

## Seminar 7

6. Write all models using DNF

→ ~~outcomes~~ using CNF

$$( \vee \quad \vee ) \wedge ( \vee \quad \vee ) \dots$$

$\underbrace{\quad}_{F}$

$$Ex \quad \delta \quad (\neg P \vee j) \wedge (\neg r \vee \neg P) \wedge (\neg r \vee \neg j)$$

$(CNF)$

$$\text{Clause } 1: (\neg P \vee j) : F$$

$$\Rightarrow \neg P = F \rightarrow P = \bar{T}$$

$$j = F$$

	P	j	r
i <sub>1</sub>	T	F	T
i <sub>2</sub>	T	F	F

Clause 2:  $(\neg R \vee \neg P)^F$

$\neg$	$R$	$\neg$	$P$	$F$	$T$	$T$	$\neg$	$T$
$\neg$	$T$	$\neg$	$T$	$F$	$i_B$	$T$	$T$	$T$
$T$	$\neg$	$T$	$\neg$	$F$	$i_u$	$T$	$T$	$T$

Clause 3:  $\neg R \vee \neg P^F$  is  $\neg T \neg \neg$

$\neg$	$R$	$\neg$	$P$	$F$	$i_6$	$F$	$T$	$T$
$\neg$	$T$	$\neg$	$T$	$F$	$i_6$	$F$	$T$	$T$
$T$	$\neg$	$T$	$\neg$	$F$	$i_6$	$F$	$T$	$T$

$$i_1 = i_u$$

$$i_3 = i_5$$

Optimaleis:  $i_1, i_2, i_3, i_6$

Q) Using the definition of deduction prove the following

$$q. \quad \vdash \vee (q \rightarrow P), \vdash q, \vdash \vdash P \\ f_1: \vdash \vee (q \rightarrow P) \equiv \vdash \rightarrow (q \rightarrow P)$$

INFERENCE RULES

$$P \rightarrow q \equiv \vdash P \vee q$$

$$U, U \rightarrow V \vdash_{mp} V$$

$$U \wedge V \vdash_{\text{simp}} U$$

$$U \wedge V \vdash_{\text{simp}} V$$

$$U, V \vdash_{\text{cop}} U \wedge V$$

$$\neg V, U \rightarrow V \vdash_{\text{mt}} \neg U$$

$$f_2: \vdash \vee q \equiv \vdash \rightarrow (q \rightarrow P)$$

$$f_3: \vdash$$

$$f_1, f_3 \vdash_{mp} q \rightarrow P = f_4$$

$$f_2, f_3 \vdash_{mp} q = f_5$$

$$f_4, f_5 \vdash_{mp} P \neq$$

Conclusion

$(f_1, \dots, f_6)$  is the deduction of the conclusion from the

finisg

10

12.  $H_1: T_{\text{sunny}} \wedge \text{color}$

$f_2: \text{min} \rightarrow \text{sunny}$

$H_3: T_{\text{sunny}} \rightarrow \text{cancel}$

$f_4: \text{cancel} \rightarrow \text{hail}$

C = hail

$f_1: T_{\text{sunny}} \wedge \text{color}$

$f_1 \xrightarrow{\text{split}} \begin{cases} T_{\text{sunny}} = f_2 \\ \text{color} = f_3 \end{cases}$

$f_4: \text{min} \rightarrow \text{sunny}$

$f_2, f_4 \xrightarrow{\text{out}} T_{\text{sunny}} = f_5$

$f_6: T_{\text{min}} \rightarrow \text{cancel}$

$f_5, f_6 \xrightarrow{\text{imp}} \text{cancel} = f_7$

$f_7, f_8 \xrightarrow{\text{imp}} \text{hail}$

$f_8: \text{cancel} \rightarrow \text{hail}$

$f_1$ : If the day doctors can introduce he  
will work

$f_2$ : The day did not work

# Predicate logic

Ex 1:

M. For any positive integer  $x$ , if  
 $x$  is ~~not~~ prime, then there  
exists a prime  $y$  such that  
 $y$  divides  $x$  and  $x \neq y$

$p_i : \mathbb{N} \rightarrow \{\text{T}, \text{F}\}$ ,  $p_i(x) = \text{T}$ , if  $x$  is a  
positive integer

$p_{\text{pr}} : \mathbb{N} \rightarrow \{\text{T}, \text{F}\}$ ,  $p_{\text{pr}}(x) = \text{T}$ , if  $x$  is prime

$d_r : \mathbb{N} \rightarrow \{\text{T}, \text{F}\}$ ,  $d_r(y, x) = \text{T}$  if  $y$  divides  $x$

$m : \mathbb{N} \rightarrow \{\text{T}, \text{F}\}$ ,  $m(y, x) = \text{T}$  if  $y < x$

$\vdash Cx \left( P_i(x) \wedge P_2(x) \rightarrow \left( \begin{array}{c} y \\ f(y) \end{array} \right) M_2(f) \right)$

$$1 \operatorname{der}(y, x) \wedge \operatorname{Im}(x, y))$$

$$D = \mathbb{Z}^+ = N \Leftrightarrow$$

$$\hookrightarrow \left( \begin{matrix} f(x) \\ x \in N \end{matrix} \right) \text{PR}(x) \rightarrow \left( \begin{matrix} y \\ f(y) \end{matrix} \right) (\text{PR}(y)) \wedge \text{def}_{f(y), x}$$

1.5. For every positive integer  $x$ , if  $x$  is  
 equal at an integer then there exists  
 an integer  $y$  s.t.  $(y+1) \cdot (y-1) = x - 1$

- at the exam, write the properties for the function as null

2.

7. Any and other bases rule only is not a consistent function.

conf : FD  $\rightarrow \{\top, \perp\}$ ,  $c(x) = \top$  if  $x$  is a conf

ant. Fcn : P  $\rightarrow \{\top, \perp\}$ ,  $mf(x) = \top$  if  $x$  is an antifunction

l-les :  $P \times FD \rightarrow \{\top, \perp\}$  an addition function  
 $l(x, y) = \top$  if  $x$  less  $y$

$(\forall x)(\forall y)$   $(D = FD)$   $(conf(y) \wedge l(x, y) \rightarrow mf(x))$

14. Walley do not like the ext functions or graphs

w : A  $\rightarrow \{\top, \perp\}$ ,  $w(x) = \top$  if  $x$  is a mult

A = animals

P = plants

$f_x : A \rightarrow \{T, F\}$ ,  $f_x(x) = T$  if  $x$  is a ~~cat~~

$f : P \rightarrow \dots$  if  $x$  is a ~~plant~~

like Eat ( $x, y$ ) :  $A \times P \rightarrow \{T, F\}$

- T if  $x$  likes to eat  $y$

$(\forall x)(\forall y)(x \in A \wedge (\forall y)(f_x(y) \vee g(y)) \rightarrow ?_{\text{like Eat}}(x, y))$

$H_1$ : if  $x$  is king and  $y$  is his oldest son, then  $y$  king

$H_2$ : if  $x$  king and  $y$  defects  $x$  then  $y$  king

$H_3$ :  $R_3$  is king

$H_4$ :  $H_2$  defects  $R_3$

$H_5$ :  $H_2$  is  $H_3$ 's oldest son

C:  $H_2$  king?

$K(x)$  " $x$  is king"

$B(x,y)$  " $y$  is oldest son of  $x$ "

$D(x,y)$  " $y$  defects  $x$ "

$$H_1: (\mathbb{H}_x)(\mathbb{H}_y)(K(G) \wedge \mathcal{O}(x,y)) \rightarrow K(\mathbb{H})$$

$$H_2: (\mathbb{H}_x)(\mathbb{H}_y)(K(G) \wedge D(x,y)) \rightarrow K(\mathbb{H})$$

$$L_3: K(R_3)$$

$$H_4: D(R_3, H_{\mathbb{H}})$$

$$H_5: \mathcal{O}(\mathbb{H}_{\mathbb{H}}, H_8)$$

$$C: K(H_8)$$

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$$H_2 \xrightarrow[\text{unisol int.}]{} (\mathbb{H}_y)(K(R_3) \wedge D(R_3, \mathbb{H})) \rightarrow K(\mathbb{H}) = H_6$$

$$x \in R_3$$

$$H_0 \xrightarrow[\text{unisol int.}]{} K(R_3) \wedge D(R_3, H_{\mathbb{H}}) \rightarrow K(H_{\mathbb{H}}) = H_{\mathbb{H}}$$

$$y \in H_{\mathbb{H}}$$

