





Optimization Algorithms

A. Comparing Optimization Algorithms

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Outline

- 1. Introduction
- 2. Performance Indicators and Time
- 3. Statistical Measures
- 4. Measures of Spread





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- For solving an optimization problem, we want to use the algorithm most suitable for it.
- What does this mean?

Performance Indicators and Time



• Key parameters^{3–6}

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 - 1. Solution quality reached after a certain runtime

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- Measure data samples A containing the results from multiple runs and estimate key parameters.

Runtime)

• What actually is runtime?

Measure the (absolute) consumed runtime of the algorithm in ms

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- ... so for research they may be less interesting, while for a specific application they do matter.

Measure the number of fully constructed and tested candidate solutions

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 - No clear relationship to real runtime
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 - 1 FE: very different costs in different situations!
- Relevant for comparing algorithms, but not so much for the practical application

Runtime

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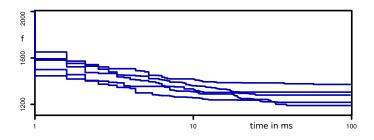
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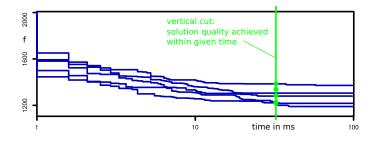
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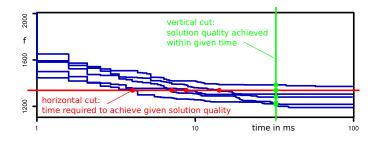
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- This question actually does not really need an answer...

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- Preferred by, e.g., the BBOB/COCO benchmark suite³:
 - Measures a time needed to reach a target function value \Rightarrow 'Algorithm A is two/ten/hundred times faster than Algorithm B in solving this problem."
 - Benchmark Perspective: No interpretable meaning to the fact that Algorithm A reaches a function value that is two/ten/hundred times smaller than the one reached by Algorithm B.

| Which Indicator is better? | |
|---------------------------------------|-------------------------------|
| Best objective function value reached | after a certain number of FEs |
| | |

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- This is the scenario in our JSSP example, too.

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- Best approach: Evaluate algorithm according to both methods.⁵⁶⁸

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 - 5. Based on known lower bounds

Statistical Measures



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- Results can be different for each run!
- Executing a randomized algorithm one time does not give reliable information.
- Statistical evaluation over a set of runs necessary.

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- How likely is it to roll a 1, 2, 3, 4, 5, or 6?



| # throws | number | f(1) | f(2) | f(3) | f(4) | f(5) | f(6) |
|----------|--------|--------|--------|--------|--------|--------|--------|
| 1 | 5 | 0.0000 | 0.0000 | 0.0000 | 0.0000 | 1.0000 | 0.0000 |



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| 3 | 1 | 0.3333 | 0.0000 | 0.0000 | 0.3333 | 0.3333 | 0.0000 |
| 4 | 4 | 0.2500 | 0.0000 | 0.0000 | 0.5000 | 0.2500 | 0.0000 |



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| 3 | 1 | 0.3333 | 0.0000 | 0.0000 | 0.3333 | 0.3333 | 0.0000 |
| 4 | 4 | 0.2500 | 0.0000 | 0.0000 | 0.5000 | 0.2500 | 0.0000 |
| 5 | 3 | 0.2000 | 0.0000 | 0.2000 | 0.4000 | 0.2000 | 0.0000 |



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| 5 | 3 | 0.2000 | 0.0000 | 0.2000 | 0.4000 | 0.2000 | 0.0000 |
| 6 | 3 | 0.1667 | 0.0000 | 0.3333 | 0.3333 | 0.1667 | 0.0000 |



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| 6 | 3 | 0.1667 | 0.0000 | 0.3333 | 0.3333 | 0.1667 | 0.0000 |
| 7 | 2 | 0.1429 | 0.1429 | 0.2857 | 0.2857 | 0.1429 | 0.0000 |
| 8 | 1 | 0.2500 | 0.1250 | 0.2500 | 0.2500 | 0.1250 | 0.0000 |
| 9 | 4 | 0.2222 | 0.1111 | 0.2222 | 0.3333 | 0.1111 | 0.0000 |
| 10 | 2 | 0.2000 | 0.2000 | 0.2000 | 0.3000 | 0.1000 | 0.0000 |



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| 9 | 4 | 0.2222 | 0.1111 | 0.2222 | 0.3333 | 0.1111 | 0.0000 |
| 10 | 2 | 0.2000 | 0.2000 | 0.2000 | 0.3000 | 0.1000 | 0.0000 |
| 11 | 6 | 0.1818 | 0.1818 | 0.1818 | 0.2727 | 0.0909 | 0.0909 |
| 12 | 3 | 0.1667 | 0.1667 | 0.2500 | 0.2500 | 0.0833 | 0.0833 |



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| 100 | | 0.1900 | 0.2100 | 0.1500 | 0.1600 | 0.1200 | 0.1700 |
| 1'000 | | 0.1700 | 0.1670 | 0.1620 | 0.1670 | 0.1570 | 0.1770 |
| 10'000 | | 0.1682 | 0.1699 | 0.1680 | 0.1661 | 0.1655 | 0.1623 |
| 100'000 | | 0.1671 | 0.1649 | 0.1664 | 0.1676 | 0.1668 | 0.1672 |
| 1'000'000 | | 0.1673 | 0.1663 | 0.1662 | 0.1673 | 0.1666 | 0.1664 |



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| 100'000'000 | | 0.1667 | 0.1666 | 0.1666 | 0.1667 | 0.1667 | 0.1667 |
| 1'000'000'000 | | 0.1667 | 0.1667 | 0.1667 | 0.1667 | 0.1667 | 0.1667 |

- Crucial Difference: distribution and sample
- A sample is what we measure.
- A distribution is the asymptotic result of the ideal process
- Statistical parameters of the distribution can be estimated from a sample
- Example: Dice Throw
- How likely is it to roll a 1, 2, 3, 4, 5, or 6?
- Never forget: All measured parameters are just estimates.



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- Never forget: All measured parameters are just estimates.
- The parameters of a random process cannot be measured directly, but only be approximated from multiple measures



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- We usually want to reduce this set of numbers to a single value which can give us an impression of what the "average outcome" (or result quality is).
- Two of the most common options for doing so, for estimating the "center" of a distribution, are to either compute the arithmetic mean or the median.

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Definition (Arithmetic Mean)

The arithmetic mean mean(A) is an estimate of the expected value of a data sample $A = (a_0, a_1, \dots, a_{n-1})$.

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$$\operatorname{mean}(A) = \frac{1}{n} \sum_{i=0}^{n-1} a_i \tag{1}$$

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$$\operatorname{med}(A) = \begin{cases} a_{\frac{n-1}{2}} & \text{if } n \text{ is odd} \\ \frac{1}{2} \left(a_{\frac{n}{2}-1} + a_{\frac{n}{2}} \right) & \text{otherwise} \end{cases}$$
 if $a_{i-1} \le a_i \ \forall i \in 1 \dots (n-1)$

(2)

• Sometimes the data contains outliers ⁹ 10.

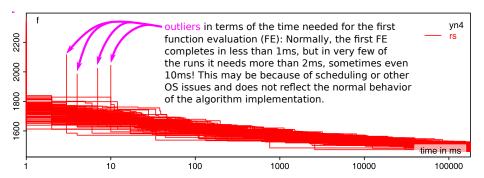
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- For example, maybe the operating system was updating itself during a run of one of our JSSP algorithms and, thus, took away much of the 3 minute computation budget.
- We can see that such odd times are possible, as our experimental data shows that there are sometimes outliers in the time it takes to create and evaluate the first candidate solution.



Example for Data Samples w/o Outlier

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- So we can conclude: It is best to have both the mean and median statistic of a given performance indicator.



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- For this, we can compute a measure of dispersion, i.e., a value that tells us whether the observations are stretched and spread far or squeezed tight around the center.

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The variance is the expectation of the squared deviation of a random variable from its mean. The variance var(A) of a data sample $A = (a_0, a_1, \dots, a_{n-1})$ with n observations can be estimated as:

$$var(A) = \frac{1}{n-1} \sum_{i=0}^{n-1} (a_i - mean(A))^2$$

Definition (Standard Deviation)

The statistical estimate $\operatorname{sd}(A)$ of the standard deviation of a data sample $A=(a_0,a_1,\ldots,a_{n-1})$ with n observations is the square root of the estimated variance $\operatorname{var}(A)$.

$$sd(A) = \sqrt{var(A)}$$

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Quantiles

Definition (Quantile)

The q-quantiles are the cut points that divide a sorted data sample $A=(a_0,a_1,\ldots,a_{n-1})$ where $a_{i-1}\leq a_i \ \forall i\in 1\ldots (n-1)$ into q-equally sized parts.

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$$\begin{array}{rcl} h & = & (n-1)\frac{k}{q} \\ \mathrm{quantile}_q^k(A) & = & \left\{ \begin{array}{ll} a_h & \text{if h is integer} \\ a_{\lfloor h \rfloor} + (h - \lfloor h \rfloor) * \left(a_{\lfloor h \rfloor + 1} - a_{\lfloor h \rfloor}\right) & \text{otherwise} \end{array} \right. \end{array}$$

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- We sometimes write things like "the 25% quantile," meaning quantile²⁵₁₀₀.

Standard Deviation: Example

ullet Two data samples A and B with $n_a=n_b=19$ values.

$$\begin{array}{rcl} A & = & (1,3,4,4,4,5,6,6,6,6,7,7,9,9,9,10,11,12,14) \\ \operatorname{mean}(A) & = & 7 \\ B & = & (1,3,4,4,4,5,6,6,6,6,7,7,9,9,9,10,11,12,10'008) \\ \operatorname{mean}(B) & = & 533 \end{array}$$

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$$var(A) = \frac{1}{19 - 1} \sum_{i=1}^{19} (a_i - mean(a))^2 = \frac{198}{18} = 11$$

$$var(B) = \frac{1}{19 - 1} \sum_{i=1}^{19} (b_i - mean(b))^2 = \frac{94'763'306}{18} \approx 5'264'628.1$$

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$$\operatorname{sd}(A) = \sqrt{\operatorname{var}A} = \sqrt{11} \approx 3.31662479$$

$$\operatorname{sd}(B) = \sqrt{\operatorname{var}B} = \sqrt{\frac{94'763'306}{18}} \approx 2294.477743$$

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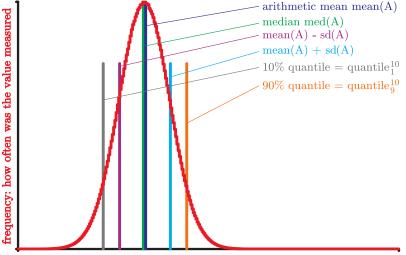
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$$quantile_4^1(A) = quantile_4^1(B) = 4.5$$

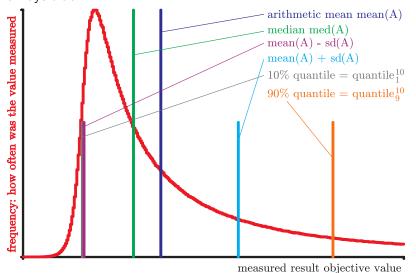
$$quantile_4^3(A) = quantile_4^3(B) = 9$$

ullet The implicit assumption that $mean \pm sd$ is a meaningful range is not always true!

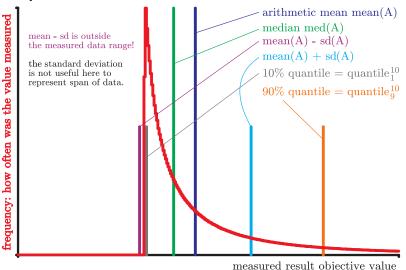


measured result objective value

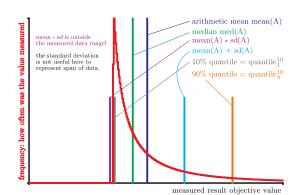
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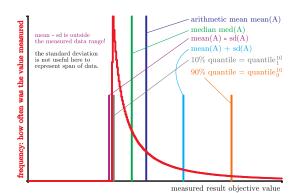
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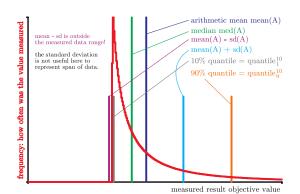
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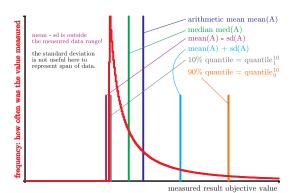
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- \bullet The implicit assumption that $\mathrm{mean} \pm \mathrm{sd}$ is a meaningful range is not always true!
- Such a shape is possible in optimization:
 - The global optimum marks a lower bound for the possible objective values.
 - A good algorithm often returns results which are close-to-optimal.
 - There may be a long tail of few but significantly worse runs.



Thank you

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