Efficient Planning for Sokoban with Answer Set Programming - A Case Study

Thomas Verweyen

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Abstract

Sokoban approach with ASP planning \rightarrow Plain Enc, IncEnc, Optimization (?)

Introduction

What do i do? I'm working on the Sokoban problem and trying to do my best to find a fast and easy solution to it. I'm not working on any level but just the ones meeting some criteria. I'm picking qualitative and quantitative characteristics and deciding on some values to examine. My encodings are supposed to work better on some of the instances and worse on others. What rules describe those fluctuations? What assumptions can be made? Is there a trend when comparing runtimes to my qualitative and quantitative characteristics? E.g. you would expect the runtime to go up when the instance's size grows or when the amount of boxes in the level grows. In this thesis I will work on optimizing ASP encodings for the game Sokoban towards different instances. First Ill briefly explain what Sokoban is and how it is played, then Ill give a short introduction on comparable/similar/other scientific studies concerning sokoban. Then Ill go over my work, ... and finally evaluation and Discussion of Results.

What is Sokoban

Sokoban is a ligostic planning puzzle. The player pushes boxes around a warehouse, trying to cover all goal fields with one of the boxes. For that the player can move one step at a time into any of the four directions of a two-dimensional grid. When the player moves onto a box's position the box is pushed one field into the direction of the move. Each box that is not pushed at step T stays at the same position for step T+1. There can never be multiple boxes on one field, neither can there be a box where there is no field. This is a very difficult planning problem for multiple reasons. There are a lot of possibilities to create a deadlock (a gamestate from which no solution can be found), the branching factor is large and the optimal solution can be very long [5].

State of the Art

- mention other works on Sokoban - what tests? - what results? - whats the connection?

My Work

what did i do? what is my baseline? what changes from encoding to encoding which of the encodings works best in which kind of instance how do i evaluate that show my experiments and their results

4.1 Encodings

4.1.1 Fact Format



```
1 #const horizon = 23.
2 #const dimX = 5.
3 #const dimY = 4.
4 init (field (1..5,1..4)).
5 init (at (0,1,2)).
6 init (at (1,1,3)).
7 init (at (2,2,2)).
8 init (at (3,3,1)).
9 init (at (4,3,3)).
10 init (at (5,4,2)).
11 init (at (6,4,4)).
```

```
12 init (target (1,4)).
13 init (target (2,1)).
14 init (target (3,4)).
15 init (target (4,1)).
16 init (target (5,3)).
17 init (target (5,4)).
```

4.1.2 Plain Encoding

```
1 time (0.. horizon).
3 dir (1,0;0,1;-1,0;0,-1).
4 at (N,X,Y,0): - init (at(N,X,Y)).
6 % pick one direction to move in
7 1 { move(DX,DY,T) : dir(DX,DY) } 1 :- time(T).
8 % move the player
9 at (0,X+DX,Y+DY,T): - at (0,X,Y,T-1), move (DX,DY,T), time (T).
10 % move box, the player collides with
11 at (N,X+DX,Y+DY,T) :- at (N,X,Y,T-1), at (0,X,Y,T), move (DX,DY,T),
                              time(T), N>0.
12
13 % leave boxes that aren't touched at their place with t+1
14 at (N, X, Y, T) :- at (N, X, Y, T-1), not at (0, X, Y, T), N>0, time (T).
15\ \% there can't be boxes, players, where there is no field
16 :- not init (field (X,Y)), at (\_,X,Y,\_).
17 % there can't be two boxes/players, at the same spot
18 :- at (N,X,Y,T), at (M,X,Y,T), N>M, time (T).
19
20 \operatorname{uncvrd}(X,Y,T) :- \operatorname{init}(\operatorname{target}(X,Y)), not \operatorname{at}(\_,X,Y,T), \operatorname{time}(T).
21
  goal(X,Y,T):- init(target(X,Y)), not uncvrd(X,Y,T),
                              not at (0, X, Y, T), time (T).
23 :- not goal(X,Y,horizon), init(target(X,Y)).
25 #show at /4.
```

The first try to construct a naive encoding to solve sokoban-instances yielded a solution that works in a common way in ASP planning. In line 4 we initialize the information from the level instance with the time 0. In line 7 we state that there is one move at minimum and at maximum for each timestep. This makes the ASP planner ground all sequences of moves, pick one for every timestep and infer every successive state from the moves it picks. In lines 9-13 we calculate the successor state from the move picked in line 7. Line 9 calculates the new player-position, line 11 calculates the new position of the box in the players way and line 13 calculates the new position of boxes not touched by the player. Additionally we need some integrity constraints to prevent illegal states. The integrity constraint in line 15 ensures that at no time a box or player has a position outside of the playing field. The integrity constraint in line 17 ensures that at no time boxes or the player share one field. In the lines 19-21 we define our win-condition to be that all targets in the instance have to be covered by boxes. The Plain Encoding has the predicates, time/1, dir/2, at/4, move/3, uncovered/3, goal/3 The encoding works in a common way in ASP planning, by guessing the sequence of actions and inferring the successor states by effect axioms. There is a time/1 predicate for every timestep and a at/4 predicate for every box and the player at every timestep.

4.1.3 Incremental Encoding

```
1 #include <incmode>.
 3 #program base.
 4 dir ( 1,0 ; 0,1 ; -1,0 ; 0,-1 ).
 5 at (N, X, Y, 0) :- init (at (N, X, Y)).
 7 #program step(t).
 8 % pick one direction to move in
 9 1 { move(DX,DY,t) : dir(DX,DY) } 1.
10 \% move the player
11 at (0, X+DX, Y+DY, t) :- at (0, X, Y, t-1), move (DX, DY, t), init (field (X+DX, Y+DY)).
12 % move box, the player collides with
13 at (N,X+DX,Y+DY,t): - at (N,X,Y,t-1), at (0,X,Y,t), move (DX,DY,t), N>0.
14 % leave boxes that aren't touched at their place with t+1
15 \ at \, (N,X,Y,t \,) \ :- \ at \, (N,X,Y,t-1) \,, \ not \ at \, (0\,,X,Y,\overset{\,\,{}_{}}{t} \,) \,, \ N>0.
16 % there can't be boxes, players, where there is no field 17 :- at(_,X,Y,_), not init(field(X,Y)).
18 % there can't be two boxes/players, at the same spot
19 :- at (N,X,Y,t), at (M,X,Y,t), N>M.
20
21 #program check(t).
22 #external query(t).
23 uncovered (X,Y,t):- init (target(X,Y)), not at (\_,X,Y,t), query (t).
24 \ \operatorname{goal}(X,Y,t) \ :- \ \operatorname{init}\left(\operatorname{target}(X,Y)\right), \ \operatorname{not} \ \operatorname{uncovered}(X,Y,t), \ \operatorname{not} \ \operatorname{at}\left(0,X,Y,t\right), \ \operatorname{query}(t).
25 : - \ not \ goal\left(X,Y,t\right), \ init\left(\,target\left(X,Y\right)\right), \ query\left(\,t\,\right).
27 #show at /4.
```

The incremental encoding is similar to the plain encoding. The difference is the partition of the encoding into three parts: the #program base in lines 3-5 which is independent of the step parameter t, the #program step in lines 7-19 which is the cumulative part, collecting knowledge and the #program check in lines 21-25 which is specific for each value of t (STEP ALSO IS THOUGH). This allows for gradually processing each time step and accumulating knowledge that persists even if the current step is not satisfiable.

4.1.4 Advanced Encoding

```
\begin{array}{lll} 1 \ time \, (0 \ldots horizon \,). \\ 2 \\ 3 \ dir \, (1 \, , 0 \, ; 0 \, , 1 \, ; -1 \, , 0 \, ; 0 \, , -1 \,). \\ 4 \ at \, (N,X,Y,0) \ : - \ init \, (at \, (N,X,Y) \,). \\ 5 \\ 6 \ forbidden \, (X,Y) \ : - \ init \, (field \, (X,Y) \, , \, not \, init \, (target \, (X,Y) \, , \\ 7 \ 3 \ \#sum \, \left\{ 3 \ : \, not \, init \, (field \, (X+DX,Y+DY) \, ) \, , \, not \, init \, (field \, (X+DY,Y+DX) \, ) \, , \, dir \, (DX,DY) \, ; \\ 8 \ 3 \ : \, not \, init \, (field \, (X+DX,Y+DY) \, ) \, , \, not \, init \, (field \, (X-DY,Y-DX) \, ) \, , \, dir \, (DX,DY) \, ; \\ 9 \ 1,1 \ : \, forbidden \, (X+DX,Y+DY) \, , \, dir \, (DX,DY) \, ; \, 1,2 \ : \, forbidden \, (X+DX,Y+DY) \, , \, dir \, (DX,DY) \, ; \\ 10 \ 1,3 \ : \, not \, \, init \, (field \, (X+DX,Y+DY) \, ) \, , \, dir \, (DX,DY) \, ; \\ 11 \ 1,4 \ : \, not \, \, init \, (field \, (X+DX,Y+DY) \, ) \, , \, not \, \, init \, (field \, (X-DX,Y-DY) \, ) \, , \, dir \, (DX,DY) \, \}. \\ 12 \ 13 \ \% \ pick \ one \ direction \ to \ move \ in \end{array}
```

```
14 1 { move(DX,DY,T) : dir(DX,DY) } 1 :- time(T).
15 \% move the player
16 at (0, X+DX, Y+DY, T) :- at (0, X, Y, T-1), move (DX, DY, T), time (T).
17 \% move box, the player collides with
18 \  \, at \left(N+1,\!X+\!DX,\!Y+\!DY,\!T\right) \  \, :- \  \, at \left(N+1,\!X,\!Y,\!T-1\right) \, , \  \, at \left(0 \, ,\!X-\!DX,\!Y-\!DY,t \, \right) \, , \  \, at \left(0 \, ,\!X,\!Y,\!T\right) \, , \  \, move \left(DX,\!DY,\!T\right) \, , \  \, time \left(T\right) \, , \  \, time \left
19 % leave boxes that aren't touched at their place with t+1
20 at (N,X,Y,T) :- at (N,X,Y,T-1), not at (0,X,Y,T), N>0, time (T).
21~\% don't take a move back on the very next step
22 :- move(DX,DY,T), move(-DX,-DY,T-1), at(0,X,Y,T), not at(-,X-DX,Y-DY,T-1), time(T).
23~\% there can't be boxes, players, where there is no field 24 :- not init(field(X,Y)), at(_,X,Y,_).
25~\% there can't be two boxes/players, at the same spot
26 :- at (N,X,Y,T), at (M,X,Y,T), N>M, time (T).
27 % don't push a box onto a forbidden field
28 :- at(N,X,Y,T), forbidden(X,Y), N>0, T>0, not at(N,X,Y,T-1).
29 dirPos(1;-1).
30 :- 3 \text{ #sum } \{1,1 : at(A,X+DX,Y,T), A>0 ; 1,1 : not init(field(X+DX,Y)) ;
                                                         1,2: at(B,X,Y+DY,T), B>0; 1,2: not init(field(X,Y+DY));
                                                           1\,,3\ :\ \mathrm{at}\left(\text{C},\text{X+DX},\text{Y+DY},\text{T}\right)\,,\ \text{C>0}\ ;\ 1\,,2\ :\ \mathrm{not\ init}\left(\,\mathrm{field}\left(\text{X+DX},\text{Y+DY}\right)\right)\ \big\}\,,
32
                                                          \label{eq:continuous_problem} \operatorname{dir} \operatorname{Pos}\left(\operatorname{D\!X}\right), \ \operatorname{dir} \operatorname{Pos}\left(\operatorname{D\!Y}\right), \ \operatorname{at}\left(\operatorname{N}, \operatorname{X}, \operatorname{Y}, \operatorname{T}\right), \ \operatorname{N} > 0, \ \operatorname{time}\left(\operatorname{T}\right).
33
34
35 \operatorname{uncovered}(X,Y,T) :- \operatorname{init}(\operatorname{target}(X,Y)), \operatorname{not} \operatorname{at}(_-,X,Y,T), \operatorname{time}(T).
36 \operatorname{goal}(X,Y,T) := \operatorname{init}(\operatorname{target}(X,Y)), \operatorname{not\ uncovered}(X,Y,T), \operatorname{not\ at}(0,X,Y,T), \operatorname{time}(T).
37 :- not goal(X,Y, horizon), init(target(X,Y)).
39 #show at /4.
```

4.2 Evaluation

4.3 Experiments and Results

Discussion of Results

compare chapter State of the Art with chapter My Work [1]

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