



CS639: Data Management for Data Science

Lecture 8: Reasoning about Scale
& The MapReduce Abstraction

Theodoros Rekatsinas

Logistics/Announcements

- Submission template for PA2
- Bonus problem for PA2
- Other questions on PA2?

Today's Lecture

1. Scalability and Algorithmic Complexity
2. Data-Parallel Algorithms
3. The MapReduce Abstraction

1. Scalability and Algorithmic Complexity

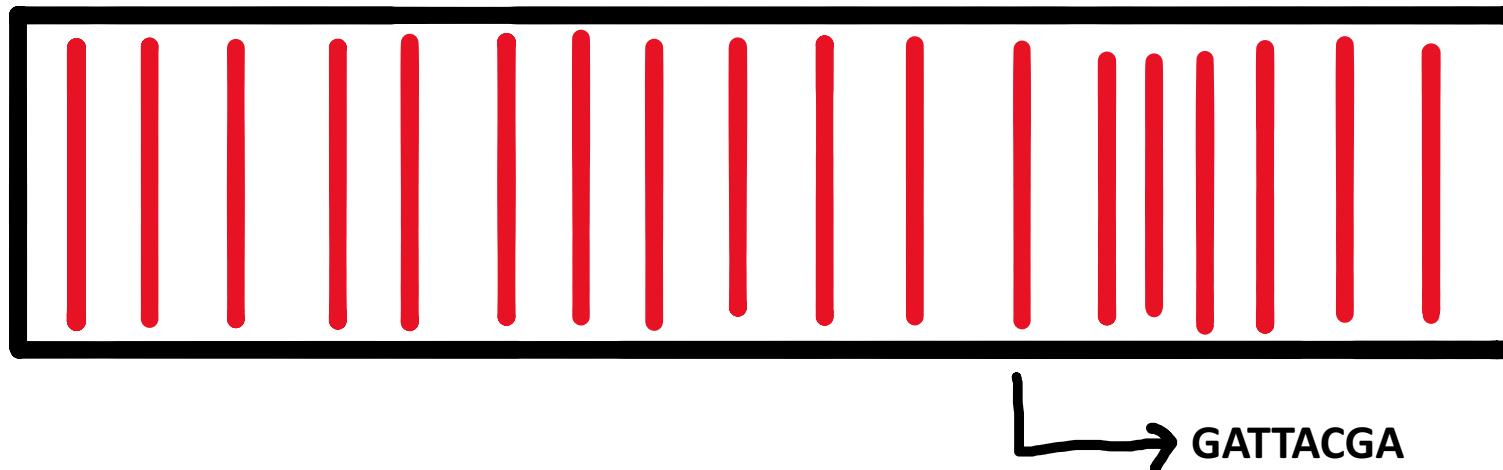
What does scalable mean?

- Operationally:
 - Works even if the data does not fit in main memory
 - Use all available resources (cores/memory) on a single node (aka **scale up**)
 - Can make use of 1000s of cheap computers (cloud) – elastic (aka **scale out**)
- Algorithmically:
 - If you have N data items you should not perform more than N^m operations (polynomial complexity)
 - In many cases it should be $N * \log(N)$ operations (streaming or too large data)
 - If you have N data items, you must do no more than N^m/k operations for some large k (k = number of cores/threads)

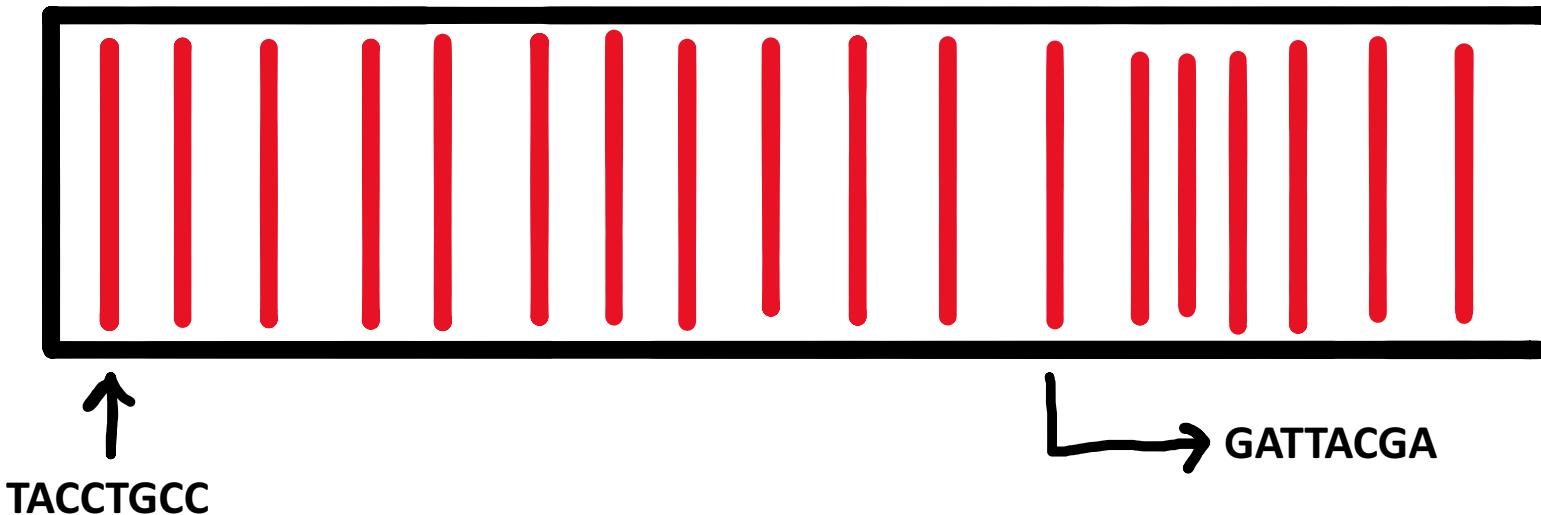
A sketch of algorithmic complexity

- Example: Find matching string sequences
- Given a set of string sequences
- Find all sequences equal to “GATTACGA”

Example: Find matching string sequences

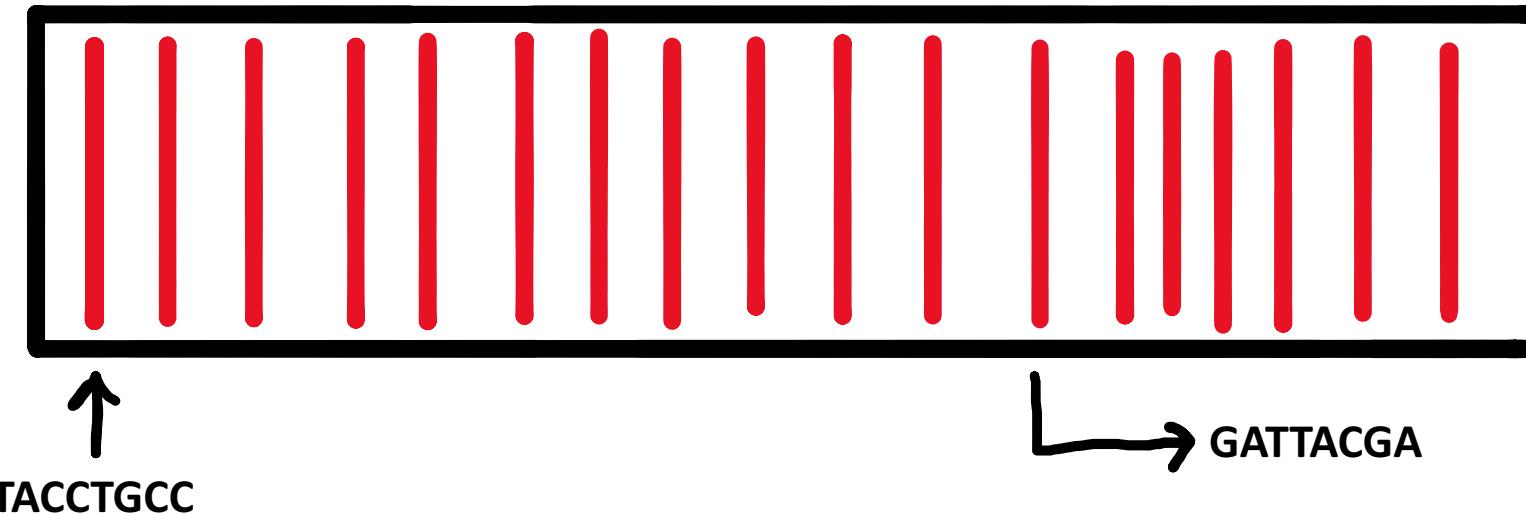


Example: Find matching string sequences



Time = 0: TACCTGCC ? GATTACGA

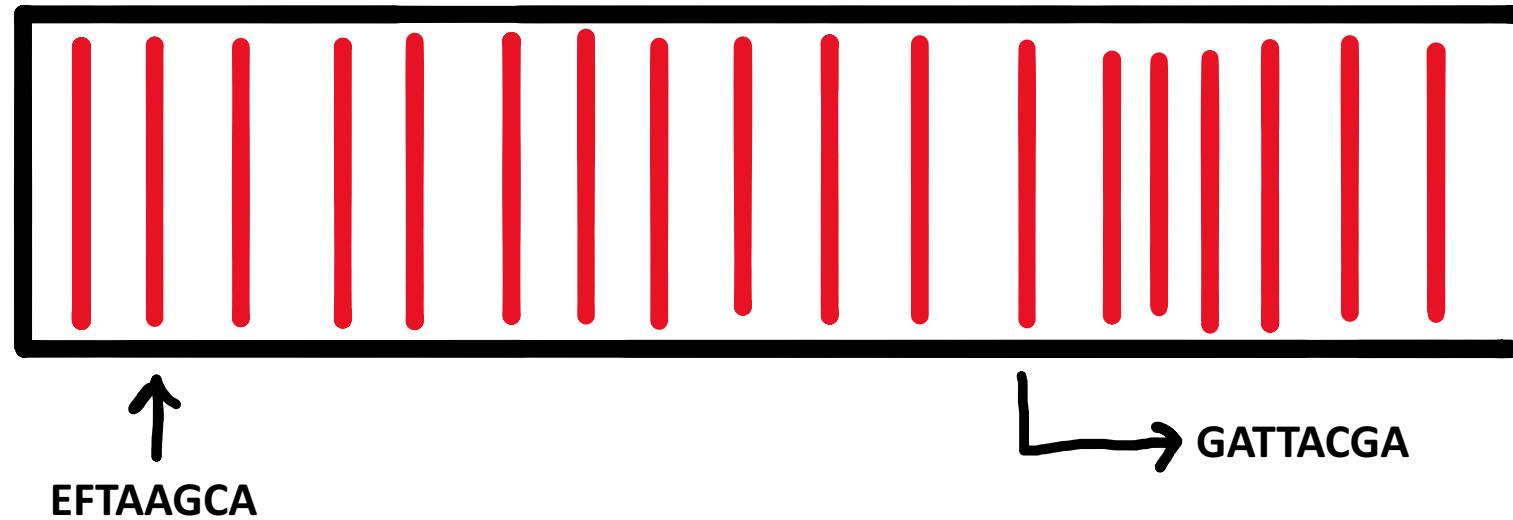
Example: Find matching string sequences



Time = 0: TACCTGCC ? GATTACGA

No move cursor to next data entry

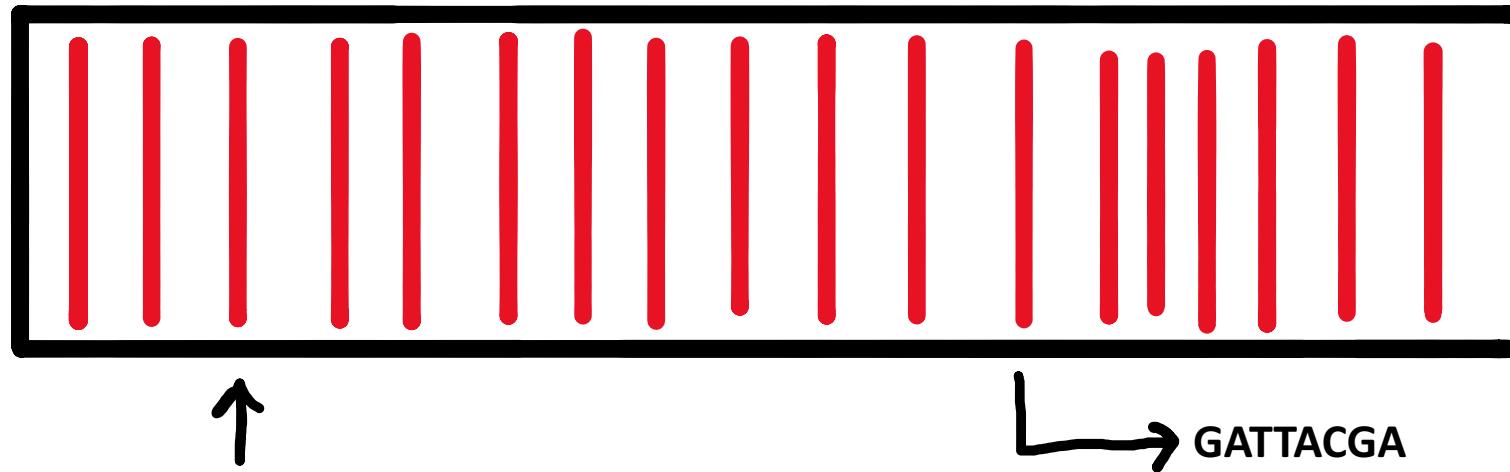
Example: Find matching string sequences



Time = 1: EFTAAGCA ? GATTACGA

No move cursor to next data entry

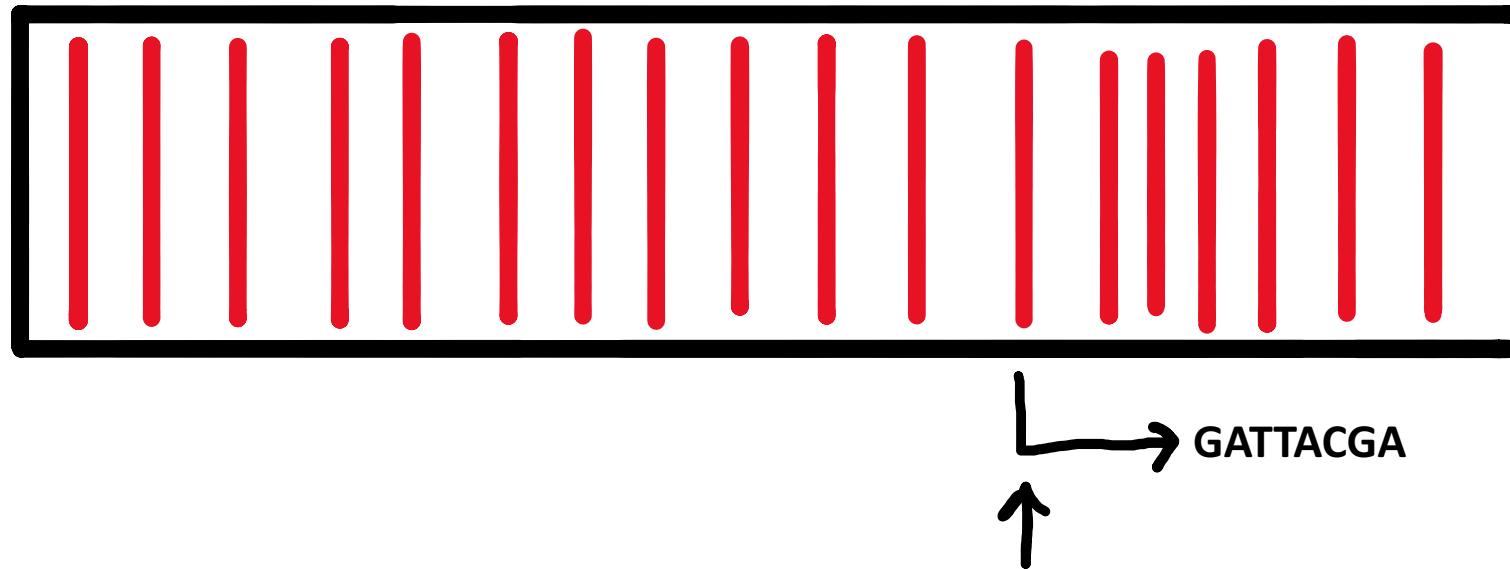
Example: Find matching string sequences



Time = 2: XXXXXXX ? GATTACGA

No move cursor to next data entry

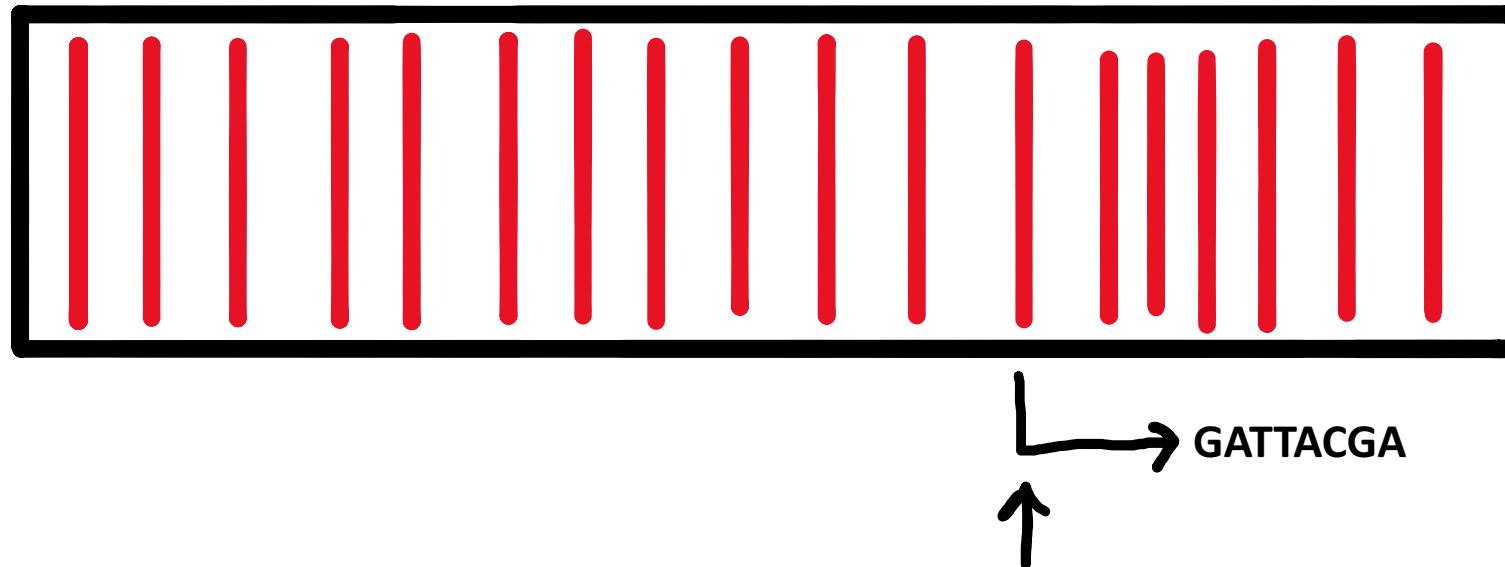
Example: Find matching string sequences



Time = n: GATTACGA ? GATTACGA

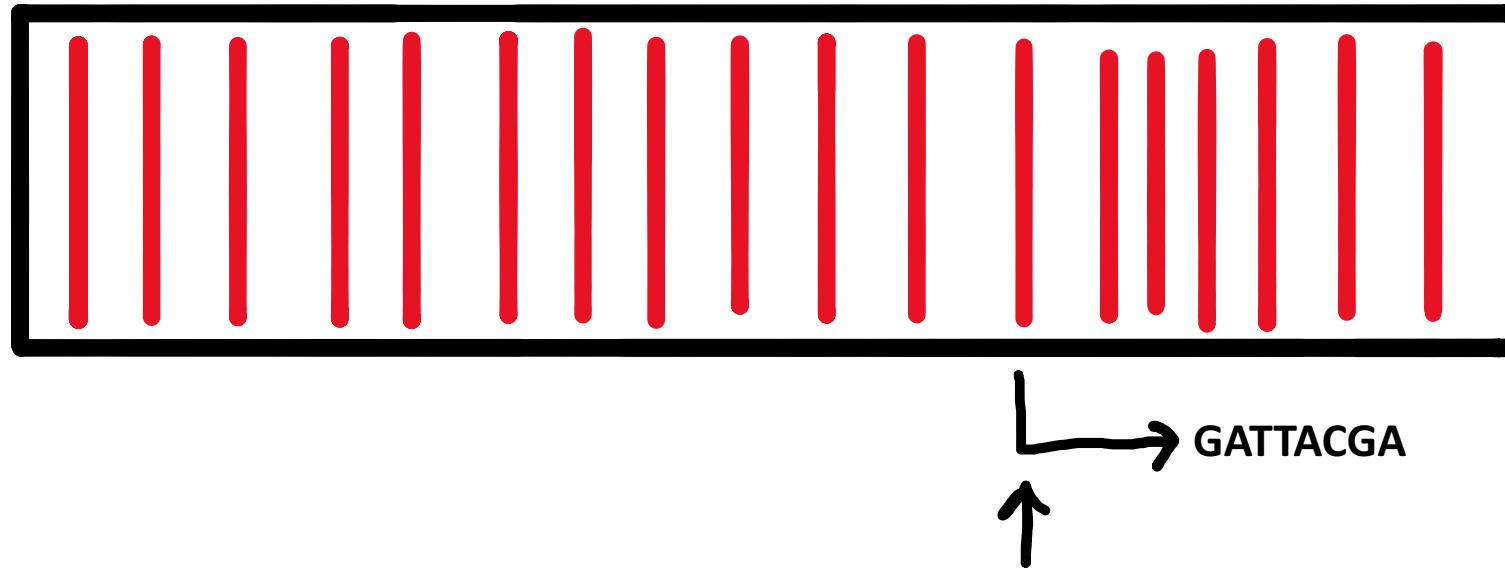
Yes! Output matching sequence

Example: Find matching string sequences



If we have 40 records we need to perform 40 comparisons

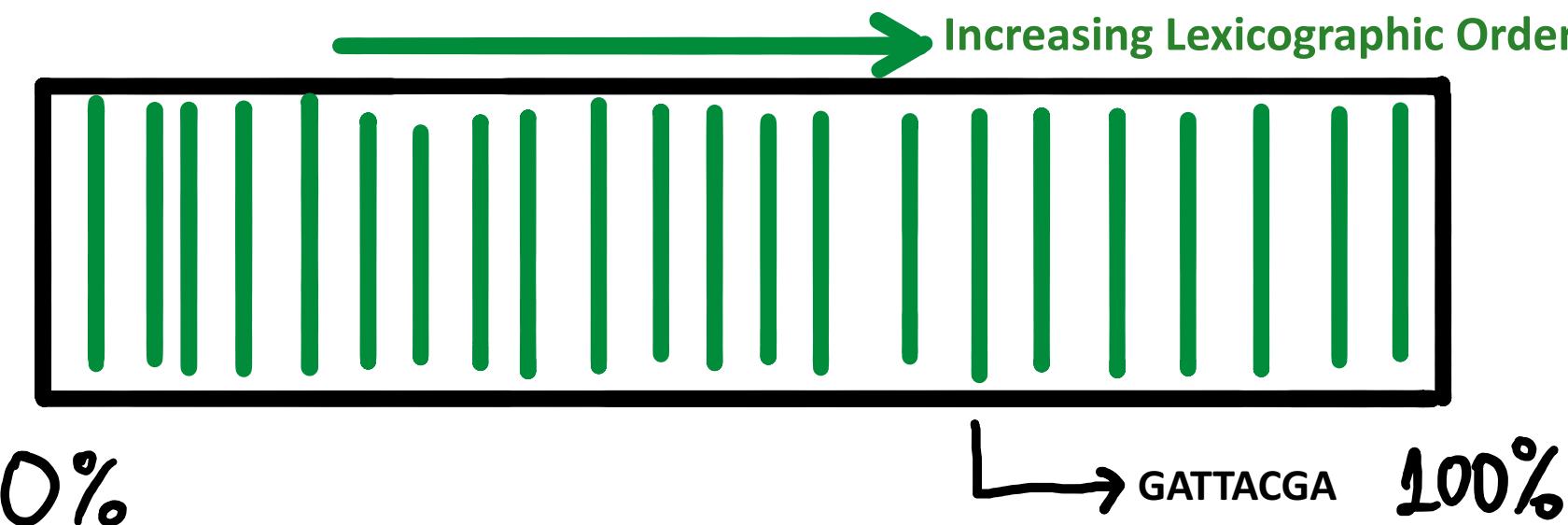
Example: Find matching string sequences



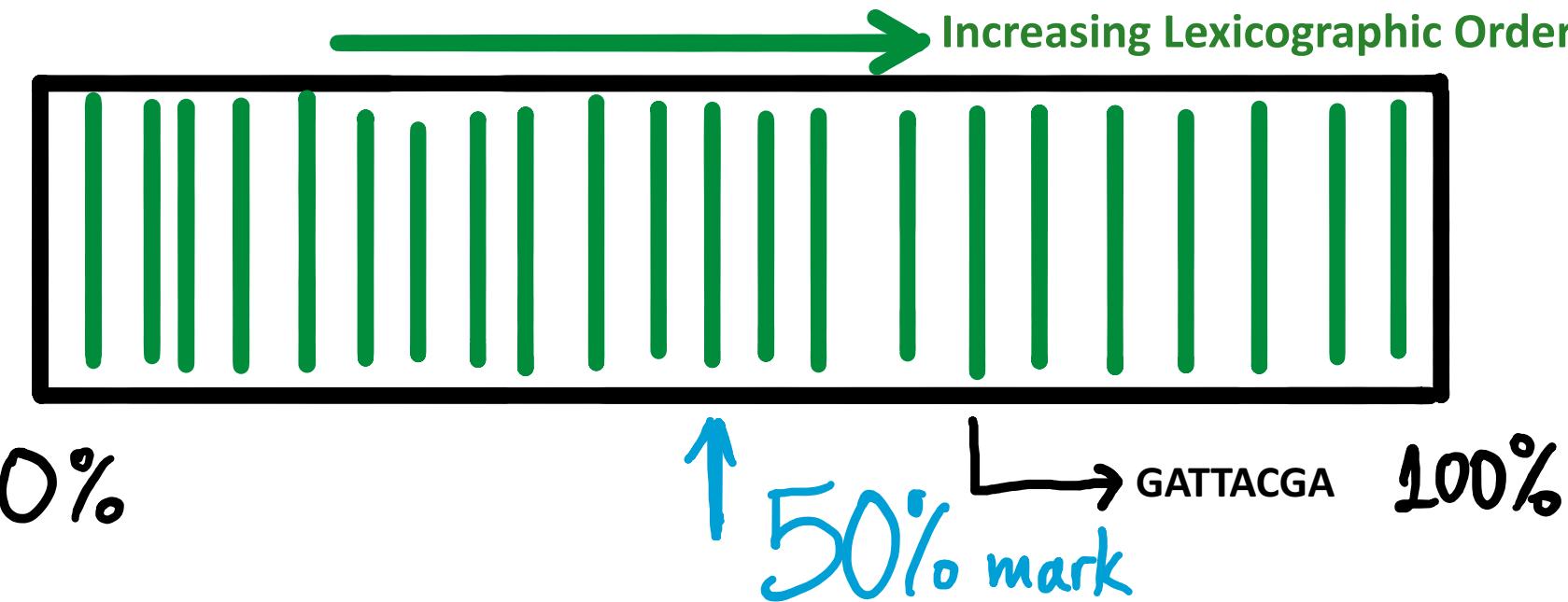
For N records we perform N comparisons

The algorithmic complexity is order N: O(N)

What if we knew the sequences are sorted

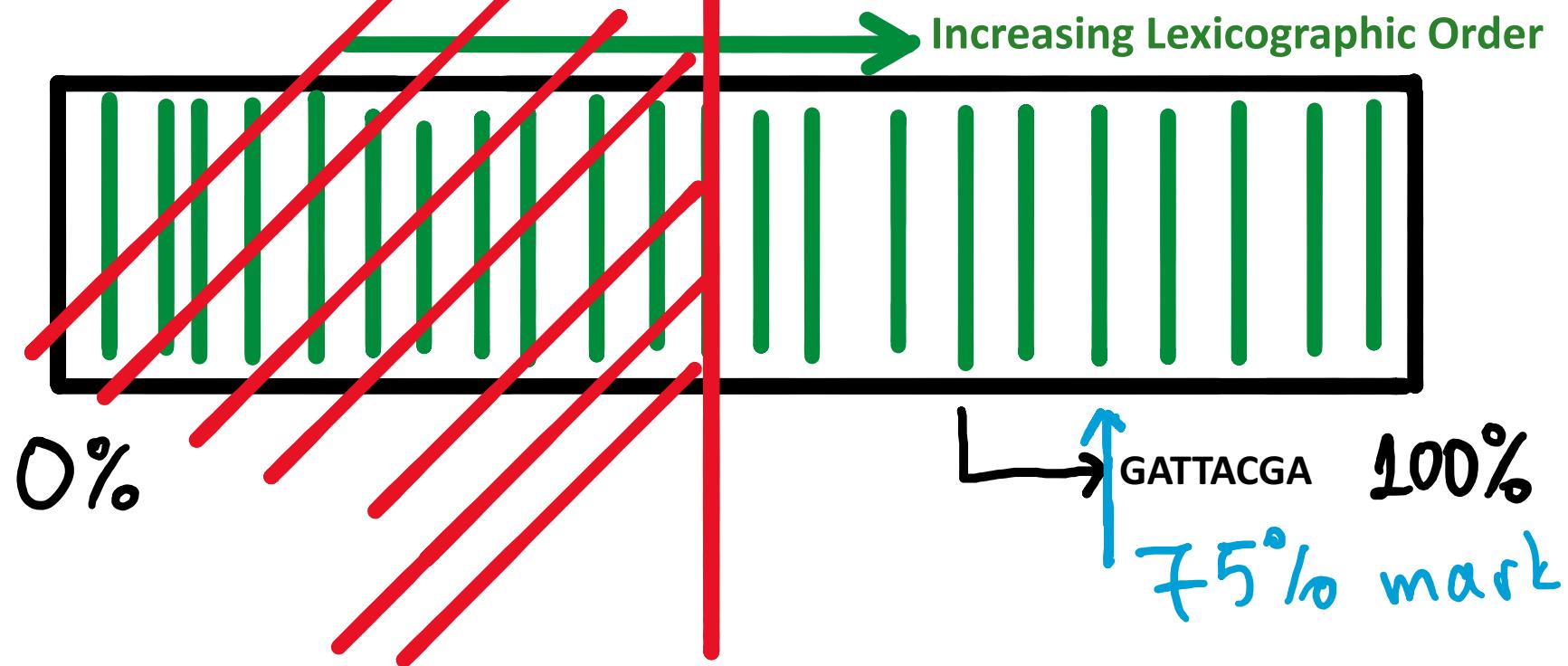


What if we knew the sequences are sorted



Time = 0: Start at 50% mark CTGTACA < GATTACGA

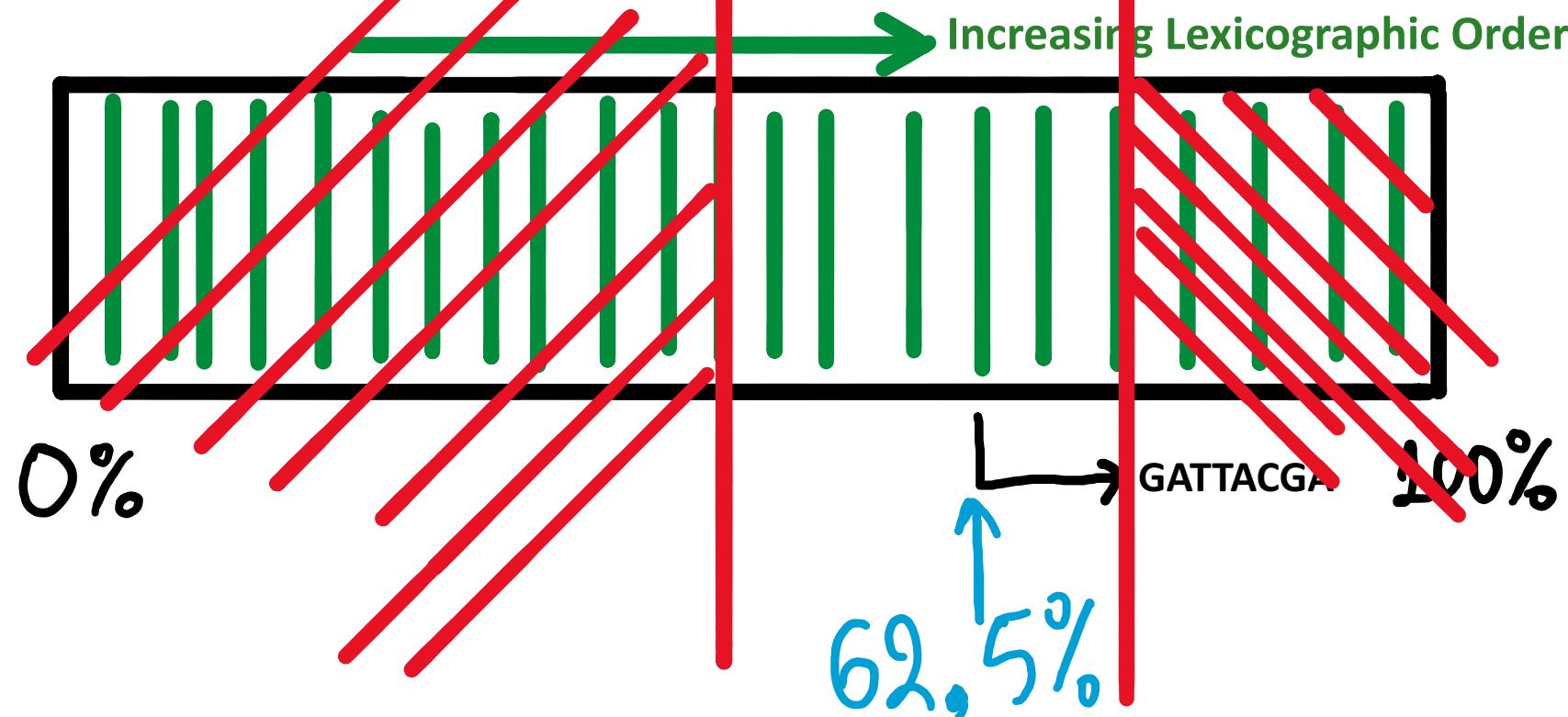
What if we knew the sequences are sorted



Time = 1: Start at 50% mark $\text{CTGTACA} < \text{GATTACGA}$

Skip to 75% mark (you know your sequence is in the second half)

What if we knew the sequences are sorted

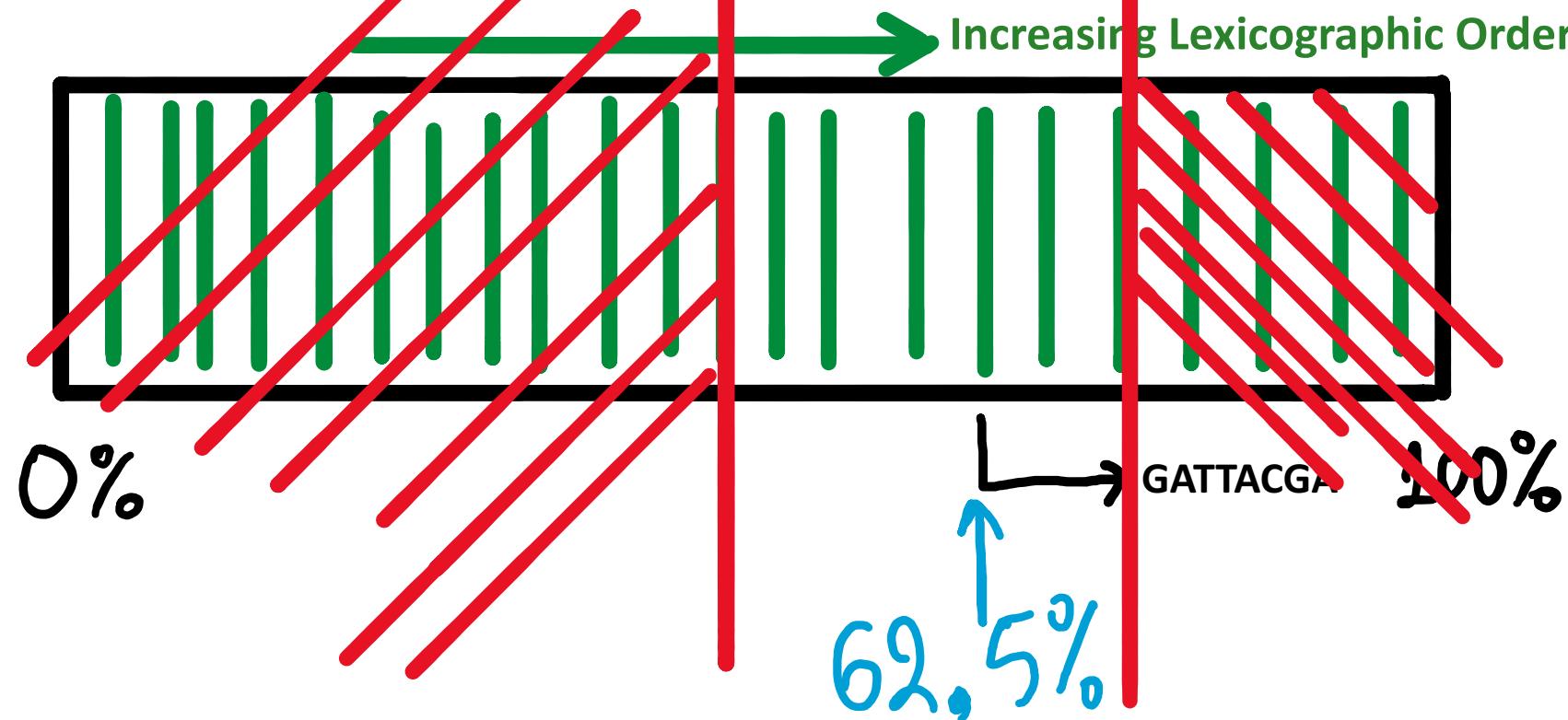


Time = 2: We are at the 75% mark $\text{TTGTCCA} > \text{GATTACGA}$

Skip to 62.5% mark Match: $\text{GATTACGA} = \text{GATTACGA}$

We find our sequence in three steps. Now we can scan entries sequentially.

What if we knew the sequences are sorted



How many comparisons?

For N records we did $\log(N)$ comparisons

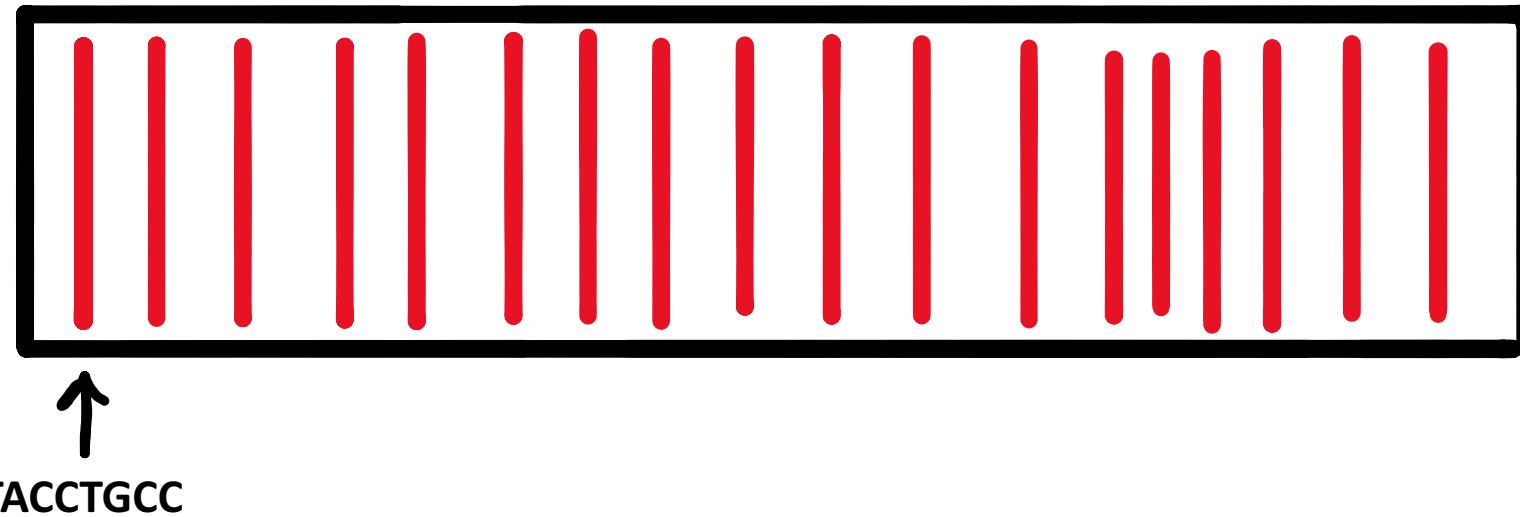
The algorithm has complexity $O(\log(N))$ – much better scalability

2. Data-Parallel Algorithms

New task: Trim string sequences

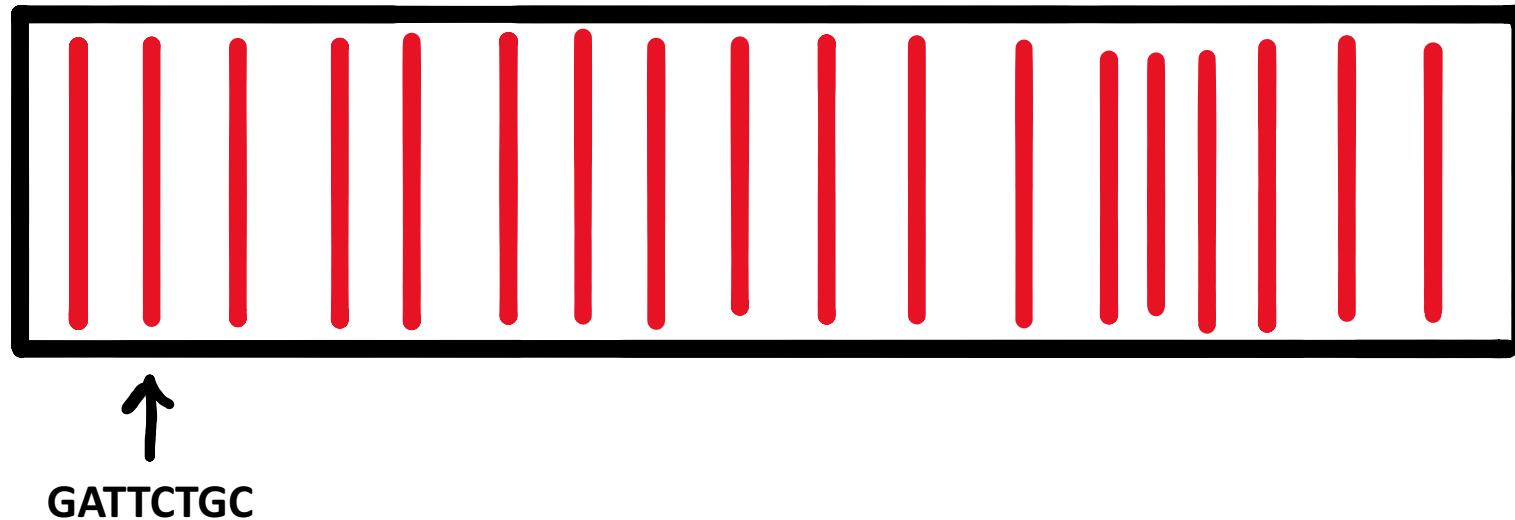
- Given a set of string sequences
- Trim the final n characters of each sequence
- Generate a new dataset

New task: Trim string sequences (last 3 chars)



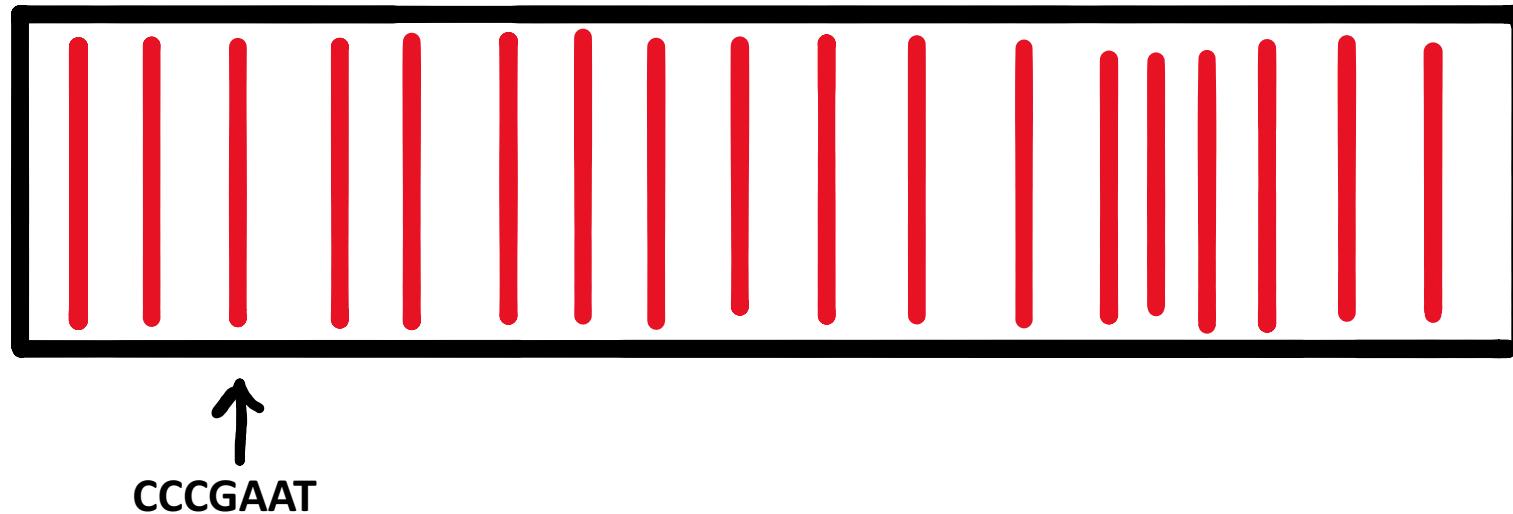
Time = 0: TACCTGCC -> TACCTG

New task: Trim string sequences (last 3 chars)



Time = 1: GATTCTGC -> GATTC

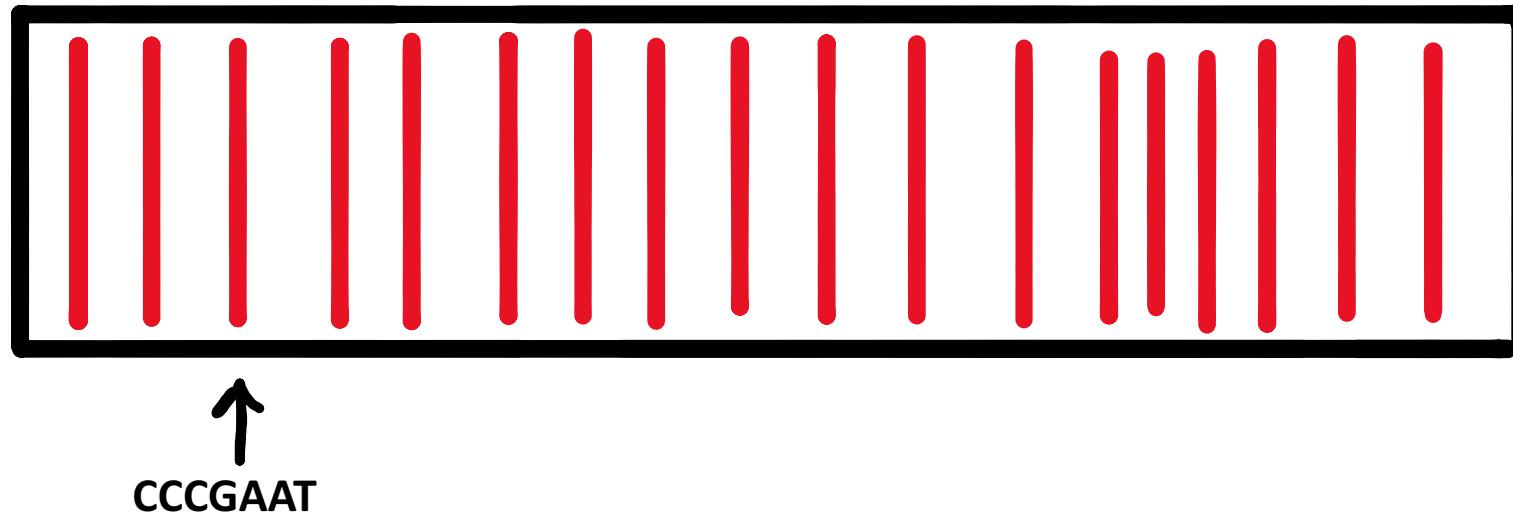
New task: Trim string sequences (last 3 chars)



Time = 2: CCCGAAT -> CCCG

Can we use a data structure to speed this operation?

New task: Trim string sequences (last 3 chars)

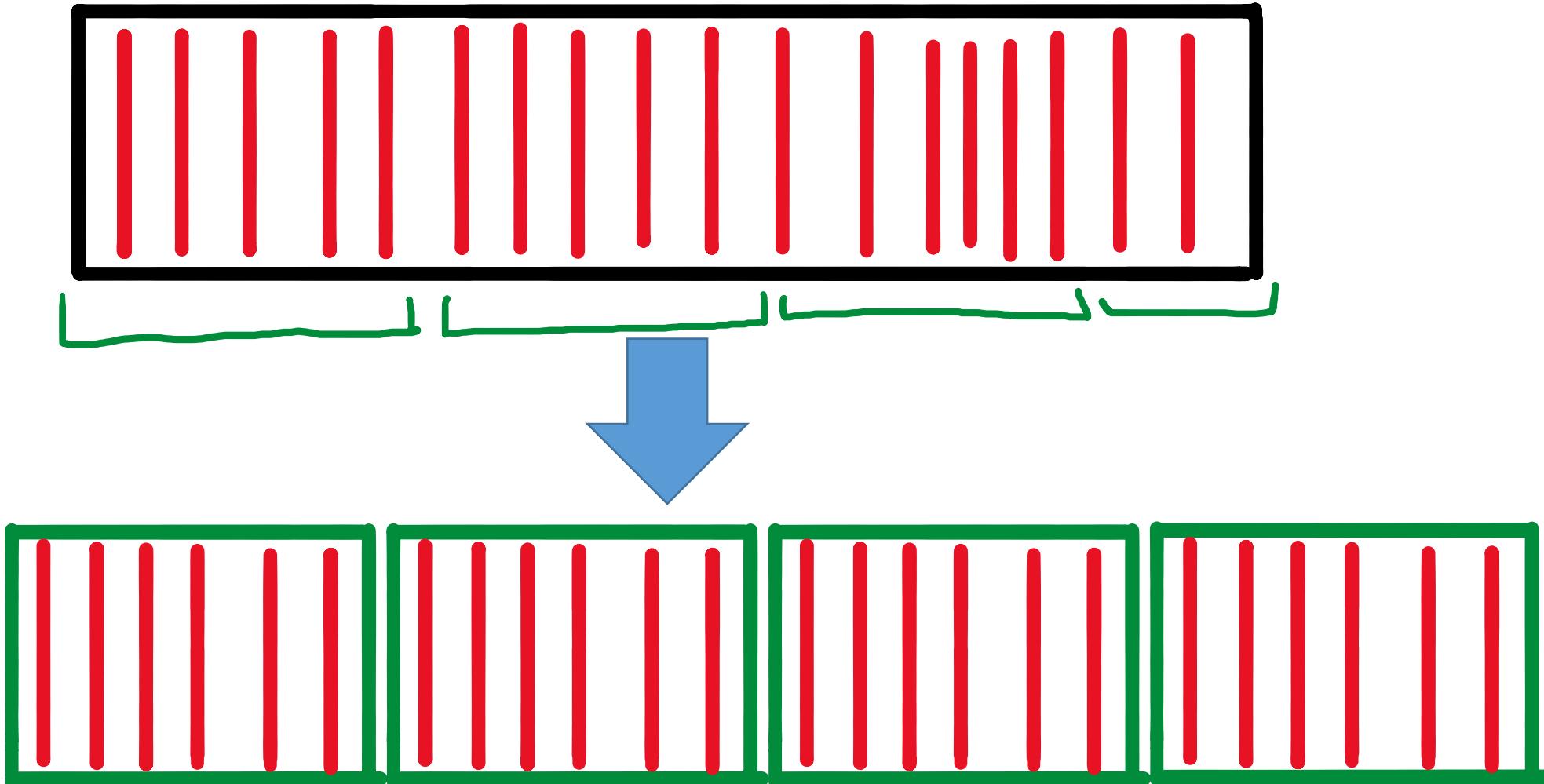


Time = 2: CCCGAAT -> CCCG

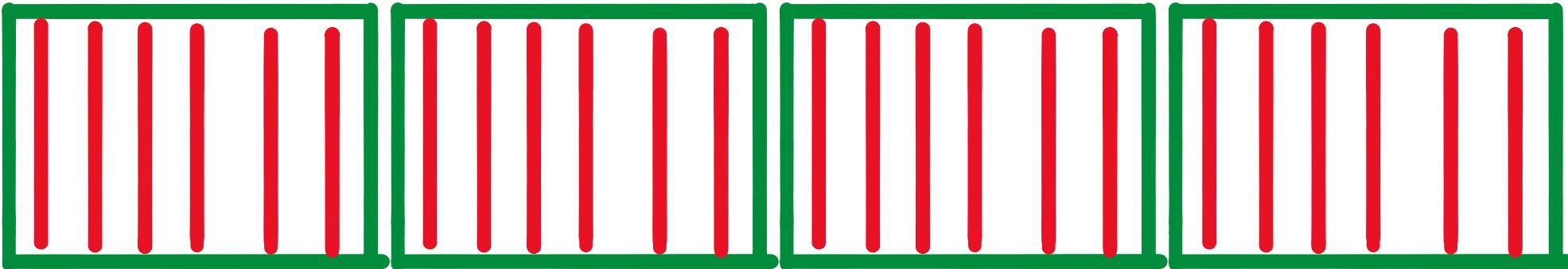
Can we use a data structure to speed this operation?

No. We have to touch every record! The task is $O(N)$.

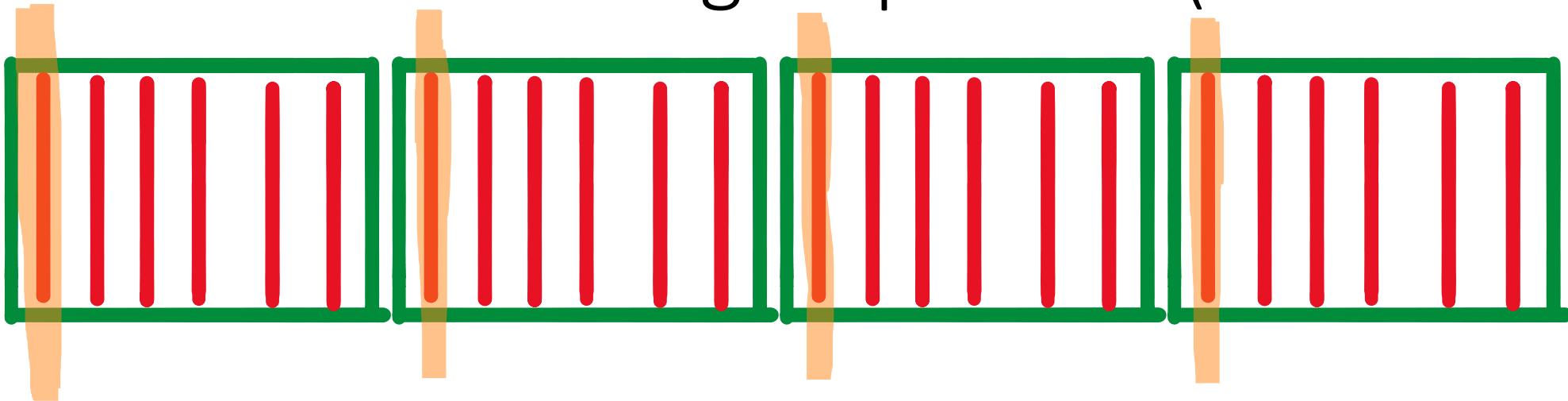
New task: Trim string sequences (last 3 chars)



New task: Trim string sequences (last 3 chars)

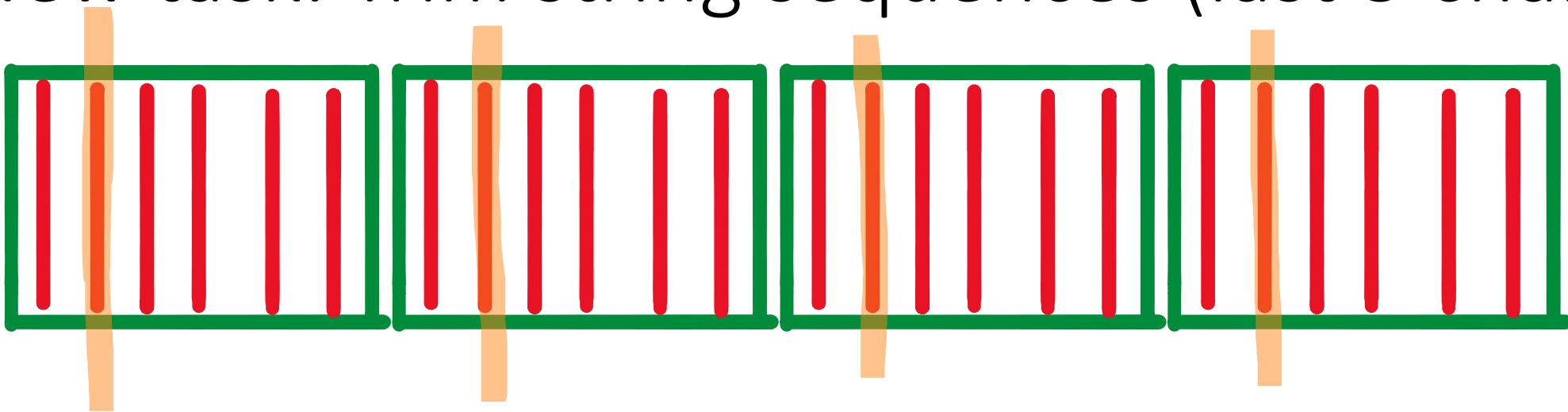


New task: Trim string sequences (last 3 chars)



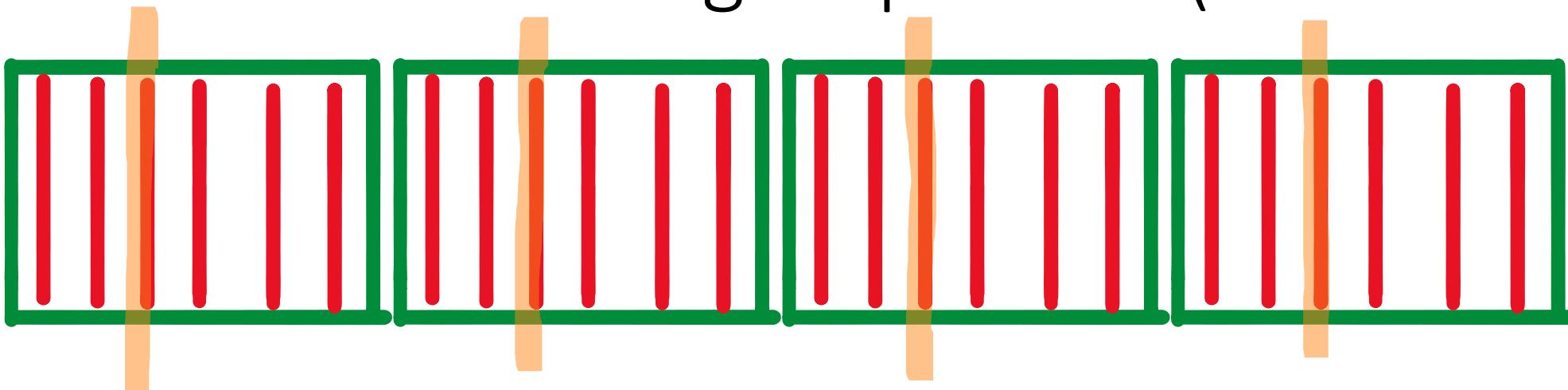
Time = 1: Process first element of each group

New task: Trim string sequences (last 3 chars)



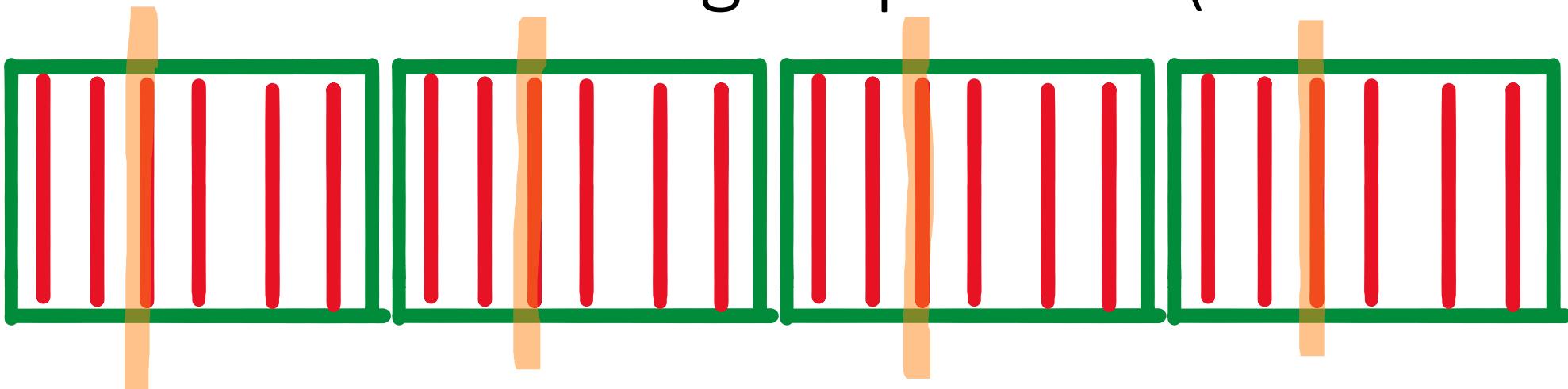
Time = 2: Process second element of each group

New task: Trim string sequences (last 3 chars)



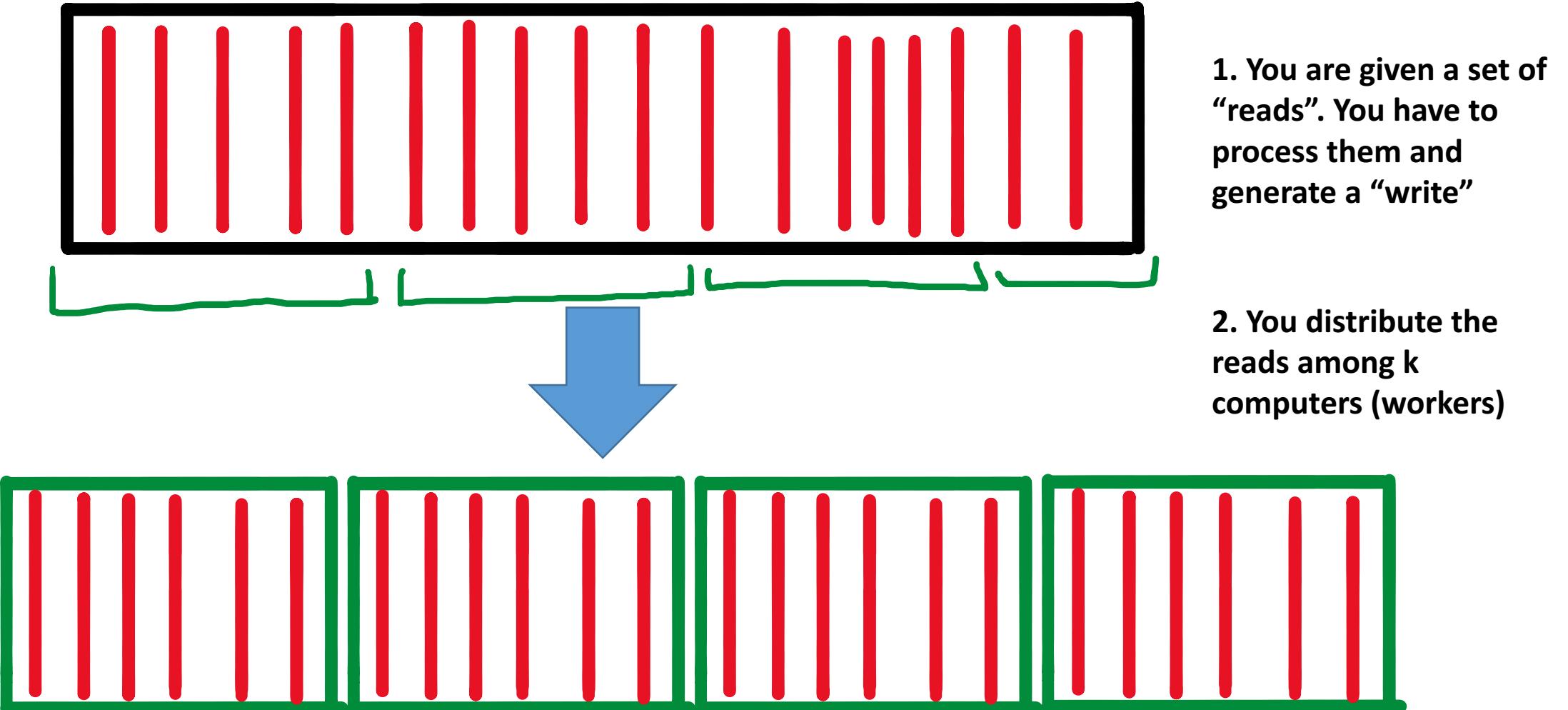
Time = 3: Process third element of each group
Etc.. How much time does this take?

New task: Trim string sequences (last 3 chars)



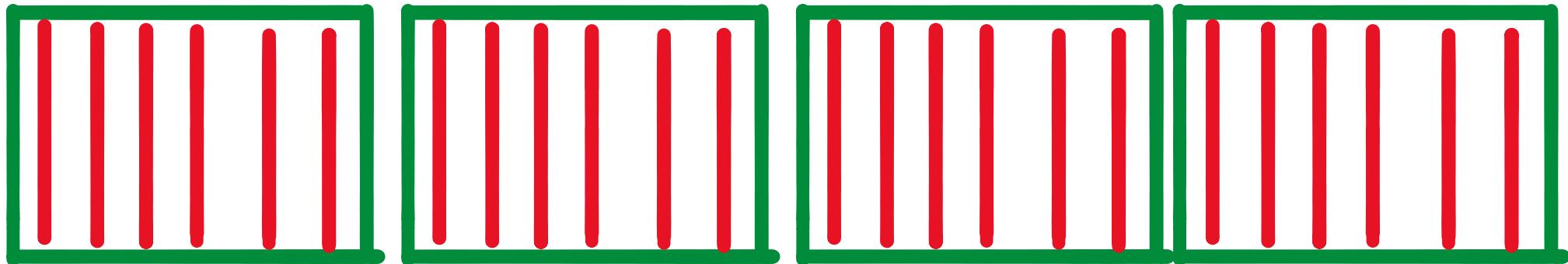
We only need $O(N/k)$ operations where k is the number of groups (workers)

Schematic of Parallel Algorithms



Schematic of Parallel Algorithms

2. You distribute the reads among k computers (workers)



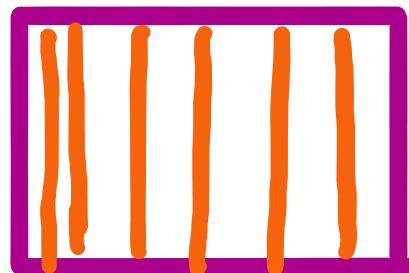
Function f

Function f

Function f

Function f

3. Apply function f to each read (for every item in each chunk)



4. Obtain a big distributed set of outputs

Applications of parallel algorithms

- Convert TIFF images to PNG
- Run thousands of simulations for different model parameters
- Find the most common word in each document
- Compute the word frequency of every word in a single document
- Etc....

Applications of parallel algorithms

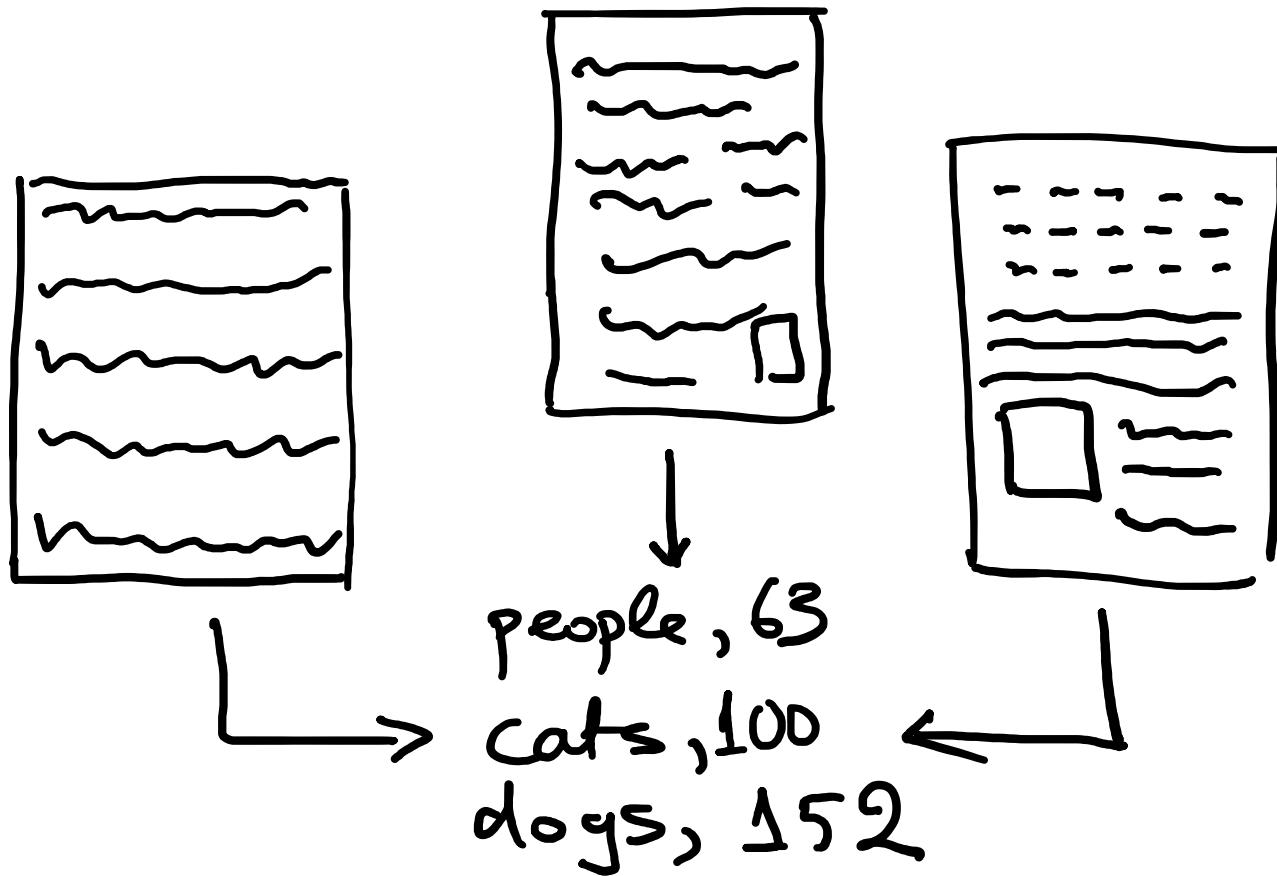
- Convert TIFF images to PNG
- Run thousands of simulations for different model parameters
- Find the most common word in each document
- Compute the word frequency of every word in a single document
- Etc....
- There is a common pattern in all these applications

Applications of parallel algorithms

- A function that *maps* a string to a trimmed string
- A function that *maps* a TIFF images to a PNG image
- A function that *maps* a set of parameters to simulation results
- A function that *maps* a document to its most common word
- A function that *maps* a document to a histogram of word frequencies

Applications of parallel algorithms

- What if we want to compute the word frequency across *all* documents?



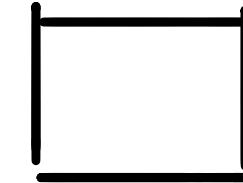
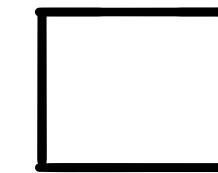
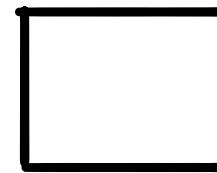
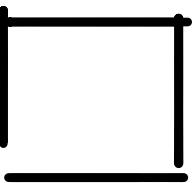
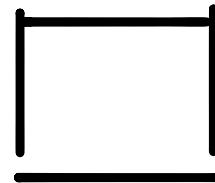
3. The MapReduce Abstraction

Compute the word frequency across 5M documents

Millions of Documents



Distribute among k workers



for each doc return map
(word,freq) pairs

|

map

|

map

|

map

|

map

|

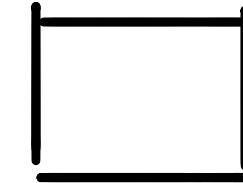
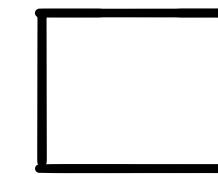
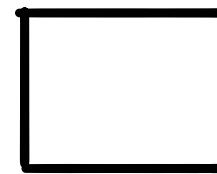
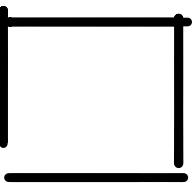
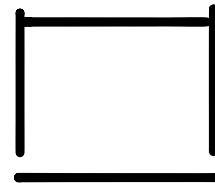
map

Compute the word frequency across 5M documents

Millions of Documents



Distribute among k workers



for each doc return map
(word,freq) pairs

map
↓

map
↓

map
↓

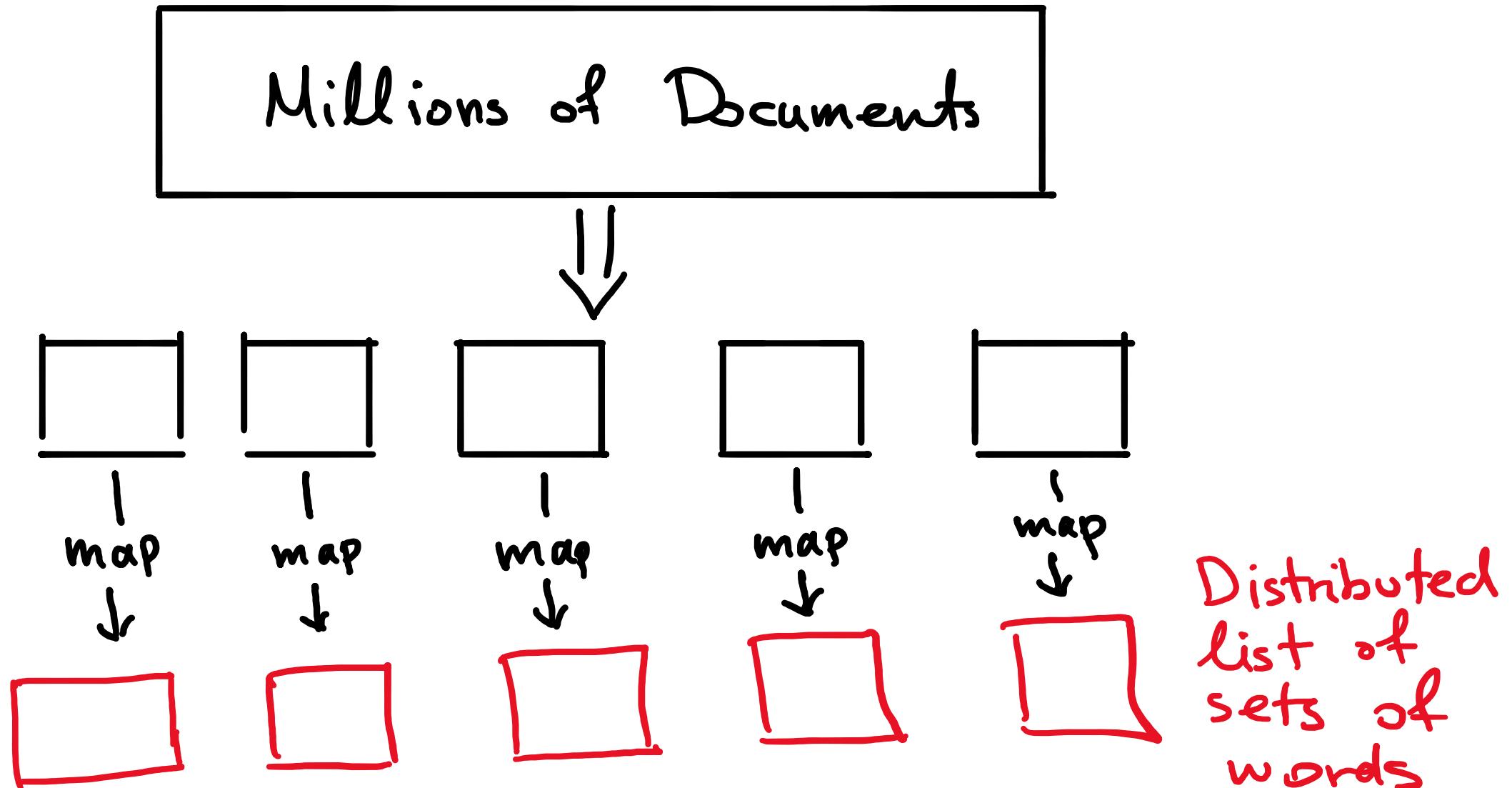
map
↓

Then what?

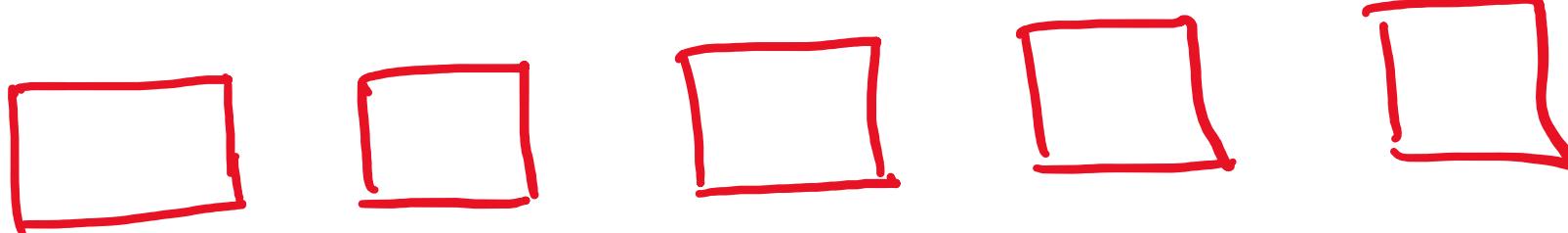
Challenge: in this task

- How can we make sure that a single computer has access to every occurrence of a given word regardless of which document it appeared in?
- *Ideas?*

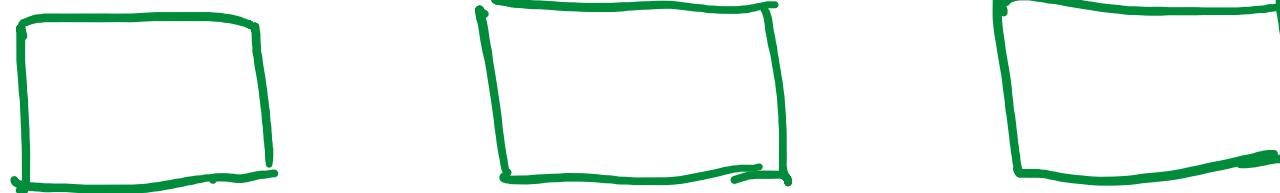
Compute the word frequency across 5M documents



Compute the word frequency across 5M documents

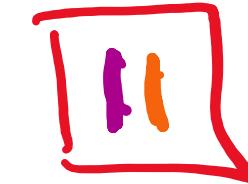
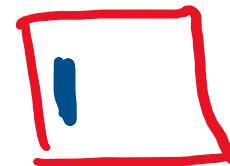
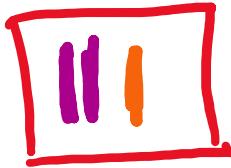
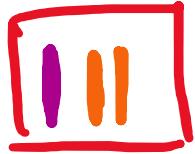
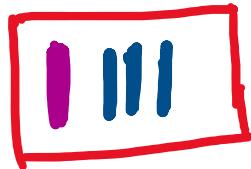


Distributed
list of
sets of
words

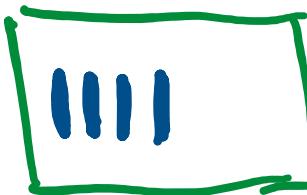
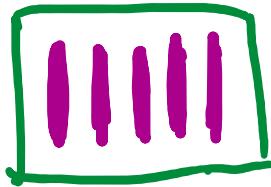


Workers to
aggregate
frequency

Compute the word frequency across 5M documents

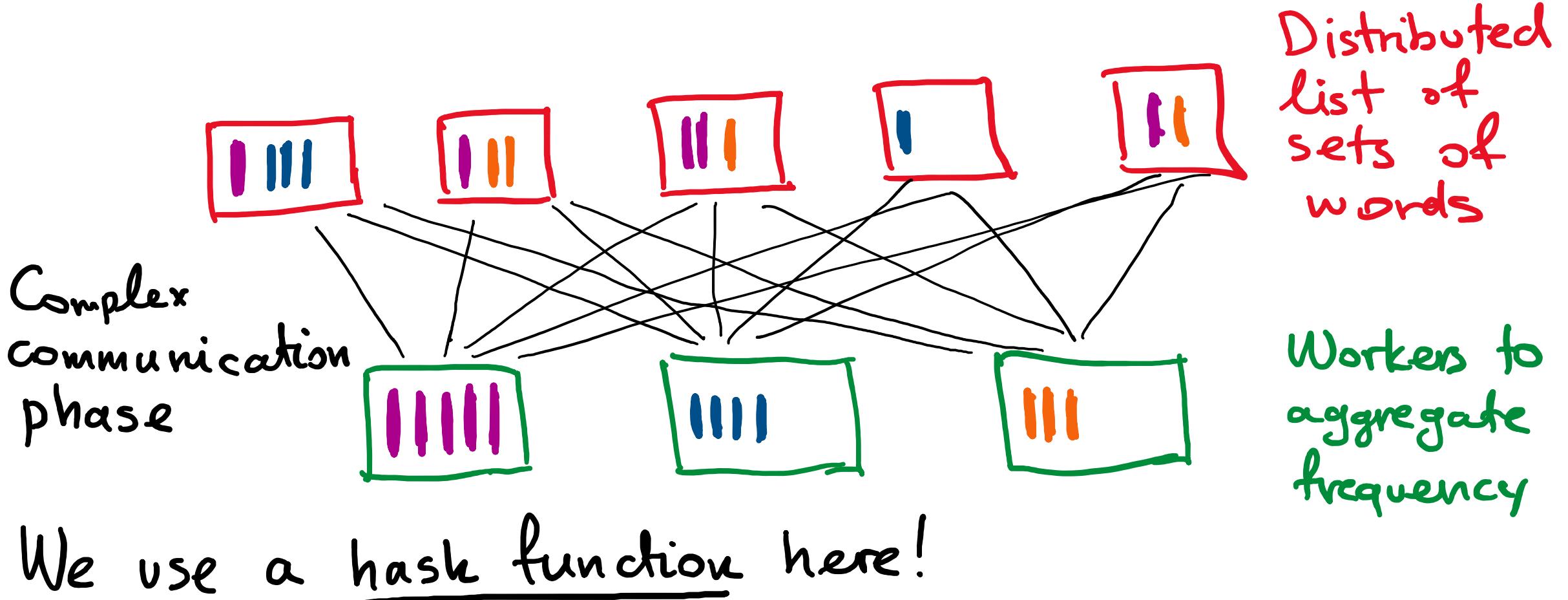


Distributed
list of
sets of
words



Workers to
aggregate
frequency

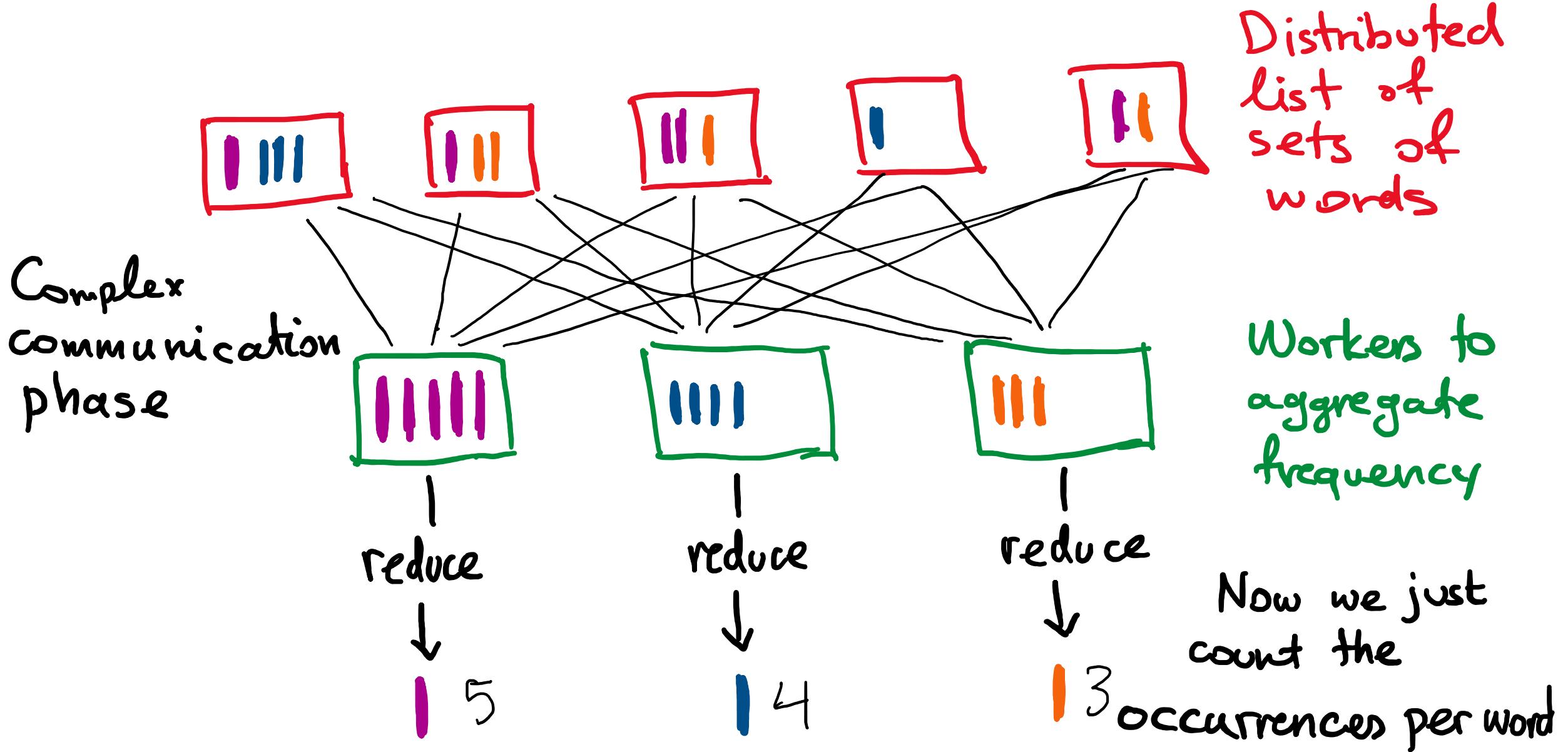
Compute the word frequency across 5M documents



We use a hash function here!

A hash function is any function that can be used to map data of arbitrary size to a data of a fixed size

Compute the word frequency across 5M documents



The Map Reduce Abstraction for Distributed Algorithms

