# From Regular Expressions To Deterministic Automata

#### **Paper**

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# Paper Overview

- A new algorithm to go from regular expressions to DFAs
- Refines Brzozowski's elegant algorithm into McNaughton and Yamada's faster algorithm

## Paper Overview

#### Brozowski's construction of a DFA

Input: Regular Expression E

Output: DFA D

- 1. The states of D are the distinct derivatives w^{-1}E, for all strings w
- 2. Construct a transition under symbol *a* from state *p* to state *q* if and only if *p* is for derivative w^{-1}E, for some *w*, and *q* is for [(wa)^(-1)]E
- 3. The state for *E* is the start state.
- 4. A state is an accepting state if and only if it is for a derivative w^{-1}E, for some w, and delta([w^(-1)]E) = 1; that is the empty string is in [L(wa)^(-1)]E.

#### McNaughton and Yamada's construction of a DFA

Input: Marked Regular Expression *E'* Output: DFA *M'* 

- 1. *M'* has a state for the continuation of each marked symbol in *E'*
- 2. Construct a transition under symbol *a* from state *p* to state *q* for the continuation of *a* if and only if *p* is for some continuation *C* and *C* can generate a string with a leading *a*
- 3. The state for the entire expression *E'* is the start state.
- A state is an accepting state if and only if it is for a continuation C and delta(C)=1

#### Paper Overview

#### Brozowski's construction of a DFA

Input: Regular Expression E

Output: DFA D

Introduces the derivative and delta of a regular expression

#### McNaughton and Yamada's construction of a DFA

Input: Marked Regular Expression E'

Output: DFA M'

Introduces marked regular expressions

#### **New Algorithm**

Input: Marked Regular Expression *E'* Output: DFA *M'* 

- Constructs transitions using first(E') and follow(b), where b is a state
- Accepting states are found by testing if a special end character is in follow(b)

# Take away from Brzozowski

#### **Delta Notation**

- $\delta(E)$  stands for 1 if L(E) contains the empty string
- $\delta(E)$  is computed recursively on E
  - $\circ$   $\delta(0) = 0$
  - $\circ$  δ(1) = 1 (where 1 is taken to mean  $\varepsilon$ )
  - $\circ \quad \delta(E_1 \cup E_2) = \delta(E_1) + \delta(E_2)$
  - $\circ \quad \delta(E_1 E_2) = \delta(E_1) \cdot \delta(E_2)$
  - $\circ \quad \delta(E^*) = 1$
- It is clear then that  $\delta(E)$  represents the number of possibilities the expression has of containing the empty string  $\varepsilon$ .

#### Take aways from McNaughton and Yamada

#### Marked Expressions

- Mark all input symbols in a regular expression
  - o  $(ab \cup b)^*ba$  becomes  $(a_1b_2 \cup b_3)^*b_4a_5$
- These subscripted symbols are distinct and will become the states of an automaton

#### First

- $first(E) = \{a : av \in L(E)\}$
- The set of possible first symbols in the language of a regular expression E

#### **Follow**

- $follow_E(a) = \{b : uabv \in L(E)\}$
- The set of possible symbols that come immediately after the symbol a in the language of E
- For a marked symbol  $a_1$ ,  $follow(a_1)$  will contain the set of marked symbols corresponding to the next possible states in the automaton

# **New Algorithm**

Input: Marked Regular Expression E'

Output: DFA M'

- 1. M' has a start state plus a state for each marked symbol a in E'
- 2. Construct a transition from the start state to the state for *a* if and only if *a* is in *first(E')*
- 3. Construct a transition from the state *b* to the state *a* if and only if *a* is in *follows(b)*
- 4. The start state is an accepting state if and only if delta(E')=1
- 5. The state for a is an accepting state if and only if ! is in follows(a)

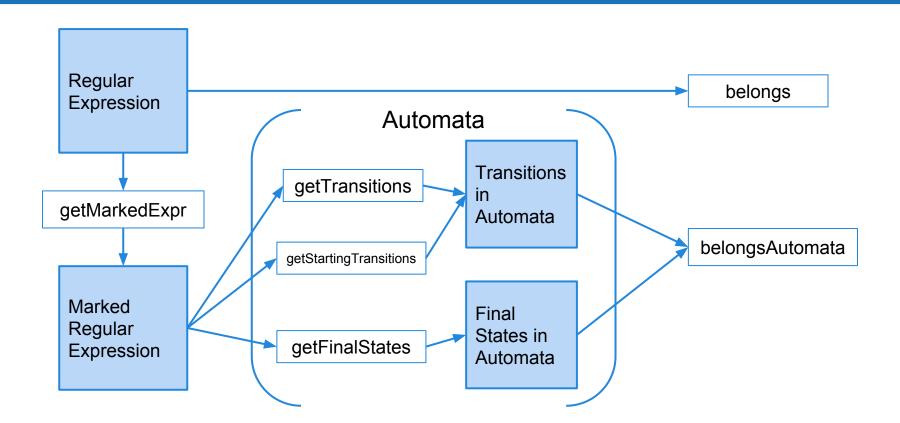
# General Idea of Using New Algorithm

- 1. Construct an automata from a marked expression
- The marks on the transitions are then erased, resulting in an NFA for the original unmarked expression
  - Note: This approach works for the usual operations of union, concatenation, and iteration, but does not work with intersection and complement. This is because marking and unmarking do not preserve the language generated by regular expressions with these operations

# **Project Implementation Overview**

 We originally intended to implement all 3 algorithms, however we only implemented the new algorithm

# What we did implement



## **Implementation**

- We used OCaml's algebraic types extensively.
  - Regex, consisting of a char, a Union, a Concatenation, a Star, or Epsilon.
  - These are all custom datatypes that use Cartesian product with regex.

```
type regex =
| Eps
| Lit of char
| Union of regex * regex
| Concat of regex * regex
| Star of regex;;
```

```
type markedRegex =

| MEps
| MLit of char * int
| MUnion of markedRegex * markedRegex
| MConcat of markedRegex * markedRegex
| MStar of markedRegex;;
```

## **Implementation**

#### Regular Expression Functions

(marked & unmarked)

- belongs
  - In: string and expr
  - Out: is the string accepted by expr
- delta
  - In: expr
  - Out: delta(expr) --> the potential for a string in the expr to include epsilon
- follow
  - In: expr and symbol
  - Out:The set of possible symbols that come immediately after the symbol a in the language of E

- first
  - In: expr
  - Out: a list of starting symbols in the expression
- last
  - o In: expr
  - Out: a list of ending symbols in the expression

## **Implementation**

#### **Automata**

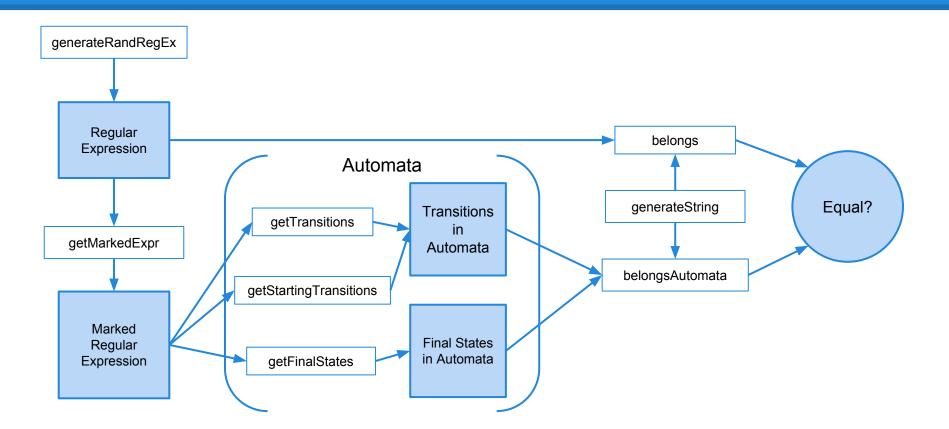
- belongs:
  - In: string, transitionlist, finalstates and currentstate
  - Out: is string accepted by automata
- getStates:
  - o In: expr
  - Out: states
- getTransitions:
  - o In: expr
  - calls two helper functions getTransitionsFromStates expr getStates
  - Out: prepares a list of transitions

- getfinalsates:
  - In: expr and states
  - Out: accepting states
- getNextStates:
  - In: transitionLists and currentstate
  - Out: possible next states
- getStartingTransitions:
  - In: expr and startingStates
  - Out: transitions from first to next state

# **Testing the Implementation**

- Created a function test
  - make an alphabet
  - randomly generates
    - string
    - regex
  - makes automata from regex
  - o calls belong on regex and automata
  - compares the two

# **Testing the Implementation**



## Script

 Script runs endlessly concatenating output to a file in which we grep errors.

```
#!/bin/bash
cat output.txt | grep "incorrect"
exit
```

#### Results

- Test cases: 25,107,109
  - (actually still running...)
    - generated regular expressions are not minimized
- Incorrect cases: 32
- Accuracy: 99.999872546%
- So we made a small mistake somewhere...

```
x [screen 0: bash] trh010@linuxremote2:/nfs/unixspace/linux/accounts...
      #ncorrect: bbbbbababbbbaa with regex: (b(b)*UbUa)*a
      incorrect: aabaabab with regex: (aeUbeUeUa)*b
      incorrect: aababbbbb with regex: (aeUeaUb)*(b)*
      incorrect: baabaaba with regex: (beaeaeUb)*(a)*
      incorrect: aabaababaaabb with regex: e(aeUbeUa)*bUaUa
      incorrect: baabaaba with regex: (beaeaeUb)*(a)*
      incorrect: aaaaababb with regex: (eaeUb(a)*)*eUeUb
     8 incorrect: bbabababba with regex: (baUbUb)*eUa
      incorrect: ababaab with regex: (ab(ea)*U(a)*)*eUeUb
    10 incorrect: aababbbbb with regex: (aeUeaUb)*(b)*
    11 incorrect: bbabababba with regex: (baUbUb)*eUa
      incorrect: baabaaba with regex: (beaeaeUb)*(a)*
      incorrect: bbbbaba with regex: (beUaeUeUb)*a
      incorrect: bbbbababbbbaa with regex: (b(b)*UbUa)*a
      incorrect: bbaa with regex: (baUeUeU(b)*)*eUbUa
      incorrect: bbaa with regex: (ebeUbUba)*bUa
      incorrect: bbaa with regex: (baUeUeU(b)*)*eUbUa
      incorrect: baababaababaaba with regex: (b(a)*e(b)*)*aUb
      incorrect: baabaaba with regex: (beaeaeUb)*(a)*
      incorrect: aaaaababb with regex: (eaeUb(a)*)*eUeUb
      incorrect: babbbbbbabba with regex: (bUbaaUb)*eUa
      incorrect: abaabababb with regex: (abUeUbUa)*(b)*
      incorrect: ababaab with regex: (ab(ea)*U(a)*)*eUeUb
      incorrect: bbabababba with regex: (baUbUb)*eUa
      incorrect: aabaababaaabb with regex: e(aeUbeUa)*bUaUa
      incorrect: bbaa with regex: (baUeUeU(b)*)*eUbUa
      incorrect: aabaababaaabb with regex: e(aeUbeUa)*bUaUa
      incorrect: aabb with regex: (eae(a)*Ub)*b
      incorrect: aabaabab with regex: (aeUbeUeUa)*b
      incorrect: aabbabaabaaabb with regex: (eUaaUea(b)*)*(b)*
      incorrect: bbbbaba with regex: (beUaeUeUb)*a
      incorrect: aababbbbb with regex: (aeUeaUb)*(b)*
                                                                 1,1
```

## We have a lot correct though...

```
screen 0: bashl trh010@linuxremote2:/nfs/unixspace/linux/accounts...
    1234698 correct: bbba with regex: (b)*
   1234699 correct: aaabb with regex: (b)*
   1234700 correct: abbabbababababa with regex: (b)*
   1234701 correct: b with regex: (b)*
   1234702 correct: bbaababbbaa with regex: (b)*
   1234703 correct: baabbaaba with regex: (b)*
   1234704 correct: aaabbbbab with regex: (b)*
   1234705 correct: bbbabbbbaab with regex: (b)*
   1234706 correct: bbabbba with regex: (b)*
   1234707 correct: baaabbbaa with regex: (b)*
   1234708 correct: aababbbbbbbbbbb with regex: (b)*
   1234709 correct: aaab with regex: (b)*
   1234710 correct: abbbaab with regex: (b)*
   1234711 correct: abbabaabba with regex: (b)*
   1234712 correct: babbb with regex: (b)*
   1234713 correct: bbab with regex: (bUbUbbUeb)*
   1234714 correct: bababbaaababbba with regex: (bUbUbbUeb)*
   1234715 correct: abaaaabaab with regex: (bUbUbbUeb)*
   1234716 correct: babbabb with regex: (bUbUbbUeb)*
   1234717 correct: baaabb with regex: aUbUeUaUe(eea)*
   1234718 correct: bbaabbaabababb with regex: aUbUeUaUe(eea)*
   1234719 correct: bb with regex: aUbUeUaUe(eea)*
    1234720 correct: bbbaababbba with regex: aUbUeUaUe(eea)*
   1234721 correct: abaaaa with regex: aUbUeUaUe(eea)*
   1234722 correct: a with regex: aUbUeUaUe(eea)*
    1234723 correct: abbbaa with regex: aUbUeUaUe(eea)*
    1234724 correct: a with regex: aUbUeUaUe(eea)*
    1234725 correct: bbaabbaabbaaab with regex: aUbUeUaUe(eea)*
   1234726 correct: bbbaaaaaaababa with regex: aUbUeUaUe(eea)*
    1234727 correct: bbbbbbbba with regex: aUbUeUaUe(eea)*
    1234728 correct: baaab with regex: aUbUeUaUe(eea)*
    1234729 ∰orrect: abaab with regex: aUbUeUaUe(eea)*
                                                                 1234729.1
```

```
screen 0: bashl trh010@linuxremote2:/nfs/unixspace/linux/accounts...
           morrect: aabbbbbbaabaaa with regex: a
           correct: bbaaababbbabbab with regex: a
           correct: aabbbbbbabaa with regex: a
           correct: abbbaabbaba with regex: a
           correct: babbb with regex: a
           correct: aa with regex: a
           correct: baabbbaaaaabba with regex: a
           correct: ababbbbab with regex: a
           correct: bbaaaabbabbbaab with regex: a
           correct: bbaabbbb with regex: a
           correct: bab with regex: a
           correct: aa with regex: a
           correct: baaabbaaaabbab with regex: a
           correct: aabab with regex: (a)*b
           correct: bab with regex: (a)*b
           correct: babababab with regex: (a)*b
           correct: baaab with regex: (a)*b
           correct: babaaaaa with regex: (a)*b
           correct: abaaaaaaa with regex: eUb
           correct: bbbb with regex: eUb
           correct: aaababbaababba with regex: eUb
           correct: bbaaabbabaa with regex: eUb
           correct: bbaaaabaaabbbba with regex: eUb
           correct: babbbbaaa with regex: eÜb
           correct: abb with regex: eUb
           correct: abb with regex: eUb
           correct: a with regex: eUb
           correct: aaaabaab with regex: eUb
           correct: baaabababbaabb with regex: eUb
           correct: bbbbbbabaaababb with regex: eUb
           correct: aabbbbba with regex: eUb
    9999978 correct: aab with regex: eUb
                                                                 9999947,1
```

#### and more...

```
screen 0: bashl trh010@linuxremote2:/nfs/unixspace/linux/accounts...
           Borrect: baaab with regex: (b)*UeUbebUbeUa
   4999975 correct: bbabbbbba with regex: (b)*UeUbebUbeUa
   4999976 correct: abababbaab with regex: (b)*UeUbebUbeUa
   4999977 correct: aaab with regex: (b)*UeUbebUbeUa
   4999978 correct: abaaaa with regex: (b)*UeUbebUbeUa
   4999979 correct: ababbbbabbb with regex: (b)*UeUbebUbeUa
   4999980 correct: aabbb with regex: (b)*UeUbebUbeUa
   4999981 correct: ab with regex: (b)*UeUbebUbeUa
   4999982 correct: aabaaabbba with regex: (b)*UeUbebUbeUa
   4999983 correct: bbabaa with regex: (b)*UeUbebUbeUa
   4999984 correct: a with regex: (b)*UeUbebUbeUa
   4999985 correct: aaabaaabaaaaaaa with regex: (b)*UeUbebUbeUa
   4999986 correct: abbaab with regex: (b)*UeUbebUbeUa
   4999987 correct: abaababbbbab with regex: (b)*UeUbebUbeUa
   4999988 correct: aabbbabbbbbbb with regex: (b)*UeUbebUbeUa
   4999989 correct: bbbaabaaababa with regex: (b)*UeUbebUbeUa
   4999990 correct: baabaaab with regex: (b)*UeUbebUbeUa
   4999991 correct: baaabbaaba with regex: eUeUaUeUa
   4999992 correct: baaababbbaaa with regex: eUeUaUeUa
    4999993 correct: bbabab with regex: eŪeUaUeUa
    4999994 correct: babbabbb with regex: eUeUaUeUa
    4999995 correct: abaaaaaabbaaaa with regex: eUeUaUeUa
    4999996 correct: bbbbaaabbbaaaaa with regex: eUeUaUeUa
   4999997 correct: baaababaaa with regex: eUeUaUeUa
   4999998 correct: aaabbaabbbbaba with regex: eUeUaUeUa
   4999999 correct: babaa with regex: eUeUaUeUa
   5000000 correct: bbabb with regex: eUeUaUeUa
   5000001 correct: ab with regex: eUeUaUeUa
   5000002 correct: abbba with regex: eUeUaUeUa
   5000003 correct: bbaaaaaaaabb with regex: eUeUaUeUa
   5000004 correct: bbbaabbaababbbb with regex: eUeUaUeUa
   5000005 correct: a with regex: eUeUaUeUa
                                                                 4999974,1
```

```
screen 0: bashl trh010@linuxremote2:/nfs/unixspace/linux/accounts...
          14999986 correct: baaaabbbb with regex: bUaUa
   14999987 correct: bbbabbbbbbbb with regex: bUaUa
  14999988 correct: bbbbbaabba with regex: bUaUa
  14999989 correct: baaabab with regex: bUaUa
  14999990 correct: ab with regex: bUaUa
  14999991 correct: bbbababbb with regex: bUaUa
  14999992 correct: bbbbaaabaab with regex: bUaUa
  14999993 correct: bab with regex: bUaŪa
  14999994 correct: bbaaaabb with regex: bUaUa
   14999995 correct: abaaaa with regex: bUaUa
  14999996 correct: abaaabbaaa with regex: bUaUa
  14999997 correct: babbabbbbabbabb with regex: bUaUa
  14999998 correct: aaaabbb with regex: bUaŪa
  14999999 correct: bbbbb with regex: bUaUa
   15000000 Խrrect: aabaab with regex: bUaUa
  15000001 correct: abababaababaa with regex: bUaUa
  15000002 correct: baaa with regex: bUaUa
   .5000003 correct: aababbaaabaaa with regex: bUaUa
   5000004 correct: aabbaabbbabb with regex: bUaUa
   5000005 correct: bbbbba with regex: bUaUa
   5000006 correct: abbbbaabab with regex: bUaUa
    5000007 correct: abbbbabbbbaa with regex: bUaUa
   5000008 correct: ababababaaabb with regex: bUaUa
    5000009 correct: bb with regex: bUaUa
   5000010 correct: aabaaabbaa with regex: bUaUa
     00011 correct: aaaabbbbababb with regex: ea
          correct: abaabaababaab with regex: ea
   5000013 correct: aabbbbbbbaabbabb with regex: ea
  15000014 correct: aaab with regex: ea
   15000015 correct: bbabbaabb with regex: ea
  15000016 correct: b with regex: ea
                                                                15000000.1
```

# Questions?