BCE 602 A4

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txerosse 1 P 1, - norm: the median of the coefficients of b. (b, b) We can see that in 2-D spale

(b, b) Chy points between (b, b) and (b, b)

will lead to a optimal solution,

which means x can be in range xb, b).

As the dimensen increase, we can see that

the median of b will difficiely lead to best result.

It could not be only one result, but the median will be definitely one

of them ls-norm: the average 1 b/m.

mis the dimension of the space. For example, in two dimension, f (x)=x and minimizell 1x-b112 means the shortest length from b to tix, Assume, y=-x+h going through point b (b,,b) we can get $h = b. + b_2$. To get the solution X, we need to solve $\begin{cases} J = X \\ J = -X + b_1 + b_2 \end{cases}$ $\Rightarrow X = y = \frac{b_1 + b_2}{2}$. In higher dismension, use the same think method. Lo-norm: the midrange point (maxbi - minbi)/2 for example, in 2-D space $b_1-X = X-b_2$ $X = b_1-b_2$ (X,K) _we assume b, >>>bz

In higher dimension, the optimal solution occur at the vertex

ca) The robust least-squares problem, with interval matrix, is winimize supated IAx-blz, which is equivalent to minimize supart 11Ax-b11,2. The latter one can be reformulated as minimize y'y subject to y = Ax-b = y, for all A 6A

We also can easily get that $\sup_{A \in A} (Ax-b) = \sum_{j \ge 1}^{n} (A_{2j} \times_{j} + R_{2j} \times_{j}) - b_{2j}$ int (Ax-b): = \(\bar{\pi} (A\forall \forall) - R\forall [\forall]

Therefore, the oxiginal problem can be written as

minimize y'y subject to Ax + Rix1 - b ≤ y AX -RIX -6 >-4

In aP torm: minimize yTy subject to Ax+Pz-b=y Ax -Pz -6 >- Y The variables are & & P, ytpm, & & P

(b) results shown in pdt file