

Trying to prove it on stainless

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Outline

LLL Algorithm

Formally verifying the LLL

Yet to prove

Conclusion

LLL Algorithm ——

Gram Schmidt

Algorithm 1 Gram-Schmidt

```
Require: v_0, \dots, v_{n-1}
Ensure: span\{v_0, \dots, v_{n-1}\} = span\{g_0, \dots, g_{n-1}\} and for i \neq j, \langle g_i, g_i \rangle = 0
 1: g_0 := v_0
 2: for i = 1 to n do
 3: g_i := b_i
 4: for j = 0 to i - 1 do
 5: \mu_{i,j} := \langle v_i, g_i \rangle / \langle g_i, g_j \rangle
 6:
       g_i := g_i - \mu_{i,i}g_i
      end for
 7:
 8: end for
 9: return g_0, \dots, g_n.
```

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α -reduced basis

Definition (α -reduced basis)

Given $\mathbf{B} = \{f_1, \dots, f_n\}$ such that $f_i \in \mathbb{Z}^m$ and Gram Schmidt basis $\mathbf{G} = \{g_1, \dots, g_n\}$. We say that it is α -reduced if:

- (size reduced) $1 \le j < i \le n$: $|\mu_{i,j}| \le 1/2$, with $\mu_{i,j}$ defined earlier
- (Lovàsz condition) For k = 1, ..., n-1: $||g_{k-1}||^2 \le \alpha ||g_k||^2$ where $\alpha \ge 1$.



Algorithm 2 The LLL Basis reduction algorithm

```
Require: A list of linearly independent vectors f_0,...,f_{n-1} \in \mathbb{Z}^m and \alpha > \frac{4}{3}
Ensure: a basis for the same lattice as f_0,...,f_{n-1} that is reduced w.r.t \alpha
1: i := 0
2: g_0, \dots, g_{n-1} := \mathsf{Gram}\text{-}\mathsf{Schmidt}(f_0, \dots, f_{n-1})
3: while i < n do
4:
           for i = i - 1 downto 0 do
5:
                 \mu_{i,j} := \langle f_i, g_i \rangle / \langle g_i, g_i \rangle
6:
                f_i := f_i - \lceil \mu_{i,j} \rfloor f_i
7:
8:
9:
10:
           end for
           if i = 0 or ||g_{i-1}||^2 \le \alpha ||g_i||^2 then
                 i := i + 1
             else
11:
                  h := g_{i-1}
12:
         g_{i-1} := g_i + \mu_{i,i-1}g_{i-1}
13:
          \mu := \langle h, g_{i-1} \rangle / \langle g_{i-1}, g_{i-1} \rangle
14:
              g_i := h - \mu g_{i-1}
15:
                   (i, f_{i-1}, f_i) := (i-1, f_i, f_{i-1})
16:
             end if
17: end while
18: return f_0,...,f_{n-1}.
```

Example

$$\begin{bmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 1 \\ 1048576 & 3294200 & 10349040 & 1724840 \end{bmatrix} \xrightarrow{LLL} \begin{bmatrix} 0 & -59 & -43 & -97 \\ 0 & -54 & 65 & 78 \\ -1 & 23 & -16 & -15 \\ 6 & 1 & -2 & 0 \\ 0 & -24 & -88 & 128 \end{bmatrix}$$

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Formally verifying the LLL

The properties, we want to verify

- Lattice preservation
- Lovàsz condition : For k = 1, ..., n-1: $||g_{k-1}||^2 \le \alpha ||g_k||^2$ where $\alpha \ge 1$. This means that vectors in the basis are not much shorter than the previous ones.
- Nearly orthogonal : $1 \le j < i \le n$: $|\mu_{i,j}| \le 1/2$. This means that f_i is component-wise the closest vector in the lattice to g_i , meaning that it is nearly orthogonal.

But we will have to assume some linear algebra properties, otherwise it will be too complicated.

Lattice preservation

```
case class Lattice(rows: BigInt, cols: BigInt, vectors: List[IntegralVector], equivalentLattice: List[IntegralVector]]

require(vectors.size==rows)
require(vectors.forall(v => v.size==cols))
require(vectors.forall(v => v.size==cols))
require(vectors.forall(v => v.size==cols))

def sum_rows(b]Idx: BigInt, alpha: BigInt, bIIdx: BigInt):Lattice = {-
}.ensuring(res => res.rows==this.rows && res.cols==this.cols)

def swap(x:BigInt, y:BigInt):Lattice = {-
}.ensuring(res => res.rows==this.rows && res.cols==this.cols)
```

Lovàsz condition

```
def lovaz_condition(vectors : List[Vector], until : BigInt, alpha : Rational, size : BigInt):Boolean={
    decreases(max(until, 0))

if(!(until<=vectors.size && until>=1) || !(vectors.forall(v => v.size==size))){
    return false
}

if(until==1){
    true
    }
}
lif(until==1) {
    true
    val g_k = vectors.head
    val g_k 1 = vectors.tail.head
    g_k.norm() <= alpha*g_k_l.norm() && lovaz_condition(vectors.tail, until-1, alpha, size)
}
}</pre>
```

Nearly orthogonal

```
def nearly orthogonal(lattice : List[IntegralVector], gramBasis : List[Vector], i : BigInt, j : BigInt, size : BigInt):Boolean = {
   decreases(max(0, i*i-i))
    if(i>lattice.size || i<1 || j<1 || j>i || (i==lattice.size && i!=j)){
     if(gramBasis.size!=lattice.size || size<1 || !(lattice.forall(v => v.size==size)) || !(gramBasis.forall(v => v.siz
     if(i==1 && i==1){
         val gramVect = get(i, gramBasis)
         conserveNorm(i, gramBasis, size)
          conserveSize(j, gramBasis, size)
          conserveSizeIntegral(i, lattice, size)
         val mu = lattice(i).dotProduct(gramVect)/gramVect.norm()
          abs(mu) <= Rational(1.2) && nearly orthogonal(lattice, gramBasis, i, i+1, size)
         assert(i==j && i>1)
         val gramVect = get(0, gramBasis)
          conserveNorm(0, gramBasis, size)
          conserveSize(0, gramBasis, size)
          conserveSizeIntegral(i-1, lattice, size)
         val mu = lattice(i-1).dotProduct(gramVect)/gramVect.norm()
          abs(mu) <= Rational(1.2) && nearly orthogonal(lattice, gramBasis, i-1, 1, size)
```

The main Loop

```
def reduced(lattice : Lattice. alpha : Rational. index : BigInt): Boolean =
 (index==0 ||
   (lovaz condition(lattice.gramSchmidt basis(), index, alpha, lattice.cols)
   && nearly orthogonal(lattice.vectors. lattice.gramSchmidt basis(), index. lattice.cols)))
def LLL(lattice : Lattice, alpha : Rational): Lattice = {
 require(alpha > Rational(4,3))
 def LLLloop(index : BigInt, lattice : Lattice): Lattice = {
   require(index>=0 && index<=lattice.rows)
   require(reduced(lattice, alpha, index))
   if(index==lattice.rows){
   }else if(index==0){
     LLLloop(index+1, lattice)
     val aprox0rtho = LLLAprox(index. index. lattice.lattice.gramSchmidt basis())
     if(lovaz condition(aproxOrtho.gramSchmidt basis(), index+1, alpha, aproxOrtho.cols)){
       LLLloop(index+1 aprox0rtho)
       val swapped = swap(index, lattice)
       assert(reduced(swapped, alpha, index-1))
       LLLloop(index-1, swapped)
 }.ensuring(res => res.sameLattice(lattice) && reduced(res, alpha, res.rows))
 LLLloop(0, lattice)
.ensuring(res => res.sameLattice(lattice) && reduced(res. alpha. res.rows))
```

Proving swap branch respects invariant

```
def swap(index : BigInt, lattice : Lattice): Lattice = {
 require(index>=1 && index<lattice.rows)
 require(reduced(lattice, alpha, index))
 val swapped = lattice.swap(index-1,index)
 if(index==1){
   swapped
   lovaz condition take(lattice.gramSchmidt basis(), index. index-1, alpha, lattice.cols)
   lattice.swap conserve gramschmit(index-1, index)
   lovaz condition append(swapped.gramSchmidt basis(), index-1, alpha, swapped.cols)
   nearly orthogonal subset(lattice.vectors, lattice.gramSchmidt basis(), lattice.cols, index, index, index-1, index-1)
   nearly orthogonal take(lattice.vectors, lattice.gramSchmidt basis(), lattice.cols, index. index-1, index-1)
   lattice.swap conserve lattice(index-1, index)
   nearly orthogonal append(swapped.vectors, swapped.gramSchmidt basis(), swapped.cols, index-1, index-1, index-1)
    swapped
}.ensuring(res => res.sameLattice(lattice) && reduced(res, alpha, index-1) )
def swap conserve gramschmit(x:BigInt, v:BigInt): Unit = {
 require(x>=0 && x<this.rows)
 require(v>=0 && v<this.rows)
}.ensuring( => take vector(x, this.gramSchmidt basis(), cols)==take vector(x, this.swap(x,y).gramSchmidt basis(), cols))
def swap conserve lattice(x:BigInt, y:BigInt): Unit = {
 require(x>=0 && x<this.rows)
 require(v>=0 && v<this.rows)
 ensuring => take integral (x, this, vectors, cols) == take integral <math>(x, this, swap(x, y), vectors, cols)
```

Proving LLLAprox preserve invariant

```
def LLLAprox(i : BigInt. i: BigInt. lattice : Lattice. gramBasis : List[Vector]):Lattice = {
 require(i>=1 && i<lattice.rows && i<=i && i>=0)
  require(lattice.rows==gramBasis.size && gramBasis.forall(v => v.size==lattice.cols) && gramBasis.forall(v => v.nonZero()))
  require(lovaz condition(lattice.gramSchmidt basis(), i, alpha, lattice.cols)
  require(nearly orthogonal(lattice.vectors, gramBasis, i (i, j), j (i, j), lattice.cols))
    val g = get(j-1, gramBasis)
    conserveSize(j-1, gramBasis, lattice.cols)
    conserveNorm(j-1, gramBasis, lattice.cols)
    val mu = lattice(i).dotProduct(g)/g.norm()
    lattice.sum conserve gramschmit basis(i-1, -round(mu), i)
    val updated lattice = lattice.sum rows(i-1, -round(mu), i)
    assert(abs(upadted mu) <= Rational(1.2))
    assert(nearly orthogonal(updated lattice.vectors, gramBasis, i (i, j), j (i, j), lattice.cols))
    LLLAprox(i, j-1, updated lattice, gramBasis)
} ensuring(res => res.rows==lattice.rows && res.cols==lattice.cols && res.sameLattice(lattice) && lovaz condition(res.gramSchmidt basis(), i, alpha, res.cols)
&& nearly orthogonal(res.vectors, gramBasis, i+1, i+1, res.cols))
def sum conserve gramschmit basis(bildx : BigInt, alpha : BigInt, bildx : BigInt):Unit={
 require(bildx<bildx)
}.ensuring( => this.sum rows(bildx, alpha, bildx).gramSchmidt basis()==this.gramSchmidt basis())
```

Yet to prove

Measure[1]

First, given B a basis of vectors in \mathbb{Z}^m and induced Gram-Schmidt basis G, we have that

$$\delta(G) = \delta(B) = \det(B^{\top}B) = \prod_{i \in n} \|g_i\|^2 \in \mathbb{N}$$

$$\Delta(G) = \Delta(B) = \prod_{i=0}^{n-1} \delta(B_i) = \prod_{i=0}^{n-1} \prod_{j=0}^{i-1} \|g_j\|^2 = \prod_{i=0}^{n-1} \|g_i\|^{2(n-i)} \in \mathbb{N}$$

Measure[2]

In our project, the LLL is handled by the function

```
LLLloop(index : BigInt, lattice : Lattice)
gramBasis ← GramSchmidt(lattice)
```

The measure that decrease, proving that the LLL finishes is given by

```
decreases(\Delta(gramBasis), n-index)
```

Measure[3]

Why does this decreases ?

- If we do not switch, then gramBasis remains unchanged and Index is increased by 1.
- If we switch, then the gramBasis is updated as follows. Given g_i, g_{i+1} , we have that:

$$g_i^* = g_{i+1} + \mu_{i+1,i}g_i$$

$$g_{i+1}^* = g_i - \langle g_i, g_i^* \rangle / \langle g_i^*, g_i^* \rangle g_i^*$$

$$\mathsf{gramBasis}^* \leftarrow \mathsf{gramBasis}[:i-1] + + [g_i^*, g_{i+1}^*] + + \mathsf{gramBasis}[i+2:]$$

Measure[4]

As $\delta(\text{gramBasis}^*) = \delta(\text{gramBasis})$, We have that

$$\|g_i^*\|^2 \|g_{i+1}^*\|^2 = \|g_i\|^2 \|g_{i+1}\|^2$$

Furthermore,

$$\|g_{i}\|^{2} > \alpha \|g_{i+1}\|^{2}$$

$$\iff \alpha^{-1} \|g_{i}\|^{2} > \|g_{i+1}\|^{2}$$

$$\iff \alpha^{-1} \|g_{i}\|^{2} + \mu_{i+1,i}^{2} \|g_{i}\|^{2} > \|g_{i+1}\|^{2} + \mu_{i+1,i}^{2} \|g_{i}\|^{2}$$

$$\iff (\alpha^{-1} + \frac{1}{4}) \|g_{i}\|^{2} > \|g_{i+1} + \mu_{i+1,i}g_{i}\|^{2}$$

$$\iff \|g_{i}\|^{2} > \|g_{i}^{*}\|^{2}$$

$$\iff \|g_{i}^{*}\|^{4} \|g_{i+1}^{*}\|^{2} < (\alpha^{-1} + \frac{1}{4}) \|g_{i}\|^{4} \|g_{i+1}\|^{2}$$

$$\Delta(\operatorname{gramBasis}^{*}) < (\alpha^{-1} + \frac{1}{4}) \Delta(\operatorname{gramBasis})$$

Conclusion

We have been a bit too zealous with the whole LLL:

- Based on many linear algebra properties, so we could not do it from the ground up.
- We have significantly improved our stainless understanding.

Thank you for your attention!

 ${\color{red} *}$ for references, please refer to our report