# CS224W-Homework 2 Solutions

## Tianjian Li

# February 2022

# 1 GCN

#### 1.1

The following two graphs are isomorphic, which we demonstrate by showing a 1-to-1 correspondence between the nodes.

#### 1.2

Assume that the two graphs both have three nodes of labelled  $v_1$ ,  $v_2$  and  $v_3$ . In both graphs, edges exist between  $v_1$ ,  $v_2$  and  $v_1$ ,  $v_3$ . Each node has feature of 1. If we use the sum and mean aggregator on both  $v_1$  nodes, the result would be the same since  $aggregate_{mean}(1,1) = 1 = aggregate_{max}(1,1)$ . However  $aggregate_{sum}(1,1) = 1 + 1 = 2$ 

#### 1.3

Proof:

Assume that the WL test cannot decide whether  $G_1$  and  $G_2$  are isomorphic at the end of K'th iteration. Let  $G_1 = (V_1, E_1)$  and  $G_2 = (V_2, E_2)$ . According to our assumption, we have:

$$\{l_v^{(K)}, \forall v \in V_1\} = \{l_v^{(K)}, \forall v \in V_2\}$$

Therefore,

$$HASH(l_v^{(k-1)}, \{l_u^{(k-1)}, \forall u \in N(v)\}), \forall v \in V_1 = HASH(l_v^{(k-1)}, \{l_u^{(k-1)}, \forall u \in N(v)\}), \forall v \in V_2 = HASH(l_v^{(k-1)}, \{l_u^{(k-1)}, \forall u \in N(v)\}), \forall v \in V_2 = HASH(l_v^{(k-1)}, \{l_u^{(k-1)}, \forall u \in N(v)\}), \forall v \in V_2 = HASH(l_v^{(k-1)}, \{l_u^{(k-1)}, \forall u \in N(v)\}), \forall v \in V_2 = HASH(l_v^{(k-1)}, \{l_u^{(k-1)}, \forall u \in N(v)\}), \forall v \in V_2 = HASH(l_v^{(k-1)}, \{l_u^{(k-1)}, \forall u \in N(v)\}), \forall v \in V_2 = HASH(l_v^{(k-1)}, \{l_u^{(k-1)}, \forall u \in N(v)\}), \forall v \in V_2 = HASH(l_v^{(k-1)}, \{l_u^{(k-1)}, \{l_u^{(k$$

The hash function is an injective function on a set, so we have

$$l_v^{(K-1)}, \forall v \in V_1 = l_v^{(K-1)}, \forall v \in V_2$$

and

$$\{l_u^{(k-1)}, \forall u \in N(v), v \in V_1\} = \{l_u^{(k-1)}, \forall u \in N(v), \forall v \in V_2\}$$

This means that

$$combine(l_v^{(K-1)}, aggregate(\{l_u^{(k-1)}, \forall u \in N(v), v \in V_1\}))$$

is equal to

$$combine(l_v^{(K-1)}, aggregate(\{l_u^{(k-1)}, \forall u \in N(v), v \in V_2\}))$$

Therefore

$$readout(\{h_v^{(K)}, \forall v \in V_1\}) = readout(\{h_v^{(K)}, \forall v \in V_2\})$$

Contradiction! Therefore the assumption is false, meaning that the Weisfeiler-Lehman test also decides that the two graphs are different.

# 2 Node Embeddings with TransE

#### 2.1

Assume there are 3 nodes of A, B and C. A is connected to B and C with the same relationship vector of l = (0,0). The embedding vector learned by TransE for A, B, C would all be the same, minimizing the loss function to zero.

$$\mathcal{L} = d(v_a + l, v_b) + d(v_a + l, v_c) = 0$$

However,  $d(v_b + l, v_c) = 0$ . B and C does not have the l relationship in between. Therefore the embeddings would be useless.

#### 2.2

The example for the previous question also works here because

$$\mathcal{L} = d(v_a + l, v_b) + d(v_a + l, v_c) - d(v_b + l, v_c) = 0$$

B and C would also have the l relationship according to the embeddings but in fact they do not have the l relationship.

#### 2.3

The algorithm would artificially increase the entity norms of the negative (corrupted samples) to minimize the loss function, resulting in the embeddings of the corrupted nodes much larger than others.

#### 2.4

The type of relations that TransE cannot model are symmetric but not reflexive. Wang et al. 2019 mentions that since

$$d(i,l,j) + d(j,l,i) \le -||e_i + 0 - e_j|| + ||e_j + 0 - e_I||$$

the TransE algorithm would push the embeddings  $e_i, e_j$  to be similar and the relation l to be zero. Howeverd(i, l, i) would also be small that we would consider the relation to be reflexive.

# 3 Expressive Power of Knowledge Graph Embeddings

#### 3.1

TransE cannot model symmetric relations because if a + l = b, we cannot make b + l = a unless l = 0.

TransE can model inverse relations since a + l = b we can make the inverse to be -l and b + (-l) = a.

TransE cannot model composition because if we make  $l_{aunt} = l_{father} + l_{sister}$ , we also have  $l_{aunt} = l_{sister} + l_{father}$ . Father's sister and sister's father would both be recognized as aunt.

#### 3.2

RotateE can model symmetric relations, we can make the rotate angle  $l = \frac{\pi}{2}$ , therefore if  $A \circ l = B$ , we have  $B \circ l = A$ 

RotateE can model inverse relations with any arbitary angle besides  $\pi$  RotateE cannot model composition because if we make  $l_{aunt} = l_{father} + l_{sister}$ , we also have  $l_{aunt} = l_{sister} + l_{father}$ . Father's sister and sister's father would both again be recognized as aunt.

## 3.3

RotateE cannot model inreflexive relations because if A and A itself have some kind of relation, so does any node with any node itself, therefore the relation would be reflexive. TransE also fails in this case.

### References

Wang, Yanjie et al. (Aug. 2019). "On Evaluating Embedding Models for Knowledge Base Completion". In: Proceedings of the 4th Workshop on Representation Learning for NLP (RepL4NLP-2019). Florence, Italy: Association for Computational Linguistics, pp. 104–112. DOI: 10.18653/v1/W19-4313. URL: https://aclanthology.org/W19-4313.