Joint Cryptography and Data Compression

João Almeida

Joint work with João Barros

Instituto de Telecomunicações Faculdade de Engenharia da Universidade do Porto, Portugal jpa@fe.up.pt



U PORTO





creating and sharing knowledge for telecommunications

- Introduction
- Source coding and security
- 3 Analysis-by-synthesis algorithms
- 4 Conclusions





Security - is it worth it?

- Well understood problem with many alternative solutions;
 - Almost all of them are modular;
- Lots of constraints
 - Hardware:
 - Time:
 - Data;
- As lightweight as possible:
 - Try to remove the extra module (go for channel coding, source coding, network coding, ...);

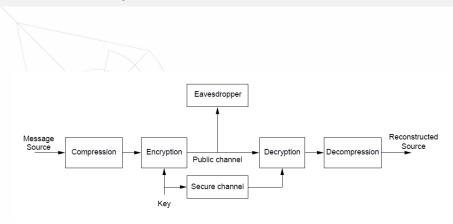




- Introduction
- Source coding and security
- 3 Analysis-by-synthesis algorithms
- 4 Conclusions

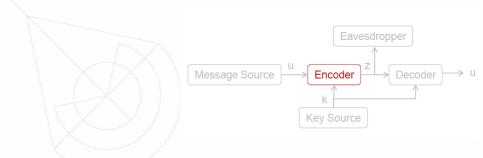


Modular security





Combining compression and encryption



- How can we design such encoder?
 - Explore knowledge given by the source and properties of source encoding algorithms;
 - Use this knowledge to compute cryptograms and key-streams;
 - Analysis-by-synthesis principle;



Prospects in variable length codes

- Catastrophic error propagation:
 - C = {A: 100, B: 0, C: 111, D: 101, E: 110};
 - BBCBECDBBB \rightarrow 001110110111101000 \rightarrow DBDDCBAB;
 - Symbol error rate depends on the codebook;
- Some effort to analyze error propagation and prevent it (reversible codes, synchronizing codewords, ...);

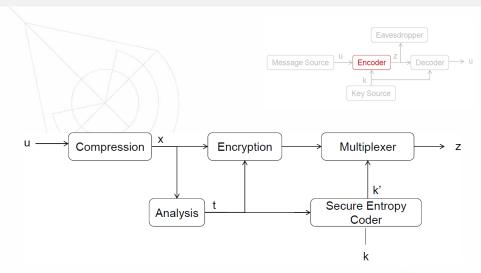


Prospects in variable length codes

- Catastrophic error propagation:
 - C = {A: 100, B: 0, C: 111, D: 101, E: 110};
 - BBCBECDBBB \rightarrow 0011101101111101000 \rightarrow DBDDCBAB:
 - Symbol error rate depends on the codebook;
- Some effort to analyze error propagation and prevent it (reversible codes, synchronizing codewords, ...);
- ... catastrophic error propagation can be useful;
 - Error patterns ↔ key-streams;
 - Carefully induce errors...
 - Unfortunately the decoder must also know where these errors have been introduced:
 - Share it or send it?



Detailed encoder

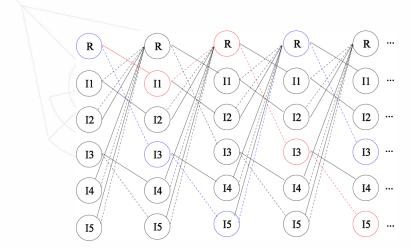




- Introduction
- 2 Source coding and security
- 3 Analysis-by-synthesis algorithms
- 4 Conclusions



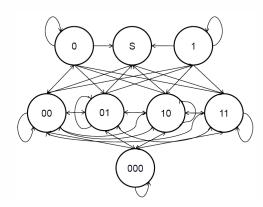
Algorithm I





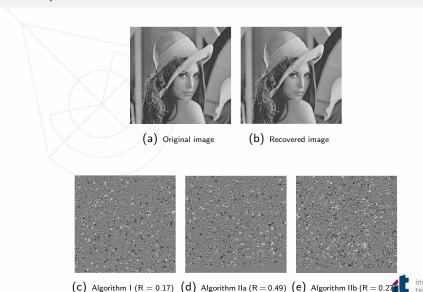
Algorithm II

s	x		t	
1	-		-	
0	0000	С	0010	Υ
0	010	R	000	Υ
1	001	Υ	010	Υ
1	011	Р	000	0
0	100	Т	000	Н
1	101	0	000	R
1	111	G	000	G
0	010	R	000	0
1	0001	Α	0000	С
1	011	Р	000	0
0	110	Н	000	G
S	001	Υ	000	Α





Example: baseline JPEG



12 / 16

Wrapping up

- Low complexity:
 - Algorithm I: $\mathcal{O}(k \cdot p)$, where k is the length of the input bit-stream and p the number of states in the trellis;
 - Algorithm II: $\mathcal{O}(n)$, where n is the size of the source message;
- Inducing loss of sync:
 - Leads to compressible key-streams;
- Codeword size diversity might become an issue:
 - Use super-symbols;





- Introduction
- 2 Source coding and security
- 3 Analysis-by-synthesis algorithms
- 4 Conclusions





Take home messages!

- Computational constraints ask for lightweight cryptography:
 - Coding schemes are good options;
- Modular approaches are insensitive to underlying communication blocks:
 - Jointly design compression and encryption stages;
- Leverage from source knowledge and encoder properties:
 - Analysis-by-synthesis encoding;
- Variable length codes:
 - Use error-patterns as key-streams (catastrophic error propagation);
 - Induce loss of synchronization;
- Further developments:
 - Secure entropy encoders;
 - Algorithms for other source encoders;





Q & A

Thank you for the attention! Questions?

