

# Is Syntactic Binding Rational?

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tag+ 2012

# Constraints in TAG and MGs

## Rational Constraints

A constraint is **rational**

- iff it can be defined by an MSO formula
  - iff it is computed by a bottom-up tree automaton
  - iff it defines a regular tree language.
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- **Most powerful class** of constraints that can be added to TAGs and MGs without increasing strong generative capacity
  - How much of syntax can be expressed in terms of rational constraints (and thus in vanilla TAGs and MGs)?  
Binding is known to be one of the hardest problems. . .

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# Necessary Restrictions

A rational theory of binding can at most

- verify the **existence of some grammatical reading**  
(but not specific readings; cf. Rogers 1998), and
- handle **syntactic binding**  
(but not discourse binding; cf. Ristad 1993).

Given these restrictions, reflexives (Principle A) are easy.  
Only bound pronouns that need to be disjoint in reference  
(Principle B) are problematic.

# Computing the Satisfiability of Principle B

## The Checkbook Version of Principle B (Simplified)

- There are **obviation domains** and **possible antecedents**.
- Obviation domains incur one point of **debt** for each syntactically bound pronoun satisfying certain conditions.
- The entire **debt** must be “paid off” by **antecedents**.

## Example

- (1) a. \* **Every patient** said that **he** should sedate **him**.  
b. **Every patient** told **some doctor** that **he** should sedate **him** in front of him.

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# Limited Obviation

The checkbook algorithm is rational iff **Limited Obviation** holds:

## Limited Obviation

There is some  $k$  such that no obviation domain's debt exceeds  $k$   
 $\Rightarrow$  only a bounded number of pronouns per obviation domain are mutually disjoint in reference.

No counterexamples to Limited Obviation in

- English,
- German,
- American Sign Language

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# References

- Ristad, Eric Sven. 1993. *The language complexity game*.  
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- Rogers, James. 1998. *A descriptive approach to language-theoretic complexity*. Stanford: CSLI.