

# Lab01-Algorithm Analysis

CS214-Algorithm and Complexity, Xiaofeng Gao, Spring 2019.

\* If there is any problem, please contact TA Mingran Peng. Also please use English in homework.

\* Name: 田雪飞 Student ID: 515030910347 Email: 13487426939@qq.com

1. **Solution.** The following is solution process:

Algorithm	Time Complexity	Space Complexity
<i>InsertionSort</i>	$O(n^2), \Omega(n)$	$\Theta(1)$
<i>CocktailSort</i>	$O(n^2), \Omega(n)$	$\Theta(1)$
<i>SelectionSort</i>	$O(n^2), \Omega(n)$	$\Theta(1)$

**CocktailSort:**

**Line 4:** The number of **comparisons** carried out by Cocktailsort in loop

$$n - 1 \quad (1)$$

and the number of **swap and assignment operation** at least(**best**)

$$0 \quad (2)$$

at most(**worst**)

$$2(n - 1) + 2(n - 3) + \cdots + 1 \quad (3)$$

**Line 9:** The number of **comparisons** carried out by in loop at most least(**best**)

$$n - 2 \quad (4)$$

and the number of **swap and assignment operation** at least(**best**)

$$0 \quad (5)$$

at most(**worst**)

$$2(n - 2) + 2(n - 4) + \cdots + 0 \quad (6)$$

**so**, the number of total executed in the loop of while at least(**best**)

$$T(n) = n - 1 + n - 2 = 2n - 3 \quad (7)$$

at most (**worst**)

$$T(n) = n - 1 + n - 2 + 2 \sum_{0}^{n-1} k = n^2 + n - 3 \quad (8)$$

**so,Time Complexity is  $O(n^2)$  and  $\Omega(n)$ ,and Space Complexity is  $\Theta(1)$ .**

**SelectionSort:**

**Line 2:** The loop executed from 1 to  $n - 1$  **so**,the number of **assignment operations** carried out by SelectionSort are:

$$2(n - 1) \quad (9)$$

**Line 4:** The comparisons of  $A[j] > max$  will be executed when

$j = 2$ ,the number of comparisons are  $n - 1$ ;

$j = 3$ ,the number of comparisons are  $n - 2$ ;

⋮

$j = n$ , the number of comparisons are 0;

so, the number of **comparisons** carried out by SelectionSort are:

$$\sum_0^{n-1} j = \frac{n(n-1)}{2} \quad (10)$$

The number of **assignment operations** carried out in second loop are at most:

$$n(n-1) \quad (11)$$

at least:

$$0 \quad (12)$$

The number of total executed in SelectionSort are at most(**worst**):

$$T(n) = 2(n-1) + \frac{n(n-1)}{2} + n(n-1) = \frac{1}{2}(3n^2 + n - 4); \quad (13)$$

The number of total executed in SelectionSort are at most(**best**):

$$T(n) = 2(n-1) + \frac{n(n-1)}{2} = \frac{1}{2}(n^2 + 3n - 4); \quad (14)$$

so, Time Complexity is  $O(n^2)$  and  $\Omega(n)$ , and Space Complexity is  $\Theta(1)$ . □

2. **Solution.** The following is solution process:

(a) **Two stack to simulated queue:**

- **stack1** is used for **enqueue**, **stack2** is used for **dequeue**;
- when pushing an element, we push it into **stack1**;
- when popping an element, if **stack2** is empty, we push all elements in **stack1** into **stack2**, if **stack2** is not empty, we pop an element in **stack2** directly.

(b) **Time Complexity(potential function):**

**Potential function:**  $\Phi(S)$  denote the num[i] of items in **stack1**.

**State:** Here state  $S_i$  refers to the state of the **stack1** after the  $i$ -th.

**Correctness:**  $\Phi(S_i) \geq 0 = \Phi(S_0)$  for any  $i$ ;

According to the definition of  $\Phi(S)$ , we know

$$\Phi(S_i) = \text{num}[i] \quad (15)$$

$$\begin{aligned} \text{PUSH: } \hat{C}_i &= C_i + \Phi(S_i) - \Phi(S_0) \\ &= 1 + \text{num}[i] - \text{num}[i-1] \\ &= 1 + \text{num}[i-1] + 1 - \text{num}[i-1] \\ &= 2 \end{aligned}$$

**POP:**

- if **stack2** is not empty:

$$\begin{aligned} \hat{C}_i &= C_i + \Phi(S_i) - \Phi(S_0) \\ &= 1 + \text{num}[i] - \text{num}[i-1] \\ &= 1 \end{aligned}$$

- if **stack2** is empty:

$$\begin{aligned} \hat{C}_i &= C_i + \Phi(S_i) - \Phi(S_0) \\ &= \text{num}[i-1] + 1 + (\text{num}[i] - \text{num}[i-1]) \\ &= \text{num}[i-1] + 1 + (0 - \text{num}[i-1]) \\ &= 1 \end{aligned}$$

**Thus**, starting from two empty stacks, any sequence of **n1 PUSH, n2 POP** operations takes at most

$$T(n) = \sum_i^n C_i \leq \sum_i^n \hat{C}_i = 2n_1 + n_2. \quad \text{Here } n = n_1 + n_2 \quad \text{and} \quad n_1 \geq n_2. \quad (16)$$

□