

**Assignment 5: Chapter 10, 11, 12****PROBLEM #1 (15 POINTS):**

Use the pumping lemma to show that this language is nonregular:

$$\{a^n b^n a^n\} = \{aba, aabbaa, aaabbbbaaa, \dots\}$$

$$\{a^n b a^n\} = \{aba, aabaa, aaabaaa, \dots\}$$

$$\{a^n b^{2n}\} = \{abb, aabbbb, aaabbbbb, \dots\}$$

**SOLUTION:**

1. First we assume that  $L$  is regular and  $n$  is the number of states.

Let  $w = a^n b^n a^n$ . Thus  $|w| = 3n \geq n$ .

By pumping lemma, let  $w = xyz$ , where  $|xy| \leq n$ .

Let  $x = a^p, y = a^q, z = a^r b^n a^n$ , where  $p + q + r = n, p \neq 0, q \neq 0, r \neq 0$ .

Let  $k = 2$ . Then  $xy^2z = a^p a^{2q} a^r b^n a^n$ .

The number of  $a$ 's is  $:(p + 2q + r) + n = (p + q + r) + q + n = 2n + q$

Hence,  $xy^2z = a^{n+q} b^n a^n$ . Since  $q \neq 0$ ,  $xy^2z$  is not of the form  $a^n b^n a^n$ .

Thus,  $xy^2z$  is not in  $L$  making  $L$  not regular.

2. First we assume that  $L$  is regular and  $n$  is the number of states.

Let  $w = a^n b a^n$ . Thus  $|w| = 2n + 1 \geq n$ .

By pumping lemma, let  $w = xyz$ , where  $|xy| \leq n$ .

Let  $x = a^p, y = a^q, z = a^r b a^n$ , where  $p + q + r = n, p \neq 0, q \neq 0, r \neq 0$ .

Let  $k = 2$ . Then  $xy^2z = a^p a^{2q} a^r b a^n$ .

The number of  $a$ 's is:  $(q + 2q + r) + n = (p + q + r) + q + n = 2n + q$

Hence,  $xy^2z = a^{n+q} b a^n$ . Since  $q \neq 0$ ,  $xy^2z$  is not of the form  $a^n b a^n$ .

Thus  $xy^2z$  is not in  $L$  making  $L$  not regular.

3. First we assume that  $L$  is regular and  $n$  is the number of states.

Let  $w = a^n b^{2n}$ . Thus  $|w| = 3n \geq n$ .

By pumping lemma, let  $w = xyz$ , where  $|xy| \leq n$ .

Let  $x = a^p, y = a^q, z = a^r b^{2n}$ , where  $p + q + r = n, p \neq 0, q \neq 0, r \neq 0$ .

Let  $k = 2$ . Then  $xy^2z = a^p a^{2q} a^r b^{2n}$ .

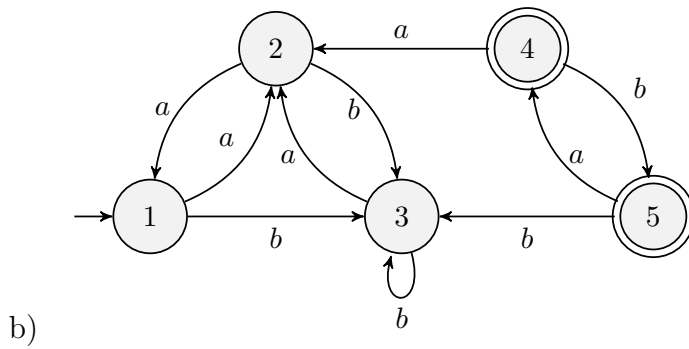
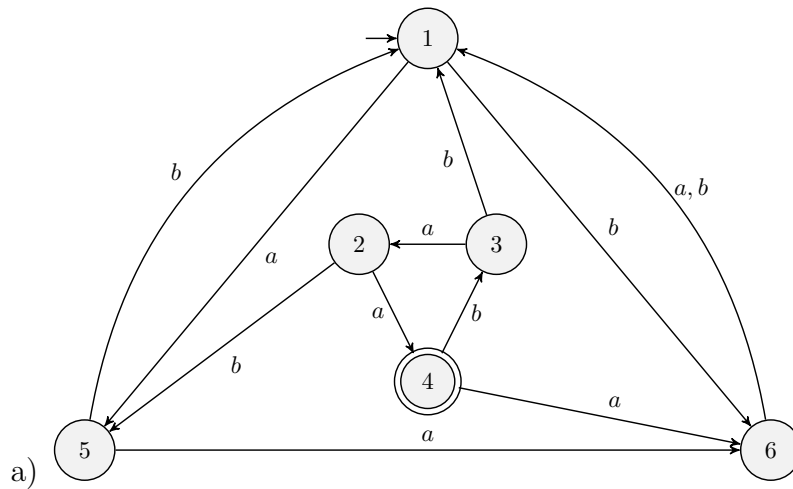
The number of  $a$ 's is:  $(q + 2q + r) = (p + q + r) + q = n + q$

Hence,  $xy^2z = a^{n+q} b^{2n}$ . Since  $q \neq 0$ ,  $xy^2z$  is not of the form  $a^n b^{2n}$ .

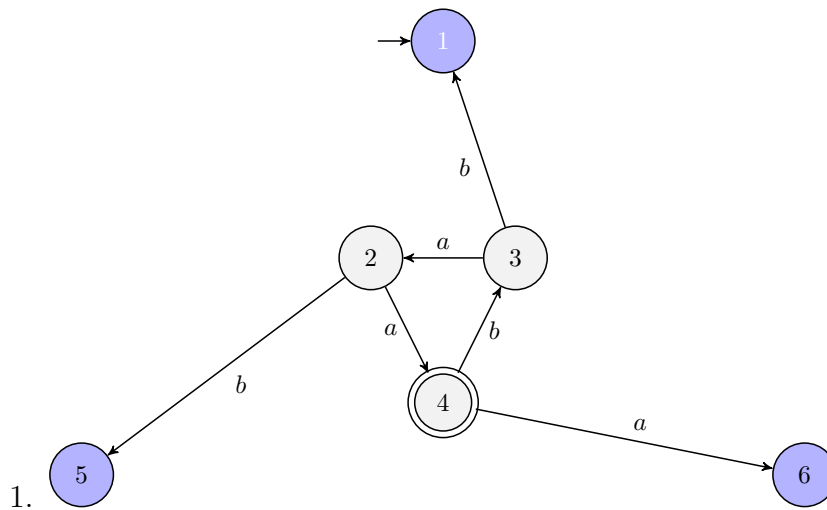
Thus  $xy^2z$  is not in  $L$  making  $L$  not regular.

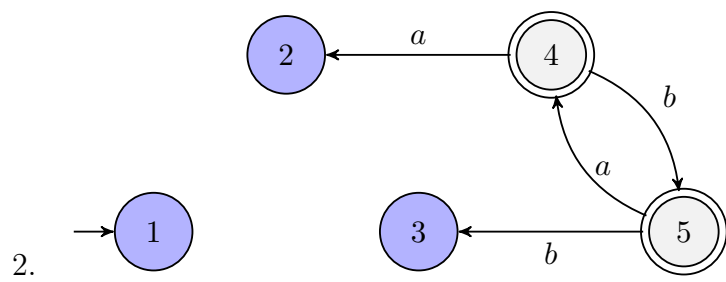
PROBLEM #2 (10 POINTS):

By using blue paint, determine which of the following FA's accept any words:



SOLUTION:





PROBLEM #3 (18 POINTS):

Describe the language generated by the following context free grammar (CFG) in English and regular expressions:

- a)  $S \rightarrow SS$   
 $S \rightarrow ZZZ$   
 $Z \rightarrow bZ$   
 $Z \rightarrow Zb$   
 $Z \rightarrow a$
- b)  $S \rightarrow aS$   
 $S \rightarrow bb$
- c)  $S \rightarrow XYX$   
 $X \rightarrow aX$   
 $X \rightarrow bX$   
 $X \rightarrow \Lambda$   
 $Y \rightarrow bbb$

SOLUTION:

- a) This CFG can be described as any number of  $b$ 's with a multiple of 3  $a$ 's.  
Regular Expression:  $(b^*ab^*ab^*ab^*)^*$
- b) This CFG can be described as any number of  $a$ 's followed by two  $b$ 's.  
Regular Expression:  $a^*bb$
- c) This CFG can be described as any string of  $a$ 's and  $b$ 's containing the string  $bbb$ .  
Regular Expression:  $(a + b)^*bbb(a + b)^*$

PROBLEM #4 (15 POINTS):

Find CFG for the following languages over the alphabet  $\Sigma = \{a, b\}$ :

- a) All words that have different first and last letters.
- b) All words in which the letter  $b$  is never tripled.
- c) All words that do not have substring  $ab$ .

SOLUTION:

a)  $S \rightarrow aXb \mid bXa$

$$X \rightarrow aX \mid bX \mid a \mid b \mid \Lambda$$

b)  $S \rightarrow \Lambda \mid b \mid bb \mid Sa \mid aS \mid Sab \mid Sabb$

c)  $S \rightarrow \Lambda \mid bS \mid bX \mid X$

$$X \rightarrow aX \mid \Lambda$$

PROBLEM #5 (10 POINTS):

Show that the CFG below is ambiguous by finding a word with two distinct syntax trees. Show both syntax trees.

a)  $S \rightarrow Sbb$   
 $S \rightarrow Sbbb$   
 $S \rightarrow b$

b)  $S \rightarrow AA$   
 $A \rightarrow AAA|a|bA|Ab$

SOLUTION:

1. The CFG is ambiguous because the word *bbbbbb* can follow:

Rule	Application	Result
Start $\rightarrow$ S	Start	S
S $\rightarrow$ Sbbb	<b>S</b>	<b>Sbbb</b>
S $\rightarrow$ Sbb	<b>Sbbb</b>	<b>Sbbbbbb</b>
S $\rightarrow$ b	<b>Sbbbbbb</b>	<b>bbbbbbb</b>
Start $\rightarrow$ S	Start	S
S $\rightarrow$ Sbb	<b>S</b>	<b>Sbb</b>
S $\rightarrow$ Sbbb	<b>Sbb</b>	<b>Sbbbbbb</b>
S $\rightarrow$ b	<b>Sbbbbbb</b>	<b>bbbbbbb</b>

2. The CFG is ambiguous because the word *aba* can follow:

Rule	Application	Result
Start $\rightarrow$ S	Start	S
S $\rightarrow$ AA	<b>S</b>	<b>AA</b>
A $\rightarrow$ Ab	<b>AA</b>	<b>AbA</b>
A $\rightarrow$ a	<b>AbA</b>	<b>abA</b>
A $\rightarrow$ a	<b>abA</b>	<b>aba</b>
Start $\rightarrow$ S	Start	S
S $\rightarrow$ AA	<b>S</b>	<b>AA</b>
A $\rightarrow$ a	<b>AA</b>	<b>aA</b>
A $\rightarrow$ bA	<b>aA</b>	<b>abA</b>
A $\rightarrow$ a	<b>aba</b>	<b>aba</b>