb{P}(a|x, \lambda) \mathbb{P}(b|y, \lambda) \tag{8}$$

This pre-established agreement is called *shared* randomness.

No-signalling condition: The fact that the statistics of a (b) should nod depend on y(x), that is, the following conditions hold for all  $a, b, x, y, \lambda$ :

$$\mathbb{P}(a|x,y,\lambda) = \mathbb{P}(a|x,\lambda) \qquad \mathbb{P}(b|x,y,\lambda) = \mathbb{P}(b|y,\lambda)$$
(9)

From the above conditions one can deduce the *no-signalling constraints*:

$$\forall x, x' \in \mathcal{Y} : \sum_{b} \mathbb{P}(a, b | x, y) = \sum_{b} \mathbb{P}(a, b | x, y') = \mathbb{P}(a | x)$$

$$\forall x, x' \in \mathcal{X} : \sum_{a} \mathbb{P}(a, b | x, y) = \sum_{a} \mathbb{P}(a, b | x', y) = \mathbb{P}(b | y)$$
(10)

### Bell inequalities as a polytope

The set  $\mathcal{L}$  of all families of probability distributions that can be obtained with LV is convex. In other words, if  $\mathcal{P}_1 \in \mathcal{L}$  and  $\mathcal{P}_2 \in \mathcal{L}$ , then  $\forall \alpha \in [0,1] : \alpha \mathcal{P}_1 + (1-\alpha)\mathcal{P}_2 \in \mathcal{L}$ . Furthermore, each deterministic local point is an extremal point of  $\mathcal{L}$ .

A polytope  $\mathcal{L}$  in  $\mathbb{R}^d$  is delimited by finitely many (d-1)-dimensional hyperplanes called facets (of the polytope).

# Link between Bell inequalities and linear programming

It is possible to state to which set  $(\mathcal{L}, \mathcal{Q}, \mathcal{NS})$  a probability distribution P belongs to. Consider the linear program given below:

$$s.t.\begin{cases} (1-q)P + q\mathbb{1} &= \sum_{\lambda} \mu_{\lambda} D_{\lambda} \\ \sum_{\lambda} \mu_{\lambda} &= 1 \\ q &\leq 1 \\ \forall \lambda, \mu_{\lambda} \geq 0, q \geq 0 \end{cases}$$
(11)

where  $\{D_{\lambda}\}_{\lambda}$  is the set of all deterministic behaviors as described in [BCP<sup>+</sup>14] and 1 is the vector  $(0.5\cdots0.5)^T$  that represents a fully random noise. If the result is q=0, then the probability vector P describes a local model, otherwise, if q>0, the probability vector describes a quantum behavior.

The dual program of the program (11) is the following:

$$\max \quad -\Gamma \cdot P + \phi - \omega \qquad \text{and expectation value can be decomposed}$$

$$s.t. \begin{cases} -\Gamma \cdot D_{\lambda} + \phi & \leq 0, \forall \lambda \in \{1, \cdots, \Delta^{2m}\} \text{ terms of measurement operators such that:} \\ \Gamma \cdot (\mathbb{1} - P) - \omega & \leq 1 \\ \forall \lambda, \gamma_{\lambda} \in \mathbb{R}, \quad \omega \geq 0, \quad \phi \in \mathbb{R} \end{cases}$$

$$\langle A_{x} \otimes B_{y} \rangle = \langle (M_{1|x} - M_{-1|x}) \otimes (M_{1|y} - M_{-1|y}) \rangle$$

$$\langle A_{x} \otimes B_{y} \rangle = \langle (M_{1|x} - M_{-1|x}) \otimes (M_{1|y} - M_{-1|x}) \rangle$$

$$\langle A_{x} \otimes B_{y} \rangle = \langle (M_{1|x} - M_{-1|x}) \otimes (M_{1|y} - M_{-1|x}) \rangle$$

$$\langle A_{x} \otimes B_{y} \rangle = \langle (M_{1|x} - M_{-1|x}) \otimes (M_{1|y} - M_{-1|x}) \rangle$$

$$\langle A_{x} \otimes B_{y} \rangle = \langle (M_{1|x} - M_{-1|x}) \otimes (M_{1|y} - M_{-1|x}) \rangle$$

$$\langle A_{x} \otimes B_{y} \rangle = \langle (M_{1|x} - M_{-1|x}) \otimes (M_{1|y} - M_{-1|x}) \rangle$$

$$\langle A_{x} \otimes B_{y} \rangle = \langle (M_{1|x} - M_{-1|x}) \otimes (M_{1|y} - M_{-1|x}) \rangle$$

$$\langle A_{x} \otimes B_{y} \rangle = \langle (M_{1|x} - M_{-1|x}) \otimes (M_{1|y} - M_{-1|x}) \rangle$$

$$\langle A_{x} \otimes B_{y} \rangle = \langle (M_{1|x} - M_{-1|x}) \otimes (M_{1|y} - M_{-1|x}) \rangle$$

$$\langle A_{x} \otimes B_{y} \rangle = \langle (M_{1|x} - M_{-1|x}) \otimes (M_{1|y} - M_{-1|x}) \rangle$$

where  $\Gamma = (\gamma_1 \cdots \gamma_{\Delta^{2m}})^T$ . We need to discuss the above program.

### Non-linear programming to determine a probability vector

Considering the linear program (11), one can transform the vector P into a vector of variables. Adding the constraints for P to be a probability vector, one can solve the associated non-linear program. It yields the following non-linear program

$$\max CHSH 
s.t. \begin{cases}
(1-q)P + q\mathbb{1} &= \sum_{\lambda} \mu_{\lambda} D_{\lambda} \\
\sum_{\lambda} \mu_{\lambda} &= 1 \\
q &\leq 1 \\
\forall \lambda, \mu_{\lambda} \geq 0, q \geq 0
\end{cases}$$
(13)

where CHSH is the CHSH inequality:

$$\langle A_0 \otimes B_0 \rangle + \langle A_1 \otimes B_0 \rangle + \langle A_0 \otimes B_1 \rangle - \langle A_1 \otimes B_1 \rangle$$
(14)

and expectation value can be decomposed in

$$\langle A_x \otimes B_y \rangle = \langle (M_{1|x} - M_{-1|x}) \otimes (M_{1|y} - M_{-1|y}) \rangle$$
(15)

and considering

$$p(ab|xy) = \langle M_{a|x} \otimes M_{b|y} \rangle \tag{16}$$

where each p(ab|xy) becomes a variable.

This can be resolved by Gurobi Solver [Gur22] even though the program found is nonconvex.

Note: It should be possible to add a constraint to ensure the resulting probability vector P stays in Q to get the maximal violation of the CHSH inequality. Otherwise, the maximal value is 4 meaning that the probability vector is in  $\mathcal{NS}$ . Considering the additional constraint  $CHSH < 2\sqrt{2}$  (that ensure that P is at most in Q), we obtain in P the probabilities that maximally violate the CHSH inequality, and q = 0.91which is consistent with what we could expect considering the program (11).

### References

- [ŠB20] Ivan Šupić and Joseph Bowles, Self-testing of quantum systems: a review, Quantum 4 (2020Sep), 337.
- [BCP+14] Nicolas Brunner, Daniel Cavalcanti, Stefano Pironio, Valerio Scarani, and Stephanie Wehner, Bell nonlocality, Reviews of Modern Physics 86 (2014Apr), no. 2, 419–478.
  - [Gur22] Gurobi Optimization, LLC, Gurobi Optimizer Reference Manual, 2022.
  - [Sca15] Valerio Scarani, The device-independent outlook on quantum physics (lecture notes on the power of bell's theorem), 2015.
- [ZCP<sup>+</sup>18] Wen-Hao Zhang, Geng Chen, Xing-Xiang Peng, Xiang-Jun Ye, Peng Yin, Ya Xiao, Zhi-Bo Hou, Ze-Di Cheng, Yu-Chun Wu, Jin-Shi Xu, Chuan-Feng Li, and Guang-Can Guo, Experimentally robust self-testing for bipartite and tripartite entangled states, Physical Review Letters 121 (2018Dec), no. 24.