

Balanced Generalized Weighing Matrices and Optimal Codes

Thomas Pender

University of Lethbridge
Department of Mathematics and Computer Science

Canadian Mathematical Society Winter Meeting
Decemeber 2021

*** Joint work with Hadi Kharaghani, Sho Suda, Vlad Zatiev

- Finite collection of “strings” of given length over a given finite alphabet.
- Finite alphabet \mathcal{A} with distinguished letter “0” called zero.
- Every $\mathcal{C} \subseteq \mathcal{A}^n$ is a code.
- Write \mathcal{C} is an $(n, M, d)_a$ -code, where
 - $M = |\mathcal{C}|$
 - $a = |\mathcal{A}|$
 - d is the Hamming distance (def. below).

Codes—Hamming Weight

- Weight of a codeword $c = c_0 \cdots c_{n-1}$.

$$\text{wt}(c) = \#\{i : c_i \neq 0\}$$

- Weight of \mathcal{C} .

$$\text{wt}(\mathcal{C}) = \min_{c \in \mathcal{C}} \text{wt}(c).$$

- \mathcal{C} is constant weight if $\#\{\text{wt}(c) : c \in \mathcal{C}\} = 1$.
- Assume \mathcal{C} is constant weight with $\text{wt}(\mathcal{C}) = w$.

Codes—Hamming Distance

- Distance between code words $c = c_0 \cdots c_{n-1}$ and $c' = c'_0 \cdots c'_{n-1}$.

$$d(c, c') = \#\{i : c_i \neq c'_i\}$$

- Distance of \mathcal{C} .

$$d(\mathcal{C}) = \min_{\substack{c, c' \in \mathcal{C}^2 \\ c \neq c'}} d(c, c')$$

- \mathcal{C} is equidistant if $\#\{d(c, c') : c, c' \in \mathcal{C}, c \neq c'\} = 1$.

Codes—Restricted Johnson Bound

- Fundamental question to maximize M given n , d , w and a .
- Denote this max by $A_a(n, d, w)$.

Restricted Johnson Bound

$$A_a(n, d, w) \leq \left\lfloor \frac{nd(a-1)}{aw^2 - 2(a-1)nw + nd(a-1)} \right\rfloor \quad (1)$$

- $\mathcal{A} = \{-, 0, 1\}$, where $-$ represents -1 .

DRAFT

- $\mathcal{A} = \{-, 0, 1\}$, where $-$ represents -1 .

1 1 1 - 0 - - 1

DRAFT

- $\mathcal{A} = \{-, 0, 1\}$, where $-$ represents -1 .

1 1 1 - 0 - - 1
- 1 1 1 - 0 - -

DRAFT

- $\mathcal{A} = \{-, 0, 1\}$, where $-$ represents -1 .

1 1 1 - 0 - - 1
- 1 1 1 - 0 - -
1 - 1 1 1 - 0 -

DRAFT

Codes—Example

- $\mathcal{A} = \{-, 0, 1\}$, where $-$ represents -1 .

1 1 1 - 0 - - 1
- 1 1 1 - 0 - -
1 - 1 1 1 - 0 -
1 1 - 1 1 1 - 0

DRAFT

Codes—Example

- $\mathcal{A} = \{-, 0, 1\}$, where $-$ represents -1 .

1 1 1 - 0 - - 1
- 1 1 1 - 0 - -
1 - 1 1 1 - 0 -
1 1 - 1 1 1 - 0
0 1 1 - 1 1 1 -

DRAFT

Codes—Example

- $\mathcal{A} = \{-, 0, 1\}$, where $-$ represents -1 .

1 1 1 - 0 - - 1
- 1 1 1 - 0 - -
1 - 1 1 1 - 0 -
1 1 - 1 1 1 - 0
0 1 1 - 1 1 1 -
1 0 1 1 - 1 1 1

Codes—Example

- $\mathcal{A} = \{-, 0, 1\}$, where $-$ represents -1 .

1 1 1 - 0 - - 1
- 1 1 1 - 0 - -
1 - 1 1 1 - 0 -
1 1 - 1 1 1 - 0
0 1 1 - 1 1 1 -
1 0 1 1 - 1 1 1
- 1 0 1 1 - 1 1

- $\mathcal{A} = \{-, 0, 1\}$, where $-$ represents -1 .

1 1 1 - 0 - - 1
- 1 1 1 - 0 - -
1 - 1 1 1 - 0 -
1 1 - 1 1 1 - 0
0 1 1 - 1 1 1 -
1 0 1 1 - 1 1 1
- 1 0 1 1 - 1 1
- - 1 0 1 1 - 1

Codes—Example

- $\mathcal{A} = \{-, 0, 1\}$, where $-$ represents -1 .

1 1 1 - 0 - - 1
- 1 1 1 - 0 - -
1 - 1 1 1 - 0 -
1 1 - 1 1 1 - 0
0 1 1 - 1 1 1 -
1 0 1 1 - 1 1 1
- 1 0 1 1 - 1 1
- - 1 0 1 1 - 1
- - - 1 0 1 1 -
1 - - - 1 0 1 1
- 1 - - - 1 0 1
- - 1 - - - 1 0
0 - - 1 - - - 1
- 0 - - 1 - - -
1 - 0 - - 1 - -
1 1 - 0 - - 1 -

- $n = 8$, $a = 3$, $d = 5$, $w = 7$.
- Extremal case of (1).
- Consequence of the following result due to Östergård, Svanström [2].

Theorem

If p is an odd prime,

$$A_3(p^m + 1, (p^m + 3)/2, p^m) = 2p^m + 2, \quad (2)$$

for $m \geq 1$.

(Balanced) Weighing Matrices—Definition

- $(-1, 0, 1)$ -matrix W of order v such that $WW^t = kI_v$, for some k .
- k is the weight of the matrix.
- Denoted $W(v, k)$.
- The $W(v, v)$ are the Hadamard matrices, and $W(v, v-1)$ are the conference matrices.
- If $W * W$ is an SBIBD, then W is balanced.
- Every conference matrix is balanced.

Weighing Matrices—Example

- The first seven codewords of the previous example form a $W(8, 7)$.

$$W = \begin{bmatrix} 1 & 1 & 1 & - & 0 & - & - & 1 \\ - & 1 & 1 & 1 & - & 0 & - & - \\ 1 & - & 1 & 1 & 1 & - & 0 & - \\ 1 & 1 & - & 1 & 1 & 1 & - & 0 \\ 0 & 1 & 1 & - & 1 & 1 & 1 & - \\ 1 & 0 & 1 & 1 & - & 1 & 1 & 1 \\ - & 1 & 0 & 1 & 1 & - & 1 & 1 \\ - & - & 1 & 0 & 1 & 1 & - & 1 \end{bmatrix}$$

- G a finite group.
- $W = [w_{ij}]$ a $(0, G)$ -matrix of order v with k nonzero entries in every row such that the multisets

$$\{w_{i\ell}w_{j\ell}^{-1} : w_{i\ell} \neq 0 \neq w_{j\ell}, 0 \leq \ell < v\}, \quad i \neq j,$$

contains each group element $\lambda/|G|$ times, for some λ .

- W is a balanced generalized weighing matrix.
- Denoted as $\text{BGW}(v, k, \lambda; G)$.

- q a prime power, $d > 1$.
- Take $K = \text{GF}(q)$, $F = \text{GF}(q^d)$.
- Relative trace $\text{Tr}_{F/K} : F \rightarrow K$ defined as

$$\text{Tr}_{F/K}(\alpha) = \alpha + \alpha^q + \cdots + \alpha^{q^{d-1}}, \quad \alpha \in F$$

BGWs—Classical Parameters cont.

- α a primitive element of F .
- $m = \frac{q^d - 1}{q - 1}$.
- Define $\omega = \alpha^{-m}$.
- Form the m dimensional vector

$$u = (\text{Tr}_{F/K}(\alpha^0), \dots, \text{Tr}_{F/K}(\alpha^{m-1})).$$

- Form the matrix W over K of order m by taking u together with its first $m - 1$ ω -shifts.
- Jungnickel, Tonchev [1] showed that W is BGW with parameters

$$\left(\frac{q^d - 1}{q - 1}, q^{d-1}, q^{d-1} - q^{d-2} \right)$$

with respect to K^* .

- Take $K = \text{GF}(7)$ and $d = 2$. Then $m = 8$.

DRAFT

BGWs—Example

- Take $K = \text{GF}(7)$ and $d = 2$. Then $m = 8$.

$$\omega^4 \quad 1 \quad \omega^4 \quad \omega^3 \quad 0 \quad \omega^5 \quad \omega^5 \quad 1$$

DRAFT

BGWs—Example

- Take $K = \text{GF}(7)$ and $d = 2$. Then $m = 8$.

$$\begin{array}{cccccccc} \omega^4 & 1 & \omega^4 & \omega^3 & 0 & \omega^5 & \omega^5 & 1 \\ \omega & \omega^4 & 1 & \omega^4 & \omega^3 & 0 & \omega^5 & \omega^5 \end{array}$$

DRAFT

BGWs—Example

- Take $K = \text{GF}(7)$ and $d = 2$. Then $m = 8$.

ω^4	1	ω^4	ω^3	0	ω^5	ω^5	1
ω	ω^4	1	ω^4	ω^3	0	ω^5	ω^5
1	ω	ω^4	1	ω^4	ω^3	0	ω^5

DRAFT

- Take $K = \text{GF}(7)$ and $d = 2$. Then $m = 8$.

ω^4	1	ω^4	ω^3	0	ω^5	ω^5	1
ω	ω^4	1	ω^4	ω^3	0	ω^5	ω^5
1	ω	ω^4	1	ω^4	ω^3	0	ω^5
1	1	ω	ω^4	1	ω^4	ω^3	0

DRAFT

BGWs—Example

- Take $K = \text{GF}(7)$ and $d = 2$. Then $m = 8$.

ω^4	1	ω^4	ω^3	0	ω^5	ω^5	1
ω	ω^4	1	ω^4	ω^3	0	ω^5	ω^5
1	ω	ω^4	1	ω^4	ω^3	0	ω^5
1	1	ω	ω^4	1	ω^4	ω^3	0
0	1	1	ω	ω^4	1	ω^4	ω^3

BGWs—Example

- Take $K = \text{GF}(7)$ and $d = 2$. Then $m = 8$.

ω^4	1	ω^4	ω^3	0	ω^5	ω^5	1
ω	ω^4	1	ω^4	ω^3	0	ω^5	ω^5
1	ω	ω^4	1	ω^4	ω^3	0	ω^5
1	1	ω	ω^4	1	ω^4	ω^3	0
0	1	1	ω	ω^4	1	ω^4	ω^3
ω^4	0	1	1	ω	ω^4	1	ω^4

BGWs—Example

- Take $K = \text{GF}(7)$ and $d = 2$. Then $m = 8$.

ω^4	1	ω^4	ω^3	0	ω^5	ω^5	1
ω	ω^4	1	ω^4	ω^3	0	ω^5	ω^5
1	ω	ω^4	1	ω^4	ω^3	0	ω^5
1	1	ω	ω^4	1	ω^4	ω^3	0
0	1	1	ω	ω^4	1	ω^4	ω^3
ω^4	0	1	1	ω	ω^4	1	ω^4
ω^5	ω^4	0	1	1	ω	ω^4	1

BGWs—Example

- Take $K = \text{GF}(7)$ and $d = 2$. Then $m = 8$.

ω^4	1	ω^4	ω^3	0	ω^5	ω^5	1
ω	ω^4	1	ω^4	ω^3	0	ω^5	ω^5
1	ω	ω^4	1	ω^4	ω^3	0	ω^5
1	1	ω	ω^4	1	ω^4	ω^3	0
0	1	1	ω	ω^4	1	ω^4	ω^3
ω^4	0	1	1	ω	ω^4	1	ω^4
ω^5	ω^4	0	1	1	ω	ω^4	1
ω	ω^5	ω^4	0	1	1	ω	ω^4

- A $\text{BGW}(8, 7, 6; \text{GF}(7)^*)$.

Codes from BGWs—Example

$\omega^4 1$	$\omega^4 \omega^3 0$	$\omega^5 \omega^5 1$	$\omega^5 \omega$	$\omega^5 \omega^4 0$	1	1	ω	1	$\omega^2 1$	$\omega^5 0$	ω	$\omega \omega^2$
$\omega \omega^4 1$	$\omega^4 \omega^3 0$	$\omega^5 \omega^5$	$\omega^2 \omega^5 \omega$	$\omega^5 \omega^4 0$	1	1		$\omega^3 1$	$\omega^2 1$	$\omega^5 0$	ω	ω
$1 \omega \omega^4 1$	$\omega^4 \omega^3 0$	ω^5	$\omega \omega^2 \omega^5 \omega$	$\omega^5 \omega^4 0$	1			$\omega^2 \omega^3 1$	$\omega^2 1$	$\omega^5 0$	ω	
$1 1 \omega \omega^4 1$	$\omega^4 \omega^3 0$		$\omega \omega \omega^2 \omega^5 \omega$	$\omega^5 \omega^4 0$				$\omega^2 \omega^2 \omega^3 1$	$\omega^2 1$	$\omega^5 0$		
$0 1 1 \omega \omega^4 1$	$\omega^4 \omega^3$		$0 \omega \omega \omega^2 \omega^5 \omega$	$\omega^5 \omega^4$				$0 \omega^2 \omega^2 \omega^3 1$	$\omega^2 1$	ω^5		
$\omega^4 0 1 1 \omega \omega^4 1$	ω^4		$\omega^5 0 \omega \omega \omega^2 \omega^5 \omega$	ω^5				$1 0 \omega^2 \omega^2 \omega^3 1$	$\omega^2 1$			
$\omega^5 \omega^4 0 1 1 \omega \omega^4 1$			$1 \omega^5 0 \omega \omega \omega^2 \omega^5 \omega$					$\omega 1 0 \omega^2 \omega^2 \omega^3 1$	ω^2			
$\omega \omega^5 \omega^4 0 1 1 \omega \omega^4$			$\omega^2 1 \omega^5 0 \omega \omega \omega^2 \omega^5$					$\omega^3 \omega 1 0 \omega^2 \omega^2 \omega^3 1$				
$\omega \omega^3 \omega 1 0 \omega^2 \omega^2 \omega^3$			$\omega^2 \omega^4 \omega^2 \omega 0 \omega^3 \omega^3 \omega^4$					$\omega^3 \omega^5 \omega^3 \omega^2 0 \omega^4 \omega^4 \omega^5$				
$\omega^4 \omega \omega^3 \omega 1 0 \omega^2 \omega^2$			$\omega^5 \omega^2 \omega^4 \omega^2 \omega 0 \omega^3 \omega^3$					$1 \omega^3 \omega^5 \omega^3 \omega^2 0 \omega^4 \omega^4$				
$\omega^3 \omega^4 \omega \omega^3 \omega 1 0 \omega^2$			$\omega^4 \omega^5 \omega^2 \omega^4 \omega^2 \omega 0 \omega^3$					$\omega^5 1 \omega^3 \omega^5 \omega^3 \omega^2 0 \omega^4$				
$\omega^3 \omega^3 \omega^4 \omega \omega^3 \omega 1 0$			$\omega^4 \omega^4 \omega^5 \omega^2 \omega^4 \omega^2 \omega 0$					$\omega^5 \omega^5 1 \omega^3 \omega^5 \omega^3 \omega^2 0$				
$0 \omega^3 \omega^3 \omega^4 \omega \omega^3 \omega 1$			$0 \omega^4 \omega^4 \omega^5 \omega^2 \omega^4 \omega^2 \omega$					$0 \omega^5 \omega^5 1 \omega^3 \omega^5 \omega^3 \omega^2$				
$\omega 0 \omega^3 \omega^3 \omega^4 \omega \omega^3 \omega$			$\omega^2 0 \omega^4 \omega^4 \omega^5 \omega^2 \omega^4 \omega^2$					$\omega^3 0 \omega^5 \omega^5 1 \omega^3 \omega^5 \omega^3$				
$\omega^2 \omega 0 \omega^3 \omega^3 \omega^4 \omega \omega^3$			$\omega^3 \omega^2 0 \omega^4 \omega^4 \omega^5 \omega^2 \omega^4$					$\omega^4 \omega^3 0 \omega^5 \omega^5 1 \omega^3 \omega^5$				
$\omega^4 \omega^2 \omega 0 \omega^3 \omega^3 \omega^4 \omega$			$\omega^5 \omega^3 \omega^2 0 \omega^4 \omega^4 \omega^5 \omega^2$					$1 \omega^4 \omega^3 0 \omega^5 \omega^5 1 \omega^3$				

Codes from BGWs—Example cont.

- $n = 8, w = 7, d = 7$, and $a = 7$.
- Equality in (1).
- $A_7(8, 7, 7) = 48$.

DRAFT

Theorem

If q is a prime power, and if $d > 1$, then



$$A_q \left(\frac{q^d - 1}{q - 1}, q^{d-1}, q^{d-1} \right) = q^d - 1.$$

Moreover, the code can be assumed to be completely generated by a single codeword and equidistant.

- If q odd, taking $d = 2$ implies (2) as a corollary.

Thank You!!!

DRAFT

-  Dieter Jungnickel and Vladimir D. Tonchev.
Perfect codes and balanced generalized weighing matrices. II.
Finite Fields Appl., 8(2):155–165, 2002.
-  Patric R. J. Östergård and Mattias Svanström.
Ternary constant weight codes.
Electron. J. Combin., 9(1):Research Paper 41, 23, 2002.