

# Lab01-Algorithm Analysis

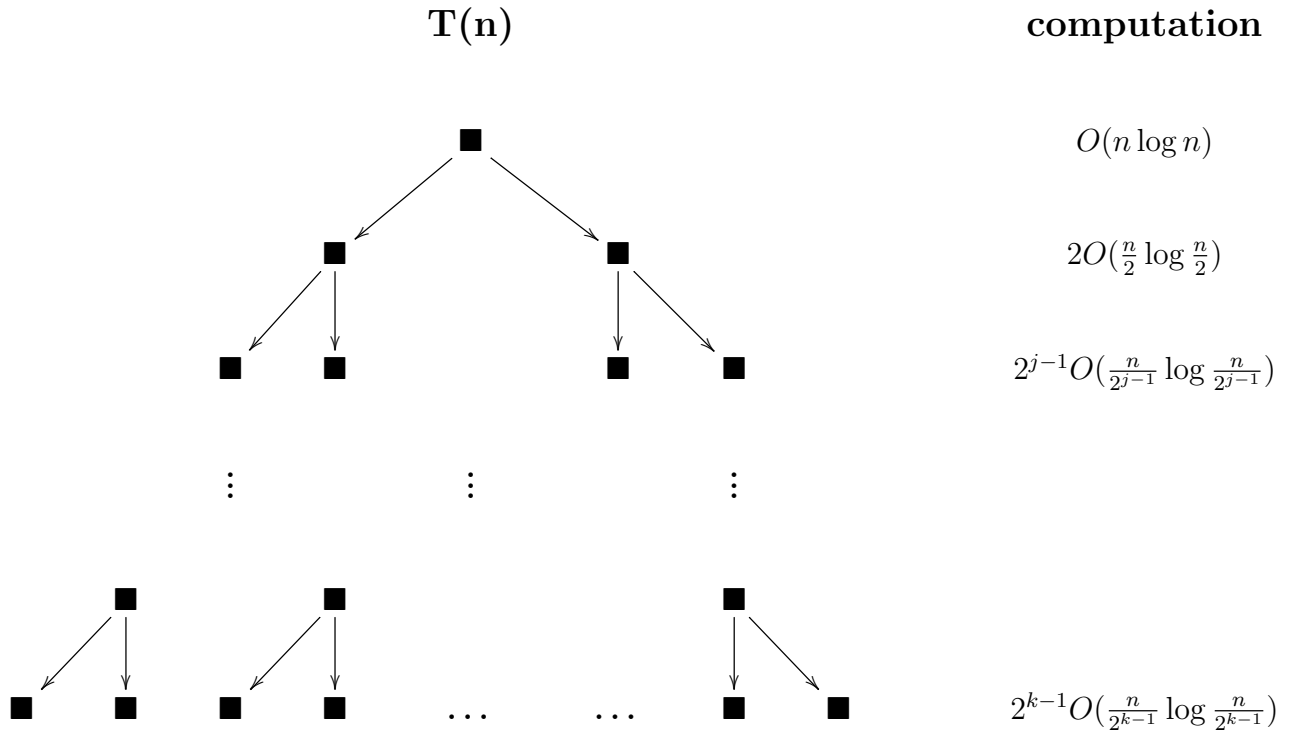
CS214-Algorithm and Complexity, Xiaofeng Gao, Spring 2019.

\* If there is any problem, please contact TA Mingran Peng. Also please use English in homework.

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1. **Solution.** The following is solution process:

(a) recurrence tree  $T(n)$ :



Complexity of  $T(n)$ =sum up all computations at each level.

(b) No,because  $O(n \log n)$  is not same as  $O(n^d)$ .

As we all know: $k = \log n$

$$T(n) = \sum_{j=0}^{k-1} 2^j O(\frac{n}{2^j} \log \frac{n}{2^j}) = \sum_{j=1}^k 2^j O(\frac{n}{2^j} \log \frac{n}{2^j}) = \sum_{j=1}^k O(n \log \frac{n}{2^j}) \quad (1)$$

$$= \sum_{j=1}^k O(n \log \frac{2^k}{2^j}) = \sum_{j=1}^k O(n \log 2^j) = O(n \log^2 n) \quad (2)$$

□

2. **Solution.** The following is solution process:

(a) The answer is in homework2.cpp

(b) recurrence  $T(n)$  by using Master Theorem:

$$T(n) = 2T(\frac{n}{2}) + O(n^2) = O(n^2) \quad (3)$$

□

3. **Solution.** The following is solution process:

(a) induction:

- $n=1$ , we do not need sort;;
- $n=2$ , obviously, we can use a comparator to sort two elements;
- we assume that we can sort  $k$  elements by using sort working;
- $n=k+1$ ; we can construct a function:

$$f(x) = \begin{cases} 0 & \text{if } x < a_{k+1} \\ 1 & \text{if } x \geq a_{k+1} \end{cases}$$

we can sort  $k$  elements to  $\{0, 0 \dots, 1, 1\}$ ;

$a_{k+1} = 1$ , and we can compare  $a_{k+1}$  with element in  $\{0, 0 \dots, 1, 1\}$  one-by-one. obviously, we can find a right index for  $a_{k+1}$  after  $k$  comparisons.

**So**, a transposition network with  $n$  input is a sorting network if and only if it sort the sequence  $\{n, n-1 \dots, 1\}$ .

(b) The answer is in homework2.py

□