

CS188 HW 188

1) a) i)  $\alpha_1 = A \vee B : \frac{3}{4}$  or  $\frac{12}{16}$  so 12 models out of 16

ii)  $\alpha_2 = (A \wedge B) \Rightarrow C : \frac{7}{8}$  or  $\frac{14}{16}$  so 14 models out of 16

iii)  $\alpha_3 = (A \wedge B) \vee (\neg C \vee D) : \frac{13}{16}$  so 13 models out of 16

A	B	C	D	$\alpha_3$
T	T	T	T	T
T	T	T	F	T
T	T	F	T	T
T	T	F	F	T
T	F	T	T	T
T	F	T	F	F
T	F	F	T	T
F	T	T	T	T
F	T	T	F	F
F	T	F	F	T
F	F	T	T	T
F	F	T	F	F
F	F	F	T	T
F	F	F	F	T

b) i) A doesn't want to sit next to B.

$$\neg(X_{A,L} \wedge X_{B,M}) \wedge \neg(X_{A,M} \wedge X_{B,L}) \wedge \neg(X_{A,M} \wedge X_{B,R}) \\ \wedge \neg(X_{A,R} \wedge X_{B,M})$$

ii) A doesn't want to sit in the left chair  
 $\neg X_{A,L}$

iii) C doesn't want to sit to the (adjacent) right of B.

$$\neg(X_{C,M} \wedge X_{B,L}) \wedge \neg(X_{C,R} \wedge X_{B,M})$$

iv) Yes, we have to specify that no 2 person can sit at the same chair. (This is going to be long :()

$$\neg(X_{A,L} \wedge (X_{B,L} \vee X_{C,L})) \wedge \neg(X_{A,M} \wedge (X_{B,M} \vee X_{C,M}))$$

$$\neg(X_{A,R} \wedge (X_{B,R} \vee X_{C,R})) \wedge \neg(X_{B,L} \wedge (X_{A,L} \vee X_{C,L}))$$

$$\neg(X_{B,M} \wedge (X_{A,M} \vee X_{C,M})) \wedge \neg(X_{B,R} \wedge (X_{A,R} \vee X_{C,R}))$$

$$\neg(X_{C,L} \wedge (X_{A,L} \vee X_{B,L})) \wedge \neg(X_{C,M} \wedge (X_{A,M} \vee X_{B,M}))$$

$$\neg(X_{C,R} \wedge (X_{A,R} \vee X_{B,R}))$$

5) No, not satisfiable since parts i and ii  $\Rightarrow X_{A,R} \wedge X_{B,L}$

only spot left for C is middle which violates iii thus not satisfiable

c) i) Initially inferred is all false, count set to # of premises for each clause and agenda empty, which will have A at the beginning

0th Iteration	Inferred:	<table border="1"> <tr> <td>A</td><td>B</td><td>C</td><td>D</td></tr> <tr> <td>F</td><td>F</td><td>F</td><td>F</td></tr> </table>	A	B	C	D	F	F	F	F	Agenda: [A]	Count	<table border="1"> <tr> <td>A</td><td>B</td><td>C</td><td>D</td></tr> <tr> <td>0</td><td>1</td><td>1</td><td>2</td></tr> </table>	A	B	C	D	0	1	1	2
A	B	C	D																		
F	F	F	F																		
A	B	C	D																		
0	1	1	2																		

Pops A and does the appropriate actions to get to

1st Iteration	Inferred:	<table border="1"> <tr> <td>A</td><td>B</td><td>C</td><td>D</td></tr> <tr> <td>T</td><td>F</td><td>F</td><td>F</td></tr> </table>	A	B	C	D	T	F	F	F	Agenda: [B,C]	Count	<table border="1"> <tr> <td>A</td><td>B</td><td>C</td><td>D</td></tr> <tr> <td>0</td><td>0</td><td>0</td><td>2</td></tr> </table>	A	B	C	D	0	0	0	2
A	B	C	D																		
T	F	F	F																		
A	B	C	D																		
0	0	0	2																		

$\Rightarrow$  Pops B and does the appropriate actions to get to

2nd Iteration	Inferred:	<table border="1"> <tr> <td>A</td><td>B</td><td>C</td><td>D</td></tr> <tr> <td>T</td><td>T</td><td>F</td><td>F</td></tr> </table>	A	B	C	D	T	T	F	F	Agenda: [C]	Count	<table border="1"> <tr> <td>A</td><td>B</td><td>C</td><td>D</td></tr> <tr> <td>0</td><td>0</td><td>0</td><td>1</td></tr> </table>	A	B	C	D	0	0	0	1
A	B	C	D																		
T	T	F	F																		
A	B	C	D																		
0	0	0	1																		

Continues on the next page (sorry)

3rd Iteration

Pops C and does the appropriate actions to get to

Inferred	A	B	C	D	Agenda: [D]	Count	A	B	C	D
	T	T	T	F			0	0	0	0

4th Iteration

Pops D and does the appropriate actions to get to

Inferred	A	B	C	D	Agenda: [-]	Count	A	B	C	D
	T	T	T	T			0	0	0	0

ii) Simply call DPLL ( $\text{KB}.$  clauses,  $\text{KB}.$  symbols,  $\text{True} \Rightarrow D$ ) or in other words use SAT solver to check if  $\text{True} \Rightarrow D$  is satisfiable in KB

iii) In CNF the clauses may be represented as:

$$(A) \wedge (\neg A \vee B) \wedge (\neg A \vee C) \wedge (\neg B \vee \neg C \vee D)$$

$\wedge (\text{False} \vee D)$  last clause comes from  $\text{True} \Rightarrow D$  or  $\models_{\text{KB}} D$

Simplifies to just D

iv) We begin with the set of clauses from part iii; since none of the base cases are satisfied we look at pure-symbols/literals

Since A is pure set  $A = \text{True}$  the simplified remaining clauses are

$$(\text{False} \vee B) \wedge (\text{False} \vee C) \wedge (\neg B \vee \neg C \vee D) \wedge (\text{False} \vee D)$$

We continue with unit-clauses which simplifies clauses to

$$(B) \wedge (C) \wedge (\neg B \vee \neg C \vee D) \wedge (D)$$
 then assign values to pure literals

to get  $B = \text{True}$   $C = \text{True}$   $D = \text{True}$  to simplify:

$(\text{False} \vee \text{False} \vee \text{True})$  which hits the base case of all statements being True thus returns true as satisfiable for  $\text{KB} \models D$