### Lattice Reduction for Modular Knapsack

# Thomas PLANTARD Willy SUSILO Zhenfei ZHANG

Centre for Computer and Information Security Research University of Wollongong

http://www.uow.edu.au/~thomaspl thomaspl@uow.edu.au

#### Outline

- Introduction
- 2 Lattice Theory
  - Lattice Basics
  - Lattice Reduction
  - LLL
- 3 LLL for Modular Knapsack Lattice
  - Modular Knapsack lattice
  - LLL for modular knapsack lattice
- A Recursive LLL
- Conclusion



#### Introduction

- Introduction
- 2 Lattice Theory
  - Lattice Basics
  - Lattice Reduction
  - LLL
- 3 LLL for Modular Knapsack Lattice
  - Modular Knapsack lattice
  - LLL for modular knapsack lattice
- 4 A Recursive LLL
- Conclusion



# Cryptography concerned by Lattice Reduction

#### **Problem**

- Shortest Vector Problem (SVP): Ajtai-Dwork, Regev, ...
- Closet Vector Problem (CVP): GGH, NTRU, ...
- Knapsack Problem
- Coding based cryptosystem
- RSA, Factorization.
- Short Integer Solution (SIS): SWIFFT, SWIFFTX, ...
- Learning With Error (LWE).
- Approximate-GCD Problem.

#### Lattice Reduction: Heuristic BUT successful

- Weeks, Month of Computation: Good Estimation.
- 2<sup>80</sup>, 2<sup>100</sup>: Unknown.

# The 2010 FHE Gentry-Halevi implementation

#### Challenge

- Find a short vector in a modular knapsack type lattice.
- Dimension, d = 2048
- Length of digits,  $\beta = 720,000$ .

#### Security based on impossibility to run LLL

- Perform a LLL reduction is enough to break challenge.
- However,  $d^3\beta^2 = (2^{11})^3(2^{19.5})^2 = 2^{72}$ .

### Lattice Theory

- Introduction
- 2 Lattice Theory
  - Lattice Basics
  - Lattice Reduction
  - LLL
- 3 LLL for Modular Knapsack Lattice
  - Modular Knapsack lattice
  - LLL for modular knapsack lattice
- A Recursive LLL
- Conclusion



#### Lattice

#### Definition of a Lattice

• All the integral combinations of  $d \le n$  linearly independent vectors over  $\mathbb R$ 

$$\mathcal{L} = \mathbb{Z} \, \mathbf{b}_1 + \dots + \mathbb{Z} \, \mathbf{b}_d = \{ \lambda_1 \mathbf{b}_1 + \dots + \lambda_d \mathbf{b}_d : \lambda_i \in \mathbb{Z} \}$$

- d dimension.
- $\mathbf{B} = (\mathbf{b}_1, \dots, \mathbf{b}_d)$  is a basis.

### An Example

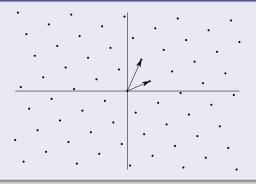
$$\mathbf{B} = \begin{pmatrix} 5 & \frac{1}{2} & \sqrt{3} \\ \frac{3}{5} & \sqrt{2} & 1 \end{pmatrix}$$

 $d = 2 \le n = 3$ 

In this work, integer Basis:  $B \in \mathbb{Z}^{d,n}$ .

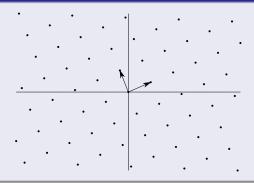
#### A lattice $\mathcal{L}$

$$\mathbf{B} = \begin{pmatrix} 8 & 5 \\ 5 & 16 \end{pmatrix}$$



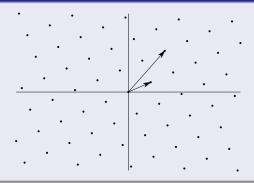
#### A lattice $\mathcal{L}$

$$\mathbf{UB} = \begin{pmatrix} 1 & 0 \\ -1 & 1 \end{pmatrix} \begin{pmatrix} 8 & 5 \\ 5 & 16 \end{pmatrix} = \begin{pmatrix} 8 & 5 \\ -3 & 11 \end{pmatrix}$$



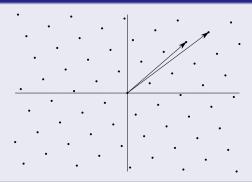
#### A lattice $\mathcal{L}$

$$\mathbf{UB} = \begin{pmatrix} 1 & 0 \\ 1 & 1 \end{pmatrix} \begin{pmatrix} 8 & 5 \\ 5 & 16 \end{pmatrix} = \begin{pmatrix} 8 & 5 \\ 13 & 21 \end{pmatrix}$$



#### A lattice $\mathcal{L}$

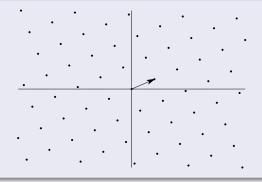
$$\mathbf{UB} = \begin{pmatrix} 3 & 1 \\ 2 & 1 \end{pmatrix} \begin{pmatrix} 8 & 5 \\ 5 & 16 \end{pmatrix} = \begin{pmatrix} 29 & 31 \\ 21 & 26 \end{pmatrix}$$



#### The Shortest Vector and The First Minima

$$\mathbf{v} = \begin{pmatrix} 8 & 5 \end{pmatrix}$$
, with  $\lambda_1 = \sqrt{8^2 + 5^2} = 9.434$ 

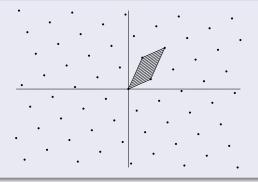
#### The Shortest Vector



#### The Determinant

$$\det \mathcal{L} = \sqrt{\det \left( \mathbf{B} \mathbf{B}^T \right)} = 103$$

### The Determinant



### Lattice Reduction Algorithm

#### Find $v \in \mathcal{L}$ smallest

- SVP is NP-Hard under randomized reduction.
- Deterministic  $O(d^{\frac{d}{2e}})$ : Kannan 1986, Hanrot and Sthele 2007.
- Probabilistic  $O(2^d)$ : AKS 2001.

#### Find $v \in \mathcal{L}$ small

- LLL: Lenstra, Lenstra and Lovasz (Poly in d).
- $DEEP_k$ : LLL with Deep Insertion (Exponential in k, Poly in d).
- $BKZ_k$ : Block Korkine Zolotaref (Exponential in k, Poly in d).
- ...



### LLL

#### LLL

- Input: a matrix  $A \in [-2^{\beta}, 2^{\beta}]^{d,n}$ .
- Output: a matrix  $B \in \mathbb{Z}^{d,n}$  with  $\|b_i\| \sim 2^{\frac{d}{2}} det^{\frac{1}{d}}$
- Shortest Basis:  $\|b_i\| \sim \sqrt{rac{d}{2\pi e}} det^{rac{1}{d}}$

#### Comparison of time complexity

| Algorithms | Time Complexity  |
|------------|--|
| LLL        | $O(d^{5+arepsilon}eta^{2+arepsilon})$                              |
| $L^2$      | $O(d^{4+\varepsilon}\beta^2+d^{5+\varepsilon}\beta)$               |
| $L^1$      | $O(d^{4+\varepsilon}\beta^{1+\varepsilon}+d^{5+\varepsilon}\beta)$ |

# LLL for Modular Knapsack Lattice

- Introduction
- 2 Lattice Theory
  - Lattice Basics
  - Lattice Reduction
  - LLL
- 3 LLL for Modular Knapsack Lattice
  - Modular Knapsack lattice
  - LLL for modular knapsack lattice
- A Recursive LLL
- Conclusion



## Modular Knapsack Basis

#### A modular knapsack basis

$$\mathbf{A} = egin{pmatrix} A_0 & 0 & 0 & 0 & 0 \ A_1 & 1 & 0 & 0 & 0 \ A_2 & 0 & 1 & 0 & 0 \ dots & 0 & 0 & \ddots & 0 \ A_{d-1} & 0 & 0 & 0 & 1 \ \end{pmatrix} ext{ with } |A_i| < 2^{eta}.$$

#### A classic format

- Natural format of lattice attack on knapsack problem.
- Use as public key as most of lattice based cryptosystem.
- Easy to compute from a random basis, using Hermite Normal Form.

### LLL for modular knapsack lattice

### Comparison of time complexity

| Algorithms                  | Time Complexity  |
|-----------------------------|--|
| LLL for knapsack            | $O(d^{4+arepsilon}eta^{2+arepsilon})$                              |
| L <sup>2</sup> for knapsack | $O(d^{3+\varepsilon}\beta^2+d^{4+\varepsilon}\beta)$               |
| $L^1$                       | $O(d^{4+\varepsilon}\beta^{1+\varepsilon}+d^{5+\varepsilon}\beta)$ |

#### Why faster than random basis: an intuition

- To reduce i + 1 vectors, LLL requires the first i vectors to be reduced.
- For modular knapsack basis, each i first vectors are a triangular matrix.

```
      (86670401
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
```

```
      86670401
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0
      0<
```

| / -3227  | -3165 | 0 | 0 | 0 | 0 | 0 | 0\ |
|----------|-------|---|---|---|---|---|----|
| -14111   | 13018 | 0 | 0 | 0 | 0 | 0 | 0  |
| 10117311 | 0     | 1 | 0 | 0 | 0 | 0 | 0  |
| 38269415 | 0     | 0 | 1 | 0 | 0 | 0 | 0  |
| 45874978 | 0     | 0 | 0 | 1 | 0 | 0 | 0  |
| 33538152 | 0     | 0 | 0 | 0 | 1 | 0 | 0  |
| 61611560 | 0     | 0 | 0 | 0 | 0 | 1 | 0  |
| 66174289 | 0     | 0 | 0 | 0 | 0 | 0 | 1/ |

| / -3227  | -3165 | 0 | 0 | 0 | 0 | 0 | 0\ |
|----------|-------|---|---|---|---|---|----|
| -14111   | 13018 | 0 | 0 | 0 | 0 | 0 | 0  |
| 10117311 | 0     | 1 | 0 | 0 | 0 | 0 | 0  |
| 38269415 | 0     | 0 | 1 | 0 | 0 | 0 | 0  |
| 45874978 | 0     | 0 | 0 | 1 | 0 | 0 | 0  |
| 33538152 | 0     | 0 | 0 | 0 | 1 | 0 | 0  |
| 61611560 | 0     | 0 | 0 | 0 | 0 | 1 | 0  |
| 66174289 | 0     | 0 | 0 | 0 | 0 | 0 | 1/ |

| / $-24$   | 153 | -215 | 0 | 0 | 0 | 0 | 0\ |
|-----------|-----|------|---|---|---|---|----|
| 183       | 242 | 76   | 0 | 0 | 0 | 0 | 0  |
| -920      | 440 | 343  | 0 | 0 | 0 | 0 | 0  |
| 38269415  | 0   | 0    | 1 | 0 | 0 | 0 | 0  |
| 45874978  | 0   | 0    | 0 | 1 | 0 | 0 | 0  |
| 33538152  | 0   | 0    | 0 | 0 | 1 | 0 | 0  |
| 61611560  | 0   | 0    | 0 | 0 | 0 | 1 | 0  |
| \66174289 | 0   | 0    | 0 | 0 | 0 | 0 | 1/ |

| / -24     | 153 | -215 | 0 | 0 | 0 | 0 | 0\ |
|-----------|-----|------|---|---|---|---|----|
| 183       | 242 | 76   | 0 | 0 | 0 | 0 | 0  |
| -920      | 440 | 343  | 0 | 0 | 0 | 0 | 0  |
| 38269415  | 0   | 0    | 1 | 0 | 0 | 0 | 0  |
| 45874978  | 0   | 0    | 0 | 1 | 0 | 0 | 0  |
| 33538152  | 0   | 0    | 0 | 0 | 1 | 0 | 0  |
| 61611560  | 0   | 0    | 0 | 0 | 0 | 1 | 0  |
| \66174289 | 0   | 0    | 0 | 0 | 0 | 0 | 1/ |

| / 8       | -27 | -66 | 42  | 0 | 0 | 0 | 0/ |
|-----------|-----|-----|-----|---|---|---|----|
| -40       | -45 | -47 | -38 | 0 | 0 | 0 | 0  |
| 0         | 126 | -18 | -23 | 0 | 0 | 0 | 0  |
| 103       | 26  | 0   | -53 | 0 | 0 | 0 | 0  |
| 45874978  | 0   | 0   | 0   | 1 | 0 | 0 | 0  |
| 33538152  | 0   | 0   | 0   | 0 | 1 | 0 | 0  |
| 61611560  | 0   | 0   | 0   | 0 | 0 | 1 | 0  |
| \66174289 | 0   | 0   | 0   | 0 | 0 | 0 | 1/ |

| / 8       | -27 | -66 | 42  | 0 | 0 | 0 | 0\ |
|-----------|-----|-----|-----|---|---|---|----|
| -40       | -45 | -47 | -38 | 0 | 0 | 0 | 0  |
| 0         | 126 | -18 | -23 | 0 | 0 | 0 | 0  |
| 103       | 26  | 0   | -53 | 0 | 0 | 0 | 0  |
| 45874978  | 0   | 0   | 0   | 1 | 0 | 0 | 0  |
| 33538152  | 0   | 0   | 0   | 0 | 1 | 0 | 0  |
| 61611560  | 0   | 0   | 0   | 0 | 0 | 1 | 0  |
| \66174289 | 0   | 0   | 0   | 0 | 0 | 0 | 1/ |

$$\begin{pmatrix} -17 & -31 & -5 & 6 & 1 & 0 & 0 & 0 \\ 24 & -20 & -6 & 4 & 7 & 0 & 0 & 0 \\ -3 & -7 & -45 & 3 & 17 & 0 & 0 & 0 \\ 8 & 4 & -13 & -14 & -36 & 0 & 0 & 0 \\ 13 & 0 & 15 & -35 & 24 & 0 & 0 & 0 \\ \hline 33538152 & 0 & 0 & 0 & 0 & 1 & 0 & 0 \\ 61611560 & 0 & 0 & 0 & 0 & 0 & 1 & 0 \\ 66174289 & 0 & 0 & 0 & 0 & 0 & 0 & 1 \end{pmatrix}$$

$$\begin{pmatrix} -17 & -31 & -5 & 6 & 1 & 0 & 0 & 0 \\ 24 & -20 & -6 & 4 & 7 & 0 & 0 & 0 \\ -3 & -7 & -45 & 3 & 17 & 0 & 0 & 0 \\ 8 & 4 & -13 & -14 & -36 & 0 & 0 & 0 \\ 13 & 0 & 15 & -35 & 24 & 0 & 0 & 0 \\ \hline 33538152 & 0 & 0 & 0 & 0 & 1 & 0 & 0 \\ \hline 61611560 & 0 & 0 & 0 & 0 & 0 & 1 & 0 \\ 66174289 & 0 & 0 & 0 & 0 & 0 & 0 & 1 \end{pmatrix}$$

$$\begin{pmatrix} -7 & -6 & -14 & 11 & 4 & -7 & 0 & 0 \\ -14 & -9 & -3 & -1 & -15 & -6 & 0 & 0 \\ -2 & 15 & 14 & 10 & -6 & 4 & 0 & 0 \\ 8 & -14 & 5 & 13 & -14 & -2 & 0 & 0 \\ -5 & 11 & -6 & -10 & -12 & 12 & 0 & 0 \\ 4 & -16 & 12 & -4 & 12 & 13 & 0 & 0 \\ \hline 61611560 & 0 & 0 & 0 & 0 & 0 & 1 & 0 \\ 66174289 & 0 & 0 & 0 & 0 & 0 & 0 & 1 \end{pmatrix}$$

$$\begin{pmatrix} -7 & -6 & -14 & 11 & 4 & -7 & 0 & 0 \\ -14 & -9 & -3 & -1 & -15 & -6 & 0 & 0 \\ -2 & 15 & 14 & 10 & -6 & 4 & 0 & 0 \\ 8 & -14 & 5 & 13 & -14 & -2 & 0 & 0 \\ -5 & 11 & -6 & -10 & -12 & 12 & 0 & 0 \\ 4 & -16 & 12 & -4 & 12 & 13 & 0 & 0 \\ \hline 66174289 & 0 & 0 & 0 & 0 & 0 & 0 & 1 \end{pmatrix}$$

$$\begin{pmatrix} 5 & -4 & -8 & 0 & 1 & 10 & -1 & 0 \\ 9 & -1 & 4 & -6 & -7 & -8 & 2 & 0 \\ 1 & -4 & 0 & 6 & -12 & -1 & 4 & 0 \\ 4 & 8 & -4 & 9 & 4 & 0 & -3 & 0 \\ 3 & -2 & -11 & -4 & -5 & -3 & -6 & 0 \\ 4 & -9 & -7 & 9 & 3 & -2 & 7 & 0 \\ 7 & -10 & 5 & 7 & -2 & -1 & -4 & 0 \\ \hline 66174289 & 0 & 0 & 0 & 0 & 0 & 1 \end{pmatrix}$$

$$\begin{pmatrix} 5 & -4 & -8 & 0 & 1 & 10 & -1 & 0 \\ 9 & -1 & 4 & -6 & -7 & -8 & 2 & 0 \\ 1 & -4 & 0 & 6 & -12 & -1 & 4 & 0 \\ 4 & 8 & -4 & 9 & 4 & 0 & -3 & 0 \\ 3 & -2 & -11 & -4 & -5 & -3 & -6 & 0 \\ 4 & -9 & -7 & 9 & 3 & -2 & 7 & 0 \\ 7 & -10 & 5 & 7 & -2 & -1 & -4 & 0 \\ 66174289 & 0 & 0 & 0 & 0 & 0 & 1 \end{pmatrix}$$

$$\begin{pmatrix} -2 & 6 & -1 & -8 & 0 & 1 & -3 & 1 \\ 2 & -3 & -8 & 1 & 3 & -1 & 4 & 1 \\ 3 & -3 & 6 & 2 & -6 & 2 & 4 & 3 \\ -2 & -1 & -6 & 4 & -6 & -3 & 0 & -3 \\ 5 & -4 & 4 & -1 & -2 & 0 & -7 & 1 \\ 4 & 2 & 3 & -1 & -4 & 1 & -1 & -7 \\ -2 & 3 & -4 & 0 & 0 & 11 & 2 & -2 \\ 5 & 11 & 1 & 3 & -2 & 3 & -2 & 4 \end{pmatrix}$$

#### A Recursive LLL

- Introduction
- 2 Lattice Theory
  - Lattice Basics
  - Lattice Reduction
  - LLL
- 3 LLL for Modular Knapsack Lattice
  - Modular Knapsack lattice
  - LLL for modular knapsack lattice
- 4 A Recursive LLL
- Conclusion



# Example Recursive LLL

```
\begin{pmatrix} 86670401 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 38009011 & 1 & 0 & 0 & 0 & 0 & 0 & 0 \\ 10117311 & 0 & 1 & 0 & 0 & 0 & 0 & 0 \\ 38269415 & 0 & 0 & 1 & 0 & 0 & 0 & 0 \\ 45874978 & 0 & 0 & 0 & 1 & 0 & 0 & 0 \\ 33538152 & 0 & 0 & 0 & 0 & 1 & 0 & 0 \\ 61611560 & 0 & 0 & 0 & 0 & 0 & 1 & 0 \\ 66174289 & 0 & 0 & 0 & 0 & 0 & 0 & 1 \end{pmatrix}
```

# Example Recursive LLL

| / -3227  | -3165 | 0 | 0 | 0 | 0 | 0 | 0\ |
|----------|-------|---|---|---|---|---|----|
| -14111   | 13018 | 0 | 0 | 0 | 0 | 0 | 0  |
| 10117311 | 0     | 1 | 0 | 0 | 0 | 0 | 0  |
| 38269415 | 0     | 0 | 1 | 0 | 0 | 0 | 0  |
| 45874978 | 0     | 0 | 0 | 1 | 0 | 0 | 0  |
| 33538152 | 0     | 0 | 0 | 0 | 1 | 0 | 0  |
| 61611560 | 0     | 0 | 0 | 0 | 0 | 1 | 0  |
| 66174289 | 0     | 0 | 0 | 0 | 0 | 0 | 1/ |

| / -3227  | -3165 | 0 | 0 | 0 | 0 | 0 | 0\ |
|----------|-------|---|---|---|---|---|----|
| -14111   | 13018 | 0 | 0 | 0 | 0 | 0 | 0  |
| 10117311 | 0     | 1 | 0 | 0 | 0 | 0 | 0  |
| 38269415 | 0     | 0 | 1 | 0 | 0 | 0 | 0  |
| 45874978 | 0     | 0 | 0 | 1 | 0 | 0 | 0  |
| 33538152 | 0     | 0 | 0 | 0 | 1 | 0 | 0  |
| 61611560 | 0     | 0 | 0 | 0 | 0 | 1 | 0  |
| 66174289 | 0     | 0 | 0 | 0 | 0 | 0 | 1/ |

| / -3227   | -3165 | 0    | 0     | 0 | 0 | 0 | 0\ |
|-----------|-------|------|-------|---|---|---|----|
| -14111    | 13018 | 0    | 0     | 0 | 0 | 0 | 0  |
| 1391      | 0     | 4036 | -1067 | 0 | 0 | 0 | 0  |
| -8121     | 0     | 3949 | -1044 | 0 | 0 | 0 | 0  |
| 45874978  | 0     | 0    | 0     | 1 | 0 | 0 | 0  |
| 33538152  | 0     | 0    | 0     | 0 | 1 | 0 | 0  |
| 61611560  | 0     | 0    | 0     | 0 | 0 | 1 | 0  |
| \66174289 | 0     | 0    | 0     | 0 | 0 | 0 | 1/ |

| / | -3227    | -3165 | 0    | 0     | 0 | 0 | 0 | 0\ |
|---|----------|-------|------|-------|---|---|---|----|
| 1 | -14111   | 13018 | 0    | 0     | 0 | 0 | 0 | 0  |
| ١ | 1391     | 0     | 4036 | -1067 | 0 | 0 | 0 | 0  |
| l | -8121    | 0     | 3949 | -1044 | 0 | 0 | 0 | 0  |
| ľ | 45874978 | 0     | 0    | 0     | 1 | 0 | 0 | 0  |
| I | 33538152 | 0     | 0    | 0     | 0 | 1 | 0 | 0  |
|   | 61611560 | 0     | 0    | 0     | 0 | 0 | 1 | 0  |
| 1 | 66174289 | 0     | 0    | 0     | 0 | 0 | 0 | 1/ |

| 1 | -3227    | -3165 | 0    | 0     | 0     | 0     | 0 | 0\ |
|---|----------|-------|------|-------|-------|-------|---|----|
|   | -14111   | 13018 | 0    | 0     | 0     | 0     | 0 | 0  |
| ı | 1391     | 0     | 4036 | -1067 | 0     | 0     | 0 | 0  |
| ı | -8121    | 0     | 3949 | -1044 | 0     | 0     | 0 | 0  |
| Ī | 1348     | 0     | 0    | 0     | 2830  | -3871 | 0 | 0  |
|   | 10894    | 0     | 0    | 0     | -2009 | 2748  | 0 | 0  |
|   | 61611560 | 0     | 0    | 0     | 0     | 0     | 1 | 0  |
|   | 66174289 | 0     | 0    | 0     | 0     | 0     | 0 | 1/ |

| / | -3227    | -3165 | 0    | 0     | 0     | 0     | 0 | 0\ |
|---|----------|-------|------|-------|-------|-------|---|----|
| 1 | -14111   | 13018 | 0    | 0     | 0     | 0     | 0 | 0  |
| İ | 1391     | 0     | 4036 | -1067 | 0     | 0     | 0 | 0  |
| ١ | -8121    | 0     | 3949 | -1044 | 0     | 0     | 0 | 0  |
| ١ | 1348     | 0     | 0    | 0     | 2830  | -3871 | 0 | 0  |
| ١ | 10894    | 0     | 0    | 0     | -2009 | 2748  | 0 | 0  |
| ľ | 61611560 | 0     | 0    | 0     | 0     | 0     | 1 | 0  |
| 1 | 66174289 | 0     | 0    | 0     | 0     | 0     | 0 | 1/ |

|   | / -3227 | -3165 | 0    | 0     | 0     | 0     | 0    | 0 \          |
|---|---------|-------|------|-------|-------|-------|------|--------------|
|   | -14111  | 13018 | 0    | 0     | 0     | 0     | 0    | 0            |
| İ | 1391    | 0     | 4036 | -1067 | 0     | 0     | 0    | 0            |
| ١ | -8121   | 0     | 3949 | -1044 | 0     | 0     | 0    | 0            |
| ١ | 1348    | 0     | 0    | 0     | 2830  | -3871 | 0    | 0            |
| ı | 10894   | 0     | 0    | 0     | -2009 | 2748  | 0    | 0            |
|   | 3       | 0     | 0    | 0     | 0     | 0     | 2248 | -2093        |
| 1 | 29437   | 0     | 0    | 0     | 0     | 0     | 29   | −27 <i>J</i> |

| /-3227 | -3165 | 0    | 0     | 0     | 0     | 0    | 0 \          |
|--------|-------|------|-------|-------|-------|------|--------------|
| -14111 | 13018 | 0    | 0     | 0     | 0     | 0    | 0            |
| 1391   | 0     | 4036 | -1067 | 0     | 0     | 0    | 0            |
| -8121  | 0     | 3949 | -1044 | 0     | 0     | 0    | 0            |
| 1348   | 0     | 0    | 0     | 2830  | -3871 | 0    | 0            |
| 10894  | 0     | 0    | 0     | -2009 | 2748  | 0    | 0            |
| 3      | 0     | 0    | 0     | 0     | 0     | 2248 | -2093        |
| 29437  | 0     | 0    | 0     | 0     | 0     | 29   | $-27$ $\int$ |

| / 8   | -27 | -66 | 42  | 0     | 0     | 0    | 0 \   |
|-------|-----|-----|-----|-------|-------|------|-------|
| -40   | -45 | -47 | -38 | 0     | 0     | 0    | 0     |
| 0     | 126 | -18 | -23 | 0     | 0     | 0    | 0     |
| 103   | 26  | 0   | -53 | 0     | 0     | 0    | 0     |
| 1348  | 0   | 0   | 0   | 2830  | -3871 | 0    | 0     |
| 10894 | 0   | 0   | 0   | -2009 | 2748  | 0    | 0     |
| 3     | 0   | 0   | 0   | 0     | 0     | 2248 | -2093 |
| 29437 | 0   | 0   | 0   | 0     | 0     | 29   | -27   |

$$\begin{pmatrix} 8 & -27 & -66 & 42 & 0 & 0 & 0 & 0 \\ -40 & -45 & -47 & -38 & 0 & 0 & 0 & 0 \\ 0 & 126 & -18 & -23 & 0 & 0 & 0 & 0 \\ 103 & 26 & 0 & -53 & 0 & 0 & 0 & 0 \\ \hline -22 & 0 & 0 & 0 & -7 & -19 & -36 & 48 \\ 93 & 0 & 0 & 0 & -25 & -2 & -5 & 23 \\ -32 & 0 & 0 & 0 & -97 & 13 & 63 & 2 \\ 1 & 0 & 0 & 0 & -25 & -111 & 93 & -13 \end{pmatrix}$$

$$\begin{pmatrix} 8 & -27 & -66 & 42 & 0 & 0 & 0 & 0 \\ -40 & -45 & -47 & -38 & 0 & 0 & 0 & 0 & 0 \\ 0 & 126 & -18 & -23 & 0 & 0 & 0 & 0 \\ 103 & 26 & 0 & -53 & 0 & 0 & 0 & 0 \\ -22 & 0 & 0 & 0 & -7 & -19 & -36 & 48 \\ 93 & 0 & 0 & 0 & -25 & -2 & -5 & 23 \\ -32 & 0 & 0 & 0 & -97 & 13 & 63 & 2 \\ 1 & 0 & 0 & 0 & -25 & -111 & 93 & -13 \end{pmatrix}$$

$$\begin{pmatrix} -2 & 6 & -1 & -8 & 0 & 1 & -3 & 1 \\ 2 & -3 & -8 & 1 & 3 & -1 & 4 & 1 \\ 3 & -3 & 6 & 2 & -6 & 2 & 4 & 3 \\ -2 & -1 & -6 & 4 & -6 & -3 & 0 & -3 \\ 5 & -4 & 4 & -1 & -2 & 0 & -7 & 1 \\ 4 & 2 & 3 & -1 & -4 & 1 & -1 & -7 \\ -2 & 3 & -4 & 0 & 0 & 11 & 2 & -2 \\ 5 & 11 & 1 & 3 & -2 & 3 & -2 & 4 \end{pmatrix}$$

#### Recursive LLL

#### **RLLL**

- If d = 2, Return LLL(A);
- If d > 2,

  - $A_0' = RLLL(A_0).$
  - $A_1'' = RLLL(A_1).$
  - Reconstruct  $A' = \left(\frac{A'_0}{A'_1}\right)$ .
  - Return LLL(A')

#### Analysis

#### Why better?

- LLL (L<sup>2</sup>) complexity is in  $O(d^4\beta^2 + d^5\beta)$ .
- If d' = (d/2) and  $\beta = 2\beta'$  therefore  $d'^4\beta'^2 + d'^5\beta' < d^4\beta^2 + d^5\beta$ .
- All preprocessing are negligible compare to last LLL.

#### Complexity of the last LLL

- Assuming uniform distribution,  $\beta' = \frac{2\beta}{d}$ .
- $O(d^4\beta'^2 + d^5\beta')$
- $O\left(d^4\left(\frac{2\beta}{d}\right)^2+d^5\frac{2\beta}{d}\right)$
- $O(d^2\beta^2 + d^4\beta)$



#### Conclusion

- Introduction
- 2 Lattice Theory
  - Lattice Basics
  - Lattice Reduction
  - LLL
- 3 LLL for Modular Knapsack Lattice
  - Modular Knapsack lattice
  - LLL for modular knapsack lattice
- A Recursive LLL
- Conclusion



#### Conclusion

#### **Improvement**

- Previous complexity of LLL for knapsack lattice:  $O(d^{3+\varepsilon}\beta^2 + d^{4+\varepsilon}\beta)$ .
- New recursive techniques:  $O(d^{2+\varepsilon}\beta^2 + d^{4+\varepsilon}\beta)$ .

#### Future Work

- Specificity of Ideal Lattice.
- For given input and a given quality, estimate  $2^x$ .