# The Definition of $\lambda_Q$

## April 9, 2021

## Contents

1	Syntax	1
2	Type Checking Rules	2
3	Operational Semantics	3
1	Syntax	

t	::=		terms:
		x	variable
		unit	constant unit
		true	constant true
		false	constant false
		$\lambda x : T.t$	function abstraction
		t $t$	function application
		(t,t)	pair
		t.1	first projection
		t.2	second projection
		if t then t else t	conditional
		run C	static lifting
		$\zeta p: W.C$	circuit abstraction
v	::=		values:
		$\lambda x : T.t$	abstraction value
		(v, v)	pair value
		unit	unit value
		true	true value
		false	false value
		$\zeta p: W.C$	circuit value
T	::=		types:
_		Unit	unit type
		C 7000	dille type

		$Bool \\ T \times T \\ T \to T \\ T \leadsto T$	boolean type product type function type circuit type
Γ	::=	Ø	contexts: empty context
		$\Gamma, x:T$	term variable binding
W	::=		wire types:
		One	wire unit type
		Bit	bit type
		Qubit	qubit type
		$W \otimes W$	wire product type
p	::=		wire patterns:
		()	empty
		w	wire variable
		$\omega$	whe variable
		(p,p)	wire variable wire pair
C	::=		
C	::=		wire pair circuits: output a pattern
C	::=	$(p,p)$ $\begin{array}{c} output \ p \\ p_2 \leftarrow gate \ g \ p_1; C \end{array}$	wire pair  circuits: output a pattern gate application
C	::=	$(p, p)$ $output p$ $p_2 \leftarrow gate g p_1; C$ $p \leftarrow C; C$	wire pair  circuits:  output a pattern  gate application circuit composition
C	::=	$(p,p)$ $output p$ $p_2 \leftarrow gate g p_1; C$ $p \leftarrow C; C$ $x \Leftarrow lift p; C$	wire pair  circuits: output a pattern gate application circuit composition dynamic lifting
C	::=	$(p, p)$ $output p$ $p_2 \leftarrow gate g p_1; C$ $p \leftarrow C; C$	wire pair  circuits:  output a pattern  gate application circuit composition
C	::=	$(p,p)$ $output p$ $p_2 \leftarrow gate g p_1; C$ $p \leftarrow C; C$ $x \Leftarrow lift p; C$	wire pair  circuits: output a pattern gate application circuit composition dynamic lifting
		$(p,p)$ $output p$ $p_2 \leftarrow gate g p_1; C$ $p \leftarrow C; C$ $x \Leftarrow lift p; C$	circuits: output a pattern gate application circuit composition dynamic lifting circuit application

#### 2 Type Checking Rules

First, we need to define what is a well-formed wire context.

**Definition 1** (Well-formed Wire Contexts). A wire context  $\Omega$  is well-formed, if there are no duplicate wire variables in it. For simplicity, we always assume the wire contexts are well-formed in the following contexts. And when we write  $\Omega_1, \Omega_2$ , we require  $\Omega_1$  and  $\Omega_2$  to be disjoint to preserve the well-formedness.

Since there are three different kinds of terms: ( $\lambda$ -)terms, wire patterns and circuits, we have three different relations for each of them to check the correctness.

- $\Omega \Rightarrow p: W$  is used to check the well-formedness of patterns.
- $\Gamma; \Omega \vdash C : W$  is the typing relation for circuits;
- $\Gamma \vdash t : T$  is the typing relation for  $(\lambda$ -)terms;

Well-formed patterns:  $\Omega \Rightarrow p:W$ 

$$\overline{\varnothing \Rightarrow (): One}$$

$$\overline{w:W\Rightarrow w:Wz}$$

$$\frac{\Omega_1 \Rightarrow p_1 : W_1 \quad \Omega_2 \Rightarrow p_2 : W_2}{\Omega_1, \Omega_2 \Rightarrow (p_1, p_2) : W_1 \otimes W_2}$$

Typing rules for circuits:  $\Gamma; \Omega \vdash C : W$ TODO: add typing rules for circuits.

Typing rules for  $(\lambda$ -)terms:  $\Gamma \vdash t : T$ 

TODO: add typing rules for terms.

#### 3 **Operational Semantics**

TODO: add semantics.