

Security Assessment & Formal Verification Final Report

Atlas

December 2024

Prepared for FastLane Labs



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© certora Project Summary

Project Scope

Project Name	Repository (link)	Latest Commit Hash	Platform
Atlas	https://github.com/FastLane-L abs/atlas	8127c0152aac7161d 8aa1e3b69d8adc24 e2755b4	EVM/Solidity 0.8

Project Overview

This document describes the specification and verification of Atlas using the Certora Prover and manual code review findings. The work was undertaken from November 19, 2024 to December 20, 2024.

The following contract list is included in our scope:

```
src/contracts/atlas/*
src/contracts/common/*
src/contracts/dapp/*
src/contracts/gasCalculator/*
src/contracts/helpers/Simulator.sol
src/contracts/helpers/Sorter.sol
src/contracts/libraries/*
src/contracts/types/*
```

The Certora Prover demonstrated that the implementation of the Solidity contracts above is correct with respect to the formal rules written by the Certora team. In addition, the team performed a manual audit of all the Solidity contracts. During the verification process and the manual audit, the Certora team discovered bugs in the Solidity contracts code, as listed on the following page.

Please note that a few more formal rules are not included in this report, as they were proven with an unreleased version of the Certora Prover. Once those rules are proven on a released version of the Certora Prover, we will add them to the next version of this document.

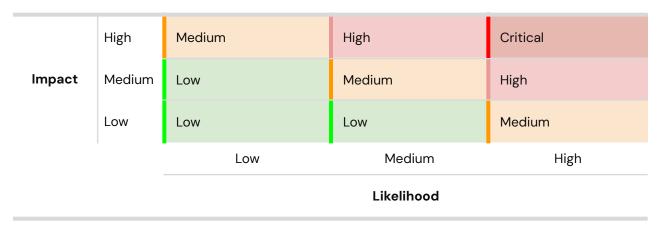


Findings Summary

The table below summarizes the findings of the review, including type and severity details.

Severity	Discovered	Confirmed	Fixed
Critical	0		
High	6	0	6
Medium	3	0	3
Low	4	1	3
Informational	6	2	4
Total	19	3	16

Severity Matrix





Detailed Findings

ID	Title	Severity	Status
H-01	Repeatable Temporary Freeze of User Assets	High	Fixed
H-02	Escrow duration can be bypassed and bonded balance can be withdrawn immediately	High	Fixed
H-03	_bundlerSurcharge will always be less than _maxBundlerRefund	High	Fixed
H-04	Bundler can get unfairly penalized and loses funds to winning solver due to gas limit	High	Fixed
H-05	Solver can frontrun bundler and make sure he becomes the winning solver by DOSing other solvers and allow MEV-like behavior	High	Fixed
H-06	Bundler-solvers can exploit users by setting unrealistic high gas limit	High	Fixed
M-01	Bundler wrongly pays Atlas surcharge when operation fails in certain situations	Medium	Fixed



M-02	Atlas might not receive storage refund	Medium	Fixed
M-03	Solver might fail when it should succeed due to incorrect bonded amount	Medium	Fixed
L-01	Remove Bundler surcharge in _updateAnalytics	Low	Fixed
L-02	Governance signatory behaves unexpectedly	Low	Fixed
L-03	False Update to Analytics	Low	Fixed
L-04	PerBlockLimit can lead to accounting exploit on Bundlers or Solvers	Low	Acknowledged



High Severity Issues

H-01 Repeatable Temporary Freeze of User Assets

Severity: High	Impact: High	Likelihood: Medium
Files: Escrow.sol	Status: Fixed	Violated Property: None

Description:

Attacker can cause a solver to not be able to withdraw his funds. By crediting the solver 1 wei with malicious bundler and update the solver's lastAccessedBlock

and the solver's timelock resets. Whenever the timelock is going to end, the attacker will give solver 1 wei with the _credit function so the solver cant ever withdraw his funds.

Recommendation:

Like in HIGH [5], remove _aData.lastAccessedBlock = uint32(block.number); from the _credit function. If the team wants to keep the functionality of coupling the solvers' winning actions, add it in a different place.

Customer's Response:

Fixed in commit fae5ba5



H-O2 Escrow duration can be bypassed and bonded balance can be withdrawn immediately

Severity: High	Impact: High	Likelihood: Medium
Files: Escrow.sol	Status: Fixed	Violated Property: None

Description:

The unbonding timelock exists to offer stability to off-chain calculations. However, it is possible to bypass this timelock, leading to off-chain calculations to fail and bundles to behave unexpectedly potentially causing a loss of funds to the Bundler because it would not get the gas refund as expected.

This is possible because a malicious actor could withdraw the atIETH without timelock by running their own protocol and bundler. By calling the borrow function as the winning solver, they can send the ETH away, and repay with the bonded amount, thus withdrawing without the unbonding timelock (and a very small fee). Usually they would need to wait or pay the atlas and bundler fees, but now they don't need to do either.

Recommendation:

Borrowed funds should be paid back on top of the bonded balance. More explicitly, the bonded balance should be considered only for the execution of the operation, not as a way to repay borrowed funds. Instead borrowed funds should be repaid through a different accounting system. This would prevent malicious actors from bypassing the unbonding period.

Customer's Response:

Fixed in commit 4d5c97b



H-03 _bundlerSurcharge will always be less than _maxBundlerRefund

Severity: High	Impact: High	Likelihood: Medium
Files: Escrow.sol	Status: Fixed	Violated Property: None

Description:

In GasAccounting.sol if (_bundlerSurcharge > _maxBundlerRefund) does not make sense during the execution flow that goes through _adjustAccountingForFees() function when there are no winning solvers. This is because _bundlerSurcharge will always be less than _maxBundlerRefund. Max bundler refund is 80% of the metacall gas cost and the _bundlerSurcharge is much less than 80% so this branch will never happen and bundlers will get this extra refund when atlas should get it.

Recommendation:

We recommend removing the if statement (_bundlerSurcharge > _maxBundlerRefund) { as it can never be reached (unless the bundler surcharge is unusually high). We propose to simply reward 20% of the original gas cost to Atlas and give the remaining to the bundler.

Customer's Response:

Fixed during the refactor of the Gas Accounting system, in commit fbalc9e



H-O4 Bundler can get unfairly penalized and loses funds to winning solver due to gas limit

Severity: High	Impact: High	Likelihood: Medium
Files: Escrow.sol	Status: Fixed	Violated Property: None

Description:

Taking in consideration the _adjustAccountingForFees() function (lines 387-404), let's assume we have 10 solvers. Each of them requests solverGasLimit as gas limit (and they have the bond to pay) but when making the actual call 9 of them revert and don't spend any gas. The last solver is successful. Now gasLeft is at least 9 * solverGasLimit but the _upperGasRemainingEstimate is 1 * solverGasLimit. The bundler will be penalized for this but it is not their fault.

Recommendation:

We recommend either dynamically deriving solverGasLimit from the total bonded among solvers or subtracting gasLeft from the gas limit of solvers that have reverted (have not spent any gas) or solvers that spent a significantly smaller percentage of it.

Customer's Response:

Fixed in commit 8de7b0d



H-05 Solver can frontrun bundler and make sure he becomes the winning solver by DOSing other solvers and allow MEV-like behavior

Severity: High	Impact: High	Likelihood: Medium
Files: Escrow.sol	Status: Fixed	Violated Property: None

Description:

In the credit function there is an update to _aData.lastAccessedBlock = uint32(block.number);.

A malicious solver can frontrun the bundler call and create a MEV-like arbitrage (with the userOp the arbitrage is big). Usually the solvers will bid the highest amount and extract the value the MEV attacker placed to create the arbitrage. But a malicious solver can frontrun the metacall and give small amounts (even 0) of gas to each solver in the bundle and make sure their next solverOp in the same block will revert.

For example by permissionlessly creating a scenario where all solvers failed the gasRefundBeneficiary gets credited.

The malicious solver can call Atlas with their own protocol and use the gasRefundBeneficiary address as the honest solver to DOS. Now in this block this solver cannot run. We can repeat this for all the solvers in the real bundle so the malicious solver is the only one left. Now the attacker can do this backrun with a very small bid (1 wei) and keep the rest of the profit as if Atlas does not exist.

This allows a **MEV-like** behavior which guarantees to the **MEV** that he will be the backrunner and thus allow him to commit the **sandwich attack**, this is exactly what Atlas should **prevent**.

Recommendation:



Remove _aData.lastAccessedBlock = uint32(block.number); from the _credit function. If you want to prevent solver from operating twice in that same block make sure that Atlas counts only operations that the solver approved.

Customer's Response:

Fixed in commit fae5ba5



H-06 Bundler-solvers can exploit users by setting unrealistic high gas limit

Severity: High	Impact: High	Likelihood: Medium
Files: Escrow.sol	Status: Fixed	Violated Property: None

Description:

Malicious bundler-solvers can send a high gas limit and because the _bundlerGasOveragePenalty goes to the winning solver there is no real penalty when the bundler is the winning solver.

Consider a bundler-solver with a high bond and the worst bid (1 wei) that sends the rest of the profit to an external address.

The attacker can give a very high gas limit for the tx. All the solvers fail and the winning solver (the bundler) is the one with the high bond. They paid the smallest fee possible and didn't get punished. Thus stealing all profits from Atlas and the user.

Recommendation:

Allow DappControl to establish a gas limit operation rather than having a gas limit that can be arbitrarily set by the bundler.

Customer's Response:

Fixed in commit 9fbf40d



Medium Severity Issues

M-O1 Bundler wrongly pays Atlas surcharge when operation fails in certain situations

Severity: Medium	Impact: High	Likelihood: Low
Files: Escrow.sol	Status: Fixed	Violated Property: None

Description:

Atlas surcharge is always included when t_writeoffs is increased.

Consequently, if the operation fails due to the bundler's fault or it fails due to solvers not having enough bonded amount, the bundler will end up paying for the Atlas surcharge. In the case of _bidFindingIteration the bundler is responsible for the onchain validation and there is a writeoff by _writeOffBidFindGasCost but Atlas still gets the surcharge on the bundler expense.

Recommendation:

In _adjustAccountingForFees() find atlasSurchargeWriteoffs from t_writeoffs then:

netAtlasGasSurcharge -= atlasSurchargeWriteoffs

adjustedWriteoffs -= atlasSurchargeWriteoffs

Customer's Response:

Fixed during the refactor of the Gas Accounting system, in commit fbalc9e



M-02 Atlas might not receive storage refund

Severity: Medium	Impact: High	Likelihood: Low
Files: Escrow.sol	Status: Fixed	Violated Property: None

Description:

If t_solverSurcharge is zero because all the solvers have failed and have bonded just exactly enough to cover their gas but nothing more, than the gas refund also goes to the bundler because all the logic that takes the refund from the bundler is inside the if block that happens when t_solverSurcharge is not zero.

Recommendation:

We believe the bonded amount should already consider the fees for the solver. Change Escrow.sol line 353 to:

```
if (_gasCost.withSurcharges(ATLAS_SURCHARGE_RATE, BUNDLER_SURCHARGE_RATE) <
    _solverBalance) {</pre>
```

Customer's Response:

Fixed during the refactor of the Gas Accounting system, in commit fbalc9e



M-03 Solver might fail when it should succeed due to incorrect bonded amount

Severity: Medium	Impact: High	Likelihood: Low
Files: Escrow.sol	Status: Fixed	Violated Property: None

Description:

When assessing whether the solver has enough balance to pay for the transaction, the surcharges for Atlas and Bundler are not accounted for. This could lead to a Solver running out of balance, even though theoretically their bonded amount should have been enough.

Recommendation:

We recommend the same fix as MEDIUM [2]:

Change Escrow.sol line 353 to:

if (_gasCost.withSurcharges(ATLAS_SURCHARGE_RATE, BUNDLER_SURCHARGE_RATE) <
 _solverBalance) {</pre>

Customer's Response:

Fixed in commit 5f53a6a



Low Severity Issues

L-01 Remove	Bundler	surcharge in	_updateAnal	vtics
_ 0111011040	Danaidi	oai oilai go ili	_ ap aato,a.	,

Severity: Low	Impact: High	Likelihood: Low
Files: GasAccounting.sol	Status: Fixed	Violated Property: None

Description:

When _updateAnalytics is called through _assign() we just need to remove the surcharge or change the name of that variable and add the Atlas surcharge.

Customer's Response:

Fixed in commit c302995



L-02 Governance signatory behaves unexpectedly

Severity: Low	Impact: High	Likelihood: Low
Files: DappIntegration.sol DappControl.sol	Status: Fixed	Violated Property: None

Description: The DappControl.sol governance can call DappIntegration.removeSignatory() and remove themselves as a signatory. This would make it impossible to transfer governance via the DappControl.acceptGovernance() function, because as the execution function reaches DappIntegration.changeDAppGovernance() it would revert, as the signatory would not longer exist. This means the governance would still be in control, even though it should not, as it has been removed as a signatory. If it is desired to completely remove governance, we suggest making it so DappControl.governance also gets changed if it removes itself as a signatory. This would prevent the case where governance falsely removes itself, in an attempt to high-jack the intended newGovernance from calling acceptGovernance(), which could happen if there is an untimely key compromise.

Customer's Response:

Fixed in commit Oe114a5



L-03 False Update to Analytics

Severity: Low	Impact: High	Likelihood: Low
Files: GasAccounting.sol	Status: Fixed	Violated Property: None

Description:

The gasRefundBeneficiary is marked as a winner in the auctionWins when all the solvers failed. If a solver is also a gasRefundBeneficiary and all the solvers fail then he gets +1 in the auctionWins and in the auctionFails. There is double counting of the totalGasValueUsed in that case because when the fail happens we increase the totalGasValueUsed for the solver run and we count that again in the _credit inside _settle .

Customer's Response:

Fixed in commit c302995



L-04 PerBlockLimit can lead to accounting exploit on Bundlers or Solvers

Severity: Low	Impact: High	Likelihood: Low
Files: Escrow.sol	Status: Acknowledged	Violated Property: None

Description:

Scenario 1: Evil Auctioneer could make Solvers lose money

An Auctioneer knowingly replays solverOps and Solver has to pay for it due to _validateSolverOpDeadline() and bundlersFault() logic.

Malicious Auctioneer could replay solverOp in the same bundle and the solver would have to pay for it according to _handleSolverAccounting() logic (plus other factors mentioned above). We believe the bundler should be the one to blame if lastAccessedBlock >= block.number(Escrow line 387). Because they should know better than to replay the same solverOps. While the Solver has no control over it. Consequently Bundlers can cause Solvers to lose money either willingly or by mistake.

Scenario 2: Evil Solver could make Bundlers lose money

Solver submits two operations, the first uses all his bonded/unbonding balance, leaving the second one with no balance to subtract from. Bundler will have to pay for this deficit, which could be big if depending on solverOps cost.

Recommendation:

Blame the bundler if there are two solverOps from the same msg.sender in the same bundle.

This aligns incentives with Bundlers and makes it so scenario 2 will not happen in practice, and if this happens it can purely be attributed to Bundler incompetence.



We can further enforce on-chain by looping over the bundle and checking each solverOp to see if there is a repeating msg.sender in it.

Customer's Response:

Acknowledged. No changes made as we still feel blaming solvers for the PerBlockLimit error is the best balance given the system's incentives and sophistication of the auctioneer, bundler, and solver parties, as well as the existing trust assumptions around the auctioneer.

To address scenario 1: We assume solvers to be sophisticated actors in this system. If an auctioneer acts dishonestly such as attempting to replay solverOps, we expect solvers to quickly notice and stop participating in those metacalls. The downside risk to solvers is also fairly low as their solverOp would not execute and the gas used would be minimal. We cannot fix this by shifting blame to the bundler as other bundlers may be coincidentally including the same solver in a different metacall in the same block, which is something even a sophisticated bundler wouldn't be able to detect, so cannot be attributed to bundler incompetence.

To address scenario 2: The auctioneer has the ability to exclude the 2nd solver from the metacall when their simulation of the solverOps shows that it would fail. The bundler has the ability to not bundle the metacall after simulating it shows that it would fail and lead to a loss. Finally, this type of attack would be costly for a solver to perform as they'd pay the gas cost + surcharges on any intentionally failed solverOps.

There is also a level of trust that solvers have in an auctioneer when submitting their solverOps, as an auctioneer could steal all MEV anyway by analysing the opportunities in the metacall bundle. With that trust assumption in place, we can assume the auctioneer will not abuse the PerBlockLimit mechanism at solvers' expense.



Informational Severity Issues

I-01. Unnecessary variable: adjustedDeposits

Description: adjustedDeposits is never used in _adjustAccountingForFees(). We recommend just using the original t_ values on GasAccounting.sol line 440 instead and save some gas by decreasing the inputs and outputs of _adjustAccountingForFees().

Customer's Response:

Fixed during the refactor of the Gas Accounting system, in commit fbalc9e

I-02. Unnecessarily setting t_withdrawal on _credit()

Description: As t_withdrawal is never used after it is increased by amount at the end of the _credit() function, we recommend removing the line in question.

Customer's Response:

Fixed in commit 57b229b

I-03. exPostBid terminology

Description: We recommend a more transparent terminology such as sortOnChain for example. As discussed with the team, this might not be possible at the moment due to backward compatibility issues, but it is something to keep in mind once backward compatibility can be overcome.

Customer's response:

Acknowledged, but exPostBids terminology left as is as some integrating partners are already familiar with the term.



I-04. postOpsCall hook is unnecessary

Description: After discussing with the team we arrived at the conclusion that postOpsCall could be combined with allocateValueCall, by making bool solverSuccessful and uint winningBid sent to the single hook and leaving it to the dapp to implement logic that is only run in the case of a winning solver in that hook.

Customer's response:

Fixed in commit 3269e9f

I-05. _contribute() in ExecutionBase_.sol could be named better

Description: Change this function name to contributeToAtlas() for better clarity and to avoid confusion with the other contribute() in GasAccounting.sol

Customer's response:

Fixed in commit ccOa4ff

I-06.Unreachable but concerning DOS of a bundle

Description: This scenario is not technically possible but we believe it still should be shared as it could be relevant for future development of the protocol.

An attacker can flashloan a lot of native tokens (ETH, POL...) and call depositAndBond then call borrow and loop this depositAndBond and borrow as much as possible. Once the solver runs successfully finishes he returns the flashloan _amountSolverPays is very high (_withdrawals is super inflated not bound by the original ETH balance of the atlas contract) due to the fact that the borrowed ETH will be repaid from the bond (at the _assign function inside _settle). In the _assign there is a cast to uint112 from the amount (uint256) to be charged from the solver. This cast may cause a revert to the metacall and fail everything. This is possible because every time the attacker deposited, they increased their bond so at the end the reconciliation will pass.



Customer's response:

Acknowledged. Agreed that this attack should not be possible due to the loops needed to overflow in a uint112 (changed to a uint128 during the refactor of the Gas Accounting system).

Mitigation Review

Project Scope

Project Name	Repository (link)	Latest Commit Hash	Platform
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Atlas	https://github.com/FastLane-L abs/atlas	288f2ca7dcb8d0636 b2e9f721e94cbfbfcef caf2	EVM/Solidity 0.8
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Project Overview

This document describes the Mitigation review of **Atlas** by manual code review findings. The work was undertaken from **March 20, 2025** to **April 10, 2025**.

The following contract list is included in our scope:

```
src/contracts/atlas/*
src/contracts/common/*
src/contracts/dapp/*
src/contracts/gasCalculator/*
src/contracts/helpers/Simulator.sol
src/contracts/helpers/Sorter.sol
src/contracts/libraries/*
src/contracts/types/*
```

The Certora Prover demonstrated that the implementation of the **Solidity** contracts above is correct with respect to the formal rules written by the Certora team. In addition, the team performed a manual audit of all the Solidity contracts. During the verification process and the manual audit, the Certora team discovered bugs in the Solidity contracts code, as listed on the following page.

Please note that a few more formal rules are not included in this report, as they were proven with an unreleased version of the Certora Prover. Once those rules are proven on a released version of the Certora Prover, we will add them to the next version of this document.

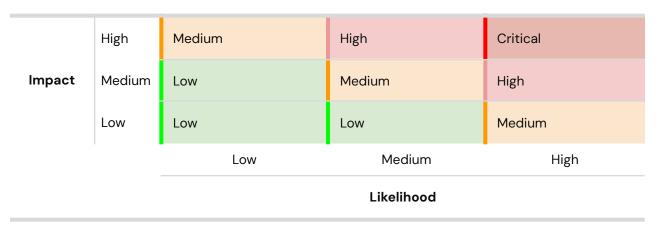


Findings Summary

The table below summarizes the findings of the review, including type and severity details.

Severity	Discovered	Confirmed	Fixed
Critical	0	0	0
High	6	0	6
Medium	2	0	2
Low	6	1	5
Informational	1	0	1
Total	15	1	14

Severity Matrix





Detailed Findings

ID	Title	Severity	Status
H-01	Bundler can make msg.data bigger than it is, inflating gas costs	High	Fixed
H-02	_netRepayments is unfairly limited to _maxRefund	High	Fixed
H-03	_gL.remainingMaxGas isn't decreased correctly	High	Fixed
H-04	unreachedSolverGas is only decreased when signatures are valid	High	Fixed
H-05	Winning solver is unfairly punished if unreached solvers can't pay their bond	High	Fixed
H-06	solverOps are not verified if they are signed when calculating unreached calldata	High	Fixed
M-01	Winning solver pays for gas costs of handling failed solvers accounting	Medium	Fixed
M-02	Winning solver/bundler will pay for running the loop for each failed solver	Medium	Fixed



L-01	Surcharge rates can be changed mid metacall	Low	Acknowledged
L-02	unreachedCalldataValuePaid is used even when there are no winning solvers	Low	Fixed
L-03	_gasWaterMark initialized in incorrect place in function	Low	Fixed
L-04	SolverTxResult is emitted with incorrect variable	Low	Fixed
L-05	_gasUsed isn't capturing the correct gas left prior to solver fail accounting	Low	Fixed
L-06	Bundler can add duplicate solver ops, forcing them to fail and charging the solver	Low	Fixed



High Severity Issues

H-01 Bundler can make msg.data bigger than it is, inflating gas costs

Severity: High	Impact: High	Likelihood: Medium
Files: Atlas.sol	Fixed	

Description:

The bundler can make the calldata bigger than it is, because of the way the EVM interprets the offsets in the calldata we can make msg.data.length to be bigger than the sum of all it's parts. The bundler can do this without changing the solverOps or anything else, by just adding zeros in the middle, because the EVM use offsets in the calldata parsing the bundler can inflate only msg.data.length this charge goes to the winning solver who will pay all of his bond to do it by doing this we can drain the winning solver.

Recommendation:

Calculate only the calldata costs for solverOps, userOp and dAppOp, instead plainly using msg.data

Customer Response:



H-02 _netRepayments is unfairly limited to _maxRefund

Severity: High	Impact: High	Likelihood: High
Files: GasAccounting.sol	Fixed	

Description:

_netRepayments represents the deposit made by the bundler and is currently capped by _maxRefund (80%) which is unfair to the bundler because this _netRepayments has the msg.value the bundler sent so capping this will cause the bundler to lose Eth.

Recommendation:

Add _netRepayments to claimsPaidToBundler at the end of the else, after the max refund logic as to not subject it to the limit.

Customer Response:

Fixed as per recommendation <u>here</u>.



H-03 _gL.remainingMaxGas isn't decreased correctly

Severity: High	Impact: High	Likelihood: High
Files: Escrow.sol	Fixed	

Description:

In the findBid iteration loop we don't decrease the _gL.remainingMaxGas like here because the solvers are not failing and we do decrease and update _gL.unreachedSolverGas here inside _validateSolverOpGasAndValue function. This will cause the later solvers to have a higher result in this function solverGasLiability here which depending on where the solvers are at line affects the result of the bundler (still should not happen even if it is the bundlers fault) and more important if we pass that check we are affected here inside reconcile function which will affect the solver's bid in the bidFinding phase and this may affect the result where the solver could have had a better bid but was forced not to by this behaviour.

Recommendation:

Adjust the _gL.unreachedSolverGas so this will simulate the proper behaviour.

Customer Response:



H-04 unreachedSolverGas is only decreased when signatures are valid

Severity: High	Impact: High	Likelihood: High
Files: Escrow.sol	Fixed	

Description:

In the knownBid path here we decrease the unreachedSolverGas but this happens only if the signature is valid. When the signature is not valid we just go to the handleFail and here we always decrease remainingMaxGas which means that is lets say the first couple of solvers just fail. In that case this will revert, the reason is that the remainingMaxGas will be smaller than unreachedSolverGas and the sub will underflow. This will cause the metacall to fail completely even if there are good solvers in the future. Also because this check for the solverGasLiability is wrong we can also (if we fail less solver so remainingMaxGas is still a bit bigger than unreachedSolverGas but less than it should) we don't really check that the solver has bond to pay we will execute and at the end there will be a deficit that the bundler will pay (no matter the result solver won or failed) this is bad for them.

Recommendation:

Decrease unreachedSolverGas even if the signature is invalid for both bid-finding execution and real execution.

Customer Response:



H-05 Winning solver is unfairly punished if unreached solvers can't pay their bond

Severity: High	Impact: High	Likelihood: High
Files: Escrow.sol	Fixed	

Description:

If exPostBids = false then any uncreached solvers must pay for the gas costs of their calldata.

If they cannot cover the cost, the deficit is taken out of unreachedCalldataValuePaid which will be paid by the winning solver, effectively punishing him for being the winner. This can be even more impactful for the winning solver if the bundler and auctioneer collude by ordering the solverOps in such a way that solvers with large calldata costs will be placed after the winning solver, punishing him even harder and draining his bond potentially.

Recommendation:

Do not punish the winning solver as he has no control over which and how many solvers are after him. Instead, add the _deficit to the bundler's writeoffs.

Customer Response:



H-06 solverOps are not verified if they are signed when calculating unreached calldata

Severity: High	Impact: High	Likelihood: Medium
Files: GasAccounting.sol	Fixed	

Description:

In the case of exPostBids = false, any unreached solvers are charged for their calldata costs, by calling _chargeUnreachedSolversForCalldata. The calldata costs are calculated and then each unreached solver is charged for his calldata. The issue is that the actual solverOp is not verified, if solverOp.from signed it, because of this the bundler + auctioneer can add fake solver ops after the winning one, making the calldata massive and drain the bond of each targeted solver.

Recommendation:

Verify each unreached solverOp to make sure it was signed by the corresponding solverOp.from

Customer Response:



Medium Severity Issues

M-01 Winning solver pay	s for gas costs of handling	failed solvers accounting
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Severity: Medium	Impact: Medium	Likelihood: Medium
Files: GasAccounting.sol	Fixed	

Description:

When a solver operation fails, it's handled through _handleSolverFailAccounting. We have two cases, if the solver op fails because it's the bundler's fault or because it's the solver's fault. The issue is that the gas costs of actually running the function are paid by the winning solver (if there is any), since none of these costs are charge neither to the bundler or the solver, unfairly punishing him for failed solver ops that he had nothing to do with. Considering the function can be run solverOps.length - 1 amount of times, the gas costs of the winning solver can become quite large..

Recommendation:

Add the gas costs of running the function to _gasUsed and then add them to the bundler's writeoffs or to the solver fault failure gas.

Customer Response:



M-O2 Winning solver/bundler will pay for running the loop for each failed solver

Severity: Medium	Impact: Medium	Likelihood: Medium
Files: Escrow.sol	Fixed	

Description:

When running _bidFindingIteration we loop through each valid solver op until we reach a winning solver. All the gas for running each _executeSolverOperation is either written off or added to the solver fault gas, but the actual gas costs of looping through each solver op isn't, which means at the end the winning solver (or bundler if all solver ops failed) will get charged, which is unfair to both parties.

Recommendation:

Remove the _gasWaterMark inside _executeSolverOperation and add a _gasWaterMark at the start of the loop which iterates over _bidsAndIndicesLastIndex.

Customer Response:



Low Severity Issues

L-01 Surcharge rates can be changed mid metacall		
Severity: Low	Impact: High	Likelihood: Low
Files: Storage.sol	Acknowledged	

Description:

Surcharge rates can be changed via setAtlasSurchargeRate during metacall if one of the solver ops is also S_surchargeRecipient which shouldn't be allowed. Consider adding _checkIfUnlocked inside setAtlasSurchargeRate.

Customer Response:

Acknowledged. As this is a permissioned operation, we believe there is very little risk of this happening as the Atlas surcharge recipient would never intentionally also be a solver. This part of the code is also refactored in the upcoming Atlas v1.6 upgrade, where this problem will be fully resolved.



L-02 unreachedCalldataValuePaid is used even when there are no winning solvers

Severity: Low	Impact: Low	Likelihood: High
Files: GasAccounting.sol	Fixed	

Description:

unreachedCalldataValuePaid is unnecessarily added to _bundlerCutBeforeLimit when we have no winning solvers. If we have no winning solvers then unreachedCalldataValuePaid is always O, so it's redundant to use it in this case.

Customer Response:



L-03 _gasWaterMark initialized in incorrect place in function

Severity: Low	Impact: Low	Likelihood: High
Files: Atlas.sol	Fixed	

Description:

Inside _bidFindingIteration, _gasWaterMark should be initialized at the start of the function to correctly capture the right amount of _gasLeft for the function.

Customer Response:



L-O4 SolverTxResult is emitted with incorrect variable Severity: Low Impact: Low Likelihood: High Files: Fixed Escrow.sol

Description:

SolverTxResult is emitted with bidAmount which is the previous bid. Use _solverTracker.bidAmount instead.

Customer Response:



Description:

Inside _handleSolverFailAccounting, _gL.remainingMaxGas should be calculated before _gasUsed in order to correctly capture the gas left prior to handling solver fail accounting.

Customer Response:



L-06 Bundler can add duplicate solver ops, forcing them to fail and charging the solver

Severity: Low	Impact: Medium	Likelihood: Low
Files: Escrow.sol	Fixed	

Description:

A malicious bundler can add duplicate solver ops to the solver0ps array, because of this, each duplicate will fail because his lastAccessedBlock == block.number which results in a fault for the solver, not the bundler. We recommend implementing a way to track solver0p.from and if a duplicate appears, charge the bundler.

Customer Response:

Acknowledged. We assume solvers are sophisticated and will quickly realise if a malicious auctioneer/bundler collusion is duplicating their solver operations maliciously. Adding a duplicate check mechanism would increase gas cost, and there are other trust assumptions around the auctioneer, so for now we rely on those trust assumptions and save gas on the honest metacalls.



Informational Severity Issues

I-01. SolverTxResult should be before _handleSolverFailAccounting

Description: Consider emitting the event prior to calling _handleSolverFailAccounting so the party at fault pays for the event emission.

Customer Response:



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