

$n \Rightarrow n!$ (fixed) \times

simple \leftarrow far out

$f(e_1, e_2) = p(e_1, e_2) \quad [t_1 = -]$

\uparrow
 \neq order

$\max_{k \in \{1, \dots, n\}} \sum_{i=1}^n p(a_i, b_{k_i})$

$\begin{bmatrix} 1 \\ 0.5 \\ 0 \end{bmatrix}$

Fixed!

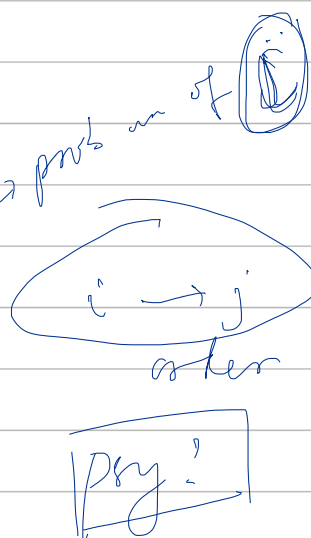
white 1rs
 vs 2.
 b

asm: $x_{ij} \in \{0, 1\}$

supply: $\sum_j x_{ij} = 1 \quad \text{for } i = 1 \rightarrow n$

dem: $\sum_i x_{ik} = 1 \quad \text{for } k = 1 \rightarrow n$

$\max \sum_{i,j} x_{ij} \cdot p(i,j) \cdot \begin{bmatrix} 1 \\ 0.5 \\ 0 \end{bmatrix}$



Order of round: fixed: 1 2 3

$p(i,j, \text{index})$

$\begin{matrix} 1 \rightarrow \text{white} \\ 2 \rightarrow \text{black} \\ \%2 \end{matrix}$

fixed (order \rightarrow color \rightarrow k)

We don't know strategy! 1st choose the best

{ 2nd }

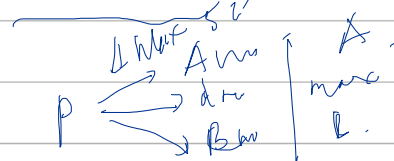
1 strategy: permutation: $\{1, 2, \dots, m\}$

isn't equilibrium

| | B | C | D | ... |
|-----|---|---|---|-----|
| A | | | | |
| T | | 1 | 1 | |
| F | | | | |
| ... | | | | |

$n \times m$

2 reverses



$A \in B = \{1\}$

Nash equilibrium opt

| A \ B | I | II | III |
|-------|--------|--------|--------|
| I | (1, 2) | (1, 1) | (1, 1) |
| II | (1, 1) | (1, 1) | (1, 1) |

best for all

A - B

Calculate

max = max

checking pure strategy

has a

Nash equilibrium:

For A: fixed $p_1, \dots, p_n \Rightarrow$

Max min (g_1, \dots, g_n)

s.t. $0 \leq p_i \leq 1$

A: $g_1 = p_1 \cdot b_{11} + p_2 \cdot b_{12} + \dots + p_n \cdot b_{1n}$

$g_2 = \dots$

what you pick p_i higher

For B: fix q_1, \dots, q_n

③ all Nash equilibrium:

Max

| | perm | p_1 |
|-------|------|-------|
| q_1 | x | |

Max of all t_k $\{t_1, \dots, t_n\}$
 Max \uparrow max(score)