

Team Notebook

HSG SQRT

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Contents

1 Combinatorics	3	3 Dynamic Programming	6	5.11 HullDiameter	11
1.1 Permutations	3	3.1 DP Convex Hull Trick	6	5.12 InsidePolygon	11
1.1.1 Factorial	3	3.2 DP Divide and Conquer	7	5.13 linearTransformation	12
1.1.2 Cycles	3	3.3 DP Knuth	7	5.14 lineDistance	12
1.1.3 Derangements	3	3.4 Subset Sum Convolution	7	5.15 lineIntersection	12
1.1.4 Burnside's lemma	3	4 Formula	7	5.16 LineProjectionReflection	12
1.2 Partitions and subsets	3	4.1 Equations	7	5.17 ManhattanMST	12
1.2.1 Partition function	3	4.2 Recurrences	8	5.18 MinimumEnclosingCircle	13
1.2.2 Lucas' Theorem	3	4.3 Trigonometry	8	5.19 OnSegment	13
1.3 General purpose numbers	3	4.4 Geometry	8	5.20 Point	13
1.3.1 Bernoulli numbers	3	4.4.1 Triangles	8	5.21 PointInsideHull	13
1.3.2 Stirling numbers of the first kind	3	4.4.2 Quadrilaterals	8	5.22 PolygonArea	13
1.3.3 Eulerian numbers	3	4.4.3 Spherical coordinates	8	5.23 PolygonCenter	14
1.3.4 Stirling numbers of the second kind	3	4.5 Sums	8	5.24 PolygonCut	14
1.3.5 Bell numbers	3	4.6 Probability theory	8	5.25 PolygonUnion	14
1.3.6 Labeled unrooted trees	3	4.6.1 Discrete distributions	8	5.26 SegmentDistance	14
1.3.7 Catalan numbers	4	5 Geometry	9	5.27 SegmentIntersection	14
2 Data Structure	4	5.1 Angle	9	5.28 sideOf	15
2.1 Gilbert Order	4	5.2 CircleIntersection	9	6 Graph	15
2.2 Li Chao Tree	4	5.3 CircleLine	9	6.1 2SAT	15
2.3 Persistent Segment Tree	4	5.4 CirclePolygonIntersection	9	6.2 Blocc Cut Tree	15
2.4 Persistent Trie	5	5.5 CircleTangents	10	6.3 DeMen Bipartite Matching	16
2.5 Segment Tree 2D	5	5.6 circumcircle	10	6.4 General Matching	16
2.6 STL	6	5.7 ClosestPair	10	6.5 Gomory Hu Tree	17
2.7 Treap	6	5.8 ConvexHull	10	6.6 Max Cost General Matching	17
		5.9 DelaunayTriangulation	10	6.7 Max Flow	19
		5.10 FastDelaunay	11	6.8 Maximum Independent Set	20
				6.9 Prüfer Decode	20

6.10	Prüfer Encode	20	7.7	Karatsuba	22	8.3	Primes	24
			7.8	Lagrange	23	8.4	Estimates	24
7	Mathematics	21	7.9	Pisano	23	8.5	Mobius Function	24
7.1	CRT	21	7.10	Rho and Miller Rabin	23			
7.2	Extended Euclid	21	7.11	Tonelli Shanks	24	9	String	25
7.3	Fast Fourier Transform - Mod	21				9.1	Aho Corasick	25
7.4	Fast Fourier Transform	21	8	Number Theory	24	9.2	Suffix Array	25
7.5	Fast Walsh–Hadamard Transform	21	8.1	Bézout’s identity	24			
7.6	Gaussian Elimination	22	8.2	Pythagorean Triples	24			

1 Combinatorics

1.1 Permutations

1.1.1 Factorial

n	1	2	3	4	5	6	7	8	9	10
$n!$	1	2	6	24	120	720	5040	40320	362880	3628800
n	11	12	13	14	15	16	17			
$n!$	4.0e7	4.8e8	6.2e9	8.7e10	1.3e12	2.1e13	3.6e14			
n	20	25	30	40	50	100	150	171		
$n!$	2e18	2e25	3e32	8e47	3e64	9e157	6e262	>DBL_MAX		

1.1.2 Cycles

Let $g_S(n)$ be the number of n -permutations whose cycle lengths all belong to the set S . Then

$$\sum_{n=0}^{\infty} g_S(n) \frac{x^n}{n!} = \exp \left(\sum_{n \in S} \frac{x^n}{n} \right)$$

1.1.3 Derangements

Permutations of a set such that none of the elements appear in their original position.

$$D(n) = (n-1)(D(n-1) + D(n-2)) = nD(n-1) + (-1)^n = \left\lfloor \frac{n!}{e} \right\rfloor$$

1.1.4 Burnside's lemma

Given a group G of symmetries and a set X , the number of elements of X up to symmetry equals

$$\frac{1}{|G|} \sum_{g \in G} |X^g|,$$

where X^g are the elements fixed by g ($g.x = x$).

If $f(n)$ counts “configurations” (of some sort) of length n , we can ignore rotational symmetry using $G = Z_n$ to get

$$g(n) = \frac{1}{n} \sum_{k=0}^{n-1} f(\gcd(n, k)) = \frac{1}{n} \sum_{k|n} f(k) \phi(n/k).$$

1.2 Partitions and subsets

1.2.1 Partition function

Number of ways of writing n as a sum of positive integers, disregarding the order of the summands.

$$p(0) = 1, \quad p(n) = \sum_{k \in \mathbb{Z} \setminus \{0\}} (-1)^{k+1} p(n - k(3k-1)/2)$$

$$p(n) \sim 0.145/n \cdot \exp(2.56\sqrt{n})$$

n	0	1	2	3	4	5	6	7	8	9	20	50	100
$p(n)$	1	1	2	3	5	7	11	15	22	30	627	$\sim 2e5$	$\sim 2e8$

1.2.2 Lucas' Theorem

Let n, m be non-negative integers and p a prime. Write $n = n_k p^k + \dots + n_1 p + n_0$ and $m = m_k p^k + \dots + m_1 p + m_0$. Then $\binom{n}{m} \equiv \prod_{i=0}^k \binom{n_i}{m_i} \pmod{p}$.

1.3 General purpose numbers

1.3.1 Bernoulli numbers

EGF of Bernoulli numbers is $B(t) = \frac{t}{e^t - 1}$ (FFT-able).

$$B[0, \dots] = [1, -\frac{1}{2}, \frac{1}{6}, 0, -\frac{1}{30}, 0, \frac{1}{42}, \dots]$$

Sums of powers:

$$\sum_{i=1}^n i^m = \frac{1}{m+1} \sum_{k=0}^m \binom{m+1}{k} B_k \cdot (n+1)^{m+1-k}$$

Euler-Maclaurin formula for infinite sums:

$$\begin{aligned} \sum_{i=m}^{\infty} f(i) &= \int_m^{\infty} f(x) dx - \sum_{k=1}^{\infty} \frac{B_k}{k!} f^{(k-1)}(m) \\ &\approx \int_m^{\infty} f(x) dx + \frac{f(m)}{2} - \frac{f'(m)}{12} + \frac{f'''(m)}{720} + O(f^{(5)}(m)) \end{aligned}$$

1.3.2 Stirling numbers of the first kind

Number of permutations on n items with k cycles.

$$\begin{aligned} c(n, k) &= c(n-1, k-1) + (n-1)c(n-1, k), \quad c(0, 0) = 1 \\ \sum_{k=0}^n c(n, k) x^k &= x(x+1) \dots (x+n-1) \\ c(8, k) &= 8, 0, 5040, 13068, 13132, 6769, 1960, 322, 28, 1 \\ c(n, 2) &= 0, 0, 1, 3, 11, 50, 274, 1764, 13068, 109584, \dots \end{aligned}$$

1.3.3 Eulerian numbers

Number of permutations $\pi \in S_n$ in which exactly k elements are greater than the previous element. k j :s s.t. $\pi(j) > \pi(j+1)$, $k+1$ j :s s.t. $\pi(j) \geq j$, k j :s s.t. $\pi(j) > j$.

$$E(n, k) = (n-k)E(n-1, k-1) + (k+1)E(n-1, k)$$

$$E(n, 0) = E(n, n-1) = 1$$

$$E(n, k) = \sum_{j=0}^k (-1)^j \binom{n+1}{j} (k+1-j)^n$$

1.3.4 Stirling numbers of the second kind

Partitions of n distinct elements into exactly k groups.

$$S(n, k) = S(n-1, k-1) + kS(n-1, k)$$

$$S(n, 1) = S(n, n) = 1$$

$$S(n, k) = \frac{1}{k!} \sum_{j=0}^k (-1)^{k-j} \binom{k}{j} j^n$$

1.3.5 Bell numbers

Total number of partitions of n distinct elements. $B(n) = 1, 1, 2, 5, 15, 52, 203, 877, 4140, 21147, \dots$. For p prime,

$$B(p^m + n) \equiv mB(n) + B(n+1) \pmod{p}$$

1.3.6 Labeled unrooted trees

on n vertices: n^{n-2}

on k existing trees of size n_i : $n_1 n_2 \dots n_k n^{k-2}$

with degrees d_i : $(n-2)! / ((d_1-1)! \dots (d_n-1)!)$

1.3.7 Catalan numbers

$$C_n = \frac{1}{n+1} \binom{2n}{n} = \binom{2n}{n} - \binom{2n}{n+1} = \frac{(2n)!}{(n+1)!n!}$$

$$C_0 = 1, C_{n+1} = \frac{2(2n+1)}{n+2} C_n, C_{n+1} = \sum C_i C_{n-i}$$

$$C_n = 1, 1, 2, 5, 14, 42, 132, 429, 1430, 4862, 16796, 58786, \dots$$

- sub-diagonal monotone paths in an $n \times n$ grid.
- strings with n pairs of parenthesis, correctly nested.
- binary trees with $n+1$ leaves (0 or 2 children).
- ordered trees with $n+1$ vertices.
- ways a convex polygon with $n+2$ sides can be cut into triangles by connecting vertices with straight lines.
- permutations of $[n]$ with no 3-term increasing subseq.

2 Data Structure

2.1 Gilbert Order

```
inline int64_t gilbertOrder(int x, int y, int pow = 21, int
    rotate = 0) {
    if (pow == 0) {
        return 0;
    }
    int hpow = 1 << (pow-1);
    int seg = (x < hpow) ? (
        (y < hpow) ? 0 : 3
    ) : (
        (y < hpow) ? 1 : 2
    );
    seg = (seg + rotate) & 3;
    const int rotateDelta[4] = {3, 0, 0, 1};
    int nx = x & (x ^ hpow), ny = y & (y ^ hpow);
    int nrot = (rotate + rotateDelta[seg]) & 3;
    int64_t subSquareSize = int64_t(1) << (2*pow - 2);
    int64_t ans = seg * subSquareSize;
    int64_t add = gilbertOrder(nx, ny, pow-1, nrot);
    ans += (seg == 1 || seg == 2) ? add : (subSquareSize - add
        - 1);
    return ans;
}
```

2.2 Li Chao Tree

```
const int N = 2e5 + 5;

struct line {
    ll a, b;
    ll val(ll x) { return a * x + b; }
};

line seg[4 * N];

void upd(int id, int l, int r, line lin) {
    if (l == r) {
        if (lin.val(l) > seg[id].val(l))
            seg[id] = lin;
        return;
    }
    int mid = (l + r) >> 1;
    if (seg[id].a > lin.a)
        swap(seg[id], lin);
    // sap xep lai do doc duong thang de tong quat hoa bai
    toan
    if (lin.val(mid) > seg[id].val(mid)) {
        swap(seg[id], lin);
        upd(id << 1, l, mid, lin);
    } else
        upd(id << 1 | 1, mid + 1, r, lin);
}

void upd_range(int id, int l, int r, int u, int v, line lin)
{
    if (l > v || r < u)
        return;
    if (l >= u && r <= v) {
        upd(id, l, r, lin);
        return;
    }
    int mid = (l + r) >> 1;
    upd_range(id << 1, l, mid, u, v, lin);
    upd_range(id << 1 | 1, mid + 1, r, u, v, lin);
}

ll get(int id, int l, int r, int pos) {
    if (l == pos && l == r) {
        return seg[id].val(pos);
    }
    int mid = (l + r) >> 1;
    if (mid >= pos)
        return max(seg[id].val(pos), get(id << 1, l, mid, pos
        ));
    return max(seg[id].val(pos), get(id << 1 | 1, mid + 1, r, pos
    ));
}
```

```
else
    return max(seg[id].val(pos), get(id << 1 | 1, mid +
        1, r, pos));
}
```

2.3 Persistent Segment Tree

```
struct Node{
    int mi;
    Node *l, *r;

    Node(int x){
        mi = x; l = nullptr; r = nullptr;
    }
    Node(Node *u, Node *v){
        l = u; r = v;
        mi = MAXA;
        if (l) mi = l -> mi;
        if (r) mi = min(mi, r -> mi);
    }
    Node(Node *u){
        mi = u -> mi;
        l = u -> l;
        r = u -> r;
    }
};

int a[N], last[N];
pair<int, int> r[N];
vector<int> v;
Node* root[N];

Node* build(int l, int r){
    if (l == r) return new Node(MAXA);

    int mid = (l + r) >> 1;
    return new Node(build(l, mid), build(mid + 1, r));
}

Node* upd(Node* node, int l, int r, int pos, int val){
    if (l == r) return new Node(val);

    int mid = (l + r) >> 1;
    if (pos > mid) return new Node(node -> l, upd(node -> r,
        mid + 1, r, pos, val));
    else return new Node(upd(node -> l, l, mid, pos, val),
        node -> r);
}
```

```

int get(Node* node, int l, int r, int u, int v){
    if (node == nullptr || l > v || r < u) return MAXA;
    if (l >= u && r <= v) return node -> mi;

    int mid = (l + r) >> 1;
    int v1 = get(node -> l, l, mid, u, v);
    int v2 = get(node -> r, mid + 1, r, u, v);
    return min(v1, v2);
}

```

2.4 Persistent Trie

```

/*
Usage: printf("%d\n", query(version[r], ~x, l));
-> Trie version[r] contains the trie for [1...r] elements
*/

```

```

struct node_t;
typedef node_t * pnode;

```

```

struct node_t {
    int time;
    pnode to[2];
    node_t() : time(0) {
        to[0] = to[1] = 0;
    }
    bool go(int l) const {
        if (!this) return false;
        return time >= l;
    }
    pnode clone() {
        pnode cur = new node_t();
        if (this) {
            cur->time = time;
            cur->to[0] = to[0];
            cur->to[1] = to[1];
        }
        return cur;
    }
};

pnode last;
pnode version[N];

void insert(int a, int time) {
    pnode v = version[time] = last = last->clone();
    for (int i = K - 1; i >= 0; --i) {
        int bit = (a >> i) & 1;
        pnode &child = v->to[bit];

```

```

        child = child->clone();
        v = child;
        v->time = time;
    }
}

int query(pnode v, int x, int l) {
    int ans = 0;
    for (int i = K - 1; i >= 0; --i) {
        int bit = (x >> i) & 1;
        if (v->to[bit]->go(l)) { // checking if this bit was
            inserted before the range
            ans |= 1 << i;
            v = v->to[bit];
        } else {
            v = v->to[bit ^ 1];
        }
    }
    return ans;
}

```

2.5 Segment Tree 2D

```

void build_y(int vx, int lx, int rx, int vy, int ly, int ry)
{
    if (ly == ry) {
        if (lx == rx)
            segtree[vx][vy] = 0;
        else
            segtree[vx][vy] = max(segtree[vx*2][vy], segtree[
                vx*2+1][vy]);
    } else {
        int my = (ly + ry) / 2;
        build_y(vx, lx, rx, vy*2, ly, my);
        build_y(vx, lx, rx, vy*2+1, my+1, ry);
        segtree[vx][vy] = max(segtree[vx][vy*2], segtree[vx][
            vy*2+1]);
    }
}

void build_x(int vx, int lx, int rx) {
    if (lx != rx) {
        int mx = (lx + rx) / 2;
        build_x(vx*2, lx, mx);
        build_x(vx*2+1, mx+1, rx);
    }
    build_y(vx, lx, rx, 1, 0, m - 1);
}

```

```

int query_y(int vx, int vy, int tly, int try_, int ly, int
    ry) {
    if (ly > ry)
        return 0;
    if (ly == tly && try_ == ry)
        return segtree[vx][vy];
    int tmy = (tly + try_) / 2;
    return max(query_y(vx, vy*2, tly, tmy, ly, min(ry, tmy))
        , query_y(vx, vy*2+1, tmy+1, try_, max(ly, tmy+1), ry)
        );
}

int query_x(int vx, int tlx, int trx, int lx, int rx, int ly
    , int ry) {
    if (lx > rx)
        return 0;
    if (lx == tlx && trx == rx)
        return query_y(vx, 1, 0, m-1, ly, ry);
    int tmx = (tlx + trx) / 2;
    return max(query_x(vx*2, tlx, tmx, lx, min(rx, tmx), ly,
        ry)
        , query_x(vx*2+1, tmx+1, trx, max(lx, tmx+1), rx, ly,
        ry));
}

void update_y(int vx, int lx, int rx, int vy, int ly, int ry
    , int x, int y, int new_val) {
    if (ly == ry) {
        if (lx == rx)
            segtree[vx][vy] = new_val;
        else
            segtree[vx][vy] = max(segtree[vx*2][vy], segtree[
                vx*2+1][vy]);
    } else {
        int my = (ly + ry) / 2;
        if (y <= my)
            update_y(vx, lx, rx, vy*2, ly, my, x, y, new_val)
                ;
        else
            update_y(vx, lx, rx, vy*2+1, my+1, ry, x, y,
                new_val);
        segtree[vx][vy] = max(segtree[vx][vy*2], segtree[vx][
            vy*2+1]);
    }
}

void update_x(int vx, int lx, int rx, int x, int y, int
    new_val) {
    if (lx != rx) {
        int mx = (lx + rx) / 2;
        if (x <= mx)

```

```

        update_x(vx*2, lx, mx, x, y, new_val);
    else
        update_x(vx*2+1, mx+1, rx, x, y, new_val);
}
update_y(vx, lx, rx, 1, 0, m-1, x, y, new_val);
}

```

2.6 STL

```

#include <bits/stdc++.h>

#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
#include <ext/rope>

using namespace __gnu_pbds;
using namespace __gnu_cxx;
using namespace std;

typedef tree<int, null_type, less<int>, rb_tree_tag,
            tree_order_statistics_node_update> ordered_set;

unsigned hash_f(unsigned x) {
    x = ((x >> 16) ^ x) * 0x45d9f3b;
    x = ((x >> 16) ^ x) * 0x45d9f3b;
    x = (x >> 16) ^ x;
    return x;
}

struct chash {
    int operator()(int x) const { return hash_f(x); }
};

ordered_set s; //ordered_set
//s.find_by_order(x)
//s.order_of_key(x)
gp_hash_table<int, int, chash> mp; //hash map
rope <int> v; //rope (almost like string...)

```

2.7 Treap

```

struct node{
    node *l, *r;
    int sz, pri;
    int val, sum;
    bool rev;

    node(){

```

```

        node(int c){
            l = r = nullptr;
            sz = 1;
            rev = false;
            pri = rng();
            val = c;
            sum = c;
        }
};

void push(node *&treap){
    if (treap == nullptr) return;
    if (!(treap -> rev)) return;
    swap(treap -> l, treap -> r);
    if (treap -> l){
        (treap -> l) -> rev ^= 1;
    }
    if (treap -> r){
        (treap -> r) -> rev ^= 1;
    }
    treap -> rev = false;
}

int get_sz(node *treap){
    if (treap == nullptr) return 0;
    else return treap -> sz;
}

int get_sum(node *treap){
    if (treap == nullptr) return 0;
    else return treap -> sum;
}

void split(node *treap, node *&l, node *&r, int k){ //[1, k]
    [k + 1, sz]
    if (treap == nullptr){
        l = r = nullptr;
        return;
    }
    push(treap);
    if (get_sz(treap -> l) < k){
        split(treap -> r, treap -> r, r, k - get_sz(treap -> l) - 1);
        l = treap;
    } else {
        split(treap -> l, l, treap -> l, k);
        r = treap;
    }
    treap -> sz = get_sz(treap -> l) + get_sz(treap -> r) + 1;
}

```

```

    treap -> sum = get_sum(treap -> l) + get_sum(treap -> r) + treap -> val;
}

void merge(node *&treap, node *l, node *r){
    if (l == nullptr){
        treap = r;
        return;
    }
    if (r == nullptr){
        treap = l;
        return;
    }
    push(l); push(r);
    if (l -> pri < r -> pri){
        merge(l -> r, l -> r, r);
        treap = l;
    } else {
        merge(r -> l, l, r -> l);
        treap = r;
    }
    treap -> sz = get_sz(treap -> l) + get_sz(treap -> r) + 1;
    treap -> sum = get_sum(treap -> l) + get_sum(treap -> r) + treap -> val;
}

void print(node *treap){
    if (treap == nullptr) return;
    push(treap);
    print(treap -> l);
    cout << treap -> val;
    print(treap -> r);
}

```

3 Dynamic Programming

3.1 DP Convex Hull Trick

```

// Decreasing Insertion, Query Min
struct CHT {
    vector<long long> a, b;

    bool cross(int i, int j, int k) {
        return 1.d*(a[j] - a[i])*(b[k] - b[i]) >= 1.d*(a[k] - a[i])*(b[j] - b[i]);
    }
}

```

```

void add(long long A, long long B) {
    a.push_back(A);
    b.push_back(B);

    while (a.size() > 2 && cross(a.size() - 3, a.size() - 2, a.size() - 1)) {
        a.erase(a.end() - 2);
        b.erase(b.end() - 2);
    }

    long long query(long long x) {
        int l = 0, r = a.size() - 1;

        while (l < r) {
            int mid = l + (r - l)/2;
            long long f1 = a[mid] * x + b[mid];
            long long f2 = a[mid + 1] * x + b[mid + 1];

            if (f1 > f2) l = mid + 1;
            else r = mid;
        }

        return a[l]*x + b[l];
    }
};

```

3.2 DP Divide and Conquer

```

void divide(int i, int L, int R, int optL, int optR) {
    if (L > R) return;
    int mid = (L+R) / 2, cut = optL;
    f[i][mid] = INF;
    for (int k = optL; k <= min(mid, optR); k++) {
        long long cur = f[i - 1][k] + Cost(k+1, mid);
        if (f[i][mid] > cur) {
            f[i][mid] = cur;
            cut = k;
        }
    }
    divide(i, L, mid - 1, optL, cut);
    divide(i, mid + 1, R, cut, optR);
}

```

3.3 DP Knuth

```
/*
```

```

* Complexity: O(N^2)
* f[i][j] = min(f[i][k] + f[k][j] + c[i][j], i < k < j)
* a[i][j] = min(k | i < k < j && f[i][j] = f[i][k] + f[k][j] + c[i][j])
* Sufficient condition: a[i][j - 1] <= a[i][j] <= a[i + 1][j] or
* c[x][z] + c[y][t] <= c[x][t] + c[y][z] (quadrangle inequality) and c[y][z] <= c[x][t] (monotonicity), x <= y <= z <= t
*/

void knuth() {
    for (int i = 1; i <= n; i++) {
        f[i][i] = 0;
        a[i][i] = i;
    }
    for (int len = 1; len <= n-1; len++)
        for (int i = 1; i <= n-len; i++)
            {
                int j = i + len;
                f[i][j] = INF;
                for (int k = a[i][j - 1]; k <= a[i + 1][j]; k++)
                {
                    if (f[i][j] > f[i][k-1] + f[k][j] + c[i][j])
                    {
                        f[i][j] = f[i][k-1] + f[k][j] + c[i][j];
                        a[i][j] = k;
                    }
                }
            }
    cout << f[1][n] << '\n';
}

```

3.4 Subset Sum Convolution

```

// Make fhat[][] = {0} and ghat[][] = {0}
for (int mask = 0; mask < (1 << N); mask++) {
    fhat[__builtin_popcount(mask)][mask] = f[mask];
    ghat[__builtin_popcount(mask)][mask] = g[mask];
}

// Apply zeta transform on fhat[][] and ghat[][]
for (int i = 0; i < N; i++) {
    for (int j = 0; j < N; j++) {
        for (int mask = 0; mask < (1 << N); mask++) {
            if ((mask & (1 << j)) != 0) {
                fhat[i][mask] += fhat[i][mask ^ (1 << j)];
                ghat[i][mask] += ghat[i][mask ^ (1 << j)];
            }
        }
    }
}

```

```

}
}
}

// Do the convolution and store into h[][] = {0}
for (int mask = 0; mask < (1 << N); mask++) {
    for (int i = 0; i < N; i++) {
        for (int j = 0; j <= i; j++) {
            h[i][mask] += fhat[j][mask] * ghat[i - j][mask];
        }
    }
}

// Apply inverse SOS dp on h[][]
for (int i = 0; i < N; i++) {
    for (int j = 0; j < N; j++) {
        for (int mask = 0; mask < (1 << N); mask++) {
            if ((mask & (1 << j)) != 0) {
                h[i][mask] -= h[i][mask ^ (1 << j)];
            }
        }
    }
}

for (int mask = 0; mask < (1 << N); mask++) fog[mask] = h[
    __builtin_popcount(mask)][mask];

```

4 Formula

4.1 Equations

$$ax^2 + bx + c = 0 \Rightarrow x = \frac{-b \pm \sqrt{b^2 - 4ac}}{2a}$$

The extremum is given by $x = -b/2a$.

$$\begin{aligned} ax + by = e &\Rightarrow \frac{ed - bf}{ad - bc} \\ cx + dy = f &\Rightarrow \frac{af - ec}{ad - bc} \end{aligned}$$

In general, given an equation $Ax = b$, the solution to a variable x_i is given by

$$x_i = \frac{\det A'_i}{\det A}$$

where A'_i is A with the i 'th column replaced by b .

4.2 Recurrences

If $a_n = c_1 a_{n-1} + \dots + c_k a_{n-k}$, and r_1, \dots, r_k are distinct roots of $x^k + c_1 x^{k-1} + \dots + c_k$, there are d_1, \dots, d_k s.t.

$$a_n = d_1 r_1^n + \dots + d_k r_k^n.$$

Non-distinct roots r become polynomial factors, e.g. $a_n = (d_1 n + d_2) r^n$.

4.3 Trigonometry

$$\sin(v+w) = \sin v \cos w + \cos v \sin w$$

$$\cos(v+w) = \cos v \cos w - \sin v \sin w$$

$$\begin{aligned} \tan(v+w) &= \frac{\tan v + \tan w}{1 - \tan v \tan w} \\ \sin v + \sin w &= 2 \sin \frac{v+w}{2} \cos \frac{v-w}{2} \\ \cos v + \cos w &= 2 \cos \frac{v+w}{2} \cos \frac{v-w}{2} \end{aligned}$$

$$(V+W) \tan(v-w)/2 = (V-W) \tan(v+w)/2$$

where V, W are lengths of sides opposite angles v, w .

$$a \cos x + b \sin x = r \cos(x - \phi)$$

$$a \sin x + b \cos x = r \sin(x + \phi)$$

where $r = \sqrt{a^2 + b^2}$, $\phi = \text{atan2}(b, a)$.

4.4 Geometry

4.4.1 Triangles

Side lengths: a, b, c

Semiperimeter: $p = \frac{a+b+c}{2}$

Area: $A = \sqrt{p(p-a)(p-b)(p-c)}$

Circumradius: $R = \frac{abc}{4A}$

Inradius: $r = \frac{A}{p}$

Length of median (divides triangle into two equal-area triangles): $m_a = \frac{1}{2} \sqrt{2b^2 + 2c^2 - a^2}$

Length of bisector (divides angles in two): $s_a =$

$$\sqrt{bc \left[1 - \left(\frac{a}{b+c} \right)^2 \right]}$$

Law of sines: $\frac{\sin \alpha}{a} = \frac{\sin \beta}{b} = \frac{\sin \gamma}{c} = \frac{1}{2R}$

Law of cosines: $a^2 = b^2 + c^2 - 2bc \cos \alpha$

Law of tangents: $\frac{a+b}{a-b} = \frac{\tan \frac{\alpha+\beta}{2}}{\tan \frac{\alpha-\beta}{2}}$

4.4.2 Quadrilaterals

With side lengths a, b, c, d , diagonals e, f , diagonals angle θ , area A and magic flux $F = b^2 + d^2 - a^2 - c^2$:

$$4A = 2ef \cdot \sin \theta = F \tan \theta = \sqrt{4e^2 f^2 - F^2}$$

For cyclic quadrilaterals the sum of opposite angles is 180° , $ef = ac + bd$, and $A = \sqrt{(p-a)(p-b)(p-c)(p-d)}$.

4.4.3 Spherical coordinates

$$\begin{aligned} x &= r \sin \theta \cos \phi & r &= \sqrt{x^2 + y^2 + z^2} \\ y &= r \sin \theta \sin \phi & \theta &= \arccos(z/\sqrt{x^2 + y^2 + z^2}) \\ z &= r \cos \theta & \phi &= \text{atan2}(y, x) \end{aligned}$$

4.5 Sums

$$c^a + c^{a+1} + \dots + c^b = \frac{c^{b+1} - c^a}{c - 1}, c \neq 1$$

4.6 Probability theory

Let X be a discrete random variable with probability $p_X(x)$ of assuming the value x . It will then have an expected value (mean) $\mu = E(X) = \sum_x x p_X(x)$ and variance $\sigma^2 =$

$V(X) = E(X^2) - (E(X))^2 = \sum_x (x - E(X))^2 p_X(x)$ where σ is the standard deviation. If X is instead continuous it will have a probability density function $f_X(x)$ and the sums above will instead be integrals with $p_X(x)$ replaced by $f_X(x)$.

Expectation is linear:

$$E(aX + bY) = aE(X) + bE(Y)$$

For independent X and Y ,

$$V(aX + bY) = a^2 V(X) + b^2 V(Y).$$

4.6.1 Discrete distributions

Binomial distribution The number of successes in n independent yes/no experiments, each which yields success with probability p is $\text{Bin}(n, p)$, $n = 1, 2, \dots$, $0 \leq p \leq 1$.

$$p(k) = \binom{n}{k} p^k (1-p)^{n-k}$$

$$\mu = np, \sigma^2 = np(1-p)$$

$\text{Bin}(n, p)$ is approximately $\text{Po}(np)$ for small p .

First success distribution The number of trials needed to get the first success in independent yes/no experiments, each which yields success with probability p is $\text{Fs}(p)$, $0 \leq p \leq 1$.

$$p(k) = p(1-p)^{k-1}, k = 1, 2, \dots$$

$$\mu = \frac{1}{p}, \sigma^2 = \frac{1-p}{p^2}$$

Poisson distribution The number of events occurring in a fixed period of time t if these events occur with a known average rate κ and independently of the time since the last event is $\text{Po}(\lambda)$, $\lambda = t\kappa$.

$$p(k) = e^{-\lambda} \frac{\lambda^k}{k!}, k = 0, 1, 2, \dots$$

$$\mu = \lambda, \sigma^2 = \lambda$$

5 Geometry

5.1 Angle

```
/**
 * Description: A class for ordering angles (as represented
 *             by int points and
 *             a number of rotations around the origin). Useful for
 *             rotational sweeping.
 *             Sometimes also represents points or vectors.
 * Usage:
 * vector<Angle> v = {w[0], w[0].t360() ...}; // sorted
 * int j = 0; rep(i,0,n) { while (v[j] < v[i].t180()) ++j; }
 * // sweeps j such that (j-i) represents the number of
 *           positively oriented triangles with vertices at 0 and i
 * Status: Used, works well
 */
#pragma once

struct Angle {
    int x, y;
    int t;
    Angle(int x, int y, int t=0) : x(x), y(y), t(t) {}
    Angle operator-(Angle b) const { return {x-b.x, y-b.y, t}; }
    int half() const {
        assert(x || y);
        return y < 0 || (y == 0 && x < 0);
    }
    Angle t90() const { return {-y, x, t + (half() && x >= 0)}; }
    Angle t180() const { return {-x, -y, t + half()}; }
    Angle t360() const { return {x, y, t + 1}; }
};

bool operator<(Angle a, Angle b) {
    // add a.dist2() and b.dist2() to also compare distances
    return make_tuple(a.t, a.half(), a.y * (11)b.x <
        make_tuple(b.t, b.half(), a.x * (11)b.y);
}

// Given two points, this calculates the smallest angle
// between
// them, i.e., the angle that covers the defined line
// segment.
pair<Angle, Angle> segmentAngles(Angle a, Angle b) {
    if (b < a) swap(a, b);
    return (b < a.t180() ?
        make_pair(a, b) : make_pair(b, a.t360()));
}

Angle operator+(Angle a, Angle b) { // point a + vector b
```

```
Angle r(a.x + b.x, a.y + b.y, a.t);
if (a.t180() < r) r.t--;
return r.t180() < a ? r.t360() : r;
}

Angle angleDiff(Angle a, Angle b) { // angle b - angle a
    int tu = b.t - a.t; a.t = b.t;
    return {a.x*b.x + a.y*b.y, a.x*b.y - a.y*b.x, tu - (b < a)
        };
}
```

5.2 CircleIntersection

```
/**
 * Description: Computes the pair of points at which two
 *             circles intersect. Returns false in case of no
 *             intersection.
 * Status: stress-tested
 */
#pragma once

#include "Point.h"

typedef Point<double> P;
bool circleInter(P a, P b, double r1, double r2, pair<P, P>* out
    ) {
    if (a == b) { assert(r1 != r2); return false; }
    P vec = b - a;
    double d2 = vec.dist2(), sum = r1+r2, dif = r1-r2,
        p = (d2 + r1*r1 - r2*r2)/(d2*2), h2 = r1*r1 - p*p*d2;
    if (sum*sum < d2 || dif*dif > d2) return false;
    P mid = a + vec*p, per = vec.perp() * sqrt(fmax(0, h2) / d2
        );
    *out = {mid + per, mid - per};
    return true;
}
```

5.3 CircleLine

```
/**
 * Description: Finds the intersection between a circle and
 *             a line.
 * Returns a vector of either 0, 1, or 2 intersection points
 * .
 * P is intended to be Point<double>.
 */
#pragma once
```

```
#include "Point.h"

template<class P>
vector<P> circleLine(P c, double r, P a, P b) {
    P ab = b - a, p = a + ab * (c-a).dot(ab) / ab.dist2();
    double s = a.cross(b, c), h2 = r*r - s*s / ab.dist2();
    if (h2 < 0) return {};
    if (h2 == 0) return {p};
    P h = ab.unit() * sqrt(h2);
    return {p - h, p + h};
}
```

5.4 CirclePolygonIntersection

```
/**
 * Description: Returns the area of the intersection of a
 *             circle with a
 *             ccw polygon.
 * Time: O(n)
 * Status: Tested on GNYR 2019 Gerrymandering, stress-tested
 */
#pragma once

#include "../content/geometry/Point.h"

typedef Point<double> P;
#define arg(p, q) atan2(p.cross(q), p.dot(q))
double circlePoly(P c, double r, vector<P> ps) {
    auto tri = [&](P p, P q) {
        auto r2 = r * r / 2;
        P d = q - p;
        auto a = d.dot(p)/d.dist2(), b = (p.dist2()-r*r)/d.dist2()
            ;
        auto det = a * a - b;
        if (det <= 0) return arg(p, q) * r2;
        auto s = max(0., -a-sqrt(det)), t = min(1., -a+sqrt(det));
        if (t < 0 || 1 <= s) return arg(p, q) * r2;
        P u = p + d * s, v = p + d * t;
        return arg(p,u) * r2 + u.cross(v)/2 + arg(v,q) * r2;
    };
    auto sum = 0.0;
    rep(i,0,sz(ps))
        sum += tri(ps[i] - c, ps[(i + 1) % sz(ps)] - c);
    return sum;
}
```

5.5 CircleTangents

```
/**
 * Description: Finds the external tangents of two circles,
 * or internal if r2 is negated.
 * Can return 0, 1, or 2 tangents -- 0 if one circle
 * contains the other (or overlaps it, in the internal
 * case, or if the circles are the same);
 * 1 if the circles are tangent to each other (in which case
 * .first = .second and the tangent line is perpendicular
 * to the line between the centers).
 * .first and .second give the tangency points at circle 1
 * and 2 respectively.
 * To find the tangents of a circle with a point set r2 to
 * 0.
 * Status: tested
 */
#pragma once

#include "Point.h"

template<class P>
vector<pair<P, P>> tangents(P c1, double r1, P c2, double r2
) {
    P d = c2 - c1;
    double dr = r1 - r2, d2 = d.dist2(), h2 = d2 - dr * dr;
    if (d2 == 0 || h2 < 0) return {};
    vector<pair<P, P>> out;
    for (double sign : {-1, 1}) {
        P v = (d * dr + d.perp() * sqrt(h2) * sign) / d2;
        out.push_back({c1 + v * r1, c2 + v * r2});
    }
    if (h2 == 0) out.pop_back();
    return out;
}
```

5.6 circumcircle

```
/**
 * Description:\\
 * The circumcircle of a triangle is the circle intersecting
 * all three vertices. ccRadius returns the radius of the
 * circle going through points A, B and C and ccCenter
 * returns the center of the same circle.
 *
 * Status: tested
 */
#pragma once
```

```
#include "Point.h"

typedef Point<double> P;
double ccRadius(const P& A, const P& B, const P& C) {
    return (B-A).dist()*(C-B).dist()*(A-C).dist()/
        abs((B-A).cross(C-A))/2;
}
P ccCenter(const P& A, const P& B, const P& C) {
    P b = C-A, c = B-A;
    return A + (b*c.dist2()-c*b.dist2()).perp()/b.cross(c)/2;
}
```

5.7 ClosestPair

```
/**
 * Source: https://codeforces.com/blog/entry/58747
 * Description: Finds the closest pair of points.
 */
#pragma once

#include "Point.h"

typedef Point<ll> P;
pair<P, P> closest(vector<P> v) {
    assert(sz(v) > 1);
    set<P> S;
    sort(all(v), [](P a, P b) { return a.y < b.y; });
    pair<ll, pair<P, P>> ret{LLONG_MAX, {P(), P()}};
    int j = 0;
    for (P p : v) {
        P d{1 + (ll)sqrt(ret.first), 0};
        while (v[j].y <= p.y - d.x) S.erase(v[j++]);
        auto lo = S.lower_bound(p - d), hi = S.upper_bound(p + d);
        for (; lo != hi; ++lo)
            ret = min(ret, {(*lo - p).dist2(), {*lo, p}});
        S.insert(p);
    }
    return ret.second;
}
```

5.8 ConvexHull

```
/**
Returns a vector of the points of the convex hull in counter
-clockwise order.
Points on the edge of the hull between two other points are
not considered part of the hull.
```

```
*/
#pragma once

#include "Point.h"

typedef Point<ll> P;
vector<P> convexHull(vector<P> pts) {
    if (sz(pts) <= 1) return pts;
    sort(all(pts));
    vector<P> h(sz(pts)+1);
    int s = 0, t = 0;
    for (int it = 2; it--; s = --t, reverse(all(pts)))
        for (P p : pts) {
            while (t >= s + 2 && h[t-2].cross(h[t-1], p) <= 0) t--;
            h[t++] = p;
        }
    return {h.begin(), h.begin() + t - (t == 2 && h[0] == h[1])
        };
}
```

5.9 DelaunayTriangulation

```
/**
 * Description: Computes the Delaunay triangulation of a set
 * of points.
 * Each circumcircle contains none of the input points.
 * If any three points are collinear or any four are on the
 * same circle, behavior is undefined.
 * Time: O(n^2)
 * Status: stress-tested
 */
#pragma once

#include "Point.h"
#include "3dHull.h"

template<class P, class F>
void delaunay(vector<P>& ps, F trifun) {
    if (sz(ps) == 3) { int d = (ps[0].cross(ps[1], ps[2]) < 0);
        trifun(0,1+d,2-d); }
    vector<P> p3;
    for (P p : ps) p3.emplace_back(p.x, p.y, p.dist2());
    if (sz(ps) > 3) for(auto t:hull3d(p3)) if ((p3[t.b]-p3[t.a
        ])).
        cross(p3[t.c]-p3[t.a]).dot(P3(0,0,1)) < 0)
        trifun(t.a, t.c, t.b);
}
```

5.10 FastDelaunay

```
/**
 * Description: Fast Delaunay triangulation.
 * Each circumcircle contains none of the input points.
 * There must be no duplicate points.
 * If all points are on a line, no triangles will be
 *   returned.
 * Should work for doubles as well, though there may be
 *   precision issues in 'circ'.
 * Returns triangles in order \{t[0][0], t[0][1], t[0][2], t
 *   [1][0], \dots\}, all counter-clockwise.
 * Time: O(n \log n)
 */
#pragma once

#include "Point.h"

typedef Point<ll> P;
typedef struct Quad* Q;
typedef __int128_t lll; // (can be ll if coords are < 2e4)
P arb(LLONG_MAX, LLONG_MAX); // not equal to any other point

struct Quad {
    Q rot, o; P p = arb; bool mark;
    P& F() { return r()->p; }
    Q& r() { return rot->rot; }
    Q prev() { return rot->o->rot; }
    Q next() { return r()->prev(); }
} *H;

bool circ(P p, P a, P b, P c) { // is p in the circumcircle?
    lll p2 = p.dist2(), A = a.dist2()-p2,
        B = b.dist2()-p2, C = c.dist2()-p2;
    return p.cross(a,b)*C + p.cross(b,c)*A + p.cross(c,a)*B >
        0;
}

Q makeEdge(P orig, P dest) {
    Q r = H ? H : new Quad{new Quad{new Quad{new Quad{0}}}};
    H = r->o; r->r()->r() = r;
    rep(i,0,4) r = r->rot, r->p = arb, r->o = i & 1 ? r : r->r
        ();
    r->p = orig; r->F() = dest;
    return r;
}

void splice(Q a, Q b) {
    swap(a->o->rot->o, b->o->rot->o); swap(a->o, b->o);
}

Q connect(Q a, Q b) {
    Q q = makeEdge(a->F(), b->p);
```

```
splice(q, a->next());
splice(q->r(), b);
return q;
}

pair<Q,Q> rec(const vector<P>& s) {
    if (sz(s) <= 3) {
        Q a = makeEdge(s[0], s[1]), b = makeEdge(s[1], s.back());
        if (sz(s) == 2) return { a, a->r() };
        splice(a->r(), b);
        auto side = s[0].cross(s[1], s[2]);
        Q c = side ? connect(b, a) : 0;
        return {side < 0 ? c->r() : a, side < 0 ? c : b->r() };
    }

#define H(e) e->F(), e->p
#define valid(e) (e->F().cross(H(base)) > 0)
    Q A, B, ra, rb;
    int half = sz(s) / 2;
    tie(ra, A) = rec({all(s) - half});
    tie(B, rb) = rec({sz(s) - half + all(s)});
    while ((B->p.cross(H(A)) < 0 && (A = A->next())) ||
        (A->p.cross(H(B)) > 0 && (B = B->r()->o)));
    Q base = connect(B->r(), A);
    if (A->p == ra->p) ra = base->r();
    if (B->p == rb->p) rb = base;

#define DEL(e, init, dir) Q e = init->dir; if (valid(e)) \
    while (circ(e->dir->F(), H(base), e->F())) { \
        Q t = e->dir; \
        splice(e, e->prev()); \
        splice(e->r(), e->r()->prev()); \
        e->o = H; H = e; e = t; \
    }
    for (;;) {
        DEL(LC, base->r(), o); DEL(RC, base, prev());
        if (!valid(LC) && !valid(RC)) break;
        if (!valid(LC) || (valid(RC) && circ(H(RC), H(LC))))
            base = connect(RC, base->r());
        else
            base = connect(base->r(), LC->r());
    }
    return { ra, rb };
}

vector<P> triangulate(vector<P> pts) {
    sort(all(pts)); assert(unique(all(pts)) == pts.end());
    if (sz(pts) < 2) return {};
    Q e = rec(pts).first;
    vector<Q> q = {e};
```

```
int qi = 0;
while (e->o->F().cross(e->F(), e->p) < 0) e = e->o;
#define ADD { Q c = e; do { c->mark = 1; pts.push_back(c->p)
    ; \
    q.push_back(c->r()); c = c->next(); } while (c != e); }
ADD; pts.clear();
while (qi < sz(q)) if (!(e = q[qi++])->mark) ADD;
return pts;
}
```

5.11 HullDiameter

```
/**
 * Description: Returns the two points with max distance on
 *   a convex hull (ccw,
 *   no duplicate/collinear points).
 */
#pragma once
#include "Point.h"

typedef Point<ll> P;
array<P, 2> hullDiameter(vector<P> S) {
    int n = sz(S), j = n < 2 ? 0 : 1;
    pair<ll, array<P, 2>> res({0, {S[0], S[0]}});
    rep(i,0,j)
        for (; j = (j + 1) % n) {
            res = max(res, {(S[i] - S[j]).dist2(), {S[i], S[j]}});
            if ((S[(j + 1) % n] - S[j]).cross(S[i + 1] - S[i]) >= 0)
                break;
        }
    return res.second;
}
```

5.12 InsidePolygon

```
/**
 * Description: Returns true if p lies within the polygon.
 *   If strict is true,
 *   it returns false for points on the boundary. The
 *   algorithm uses
 *   products in intermediate steps so watch out for overflow.
 * Time: O(n)
 * Usage:
 * vector<P> v = {P{4,4}, P{1,2}, P{2,1}};
 * bool in = inPolygon(v, P{3, 3}, false);
 */
#pragma once
```

```
#include "Point.h"
#include "OnSegment.h"
#include "SegmentDistance.h"

template<class P>
bool inPolygon(vector<P> &p, P a, bool strict = true) {
    int cnt = 0, n = sz(p);
    rep(i,0,n) {
        P q = p[(i + 1) % n];
        if (onSegment(p[i], q, a)) return !strict;
        //or: if (segDist(p[i], q, a) <= eps) return !strict;
        cnt ^= ((a.y<p[i].y) - (a.y<q.y)) * a.cross(p[i], q) > 0;
    }
    return cnt;
}
```

5.13 linearTransformation

```
/**
 * Apply the linear transformation (translation, rotation and
 * scaling) which takes line p0-p1 to line q0-q1 to point
 * r.
 */
#pragma once

#include "Point.h"

typedef Point<double> P;
P linearTransformation(const P& p0, const P& p1,
    const P& q0, const P& q1, const P& r) {
    P dp = p1-p0, dq = q1-q0, num(dp.cross(dq), dp.dot(dq));
    return q0 + P((r-p0).cross(num), (r-p0).dot(num))/dp.dist2
        ();
}
```

5.14 lineDistance

```
/**
 * Returns the signed distance between point p and the line
 * containing points a and b. Positive value on left side
 * and negative on right as seen from a towards b. a==b
 * gives nan. P is supposed to be Point<T> or Point3D<T>
 * where T is e.g. double or long long. It uses products
 * in intermediate steps so watch out for overflow if
 * using int or long long. Using Point3D will always give
```

```
a non-negative distance. For Point3D, call .dist on the
result of the cross product.
 * Status: tested
 */
#pragma once

#include "Point.h"

template<class P>
double lineDist(const P& a, const P& b, const P& p) {
    return (double)(b-a).cross(p-a)/(b-a).dist();
}
```

5.15 lineIntersection

```
/**
 * If a unique intersection point of the lines going through s1
 * ,e1 and s2,e2 exists \{1, point\} is returned.
 * If no intersection point exists \{0, (0,0)\} is returned and
 * if infinitely many exists \{-1, (0,0)\} is returned.
 * The wrong position will be returned if P is Point<ll> and
 * the intersection point does not have integer
 * coordinates.
 * Products of three coordinates are used in intermediate steps
 * so watch out for overflow if using int or ll.
 * Usage:
 * auto res = lineInter(s1,e1,s2,e2);
 * if (res.first == 1)
 *     cout << "intersection point at " << res.second << endl;
 * Status: stress-tested, and tested through half-plane
 * tests
 */
#pragma once

#include "Point.h"

template<class P>
pair<int, P> lineInter(P s1, P e1, P s2, P e2) {
    auto d = (e1 - s1).cross(e2 - s2);
    if (d == 0) // if parallel
        return {-(s1.cross(e1, s2) == 0), P(0, 0)};
    auto p = s2.cross(e1, e2), q = s2.cross(e2, s1);
    return {1, (s1 * p + e1 * q) / d};
}
```

5.16 LineProjectionReflection

```
/**
 * Description: Projects point p onto line ab. Set refl=true
 * to get reflection
 * of point p across line ab insted. The wrong point will be
 * returned if P is
 * an integer point and the desired point doesn't have
 * integer coordinates.
 * Products of three coordinates are used in intermediate
 * steps so watch out
 * for overflow.
 */
#pragma once

#include "Point.h"

template<class P>
P lineProj(P a, P b, P p, bool refl=false) {
    P v = b - a;
    return p - v.perp()*(1+refl)*v.cross(p-a)/v.dist2();
}
```

5.17 ManhattanMST

```
/**
 * Description: Given N points, returns up to 4*N edges,
 * which are guaranteed
 * to contain a minimum spanning tree for the graph with
 * edge weights w(p, q) =
 * |p.x - q.x| + |p.y - q.y|. Edges are in the form (
 * distance, src, dst). Use a
 * standard MST algorithm on the result to find the final
 * MST.
 * Time: O(N \log N)
 */
#pragma once
#include "Point.h"

typedef Point<int> P;
vector<array<int, 3>> manhattanMST(vector<P> ps) {
    vi id(sz(ps));
    iota(all(id), 0);
    vector<array<int, 3>> edges;
    rep(k,0,4) {
        sort(all(id), [&](int i, int j) {
            return (ps[i]-ps[j]).x < (ps[j]-ps[i]).y;});
        map<int, int> sweep;
        for (int i : id) {
            for (auto it = sweep.lower_bound(-ps[i].y);
```

```

        it != sweep.end(); sweep.erase(it++)) {
    int j = it->second;
    P d = ps[i] - ps[j];
    if (d.y > d.x) break;
    edges.push_back({d.y + d.x, i, j});
}
sweep[-ps[i].y] = i;
}
for (P& p : ps) if (k & 1) p.x = -p.x; else swap(p.x, p.y)
;
}
return edges;
}

```

5.18 MinimumEnclosingCircle

```

/**
 * Description: Computes the minimum circle that encloses a
 * set of points.
 * Time: expected O(n)
 * Status: stress-tested
 */
#pragma once

#include "circumcircle.h"

pair<P, double> mec(vector<P> ps) {
    shuffle(all(ps), mt19937(time(0)));
    P o = ps[0];
    double r = 0, EPS = 1 + 1e-8;
    rep(i,0,sz(ps)) if ((o - ps[i]).dist() > r * EPS) {
        o = ps[i], r = 0;
        rep(j,0,i) if ((o - ps[j]).dist() > r * EPS) {
            o = (ps[i] + ps[j]) / 2;
            r = (o - ps[i]).dist();
            rep(k,0,j) if ((o - ps[k]).dist() > r * EPS) {
                o = ccCenter(ps[i], ps[j], ps[k]);
                r = (o - ps[i]).dist();
            }
        }
    }
    return {o, r};
}

```

5.19 OnSegment

```

/**

```

```

 * Description: Returns true iff p lies on the line segment
 * from s to e.
 * Use \texttt{(segDist(s,e,p)<=epsilon)} instead when using
 * Point<double>.
 * Status:
 */
#pragma once

#include "Point.h"

template<class P> bool onSegment(P s, P e, P p) {
    return p.cross(s, e) == 0 && (s - p).dot(e - p) <= 0;
}

```

5.20 Point

```

/**
 * Description: Class to handle points in the plane.
 * T can be e.g. double or long long. (Avoid int.)
 * Status: Works fine, used a lot
 */
#pragma once

template <class T> int sgn(T x) { return (x > 0) - (x < 0); }

template<class T>
struct Point {
    typedef Point P;
    T x, y;
    explicit Point(T x=0, T y=0) : x(x), y(y) {}
    bool operator<(P p) const { return tie(x,y) < tie(p.x,p.y); }
    bool operator==(P p) const { return tie(x,y)==tie(p.x,p.y); }

    P operator+(P p) const { return P(x+p.x, y+p.y); }
    P operator-(P p) const { return P(x-p.x, y-p.y); }
    P operator*(T d) const { return P(x*d, y*d); }
    P operator/(T d) const { return P(x/d, y/d); }
    T dot(P p) const { return x*p.x + y*p.y; }
    T cross(P p) const { return x*p.y - y*p.x; }
    T cross(P a, P b) const { return (a-*this).cross(b-*this); }

    T dist2() const { return x*x + y*y; }
    double dist() const { return sqrt((double)dist2()); }
    // angle to x-axis in interval [-pi, pi]
    double angle() const { return atan2(y, x); }
    P unit() const { return *this/dist(); } // makes dist()=1
    P perp() const { return P(-y, x); } // rotates +90 degrees
    P normal() const { return perp().unit(); }
}

```

```

// returns point rotated 'a' radians ccw around the origin
P rotate(double a) const {
    return P(x*cos(a)-y*sin(a),x*sin(a)+y*cos(a)); }
friend ostream& operator<<(ostream& os, P p) {
    return os << "(" << p.x << ", " << p.y << ")"; }
};

```

5.21 PointInsideHull

```

/**
 * Description: Determine whether a point t lies inside a
 * convex hull (CCW
 * order, with no collinear points). Returns true if point
 * lies within
 * the hull. If strict is true, points on the boundary aren'
 * t included.
 * Usage:
 * Status: stress-tested
 * Time: O(\log N)
 */
#pragma once

#include "Point.h"
#include "sideOf.h"
#include "OnSegment.h"

typedef Point<ll> P;

bool inHull(const vector<P>& l, P p, bool strict = true) {
    int a = 1, b = sz(l) - 1, r = !strict;
    if (sz(l) < 3) return r && onSegment(l[0], l.back(), p);
    if (sideOf(l[0], l[a], l[b]) > 0) swap(a, b);
    if (sideOf(l[0], l[a], p) >= r || sideOf(l[0], l[b], p) <= -
        r)
        return false;
    while (abs(a - b) > 1) {
        int c = (a + b) / 2;
        (sideOf(l[0], l[c], p) > 0 ? b : a) = c;
    }
    return sgn(l[a].cross(l[b], p)) < r;
}

```

5.22 PolygonArea

```

/**
 * Description: Returns twice the signed area of a polygon.

```

```

* Clockwise enumeration gives negative area. Watch out for
  overflow if using int as T!
* Status: Stress-tested and tested on kattis:polygonarea
*/
#pragma once

#include "Point.h"

template<class T>
T polygonArea2(vector<Point<T>>& v) {
    T a = v.back().cross(v[0]);
    rep(i,0,sz(v)-1) a += v[i].cross(v[i+1]);
    return a;
}

```

5.23 PolygonCenter

```

/**
 * Description: Returns the center of mass for a polygon.
 * Time: O(n)
 * Status: Tested
 */
#pragma once

#include "Point.h"

typedef Point<double> P;
P polygonCenter(const vector<P>& v) {
    P res(0, 0); double A = 0;
    for (int i = 0, j = sz(v) - 1; i < sz(v); j = i++) {
        res = res + (v[i] + v[j]) * v[j].cross(v[i]);
        A += v[j].cross(v[i]);
    }
    return res / A / 3;
}

```

5.24 PolygonCut

```

/**
 Returns a vector with the vertices of a polygon with
 everything to the left of the line going from s to e
 cut away.
 * Usage:
 * vector<P> p = ...;
 * p = polygonCut(p, P(0,0), P(1,0));
 */
#pragma once

```

```

#include "Point.h"
#include "lineIntersection.h"

typedef Point<double> P;
vector<P> polygonCut(const vector<P>& poly, P s, P e) {
    vector<P> res;
    rep(i,0,sz(poly)) {
        P cur = poly[i], prev = i ? poly[i-1] : poly.back();
        bool side = s.cross(e, cur) < 0;
        if (side != (s.cross(e, prev) < 0))
            res.push_back(lineInter(s, e, cur, prev).second);
        if (side)
            res.push_back(cur);
    }
    return res;
}

```

5.25 PolygonUnion

```

/**
 * Description: Calculates the area of the union of $n$
   polygons (not necessarily
   * convex). The points within each polygon must be given in
   CCW order.
 * (Epsilon checks may optionally be added to sideOf/sgn,
   but shouldn't be needed.)
 * Time: $O(N^2)$, where $N$ is the total number of points
 * Status: stress-tested, Submitted on ECNA 2017 Problem A
 */
#pragma once

#include "Point.h"
#include "sideOf.h"

typedef Point<double> P;
double rat(P a, P b) { return sgn(b.x) ? a.x/b.x : a.y/b.y; }

double polyUnion(vector<vector<P>>& poly) {
    double ret = 0;
    rep(i,0,sz(poly)) rep(v,0,sz(poly[i])) {
        P A = poly[i][v], B = poly[i][(v + 1) % sz(poly[i])];
        vector<pair<double, int>> segs = {{0, 0}, {1, 0}};
        rep(j,0,sz(poly)) if (i != j) {
            rep(u,0,sz(poly[j])) {
                P C = poly[j][u], D = poly[j][(u + 1) % sz(poly[j])];
                int sc = sideOf(A, B, C), sd = sideOf(A, B, D);
                if (sc != sd) {
                    double sa = C.cross(D, A), sb = C.cross(D, B);

```

```

                    if (min(sc, sd) < 0)
                        segs.emplace_back(sa / (sa - sb), sgn(sc - sd));
                } else if (!sc && !sd && j < i && sgn((B-A).dot(D-C)) > 0) {
                    segs.emplace_back(rat(C - A, B - A), 1);
                    segs.emplace_back(rat(D - A, B - A), -1);
                }
            }
        }
        sort(all(segs));
        for (auto& s : segs) s.first = min(max(s.first, 0.0), 1.0);
        double sum = 0;
        int cnt = segs[0].second;
        rep(j,1,sz(segs)) {
            if (!cnt) sum += segs[j].first - segs[j - 1].first;
            cnt += segs[j].second;
        }
        ret += A.cross(B) * sum;
    }
    return ret / 2;
}

```

5.26 SegmentDistance

```

/**
 Returns the shortest distance between point p and the line
 segment from point s to e.
 * Usage:
 * Point<double> a, b(2,2), p(1,1);
 * bool onSegment = segDist(a,b,p) < 1e-10;
 * Status: tested
 */
#pragma once

#include "Point.h"

typedef Point<double> P;
double segDist(P& s, P& e, P& p) {
    if (s==e) return (p-s).dist();
    auto d = (e-s).dist2(), t = min(d,max(.0,(p-s).dot(e-s)));
    return ((p-s)*d-(e-s)*t).dist()/d;
}

```

5.27 SegmentIntersection

```

/**

```


If a unique intersection point between the line segments going from s_1 to e_1 and from s_2 to e_2 exists then it is returned.

If no intersection point exists an empty vector is returned. If infinitely many exist a vector with 2 elements is returned, containing the endpoints of the common line segment.

The wrong position will be returned if P is `Point<ll>` and the intersection point does not have integer coordinates.

Products of three coordinates are used in intermediate steps so watch out for overflow if using `int` or `long long`.

Usage:

```
* vector<P> inter = segInter(s1,e1,s2,e2);
* if (sz(inter)==1)
*   cout << "segments intersect at " << inter[0] << endl;
* Status: stress-tested, tested on kattis:intersection
*/
```

```
#pragma once

#include "Point.h"
#include "OnSegment.h"

template<class P> vector<P> segInter(P a, P b, P c, P d) {
    auto oa = c.cross(d, a), ob = c.cross(d, b),
          oc = a.cross(b, c), od = a.cross(b, d);
    // Checks if intersection is single non-endpoint point.
    if (sgn(oa) * sgn(ob) < 0 && sgn(oc) * sgn(od) < 0)
        return {(a * ob - b * oa) / (ob - oa)};
    set<P> s;
    if (onSegment(c, d, a)) s.insert(a);
    if (onSegment(c, d, b)) s.insert(b);
    if (onSegment(a, b, c)) s.insert(c);
    if (onSegment(a, b, d)) s.insert(d);
    return {all(s)};
}
```

5.28 sideOf

```
/**
 * Description: Returns where $p$ is as seen from $s$
 * towards $e$. 1/0/-1 $\rightarrow$ left/on line/
 * right. If the optional argument $eps$ is given 0 is
 * returned if $p$ is within distance $eps$ from the line.
 * $P$ is supposed to be Point<T> where  $T$  is e.g. double or
 * long long. It uses products in intermediate steps so
 * watch out for overflow if using int or long long.
 * Usage:
 * bool left = sideOf(p1,p2,q)==1;
```

```
* Status: tested
*/
#pragma once

#include "Point.h"

template<class P>
int sideOf(P s, P e, P p) { return sgn(s.cross(e, p)); }

template<class P>
int sideOf(const P& s, const P& e, const P& p, double eps) {
    auto a = (e-s).cross(p-s);
    double l = (e-s).dist()*eps;
    return (a > l) - (a < -l);
}
```

6 Graph

6.1 2SAT

```
//task : n people, each people have 2 request : + x or - x
//Ask : Is there a way build array m elements that for each
//       people,
//at least one of two request is satisfied. If yes, print it
//       .

void tarjan(int u) {
    st[++top] = u;
    in[u] = low[u] = ++tme;
    was_tarjan[u] = true;
    for (int v : g[u]) {
        if (!was_tarjan[v]) {
            tarjan(v);
            low[u] = min(low[u], low[v]);
        }
        else {
            low[u] = min(low[u], in[v]);
        }
    }

    if (low[u] == in[u]) {
        ++scc_cnt;
        while (st[top] != u) {
            scc_id[st[top]] = scc_cnt;
            in[st[top]] = low[st[top]] = N;
            --top;
        }
        scc_id[st[top--]] = scc_cnt;
    }
```

```
        in[u] = low[u] = N;
    }
}

void compress_scc_to_dag() {
    for (int u = 1; u <= m; ++u)
        for (int v : g[u]) if (scc_id[v] != scc_id[u])
            g2[scc_id[u]].emplace_back(scc_id[v]);
}

void dfs(int u) { // toposort
    was[u] = true;
    for (int v : g2[u]) if (!was[v]) dfs(v);

    topo[u] = top--;
}

void toposort_g2() {
    top = scc_cnt;
    fill(was + 1, was + scc_cnt + 1, false);
    for (int i = 1; i <= scc_cnt; ++i) if (!was[i]) dfs(i);
}

int main() {
    //io()
    m = n << 1;
    for (int i = 1; i <= m; ++i) if (!was_tarjan[i]) tarjan(i);

    for (int i = 1; i <= n; ++i) if (scc_id[i] == scc_id[i + n]) {
        cout << "IMPOSSIBLE";
        return 0;
    }

    compress_scc_to_dag();
    toposort_g2();

    for (int i = 1; i <= n; ++i) {
        cout << (topo[scc_id[i]] > topo[scc_id[i + n]] ? '+'
                : '-') << ' ';
    }
}
```

6.2 Blocc Cut Tree

```
struct block_cut_tree {
    int h[maxn], p[20][maxn], in[maxn], low[maxn], cnt = 0,
        node;
```

```

vector<int> adj2[maxn], st;
void DFS(int u = 1, int p = 0)
{
    in[u] = low[u] = ++cnt;
    st.emplace_back(u);
    for (int v: adj[u])
    {
        if (v == p) continue;
        if (in[v]) low[u] = min(low[u], in[v]);
        else
        {
            DFS(v, u);
            low[u] = min(low[u], low[v]);
            if (low[v] >= in[u])
            {
                node++;
                int x;
                do
                {
                    x = st.back(), st.pop_back();
                    adj2[node].emplace_back(x),
                    adj2[x].emplace_back(node);
                } while (x != v);
                adj2[node].emplace_back(u);
                adj2[u].emplace_back(node);
            }
        }
    }
}

void DFS2(int u = 1, int pa = 0)
{
    for (int v: adj2[u])
    {
        if (v == pa) continue;
        h[v] = h[u] + 1;
        p[0][v] = u;
        DFS2(v, u);
    }
}

int lca(int u, int v)
{
    if (h[u] > h[v]) swap(u, v);
    int dis = h[v] - h[u];
    for (int i=19; i>=0; i--) if ((dis>>i)&1) v = p[i][v];
    if (v == u) return u;
    for (int i=19; i>=0; i--) if (p[i][u] != p[i][v]) u = p[i][u], v = p[i][v];
    return p[0][u];
}

```

```

int dis(int u, int v) {return h[u] + h[v] - 2 * h[lca(u, v)];}
int query(int u, int v)
{
    return dis(u, v)/2 + 1;
}

void init()
{
    memset(p, -1, sizeof p);
    memset(in, 0, sizeof in);
    memset(low, 0, sizeof low);
    memset(h, 0, sizeof h);
    st.clear();
    for (int i=1; i<=n; i++) adj2[i].clear();
    node = n;
    DFS();
    DFS2();
    for (int i=1; (1<<i) <= node; i++)
        for (int j=1; j<=node; j++)
            if (p[i-1][j] != -1) p[i][j] = p[i-1][p[i-1][j]];
}

} bctree;

```

6.3 DeMen Bipartite Matching

```

vector<int> a[N];
int mr[N], cttme;
bool ml[N];
char f[N];

bool dfs(int u){
    if (f[u] == cttme) return false;
    f[u] = cttme;
    for(int i : a[u]){
        if (!mr[i] || dfs(mr[i])){
            mr[i] = u;
            return true;
        }
    }
    return false;
}

int maximum_matching(){
    int cnt = 0;
    for(bool run = true; run;){
        cttme++;
        run = false;
    }
}

```

```

for(int i = 1; i <= n; ++i){
    if (ml[i]) continue;
    if (dfs(i)){
        ml[i] = run = true;
        ++cnt;
    }
}
return cnt;
}

```

6.4 General Matching

```

struct GeneralMatching {
    GeneralMatching(int _n) : n(_n), match(_n, -1), g(_n),
        timer(-1), label(_n), parent(_n), orig(_n), aux(_n, -1) {}

    void add_edge(int u, int v) {
        g[u].push_back(v);
        g[v].push_back(u);
    }

    int get_match() {
        for (int i = 0; i < n; i++) {
            if (match[i] == -1) bfs(i);
        }
        int res = 0;
        for (int i = 0; i < n; i++) {
            if (match[i] >= 0) ++res;
        }
        return res / 2;
    }

    int n;
    vector<int> match;

private:
    int lca(int x, int y) {
        for (timer++; ; swap(x, y)) {
            if (x == -1) continue;
            if (aux[x] == timer) return x;
            aux[x] = timer;
            x = (match[x] == -1 ? -1 : orig[parent[match[x]]]);
        }
    }

    void blossom(int v, int w, int a) {

```



```

while (orig[v] != a) {
    parent[v] = w;
    w = match[v];
    if (label[w] == 1) {
        label[w] = 0;
        q.push_back(w);
    }
    orig[v] = orig[w] = a;
    v = parent[w];
}

void augment(int v) {
    while (v != -1) {
        int pv = parent[v], nv = match[pv];
        match[v] = pv; match[pv] = v; v = nv;
    }
}

int bfs(int root) {
    fill(label.begin(), label.end(), -1);
    iota(orig.begin(), orig.end(), 0);
    q.clear();
    label[root] = 0;
    q.push_back(root);
    for (int i = 0; i < (int) q.size(); ++i) {
        int v = q[i];
        for (auto x : g[v]) {
            if (label[x] == -1) {
                label[x] = 1;
                parent[x] = v;
                if (match[x] == -1) {
                    augment(x);
                    return 1;
                }
            }
            label[match[x]] = 0;
            q.push_back(match[x]);
        }
        else if (label[x] == 0 && orig[v] != orig[x]) {
            int a = lca(orig[v], orig[x]);
            blossom(x, v, a);
            blossom(v, x, a);
        }
    }
    return 0;
}

private:
vector<vector<int>> g;

```

```

int timer;
vector<int> label, parent, orig, aux, q;
};

```

6.5 Gomory Hu Tree

```

//0-based

struct PushRelabel {
    struct Edge {
        int dest, back;
        int f, c;
    };
    vector<vector<Edge>> g;
    vector<int> ec;
    vector<Edge*> cur;
    vector<vector<int>> hs; vector<int> H;
    PushRelabel(int n) : g(n), ec(n), cur(n), hs(2*n), H(n) {}

    void addEdge(int s, int t, int cap, int rcap=0) {
        if (s == t) return;
        g[s].push_back({t, (int)g[t].size(), 0, cap});
        g[t].push_back({s, (int)g[s].size()-1, 0, rcap});
    }

    void addFlow(Edge& e, int f) {
        Edge &back = g[e.dest][e.back];
        if (!ec[e.dest] && f) hs[H[e.dest]].push_back(e.dest);
        e.f += f; e.c -= f; ec[e.dest] += f;
        back.f -= f; back.c += f; ec[back.dest] -= f;
    }

    int calc(int s, int t) {
        int v = g.size(); H[s] = v; ec[t] = 1;
        vector<int> co(2*v); co[0] = v-1;
        for(int i = 0; i < v; ++i) cur[i] = g[i].data();
        for (Edge& e : g[s]) addFlow(e, e.c);

        for (int hi = 0;;) {
            while (hs[hi].empty()) if (!hi--) return -ec[s];
            int u = hs[hi].back(); hs[hi].pop_back();
            while (ec[u] > 0) // discharge u
                if (cur[u] == g[u].data() + g[u].size()) {
                    H[u] = 1e9;
                    for (Edge& e : g[u]) if (e.c && H[u] > H[e.dest]+1)
                        H[u] = H[e.dest]+1, cur[u] = &e;
                    if (++co[H[u]], !--co[hi] && hi < v)
                        for(int i = 0; i < v; ++i) if (hi < H[i] && H[i] < v)
                            --co[H[i]], H[i] = v + 1;
                    hi = H[u];
                }
        }
    }
}

```

```

} else if (cur[u]->c && H[u] == H[cur[u]->dest]+1)
    addFlow(*cur[u], min(ec[u], cur[u]->c));
else ++cur[u];
}
}

bool leftOfMinCut(int a) { return H[a] >= (int)g.size(); }
};

vector<array<int, 3>> gomoryHu(int N, vector<array<int, 3>>
    ed) {
    vector<array<int, 3>> tree;
    vector<int> par(N);
    for(int i = 0; i < N; ++i) {
        PushRelabel D(N); // Dinic also works
        for (array<int, 3> t : ed) D.addEdge(t[0], t[1], t[2], t[2]);

        int flow = D.calc(i, par[i]);
        if (flow != -1) tree.push_back({i, par[i], flow});
        for(int j = i + 1; j < N; ++j)
            if (par[j] == par[i] && D.leftOfMinCut(j)) par[j] = i;
    }
    return tree;
}

```

6.6 Max Cost General Matching

```

#include <bits/stdc++.h>
using namespace std;
// Cap ghép có trọng số lớn nhất - Thái Hoàng
#define MCM MaxCostMatching
namespace MaxCostMatching {
#define dist(e) (lab[e.u] + lab[e.v] - g[e.u][e.v].w * 2)
    const int maxn = 1e3 + 5;
    const int oo = (int) 1e9;
    struct Edge {
        int u, v, w;
    } g[maxn][maxn];

    int n, m, n_x;
    int lab[maxn], match[maxn], slack[maxn], st[maxn], pa[
        maxn];
    int flower_from[maxn][maxn], s[maxn], vis[maxn];
    vector<int> flower[maxn];
    deque<int> q;

    void init(int _n) {
        n = _n;
        for (int u = 1; u <= n; u++) {

```

```

    for (int v = 1; v <= n; v++) {
        g[u][v] = Edge{u, v, 0};
    }
}

void add(int u, int v, int w) {
    g[u][v].w = max(g[u][v].w, w);
    g[v][u].w = max(g[v][u].w, w);
}

void update_slack(int u, int x) {
    if (!slack[x] || dist(g[u][x]) < dist(g[slack[x]][x]))
        slack[x] = u;
}

void set_slack(int x) {
    slack[x] = 0;
    for (int u = 1; u <= n; u++) {
        if (g[u][x].w > 0 && st[u] != x && s[st[u]] == 0)
            update_slack(u, x);
    }
}

void q_push(int x) {
    if (x <= n) return q.push_back(x);
    for (int i = 0; i < flower[x].size(); i++) q_push(
        flower[x][i]);
}

void set_st(int x, int b) {
    st[x] = b;
    if (x <= n) return;
    for (int i = 0; i < flower[x].size(); i++) set_st(
        flower[x][i], b);
}

int get_pr(int b, int xr) {
    int pr = find(flower[b].begin(), flower[b].end(), xr)
        - flower[b].begin();
    if (pr % 2 == 1) {
        reverse(flower[b].begin() + 1, flower[b].end());
        return (int) flower[b].size() - pr;
    }
    else {
        return pr;
    }
}

void set_match(int u, int v) {
    match[u] = g[u][v].v;
    if (u <= n) return;
    Edge e = g[u][v];
    int xr = flower_from[u][e.u], pr = get_pr(u, xr);
    for (int i = 0; i < pr; i++) set_match(flower[u][i],
        flower[u][i ^ 1]);
    set_match(xr, v);
}

```

```

    rotate(flower[u].begin(), flower[u].begin() + pr,
        flower[u].end());
}

void augment(int u, int v) {
    int xnv = st[match[u]];
    set_match(u, v);
    if (!xnv) return;
    set_match(xnv, st[pa[xnv]]);
    augment(st[pa[xnv]], xnv);
}

int get_lca(int u, int v) {
    static int t = 0;
    for (t++; u || v; swap(u, v)) {
        if (u == 0) continue;
        if (vis[u] == t) return u;
        vis[u] = t;
        u = st[match[u]];
        if (u) u = st[pa[u]];
    }
    return 0;
}

void add_blossom(int u, int lca, int v) {
    int b = n + 1;
    while (b <= n_x && st[b]) b++;
    if (b > n_x) n_x++;
    lab[b] = 0, s[b] = 0;
    match[b] = match[lca];
    flower[b].clear();
    flower[b].push_back(lca);
    for (int x = u, y; x != lca; x = st[pa[y]]) {
        flower[b].push_back(x), flower[b].push_back(y =
            st[match[x]]), q_push(y);
    }
    reverse(flower[b].begin() + 1, flower[b].end());
    for (int x = v, y; x != lca; x = st[pa[y]]) {
        flower[b].push_back(x), flower[b].push_back(y =
            st[match[x]]), q_push(y);
    }
    set_st(b, b);
    for (int x = 1; x <= n_x; x++) g[b][x].w = g[x][b].w
        = 0;
    for (int x = 1; x <= n; x++) flower_from[b][x] = 0;
    for (int i = 0; i < flower[b].size(); i++) {
        int xs = flower[b][i];
        for (int x = 1; x <= n_x; x++) {
            if (g[b][x].w == 0 || dist(g[xs][x]) < dist(g[
                b][x])) {
                g[b][x] = g[xs][x], g[x][b] = g[x][xs];
            }
        }
    }
}

```

```

    for (int x = 1; x <= n; x++) {
        if (flower_from[xs][x]) flower_from[b][x] = xs
            ;
    }
    set_slack(b);
}

void expand_blossom(int b) {
    for (int i = 0; i < flower[b].size(); i++) {
        set_st(flower[b][i], flower[b][i]);
    }
    int xr = flower_from[b][g[b][pa[b]].u], pr = get_pr(b,
        xr);
    for (int i = 0; i < pr; i += 2) {
        int xs = flower[b][i], xns = flower[b][i + 1];
        pa[xs] = g[xns][xs].u;
        s[xs] = 1, s[xns] = 0;
        slack[xs] = 0, set_slack(xns);
        q_push(xns);
    }
    s[xr] = 1, pa[xr] = pa[b];
    for (int i = pr + 1; i < flower[b].size(); i++) {
        int xs = flower[b][i];
        s[xs] = -1, set_slack(xs);
    }
    st[b] = 0;
}

int on_found_Edge(const Edge &e) {
    int u = st[e.u], v = st[e.v];
    if (s[v] == -1) {
        pa[v] = e.u, s[v] = 1;
        int nu = st[match[v]];
        slack[v] = slack[nu] = 0;
        s[nu] = 0, q_push(nu);
    }
    else if (s[v] == 0) {
        int lca = get_lca(u, v);
        if (!lca) return augment(u, v), augment(v, u), 1;
        else add_blossom(u, lca, v);
    }
    return 0;
}

int matching() {
    fill(s, s + n_x + 1, -1), fill(slack, slack + n_x +
        1, 0);
    q.clear();
    for (int x = 1; x <= n_x; x++) {
        if (st[x] == x && !match[x]) pa[x] = 0, s[x] = 0,
            q_push(x);
    }
}

```

```

if (q.empty()) return 0;
while (1) {
    while (q.size()) {
        int u = q.front();
        q.pop_front();
        if (s[st[u]] == 1) continue;
        for (int v = 1; v <= n; v++) {
            if (g[u][v].w > 0 && st[u] != st[v]) {
                if (dist(g[u][v]) == 0) {
                    if (on_found_Edge(g[u][v])) return 1;
                }
                else update_slack(u, st[v]);
            }
        }
    }
    int d = oo;
    for (int b = n + 1; b <= n_x; b++) {
        if (st[b] == b && s[b] == 1) d = min(d, lab[b] / 2);
    }
    for (int x = 1; x <= n_x; x++) {
        if (st[x] == x && slack[x]) {
            if (s[x] == -1) d = min(d, dist(g[slack[x]][x]));
            else if (s[x] == 0) d = min(d, dist(g[slack[x]][x]) / 2);
        }
    }
    for (int u = 1; u <= n; u++) {
        if (s[st[u]] == 0) {
            if (lab[u] <= d) return 0;
            lab[u] -= d;
        }
        else if (s[st[u]] == 1) lab[u] += d;
    }
    for (int b = n + 1; b <= n_x; b++) {
        if (st[b] == b) {
            if (s[st[b]] == 0) lab[b] += d * 2;
            else if (s[st[b]] == 1) lab[b] -= d * 2;
        }
    }
    q.clear();
    for (int x = 1; x <= n_x; x++) {
        if (st[x] == x && slack[x] && st[slack[x]] != x && dist(g[slack[x]][x]) == 0) {
            if (on_found_Edge(g[slack[x]][x])) return 1;
        }
    }
}

```

```

    for (int b = n + 1; b <= n_x; b++) {
        if (st[b] == b && s[b] == 1 && lab[b] == 0)
            expand_blossom(b);
    }
    return 0;
}

int maxcost() {
    fill(match, match + n + 1, 0);
    n_x = n;
    int tot_weight = 0;
    int n_matches = 0;
    for (int u = 0; u <= n; u++) st[u] = u, flower[u].clear();
    int w_max = 0;
    for (int u = 1; u <= n; u++) {
        for (int v = 1; v <= n; v++) {
            flower_from[u][v] = (u == v ? u : 0);
            w_max = max(w_max, g[u][v].w);
        }
    }
    for (int u = 1; u <= n; u++) lab[u] = w_max;
    while (matching()) n_matches++;
    for (int u = 1; u <= n; u++) {
        if (match[u] && match[u] < u) {
            tot_weight += g[u][match[u]].w;
        }
    }
    return tot_weight;
}

int main() {
    MCM::init(4);
    MCM::add(1, 2, 5);
    MCM::add(2, 3, 10);
    MCM::add(3, 4, 2);
    cout << MCM::maxcost() << "\n";
    return 0;
}

```

6.7 Max Flow

```

struct FlowEdge {
    int v, u;
    long long cap, flow = 0;
    FlowEdge(int v, int u, long long cap) : v(v), u(u), cap(cap) {}
};

```

```

struct Dinic {
    const long long flow_inf = 1e18;
    vector<FlowEdge> edges;
    vector<vector<int>>> adj;
    int n, m = 0;
    int s, t;
    vector<int> level, ptr;
    queue<int> q;

    Dinic(int n, int s, int t) : n(n), s(s), t(t) {
        adj.resize(n);
        level.resize(n);
        ptr.resize(n);
    }

    void add_edge(int v, int u, long long cap) {
        edges.emplace_back(v, u, cap);
        edges.emplace_back(u, v, 0);
        adj[v].push_back(m);
        adj[u].push_back(m + 1);
        m += 2;
    }

    bool bfs() {
        while (!q.empty()) {
            int v = q.front();
            q.pop();
            for (int id : adj[v]) {
                if (edges[id].cap - edges[id].flow < 1)
                    continue;
                if (level[edges[id].u] != -1)
                    continue;
                level[edges[id].u] = level[v] + 1;
                q.push(edges[id].u);
            }
        }
        return level[t] != -1;
    }

    long long dfs(int v, long long pushed) {
        if (pushed == 0)
            return 0;
        if (v == t)
            return pushed;
        for (int& cid = ptr[v]; cid < (int)adj[v].size(); cid++) {
            int id = adj[v][cid];
            int u = edges[id].u;

```

```

    if (level[v] + 1 != level[u] || edges[id].cap -
        edges[id].flow < 1)
        continue;
    long long tr = dfs(u, min(pushed, edges[id].cap -
        edges[id].flow));
    if (tr == 0)
        continue;
    edges[id].flow += tr;
    edges[id ^ 1].flow -= tr;
    return tr;
}
return 0;
}

long long flow() {
    long long f = 0;
    while (true) {
        fill(level.begin(), level.end(), -1);
        level[s] = 0;
        q.push(s);
        if (!bfs())
            break;
        fill(ptr.begin(), ptr.end(), 0);
        while (long long pushed = dfs(s, flow_inf)) {
            f += pushed;
        }
    }
    return f;
}
};

```

6.8 Maximum Independent Set

```

void cal(int node, bool col){
    if(col)
    {
        if(visitb[node]) return;
    }
    else
    {
        if(visita[node]) return;
    }
    if(col)
        visitb[node] = 1;
    else
        visita[node] = 1;
    if(col){ // node from the right side, can only
        traverse matched edge
        cal(pb[node], col ^ 1);
    }
}

```

```

        return;
    }
    for(auto nx:adj[node]){
        //cout << nx << ' ' << pa[node] << '\n';
        if(nx==pa[node])continue;
        cal(nx, col ^ 1);
    }
}
int main(){
    DemenMatching();
    for(int i=1;i<=n;i++){
        if(pa[i])continue; // matched node from the left
        side
        cal(i,0);
    }
    vector<int> MaxISa, MaxISb, MVCa, MVCb; // find max cover
    and minimum cover
    for(int i=1;i<=n;i++){
        if(visita[i]) MaxISa.pb(i); // Minimum independent set
        is visted on the left
        else MVCa.pb(i); // Max vertex cover is not visited
        on left
    }
    for(int i=1;i<=n;i++){
        if(!visitb[i])MaxISb.pb(i); // Minimum independent set
        is not visted on the right
        else MVCb.pb(i); // Max vertex cover is visited on
        right
    }
}
}

```

6.9 Prüfer Decode

```

vector<pair<int, int>> pruefer_decode(vector<int> const&
    code) {
    int n = code.size() + 2;
    vector<int> degree(n, 1);
    for (int i : code)
        degree[i]++;

    int ptr = 0;
    while (degree[ptr] != 1)
        ptr++;
    int leaf = ptr;

    vector<pair<int, int>> edges;
    for (int v : code) {
        edges.emplace_back(leaf, v);
        if (--degree[v] == 1 && v < ptr) {

```

```

            leaf = v;
        } else {
            ptr++;
            while (degree[ptr] != 1)
                ptr++;
            leaf = ptr;
        }
    }
    edges.emplace_back(leaf, n-1);
    return edges;
}

```

6.10 Prüfer Encode

```

vector<vector<int>> adj;
vector<int> parent;

void dfs(int v) {
    for (int u : adj[v]) {
        if (u != parent[v]) {
            parent[u] = v;
            dfs(u);
        }
    }
}

vector<int> pruefer_code() {
    int n = adj.size();
    parent.resize(n);
    parent[n-1] = -1;
    dfs(n-1);

    int ptr = -1;
    vector<int> degree(n);
    for (int i = 0; i < n; i++) {
        degree[i] = adj[i].size();
        if (degree[i] == 1 && ptr == -1)
            ptr = i;
    }

    vector<int> code(n - 2);
    int leaf = ptr;
    for (int i = 0; i < n - 2; i++) {
        int next = parent[leaf];
        code[i] = next;
        if (--degree[next] == 1 && next < ptr) {
            leaf = next;
        } else {
            ptr++;
        }
    }
}

```

```

        while (degree[ptr] != 1)
            ptr++;
        leaf = ptr;
    }
}

return code;
}

```

7 Mathematics

7.1 CRT

```

//Given m1, m2, r1, r2, find x :
// - x mod m1 = r1
// - x mod m2 = r2

int cal_crt(int r1, int r2, int m1, int m2){
    ii ExEuclid(m1, m2);
    bool f = false;
    int g = __gcd(m1, m2);
    if ((r2 - r1) % g) return 1e18; //No solution

    int k = ans.x * ((r2 - r1) / g);
    k %= (m2 / g);
    return (r1 + k * m1) % lc;
    //all_ans = {ans + k*LCM(m1, m2)}
}

```

7.2 Extended Euclid

```

ii ExEuclid(int a, int b){
    int x0 = 1, y0 = 0;
    int x1 = 0, y1 = 1;
    int x2, y2;
    while (b){
        int q = a / b;
        int r = a % b;
        a = b; b = r;
        x2 = x0 - q * x1;
        y2 = y0 - q * y1;
        x0 = x1; y0 = y1;
        x1 = x2; y1 = y2;
    }
    return {x0, y0};
}

```

7.3 Fast Fourier Transform - Mod

```

/**
 * Description: Higher precision FFT, can be used for
 *               convolutions modulo arbitrary integers
 * as long as  $N \log_2 N \cdot \text{mod} < 8.6 \cdot 10^{14}$ 
 * (in practice  $10^{16}$  or higher).
 * Inputs must be in  $[0, \text{mod})$ .
 * Time:  $O(N \log N)$ , where  $N = |A| + |B|$  (twice as slow as
 *       NTT or FFT)
 */

typedef vector<ll> vl;
template<int M> vl convMod(const vl &a, const vl &b) {
    if (a.empty() || b.empty()) return {};
    vl res(sz(a) + sz(b) - 1);
    int B=32-__builtin_clz(sz(res)), n=1<<B, cut=int(sqrt(M));
    vector<C> L(n), R(n), outs(n), outl(n);
    rep(i,0,sz(a)) L[i] = C((int)a[i] / cut, (int)a[i] % cut);
    rep(i,0,sz(b)) R[i] = C((int)b[i] / cut, (int)b[i] % cut);
    fft(L), fft(R);
    rep(i,0,n) {
        int j = -i & (n - 1);
        outl[j] = (L[i] + conj(L[j])) * R[i] / (2.0 * n);
        outs[j] = (L[i] - conj(L[j])) * R[i] / (2.0 * n) / 1i;
    }
    fft(outl), fft(outs);
    rep(i,0,sz(res)) {
        ll av = ll(real(outl[i])+.5), cv = ll(imag(outs[i])+.5);
        ll bv = ll(imag(outl[i])+.5) + ll(real(outs[i])+.5);
        res[i] = ((av % M * cut + bv) % M * cut + cv) % M;
    }
    return res;
}

```

7.4 Fast Fourier Transform

```

/**
 * Time:  $O(N \log N)$  with  $N = |A| + |B|$  ( $\tilde{O}(N)$  for  $N = 2^{22}$ )
 */

typedef complex<double> C;
typedef vector<double> vd;
void fft(vector<C> &a) {
    int n = sz(a), L = 31 - __builtin_clz(n);
    static vector<complex<long double>> R(2, 1);
    static vector<C> rt(2, 1); // (~ 10% faster if double)
    for (static int k = 2; k < n; k *= 2) {

```

```

        R.resize(n); rt.resize(n);
        auto x = polar(1.0L, acos(-1.0L) / k);
        rep(i,k,2*k) rt[i] = R[i] = i&1 ? R[i/2] * x : R[i/2];
    }
    vi rev(n);
    rep(i,0,n) rev[i] = (rev[i / 2] | (i & 1) << L) / 2;
    rep(i,0,n) if (i < rev[i]) swap(a[i], a[rev[i]]);
    for (int k = 1; k < n; k *= 2)
        for (int i = 0; i < n; i += 2 * k) rep(j,0,k) {
            // C z = rt[j+k] * a[i+j+k]; // (25% faster if hand-rolled)
            // include-line
            auto x = (double *)&rt[j+k], y = (double *)&a[i+j+k];
            // exclude-line
            C z(x[0]*y[0] - x[1]*y[1], x[0]*y[1] + x[1]*y[0]);
            // exclude-line
            a[i + j + k] = a[i + j] - z;
            a[i + j] += z;
        }
    }
    vd conv(const vd& a, const vd& b) {
        if (a.empty() || b.empty()) return {};
        vd res(sz(a) + sz(b) - 1);
        int L = 32 - __builtin_clz(sz(res)), n = 1 << L;
        vector<C> in(n), out(n);
        copy(all(a), begin(in));
        rep(i,0,sz(b)) in[i].imag(b[i]);
        fft(in);
        for (C& x : in) x *= x;
        rep(i,0,n) out[i] = in[-i & (n - 1)] - conj(in[i]);
        fft(out);
        rep(i,0,sz(res)) res[i] = imag(out[i]) / (4 * n);
        return res;
    }
}

```

7.5 Fast Walsh–Hadamard Transform

```

#include <bits/stdc++.h>
using namespace std;

int fpow(int n, long long k, int p = (int) 1e9 + 7) {
    int r = 1;
    for (; k; k >>= 1) {
        if (k & 1) r = (long long) r * n % p;
        n = (long long) n * n % p;
    }
    return r;
}

/*

```

```

* matrix:
* +1 +1
* +1 -1
*/
void XORFFT(int a[], int n, int p, int invert) {
    for (int i = 1; i < n; i <= 1) {
        for (int j = 0; j < n; j += i << 1) {
            for (int k = 0; k < i; k++) {
                int u = a[j + k], v = a[i + j + k];
                a[j + k] = u + v;
                if (a[j + k] >= p) a[j + k] -= p;
                a[i + j + k] = u - v;
                if (a[i + j + k] < 0) a[i + j + k] += p;
            }
        }
    }
    if (invert) {
        long long inv = fpow(n, p - 2, p);
        for (int i = 0; i < n; i++) a[i] = a[i] * inv % p;
    }
}
/*
* Matrix:
* +1 +1
* +1 +0
*/
void ORFFT(int a[], int n, int p, int invert) {
    for (int i = 1; i < n; i <= 1) {
        for (int j = 0; j < n; j += i << 1) {
            for (int k = 0; k < i; k++) {
                int u = a[j + k], v = a[i + j + k];
                if (!invert) {
                    a[j + k] = u + v;
                    a[i + j + k] = u;
                    if (a[j + k] >= p) a[j + k] -= p;
                }
                else {
                    a[j + k] = v;
                    a[i + j + k] = u - v;
                    if (a[i + j + k] < 0) a[i + j + k] += p;
                }
            }
        }
    }
}
/*
* matrix:
* +0 +1
* +1 +1
*/

```

```

void ANDFFT(int a[], int n, int p, int invert) {
    for (int i = 1; i < n; i <= 1) {
        for (int j = 0; j < n; j += i << 1) {
            for (int k = 0; k < i; k++) {
                int u = a[j + k], v = a[i + j + k];
                if (!invert) {
                    a[j + k] = v;
                    a[i + j + k] = u + v;
                    if (a[i + j + k] >= p) a[i + j + k] -= p;
                }
                else {
                    a[j + k] = v - u;
                    if (a[j + k] < 0) a[j + k] += p;
                    a[i + j + k] = u;
                }
            }
        }
    }
}

```

7.6 Gaussian Elimination

```

const double EPS = 1e-9;
const int INF = 2; // it doesn't actually have to be
                    // infinity or a big number

int gauss (vector < vector<double> > a, vector<double> & ans
) {
    int n = (int) a.size();
    int m = (int) a[0].size() - 1;

    vector<int> where (m, -1);
    for (int col=0, row=0; col<m && row<n; ++col) {
        int sel = row;
        for (int i=row; i<n; ++i)
            if (abs (a[i][col]) > abs (a[sel][col]))
                sel = i;
        if (abs (a[sel][col]) < EPS)
            continue;
        for (int i=col; i<=m; ++i)
            swap (a[sel][i], a[row][i]);
        where[col] = row;

        for (int i=0; i<n; ++i)
            if (i != row) {
                double c = a[i][col] / a[row][col];
                for (int j=col; j<=m; ++j)
                    a[i][j] -= a[row][j] * c;
            }
    }
}

```

```

        ++row;
    }

    ans.assign (m, 0);
    for (int i=0; i<m; ++i)
        if (where[i] != -1)
            ans[i] = a[where[i]][m] / a[where[i]][i];
    for (int i=0; i<n; ++i) {
        double sum = 0;
        for (int j=0; j<m; ++j)
            sum += ans[j] * a[i][j];
        if (abs (sum - a[i][m]) > EPS)
            return 0;
    }

    for (int i=0; i<m; ++i)
        if (where[i] == -1)
            return INF;
    return 1;
}

```

7.7 Karatsuba

```

void convo(int a[], int b[], int res[], int h1, int h2, int
n){
    if (n <= 8){
        for (int i = h1; i < h1 + n; i++) {
            for (int j = h1; j < h1 + n; j++) {
                add(res[i + j], prod(a[i], b[j]));
            }
        }
    }
    else {
        const int mid = n >> 1;
        int atmp[mid], btmp[mid], E[n + 1];
        memset(E, 0, sizeof E);
        for(int i = h1; i < h1 + mid; ++i){
            atmp[i - h1] = sum(a[i], a[i + mid]);
            btmp[i - h1] = sum(b[i], b[i + mid]);
        }

        convo(atmp, btmp, E, 0, 0, mid);
        convo(a, b, res, h1, h2, mid);
        convo(a, b, res, h1 + mid, h2 + n, mid);

        for(int i = h2; i < h2 + mid; ++i){
            const int tmp = res[i + mid];
            add(res[i + mid], E[i - h2] - res[i] - res[i + 2 * mid]);
            add(res[i + 2 * mid], E[i - h2 + mid] - tmp - res[i + 3 *
mid]);
        }
    }
}

```

```

    }
}

```

7.8 Lagrange

```

long long n,k,a[maxn],fac[maxn],ifac[maxn],prf[maxn],suf[
    maxn];
void build()
{
    fac[0] = ifac[0] = 1;
    for (int i = 1; i < maxn; i++)
    {
        fac[i] = fac[i - 1] * i % MOD;
        ifac[i] = binPow(fac[i], MOD - 2);
    }
    //Calculate P(x) of degree k - 1, k values form 1 to k
    //P(i) = a[i]
    long long calc(long long x, long long k)
    {
        if(x <= k)
        {
            return a[x];
        }
        prf[0] = suf[k + 1] = 1;
        for (long long i = 1; i <= k; i++) {
            prf[i] = prf[i - 1] * (x - i + MOD) % MOD;
        }
        for (long long i = k; i >= 1; i--) {
            suf[i] = suf[i + 1] * (x - i + MOD) % MOD;
        }
        long long res = 0;
        for (long long i = 1; i <= k; i++) {
            if (((k - i) & 1)) {
                res = (res + prf[i - 1] * suf[i + 1] % MOD
                    * ifac[i - 1] % MOD * ifac[k - i] % MOD
                    * a[i]) % MOD;
            }
            else {
                res = (res - prf[i - 1] * suf[i + 1] % MOD
                    * ifac[i - 1] % MOD * ifac[k - i] % MOD
                    * a[i] % MOD + MOD) % MOD;
            }
        }
        return res;
    }
}
void solve()
{

```

```

    cin >> n >> k;
    build();
    for(int i = 1; i <= k+2; i++)
        a[i] = (a[i-1]+binPow(i,k))%MOD;
    cout << calc(n,k+2);
}

```

7.9 Pisano

```

bool check(int r, int p)
{
    if (p == 1) return false;
    swap(m, r);
    matrix q; q.build1();
    matrix dv = q;
    q = bpow(q, p + 1);
    for(int i = 0; i < 2; ++i){
        for(int j = 0; j < 2; ++j){
            if (q(i, j) != dv(i, j)) {swap(r, m); return
                false;}
        }
    }
    swap(r, m);
    return true;
}

int pisano(int m)
{
    int p = 1;
    int tmp = m;
    for (int i : prime_factors(m)) {
        int expo = 0, fact = 1;
        for (; tmp % i == 0; ++expo, fact *= i)
            tmp /= i;

        int q = 1;
        if (i == 2) {
            q = (int) fact / 2 * 3;
        }
        else if (i == 5) {
            q = (int) fact * 4;
        }
        else {
            vector<int> cands;
            if (i % 10 == 1 || i % 10 == 9)
                cands = factorize(i - 1);
            else
                cands = factorize(2 * (i + 1));

```

```

            for (int x : cands)
                if (check(i, x)) {
                    q = x;
                    break;
                }
            q = (int) fact / i * q;
        }
        p = lcm(p, q);
    }
    return p;
}

```

7.10 Rho and Miller Rabin

```

mt19937 rng(chrono::steady_clock::now().time_since_epoch().
    count());
vector<int> b;
ll binpower(ll base, ll e, ll mod) {
    ll result = 1;
    base %= mod;
    while (e){
        if (e & 1)result = (ll)result * base % mod;
        base = (ll)base * base % mod;e >>= 1;
    }
    return result;
}

bool check_composite(ll n, ll a, ll d, int s){
    ll x = binpower(a, d, n);
    if (x == 1 or x == n - 1)return false;
    for (int r = 1; r < s; r++) {
        x = (ll)x * x % n;
        if (x == n - 1)return false;
    }return true;
}

bool MillerRabin_checkprime(ll n) {
    if (n < 2)return false;
    int r = 0;
    ll d = n - 1;
    while ((d & 1) == 0) {
        d >>= 1;
        r++;
    }
    for (int a : {2, 3, 5, 7, 11, 13, 17, 19, 23, 29, 31,
        37}) {
        if (n == a)return true;
        if (check_composite(n, a, d, r))return false;
    }
    return true;
}

```

```

}

int mult(int a, int b, int mod) {
    return (ll)a * b % mod;
}

int F(int x, int c, int mod) {
    return (mult(x, x, mod) + c) % mod;
}

int rho(int n, int x0=2, int c=1) {
    int x = x0;
    int y = x0;
    int g = 1;
    if (n % 2 == 0) return 2;
    while (g == 1) {
        x = F(x, c, n);
        y = F(y, c, n);
        y = F(y, c, n);
        g = gcd(abs(x - y), n);
    }
    return g;
}

void Rho_factorization(int n){
    set<int> s;
    if (n == 1) {return;}
    while (!MillerRabin_checkprime(n)){
        ll k;
        while (1){
            int p = (rng()%(n-2))+2, q = (rng()%(n-1))+1;
            k = rho(n,p,q);
            if (MillerRabin_checkprime(k)) break;
        }
        s.insert(k);
        n/=k;
    }
    s.insert(n);
    fora(i,s) b.pb(i);
}

```

7.11 Tonelli Shanks

```

mt19937 rng(std::chrono::system_clock::now().
    time_since_epoch().count());
// sqrt mod p
int Tonelli_Shanks(int n, int p){
    int q = p - 1, s = 0;
    while (q % 2 == 0){
        ++ s,

```

```

        q >>= 1;
    }
    int z;
    do {
        z = rng() % p;
    } while(binpow(z, (p-1)/2, p) != p-1);

    int m = s,
        c = binpow(z, q, p),
        t = binpow(n, q, p),
        r = binpow(n, (q+1)/2, p);

    while (t > 1) {
        i = 0;
        do {
            i ++;
            t = 1ll * t * t % p;
        } while (t != 1);

        if (t == m)
            return -1;

        int b = binpow(c, (1<<(m-i-1)), p);
        m = i;
        c = binpow(b, 2, p);
        t = 1ll * t * c % p;
        r = 1ll * r * b;
    }

    if (t == 0)
        return 0;
    else
        return r;
}

```

8 Number Theory

8.1 Bézout's identity

For $a \neq 0$, $b \neq 0$, then $d = \gcd(a, b)$ is the smallest positive integer for which there are integer solutions to

$$ax + by = d$$

If (x, y) is one solution, then all solutions are given by

$$\left(x + \frac{kb}{\gcd(a, b)}, y - \frac{ka}{\gcd(a, b)}\right), \quad k \in \mathbb{Z}$$

8.2 Pythagorean Triples

The Pythagorean triples are uniquely generated by

$$a = k \cdot (m^2 - n^2), \quad b = k \cdot (2mn), \quad c = k \cdot (m^2 + n^2),$$

with $m > n > 0$, $k > 0$, $m \perp n$, and either m or n even.

8.3 Primes

$p = 962592769$ is such that $2^{21} \mid p - 1$, which may be useful. For hashing use 970592641 (31-bit number), 31443539979727 (45-bit), 3006703054056749 (52-bit). There are 78498 primes less than 1 000 000.

Primitive roots exist modulo any prime power p^a , except for $p = 2, a > 2$, and there are $\phi(\phi(p^a))$ many. For $p = 2, a > 2$, the group $\mathbb{Z}_{2^a}^\times$ is instead isomorphic to $\mathbb{Z}_2 \times \mathbb{Z}_{2^{a-2}}$.

8.4 Estimates

$$\sum_{d|n} d = O(n \log \log n).$$

The number of divisors of n is at most around 100 for $n < 5e4$, 500 for $n < 1e7$, 2000 for $n < 1e10$, 200 000 for $n < 1e19$.

8.5 Mobius Function

$$\mu(n) = \begin{cases} 0 & n \text{ is not square free} \\ 1 & n \text{ has even number of prime factors} \\ -1 & n \text{ has odd number of prime factors} \end{cases}$$

Mobius Inversion:

$$g(n) = \sum_{d|n} f(d) \Leftrightarrow f(n) = \sum_{d|n} \mu(d)g(n/d)$$

Other useful formulas/forms:

$$\begin{aligned} \sum_{d|n} \mu(d) &= [n = 1] \text{ (very useful)} \\ g(n) &= \sum_{n|d} f(d) \Leftrightarrow f(n) = \sum_{n|d} \mu(d/n)g(d) \\ g(n) &= \sum_{1 \leq m \leq n} f\left(\left\lfloor \frac{n}{m} \right\rfloor\right) \Leftrightarrow f(n) = \sum_{1 \leq m \leq n} \mu(m)g\left(\left\lfloor \frac{n}{m} \right\rfloor\right) \end{aligned}$$

9 String

9.1 Aho Corasick

```
const int K = 26;

struct Vertex {
    int next[K];
    bool leaf = false;
    int p = -1;
    char pch;
    int link = -1;
    int go[K];

    Vertex(int p=-1, char ch='$') : p(p), pch(ch) {
        fill(begin(next), end(next), -1);
        fill(begin(go), end(go), -1);
    }
};

vector<Vertex> t(1);

void add_string(string const& s) {
    int v = 0;
    for (char ch : s) {
        int c = ch - 'a';
        if (t[v].next[c] == -1) {
            t[v].next[c] = t.size();
            t.emplace_back(v, ch);
        }
        v = t[v].next[c];
    }
    t[v].leaf = true;
}

int go(int v, char ch);

int get_link(int v) {
    if (t[v].link == -1) {
        if (v == 0 || t[v].p == 0)
            t[v].link = 0;
        else
            t[v].link = go(get_link(t[v].p), t[v].pch);
    }
    return t[v].link;
}
```

```
}

int go(int v, char ch) {
    int c = ch - 'a';
    if (t[v].go[c] == -1) {
        if (t[v].next[c] != -1)
            t[v].go[c] = t[v].next[c];
        else
            t[v].go[c] = v == 0 ? 0 : go(get_link(v), ch);
    }
    return t[v].go[c];
}
```

9.2 Suffix Array

```
string s;
vector<int> p(400007), c(400007), lcp(400007);
int n, k;

void build(int n){
    vector<pair<int,int>> a(n);
    for(int i = 0; i < n; ++i) a[i] = {s[i], i};
    sort(a.begin(), a.end());
    for(int i = 0; i < n; ++i) p[i] = a[i].y;
    c[p[0]] = 0;
    for(int i = 1; i < n; ++i){
        c[p[i]] = c[p[i-1]];
        if (a[i].x != a[i-1].x) c[p[i]] += 1;
    }
    k = 0;
    while ((1 << k) < n){
        vector<pair<pair<int,int>,int>> a(n);
        for(int i = 0; i < n; ++i){
            a[i] = {{c[i], c[(i + (1 << k)) % n]}, i};
        }
        //Radix sort
        vector<int> cnt(n);
        for(auto i : a){
            cnt[i.x.y]++;
        }
        vector<pair<pair<int,int>,int>> b(n);
        vector<int> pos(n);
        pos[0] = 0;
        for(int i = 1; i < n; ++i) pos[i] = pos[i-1] + cnt[i-1];
    }
```

```
for(auto i : a){
    b[pos[i.x.y]] = i;
    pos[i.x.y]++;
}
a=b;
//////////
vector<int> cnt2(n);
for(auto i : a){
    cnt2[i.x.x]++;
}
vector<pair<pair<int,int>,int>> f(n);
vector<int> pos2(n);
pos2[0] = 0;
for(int i = 1; i < n; ++i) pos2[i] = pos2[i-1] + cnt2[i-1];
for(auto i : a){
    f[pos2[i.x.x]] = i;
    pos2[i.x.x]++;
}
a = f;
//////////
for(int i = 0; i < n; ++i) p[i] = a[i].y;
c[p[0]] = 0;
for(int i = 1; i < n; ++i){
    c[p[i]] = c[p[i-1]];
    if (a[i].x != a[i-1].x) c[p[i]]++;
}
k++;
}

void buildlcp(int n){
    k=0;
    for(i = 0; i < n - 1; ++i){
        k=max(0,k - 1);
        lcp[c[i]] = k;
        int s1 = i, s2 = p[c[i]-1];
        for(int j = k; j <= n - i + 1; ++j){
            if (s[s1 + j] == s[s2 + j]){
                k++;
                lcp[c[i]] = k;
            } else break;
        }
    }
}
```