Efficient Planning for Sokoban with Answer Set Programming - A Case Study

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Abstract

Sokoban approach with ASP planning → PlainEnc, IncEnc, Optimization (?)

1 Introduction

TODO: REWORK THIS What do i do? I'm working on the Sokoban problem and trying to do my best to find a fast and easy solution to it. I'm not working on any level but just the ones meeting some criteria. I'm picking qualitative and quantitative characteristics and deciding on some values to examine. My encodings are supposed to work better on some of the instances and worse

on others. What rules describe those fluctuations? What assumptions can be made? Is there a trend when comparing runtimes to my qualitative and quantitative characteristics? E.g. you would expect the runtime to go up when the instance's size grows or when the amount of boxes in the level grows. In this thesis I will work on optimizing ASP encodings for the game Sokoban towards different instances. First Ill briefly explain what Sokoban is and how it is played, then Ill give a short introduction on comparable/similar/other scientific studies concerning sokoban. Then Ill go over my work, ... and finally evaluation and Discussion of Results.

2 What is Sokoban

2.1 Background

Sokoban is a logistic planning puzzle game from Japan. Sokoban was apparently created by Hiroyuki Imabayashi in 1981, and published by *Thinking Rabbit* in 1982. Sokoban means warehouse person in japanese. The game has very simple rules but shows to be very complex for humans and for AI. This makes Sokoban a very interesting subject for research regarding planning problems.

2.2 Rules

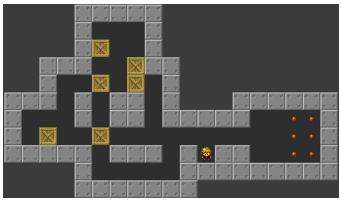


Fig.1: First of the Classical Sokoban Levels

Figure 1 illustrates the first Sokoban level of the original set of levels. The game is played on a board that can be modelled as a two-dimensional grid. The squares can either be floor or wall squares. Some floor squares are marked as goals and the same number of floor squares contain boxes. The player pushes boxes around this modelled warehouse, trying to cover the marked goal fields (targets) with one of the boxes. For that the player can move one step at a time into one of the four directions up, down, left and right. When the player moves onto a box's position the box is pushed one field into the direction of the move. There can never be multiple boxes on one field, neither can there be a box where there is no field. A solution to a level is a sequence of moves that leaves a box

on each target. In many of the levels created to this date, the warehouse has a lot of small passageways and very intricate designs. In this thesis I will look at more open level designs, e.g. rectangular levels. Sokoban is a very difficult planning problem for multiple reasons. There are a lot of possibilities to create a deadlock (a gamestate from which no solution can be found), the branching factor is large and the optimal solution can be very long [5].

3 State of the Art

Gordon Wilfong published 1991 that all motion planning problems with moving obstacles are NP-hard [12] by reducing 3Sat to such a problem. He didnt answer the question whether Sokoban is also PSPACE-complete. That was answered by Joseph C Culberson [9] in 1997. He showed that it is possible to emulate a linear bounded automaton by constructing a level. The solution to this constructed level, if there is any, will have a length in $\Theta(|w|+t(|w|))$, where w is the input and t(|w|) is the amount of transitions the LBA made. This shows PSPACE-completeness

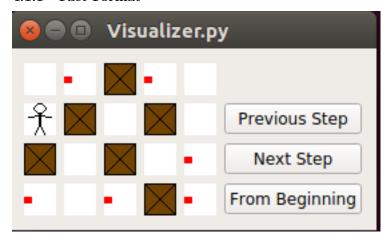
Heuristics doesnt seem to do a great deal for sokoban [5]. It helped a lot with chess e.g., but doesnt cut it for sokoban. Approach the hard domain of sokoban in another way. Macromoves a way that really works well and is used in practice, also deadlock detection (WHAT IS A DEADLOCK - stone in corner and no goal -¿ no solutions) (obviously best if you know the specific level). Reference Rolling stone (source not available (?)). Furthermore Sokoban can be abstracted into rooms and tunnels, making it into subproblems which should be easier to solve, tunnels for example, are mostly immediately solved. Similarly you can also use relevance cutting. This means creating a measurement for relevance in between moves and picking moves that are relevant to prior moves (few exceptions in special cases). Backwards Approach and LexOrdering. Conclusion

4 My Work

what did i do? what is my baseline? what changes from encoding to encoding which of the encodings works best in which kind of instance how do i evaluate that show my experiments and their results

4.1 Encodings

4.1.1 Fact Format



```
1 \# const horizon = 23.
2 \# const dim X = 5.
3 \# const dim Y = 4.
4 init (field (1..5,1..4)).
5 init (at (0,1,2)).
6 init(at(1,1,3)).
7 init (at (2,2,2)).
8 init(at(3,3,1)).
9 init(at(4,3,3)).
10 init (at(5,4,2)).
11 init(at(6,4,4)).
12 init (target (1,4)).
13 init (target(2,1)).
14 init(target(3,4)).
15 init (target(4,1)).
16 init(target(5,3)).
17 init (target (5,4)).
```

4.1.2 Plain Encoding

```
\begin{array}{c} 1 \  \, \mathrm{time}\,(\,0\,..\,\mathrm{horizon}\,)\,.\\ 2\\ 3 \  \, \mathrm{dir}\,(\,1\,,0\,;0\,,1\,;-\,1\,,0\,;0\,,-\,1\,)\,.\\ 4 \  \, \mathrm{at}\,(N,X,Y,0)\  \, :-\  \, \mathrm{init}\,(\,\mathrm{at}\,(N,X,Y)\,)\,.\\ 5\\ 6 \  \, \% \  \, \mathrm{pick} \  \, \mathrm{one} \  \, \mathrm{direction} \  \, \mathrm{to} \  \, \mathrm{move} \  \, \mathrm{in}\\ 7 \  \, 1 \  \, \{\  \, \mathrm{move}(\mathrm{DX},\mathrm{DY},\mathrm{T})\  \, :-\  \, \mathrm{dir}\,(\mathrm{DX},\mathrm{DY})\  \, \}\  \, 1\  \, :-\  \, \mathrm{time}(\mathrm{T})\,.\\ 8 \  \, \% \  \, \mathrm{move} \  \, \mathrm{the} \  \, \mathrm{player}\\ 9 \  \, \mathrm{at}\,(\,0\,,X\!\!+\!\!\mathrm{DX},Y\!\!+\!\!\mathrm{DY},\mathrm{T})\  \, :-\  \, \mathrm{at}\,(\,0\,,X,Y,T\!\!-\!1)\,,\  \, \mathrm{move}\,(\mathrm{DX},\mathrm{DY},\mathrm{T})\,,\  \, \mathrm{time}(\mathrm{T})\,.\\ 10 \  \, \% \  \, \mathrm{move} \  \, \mathrm{box}\,,\  \, \mathrm{the} \  \, \mathrm{player}\  \, \mathrm{collides}\  \, \mathrm{with}\\ 11 \  \, \mathrm{at}\,(\,N,X\!\!+\!\!\mathrm{DX},Y\!\!+\!\!\mathrm{DY},\mathrm{T})\  \, :-\  \, \mathrm{at}\,(\,N,X,Y,T\!\!-\!1)\,,\  \, \mathrm{at}\,(\,0\,,X,Y,\mathrm{T})\,,\  \, \mathrm{move}(\mathrm{DX},\mathrm{DY},\mathrm{T})\,,\\ 12 \  \, \mathrm{time}\,(\mathrm{T})\,,\  \, N\!\!>\!0\,.\\ 13 \  \, \% \  \, \mathrm{leave}\  \, \mathrm{boxes}\  \, \mathrm{that}\  \, \mathrm{aren}\,\, {}^{\prime}\,\mathrm{t}\,\, \mathrm{touched}\,\, \mathrm{at}\,\, \mathrm{their}\,\, \mathrm{place}\,\, \mathrm{with}\,\, t\!\!+\!\!1\\ 14 \  \, \mathrm{at}\,(\,N,X,Y,\mathrm{T})\,  \, :-\  \, \mathrm{at}\,(\,N,X,Y,T\!\!-\!1)\,,\  \, \mathrm{not}\,\, \mathrm{at}\,(\,0\,,X,Y,\mathrm{T})\,,\  \, N\!\!>\!\!0\,,\  \, \mathrm{time}\,(\mathrm{T})\,. \end{array}
```

```
15 % there can't be boxes, players, where there is no field 16 :- not init(field(X,Y)), at(_,X,Y,_). 17 % there can't be two boxes/players, at the same spot 18 :- at(N,X,Y,T), at(M,X,Y,T), N>M, time(T). 19 20 uncvrd(X,Y,T) :- init(target(X,Y)), not at(_,X,Y,T), time(T). 21 goal(X,Y,T) :- init(target(X,Y)), not uncvrd(X,Y,T), 22 not at(0,X,Y,T), time(T). 23 :- not goal(X,Y,horizon), init(target(X,Y)). 24 25 #show at/4.
```

The first try to construct a naive encoding to solve sokoban-instances yielded a solution that works in a common way in ASP planning. In line 4 we initialize the information from the level instance with the time 0. In line 7 we state that there is one move at minimum and at maximum for each timestep. This makes the ASP planner ground all sequences of moves, pick one for every timestep and infer every successive state from the moves it picks. In lines 9-13 we calculate the successor state from the move picked in line 7. Line 9 calculates the new player-position, line 11 calculates the new position of the box in the players way and line 13 calculates the new position of boxes not touched by the player. Additionally we need some integrity constraints to prevent illegal states. The integrity constraint in line 15 ensures that at no time a box or player has a position outside of the playing field. The integrity constraint in line 17 ensures that at no time boxes or the player share one field. In the lines 19-21 we define our win-condition to be that all targets in the instance have to be covered by boxes. The Plain Encoding has the predicates, time/1, dir/2, at/4, move/3, uncovered/3, goal/3 The encoding works in a common way in ASP planning, by guessing the sequence of actions and inferring the successor states by effect axioms. There is a time/1 predicate for every timestep and a at/4 predicate for every box and the player at every timestep.

4.1.3 Incremental Encoding

```
1 #include <incmode>.
3 #program base.
4 \ \mathrm{dir} \left( \ 1 \, , 0 \ ; \ 0 \, , 1 \ ; \ -1 \, , 0 \ ; \ 0 \, , -1 \ \right).
  at(N,X,Y,0): - init(at(N,X,Y)).
7 #program step(t).
8 % pick one direction to move in
9 1 { move(DX,DY,t) : dir(DX,DY) } 1.
10 % move the player
11 at (0, X+DX, Y+DY, t) :- at (0, X, Y, t-1), move (DX, DY, t), init (field (X+DX, Y+DY)).
12 % move box, the player collides with
13 at (N,X+DX,Y+DY,t) :- at (N,X,Y,t-1), at (0,X,Y,t), move (DX,DY,t), N>0.
14 \% leave boxes that aren't touched at their place with t+1
15 at (N, X, Y, t): - at (N, X, Y, t - 1), not at (0, X, Y, t), N > 0.
16\ \% there can't be boxes, players, where there is no field
17 :- at(-,X,Y,-), not init(field(X,Y)).
18 % there can't be two boxes/players, at the same spot
19 :- at (N,X,Y,t), at (M,X,Y,t), N>M.
20
```

```
21 #program check(t).
22 #external query(t).
23 uncovered(X,Y,t) :- init(target(X,Y)), not at(_,X,Y,t), query(t).
24 goal(X,Y,t) :- init(target(X,Y)), not uncovered(X,Y,t), not at(0,X,Y,t), query(t).
25 :- not goal(X,Y,t), init(target(X,Y)), query(t).
26
27 #show at/4.
```

The incremental encoding is similar to the plain encoding. The difference is the partition of the encoding into three parts: the #program base in lines 3-5 which is independent of the step parameter t, the #program step in lines 7-19 which is the cumulative part, collecting knowledge and the #program check in lines 21-25 which is specific for each value of t (STEP ALSO IS THOUGH). This allows for gradually processing each time step and accumulating knowledge that persists even if the current step is not satisfiable.

4.1.4 Advanced Encoding

```
1 time (0..horizon).
 3 dir (1,0;0,1;-1,0;0,-1).
 4 \operatorname{at}(N,X,Y,0) :- \operatorname{init}(\operatorname{at}(N,X,Y)).
 6 forbidden (X,Y): - init (field(X,Y)), not init (target(X,Y)),
 7 3 #sum {3 : not init(field(X+DX,Y+DY)), not init(field(X+DY,Y+DX)), dir(DX,DY);
    3 : not init (field (X+DX,Y+DY)), not init (field (X-DY,Y-DX)), dir (DX,DY);
    1,1: forbidden(X+DX,Y+DY), dir(DX,DY); 1,2: forbidden(X+DX,Y+DY), dir(DX,DY), forbidden(X-D
    \begin{array}{l} 1,3 \ : \ not \ init ( \ field (X+DX,Y+DY)) \, , \ dir (DX,DY); \\ 1,4 \ : \ not \ init ( \ field (X+DX,Y+DY)) \, , \ not \ init ( \ field (X-DX,Y-DY)) \, , \ dir (DX,DY) \}. \end{array}
11
13 % pick one direction to move in
14 1 { move(DX,DY,T) : dir(DX,DY) } 1 :- time(T).
15 % move the player
16 at (0, X+DX, Y+DY, T) :- at (0, X, Y, T-1), move (DX, DY, T), time (T).
17 \% move box, the player collides with
18 \ \ \text{at} \ (N+1,X+DX,Y+DY,T) \ :- \ \ \text{at} \ (N+1,X,Y,T-1) \ , \ \ \text{at} \ (0 \ ,X-DX,Y-DY,t) \ , \ \ \text{at} \ (0 \ ,X,Y,T) \ , \ \ \text{move} \ (DX,DY,T) \ , \ \ \text{time} \ (T)
19 % leave boxes that aren't touched at their place with t+1 20 at(N,X,Y,T) :- at(N,X,Y,T-1), not at(0,X,Y,T), N>0, time(T).
21 % don't take a move back on the very next step
22 :- move(DX,DY,T), move(-DX,-DY,T-1), at(0,X,Y,T), not at(-,X-DX,Y-DY,T-1), time(T).
23~\% there can't be boxes, players, where there is no field
24 :- not init(field(X,Y)), at(_{-},X,Y,_{-}).
25 % there can't be two boxes/players, at the same spot
26 : - \text{ at } (N,X,Y,T) \;, \; \text{ at } (M,X,Y,T) \;, \; N \!\!>\!\! M, \; \text{ time } (T) \;.
27 % don't push a box onto a forbidden field
28 :- \ at (N,X,Y,T) \,, \ forbidden \, (X,Y) \,, \ N>0, \ T>0, \ not \ at \, (N,X,Y,T-1) \,.
29 dirPos(1;-1).
30 :- \ 3 \ \#sum \ \left\{1,1 : \ at\left(A,X\!\!+\!\!DX,Y,T\right), \ A\!\!>\!\! 0 \ ; \ 1,1 : \ not \ init\left(\,field\left(X\!\!+\!\!DX,Y\right)\right) \right. ;
31
                      1,2: at(B,X,Y+DY,T), B>0; 1,2: not init(field(X,Y+DY))
32
                      1,3: at(C,X+DX,Y+DY,T), C>0; 1,2: not init(field(X+DX,Y+DY))},
33
                      dirPos(DX), dirPos(DY), at(N,X,Y,T), N>0, time(T).
35\ \operatorname{uncovered}\left(X,Y,T\right)\ :-\ \operatorname{init}\left(\operatorname{target}\left(X,Y\right)\right),\ \operatorname{not}\ \operatorname{at}\left({}_{-},X,Y,T\right),\ \operatorname{time}\left(T\right).
36 \ \operatorname{goal}\left(X,Y,T\right) \ :- \ \operatorname{init}\left(\operatorname{target}\left(X,Y\right)\right), \ \operatorname{not} \ \operatorname{uncovered}\left(X,Y,T\right), \ \operatorname{not} \ \operatorname{at}\left(0,X,Y,T\right), \ \operatorname{time}\left(T\right).
37 :- not goal(X,Y, horizon), init(target(X,Y)).
38
39 #show at /4.
```

4.2 Evaluation

4.3 Experiments and Results

5 Discussion of Results

compare chapter State of the Art with chapter My Work [1]

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