

Assignment 2: Chapter 4

PROBLEM #1 (15 POINTS):

Construct a regular expression defining each of the following languages over the alphabet $\Sigma = \{a, b\}$:

- a) The language of all words that do not begin with ba .
- b) The language of all words in which the total number of b's is divisible by 3 no matter how they are distributed, such as $bbabbaabab$.
- c) All words that contain exactly 2 b's or exactly 3 b's, not more.
- d) All words in which a appears tripled, if at all. This means that every clump of a's contains 3 or 6 or 9 or 12 ... a's.
- e) All words that contain at least one of the strings S_1, S_2, S_3 , or S_4 .

SOLUTION:

- a) a^*b^*

Contains only a's or only b's or a's followed by b's, therefore the string ba should never appear.

- b) $(a^*ba^*ba^*ba^*)^*$

- c) $a^*(b + \Lambda)a^*ba^*ba^*$

- d) $b^*(aaa)^*b^*$

Nothing was said about there being no a's so $(aaa)^*$ allows for only strings of b's or an empty string.

- e) $(S_1 + S_2 + S_3 + S_4)(S_1 + S_2 + S_3 + S_4)^*$

PROBLEM #2 (15 POINTS):

Construct a regular expression defining each of the following languages over the alphabet $\Sigma = \{a, b\}$:

- a) All strings that end in a double letter.
- b) All strings that do not end in a double letter.
- c) All strings that have exactly one double letter in them.
- d) All words in which the letter b is never tripled. This means that no word contains the substring bbb .
- e) All words in which a is tripled or b is tripled, but not both. This means each word contains the substring aaa or the substring bbb but not both.

SOLUTION:

a) $(a + b)^*(aa + bb)$

b) $(a + b)^*(ab + ba) + a + b + \Lambda$

c) $(b + \Lambda)(ab)^*aa(ba)^*(b + \Lambda) + (a + \Lambda)(ba)^*bb(ab)^*(a + \Lambda)$

d) $(\Lambda + b + bb)(a + ab + abb)^*$

e) $(\Lambda + b + bb)(a + ab + abb)^*aaa(\Lambda + b + bb)(a + ab + abb)^* +$
 $(\Lambda + a + aa)(b + ba + baa)^*bbb(\Lambda + a + aa)(b + ba + baa)^*$

PROBLEM #3 (10 POINTS):

Let us consider the regular expression

$$(a + b)^* a (a + b)^* b (a + b)$$

Show that this is equivalent to

$$(a + b)^* ab (a + b)^*$$

In the sense that they define the same language.

SOLUTION:

$$(a + b)^* = \{\Lambda, a, b, aa, ab, bb, aaa, aab, aba, abb, bbb, \dots\}$$

$$a(a + b)^* = \{a, aa, ab, aaa, aab, aba, \dots\}$$

$$b(a + b) = \{ba, bb\}$$

$$(a + b)^* = \{\Lambda, a, b, aa, ab, bb, aaa, aab, aba, abb, bbb, \dots\}$$

$$ab(a + b)^* = \{ab, aba, abb, abaa, abab, abbb, \dots\}$$

$$a(a + b)^* b(a + b) = ab(a + b)^*$$

PROBLEM #4 (5 POINTS):

If the only difference between L and L^* is the word Λ , is the only difference between L^2 and L^* the word Λ ? Show by example.

SOLUTION:

No, if $L = \{a, b, aa, ab, ba, bb, \dots\}$ then $L^* = \{\Lambda, a, b, aa, ab, ba, bb, \dots\}$.

But $L^2 = \{aa, bb, aaaa, abab, baba, bbbb\}$ which does not include $\{a, b\}$.

PROBLEM #5 (12 POINTS):

Describe in English phrases the languages associated with the following regular expressions.

a $(a + b)^*a(\Lambda + bbbb)$

b $a(a + bb)^*$

c $a(aa)^*b(bb)^*$

d $((a + b)a)^*$

SOLUTION:

- a The set of strings over the alphabet $\{a, b\}$ that ends with a or $bbbb$ and must have at least one a .
- b The set of strings over the alphabet $\{a, b\}$ with any number of a 's followed by none or even numbers of b 's.
- c The set of strings over the alphabet $\{a, b\}$ that contains only odd numbers of a 's followed by odd numbers of b 's
- d The set of strings over the alphabet $\{a, b\}$ with any number of a 's or b 's ending with an a (Λ included)