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Benchmarking Android Data Leak Detection Tools

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Abstract

In 2017, Android hit a global mobile market share of 88% which makes it the most popular mobile platform. Application stores, such as the Google Play Store, are offering millions of mobile applications to the consumers, which are installed and updated on a daily basis. However, the security of those applications is a major concern. A thorough security analysis before the publication of each application is time and resource consuming. Hence, platform providers cannot and do not manually vet every application handed in for publication. Consequently, many malicious and vulnerable applications find their way to the app stores and through there to the end users devices. Those applications exhibit serious security issues, such as leaking of sensitive information.

During the previous years, researchers have been proposing a myriad of techniques and tools to detect such issues. There also exist large scale taxonomies classifying such tools into different categories. However, it is unclear how these tools perform compared to each other. Such a comparison is almost infeasible, since most tools are no longer available or cannot be set up any more.

In this work, we review static analysis tools for detecting data leaks in Android applications. Out of 87 tools in the vulnerability detection domain, we are able to obtain 22 tools. We then identify 5 tools in the data leak detection domain and run them. We run them on a given data set with known data leak vulnerabilities and compare their performance. Furthermore, we run the tools on a larger set of real-world applications to assess the prevalence of data leak issues in open-source Android applications.

We propose our own approach to compare security analysis tools by normalising their interfaces. This simplifies result reproduction and extension to other security vulnerability domains. In addition, the user experience and usability is highly improved.

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1

Introduction

In this chapter, we provide an overview of the current state of Android security. We point out what major security challenges are still prevalent and how platform providers such as Google are trying to tackle them. We briefly present the efforts done by the research community and state the goal of this work. Finally, we formulate the research questions and provide an overview of the rest of this thesis.

1.1 The State Of Android Security

There are over 2 billion devices worldwide running on the Android platform¹. With a market share of 87.7% (Q2, 2017²), Android is the most popular mobile platform worldwide. However, since the platform exhibited several security breaches and vulnerabilities during the past few years³, it has a poor reputation in the public eye.

As of spring 2018, there are nine major versions of Android in circulation, ranging from Android Gingerbread, released in 2010, to Oreo, released in 2017. Oreo—Google’s latest Android version—is only running on 1.1% of

¹https://source.android.com/security/reports/Google_Android_Security_2017_Report_Final.pdf

²<https://www.statista.com/statistics/266136/global-market-share-held-by-smartphone-operating-systems>

³<https://www.scmagazineuk.com/android-tops-2016-vulnerability-list-security-industry-says-meh/article/630059/>

all Android devices (February 2018). Together with its predecessor Nougat, they run on about 30% of Android devices, followed by Marshmallow (2015) with 28.1%. Android Lollipop (2014) and KitKat (2013) are installed on 36.6% of all Android devices.

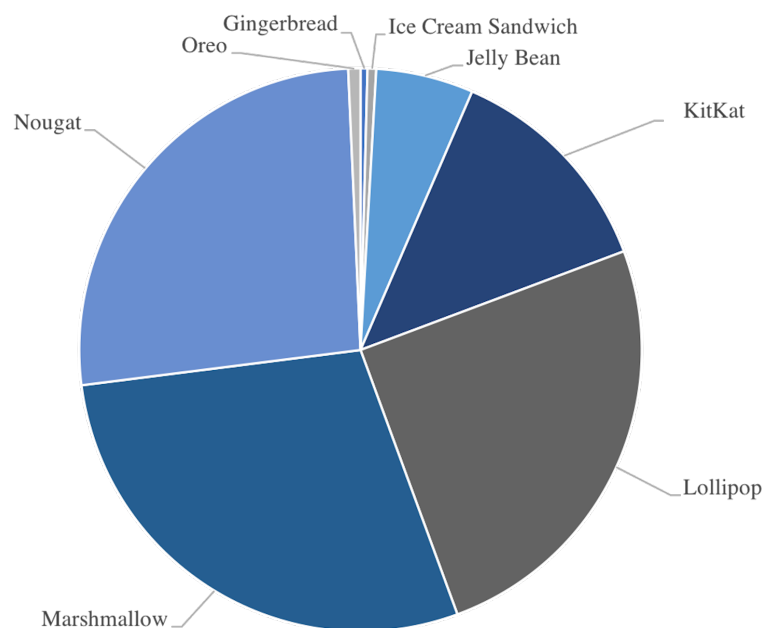


Figure 1.1: Overview of Android OS versions in the field (February 2018)⁴

Over 70%—or more than 1.4 billion Android devices—are more than two years out of date, posing several challenges for both Google and the Android ecosystem.

A major challenge is fixing security vulnerabilities. Due to the huge fragmentation of Android versions, security patches cannot be rolled out in a straightforward way. They need to be compatible with several Android versions. With its latest OS version—Android Oreo—Google tries to improve the state of security and to decrease fragmentation. Initiatives such as Project Treble⁵ or Google Play Protect⁶ shall help tackle these challenges.

⁴<https://c.mi.com/forum.php?mod=viewthread&tid=759720&aid=1668480&from=album&page=1>

⁵<https://source.android.com/devices/architecture/treble>

⁶https://www.android.com/intl/de_de/play-protect/

1.2 Analysis Tools

Not only Google is making a leap forward to improve the state of Android security and to remove vulnerable and malicious applications from online app stores. The research community has proposed a myriad of techniques and approaches over the past years to detect such applications.

Furthermore, there exist large scale taxonomies aiming to list and classify existing tools and approaches on the Android Security domain [33, 78, 80, 90]. However, these taxonomies lack comparisons among similar tools.

To provide such a comparison, we need to obtain the tools. However, we observe that most tools mentioned in literature are not available to the public anymore. One of the root causes is that the fragmentation forces researchers to update their tools on a frequent basis. We note, that this is usually not the case and that most tools are as outdated as the Android version they are built upon.

1.3 Research Questions

The aim of this work is to compare a set of analysis tools on the domain of data leak vulnerabilities for a common set of applications. Unfortunately, during our project, we came to the believe, that research papers, even if they present and discuss a particular tool, are rarely accompanied by an artefact for reproducing the results. Thus, a comparison with other tools is not feasible. This yields in the following research question:

RQ1 *To what degree are security analysis approaches and results in the information disclosure domain in Android applications reproducible?*

To answer this question, we review existing approaches proposed by the research community and try to obtain the corresponding tools through multiple channels (*e.g.*, using search engines or by contacting the authors of the papers). Once a tool is obtained, we try to set it up based on the documentation.

We report findings similar to what Amann *et al.* report in the domain of application programming interface (API) misuse analysis [4]. In their work, the authors ended up evaluating their benchmark on just five tools.

The second part of this paper presents the evaluation and comparison of five static Android security analysis tools on the domain of information disclosure vulnerabilities. To facilitate the comparison and reproduction of results, we present our own implementation of a flexible

benchmark suite that allows others to easily add their tool to the suite. The analysis will then be executed automatically, and the results consolidated with the other tools. Thanks to our implementation, we can run all tools on a common set of applications and gather unified results. This is not only useful for the evaluation, but can also serve as a basis for an analysis frontend that provides normalised reports to the user.

Using the benchmark enables investigation into the second research question, which we state as follows.

RQ2 *How do information disclosure analysis tools perform (individually and compared to each other) on a common set of applications?*

To answer this question, we run the five obtained artefacts on a common set of applications with known vulnerabilities from a given dataset. Then we report on precision and recall for each tool. We hypothesize, that tools with high agreement tend to report the same issues when configured to use the same list of sources and sinks. Thus, we formulate the final research question.

RQ3 *To what extent do reported issues from various approaches overlap?*

We perform a pairwise analysis of the tools' results after running them on the set of applications with known vulnerabilities. Afterwards, we hypothesize which tools tend to report the same findings and evaluate our assumptions on a larger scale with real-world applications.

Answering these questions allows us to identify common weaknesses amongst the tools, which helps us to point out areas in need of further research.

1.4 Outlook

The remainder of this paper is organized as follows. Chapter 2 gives an overview of relevant work. This includes an overview of data leak analysis in Android and related work to analysing and comparing tools.

In Chapter 3 we elaborate our course of action for obtaining, reviewing, and classifying analysis tools for data leak detection. We show which tools were initially considered for the benchmark and how the process of elimination yielded a set of just five tools.

We then explain our benchmarking concept and present our own benchmark implementation in Chapter 4. This implementation aims to facilitate the comparison of tools and increase reproducibility.

We proceed in showing our experimental setup used for evaluating the selected tools in Chapter 5, before presenting the obtained results in Chapter 6.

Finally, we state threats to validity in Chapter 7 and conclude the thesis in Chapter 8.

2

Related Work

Android security analysis is a popular topic. Over the past few years, the research community has done extensive work in this field. Considerable efforts have been devoted in detecting vulnerabilities in Android applications and in the Android platform itself. In most cases, the approach is evaluated using a prototype tool. Thus, the range of existing related works and approaches is exhaustive.

In this chapter, we give an overview of the works that served as a basis for this thesis. It is organised as follows. First, we cover the papers that provide an overview of the state of the art of Android security analysis in general. Second, we give a brief overview of the data leak vulnerability detection and taint analysis tools that are evaluated in this thesis. However, we do this without going into much details and instead refer to Section 3.4 for a detailed presentation of the selected tools. Finally, we point out some existing benchmarking works on comparing software artefacts, particularly for the security domain.

2.1 Android Security Analysis

There exists a vast amount of published work on the topic of Android security analysis. First, it is crucial for us to get a better understanding on what kinds of vulnerabilities to expect in Android applications. Ghafari *et al.* present a study on security smells in Android [45]. They list 28 code smells that may

lead to vulnerabilities in applications and show their prevalence in real-world Android applications. The code smells also cover data leak vulnerabilities and confirm that these issues are in fact common among Android applications.

Second, we focus on the state-of-the-art of Android security analysis. Fortunately, there exist excellent, recent taxonomies that cover a wide range of static, dynamic, and hybrid approaches to Android security analysis.

One of the latest and most extensive taxonomies is presented by Sadeghi, which provides a large-scale overview of Android security analysis in general [80]. Their paper cites more than 500 works in the domain of Android security analysis and provides a detailed categorization of approaches and tools in that domain. Among other lists, the paper provides a list of tools with the analysis objective of vulnerability detection [80, p. 12, Table 3].

This list served as a starting point for our review. In fact, we were unable to find additional approaches that were not already covered by this overview. Moreover, we also cross referenced the list with other published taxonomies such as the works by Sufatrio *et al.* [89] or Reaves *et al.* [78]. All taxonomies presented contain similar lists of references.

2.2 Data Leak Vulnerabilities

In general, data leaks describe the practice to release potentially private or sensitive data. For example, a device’s unique identifier (UID). In this thesis, we are interested in the release of data to unauthorized clients or otherwise to unintended places.

Many approaches, including those used in our evaluation, use taint analysis to determine locations in the code where possibly sensitive data is leaked. In taint analysis, one tries to determine whether sensitive data can flow from a source to a sink.

2.2.1 Sources and Sinks

We define sources and sinks similar to Rasthofer *et al.* in their work on the *SuSi* tool [77].

Definition (Source): Sources are calls into resource methods returning non-constant values into the application code.
(Rasthofer *et al.* [77])

Resource methods in this context are predefined API calls provided by the Android operating system to access shared resources outside the application’s

address space. Therefore, a source is a location within the application's source code, where a resource method is called returning a non-constant value from a shared resource.

A typical example for a source is the *getDeviceID()* method returning a unique identifier of the underlying device such as the *IMEI* (International Mobile Equipment Identity). Such globally unique identifiers can be used for tracking and physical device association. An attacker could for example use the *IMEI* to perform a remote SIM card rooting¹.

Sinks, on the other hand, are locations in the source code, where sensitive information is leaving the current component. Rasthofer *et al.* define it as follows.

Definition (Sink): Sinks are calls into resource methods accepting at least one non-constant data value from the application code as parameter, if and only if a new value is written or an existing one is overwritten on the resource.
(Rasthofer *et al.* [77])

Therefore, a sink is considered when using the Android API calls within the application's source code to write to shared resources.

A typical example of a sink is the *sendTextMessage(String message)* method with the message text being non-constant. However, a sink is only problematic if the values passed as arguments are actually sensitive information, such as a unique identifier. In the *sendTextMessage(String message)* example above, we could also pass a random String message not containing any sensitive information. In that case, the sink would not be leaking sensitive information and therefore is no data leak.

Taint analysis requires a list of sources and sinks to check possible data flows. *SuSi* is an approach to automatically generate such lists based on an *android.jar* artefact [77].

The *android.jar* file is a compiled version of the Android SDK. It mainly serves as a compile-time library for 3rd party applications and is available for different API levels. We use the *SuSi* tool in our benchmark to obtain a neutral configuration and to reduce bias stemming from individual lists shipped with the respective tools.

Employed techniques for taint analysis include, for example, static data-flow and control-flow analysis. Besides approaches in taint analysis, there

¹<https://threatpost.com/mobile-applications-leak-device-location-data/121392/>

exists a wide range of approaches and tools that aim to detect data leak vulnerabilities.

We abstain from discussing individual approaches here. Instead, we refer to the work we considered for the evaluation later in this thesis and point the interested reader to the exhaustive, generic taxonomies on Android security analysis [80, 89, 78]. In Section 3.4, we also present the five tools included in this benchmark.

2.3 Comparing Software Artefacts

There are several pieces of relevant work related to analysing and benchmarking software artefacts.

In 2016, Amann *et al.* have analysed artefacts for detecting API misuse violations [3]. Their approach is similar to ours in that they developed a framework for comparing such tools. Furthermore, the authors first reviewed corresponding literature and obtained artefacts in a similar fashion as we did.

Hoffmann *et al.* have evaluated several static and dynamic analysis tools with regard to obfuscation [50]. They use synthetic sample applications and automatically obfuscate them before running the analysis. In contrast, we do not provide our own samples; instead, we rely on *DroidBench*², a database of synthetic applications with known vulnerabilities. In addition, we evaluate the artefacts on real-world applications.

DroidBench is a collection of synthetic applications intended for evaluating analysis tools. Because vulnerabilities are documented in these applications, they are well-suited for the qualitative analysis we present in this work. Using this benchmark, not only can we identify true positives easily, but, due to the limited size of the synthetic programs, we can determine false positives as well.

The work by Reaves *et al.* comes closest to ours. In their study, they use *DroidBench* to analyse results obtained from *Amandroid*, *AppAudit*, *DroidSafe*, *Epicc*, *FlowDroid*, *MalloDroid*, and *TaintDroid* [78]. They do not, however, cover *IC3*, *COVERT*, *IccTA*, or *HornDroid*. In contrast to our work, the evaluation is lacking a comparison of tools amongst each other. Furthermore, the authors do not report on the amount of vulnerabilities that were detected.

²<https://github.com/secure-software-engineering/DroidBench>

3

Classification and Selection

In this chapter, we state hypotheses about the outcome of our benchmark. Furthermore, we explain how we selected the tools included in this work and what questions and limitations led to our choice. In addition, we present the five selected data leak detection tools.

3.1 Hypotheses

Based on the number of true positives, false positives, true negatives, and false negatives in a tool's report, we can formulate several hypotheses about the tools and the quality of their results.

- H1** *As more tools report the same vulnerability, the likelihood of it being a true positive increases.*
- H2** *The less tools report a certain vulnerability, the more likely is a false positive*
- H3** *If a tool reports much more data leaks than all the other tools, then it is likely to report more false positives than the others.*

We validate our hypotheses in Section 6.4.

3.2 Preliminary Questions

There are two preliminary questions we had to answer before starting with the selection of the tools.

Q1 *On which domain of Android security analysis do we want to focus?*

Q2 *On what level do we want to compare the tools and in what metrics are we interested?*

For the first question, we checked different domains of Android Security analysis. In general, our work is not limited to a single domain and could be easily extended to other domains. However, due to the large number of published works for static approaches and to ensure comparability of the tools, we decided to focus on the domain of data leak vulnerabilities.

To answer the second question, we refer to the literature, where common metrics in the field of software comparison are precision and recall. These metrics allow us to draw conclusions on a tool's performance and to easily compare them amongst each other. In order to calculate the precision and recall, we require the number of true positives/negatives and false positives/negatives. So, the number of reported vulnerabilities of a tool plays a key role in this benchmarking.

The identification of true/false positives and negatives also has implications on the dataset of applications used in this work, since we need to be able to identify them within the applications source code. The selection of the dataset of applications is further discussed in Section 5.1.

3.3 Obtaining Tools

In this benchmark, we take advantage of two recent taxonomies in the security analysis domain.

Sadeghi *et al.* have collected 517 references to tools and approaches in their taxonomy on security related Android analysis software [80]. Within this paper, the authors classify the tools and approaches according to different categories, such as the type of sensitivity or the type of program analysis (static, dynamic, hybrid, formal, machine learning). We use Sadeghi's work as a starting point for our tool selection. Hence, for the initial selection of tools, we chose the 87 references to vulnerability detection tools presented by the paper [80, p. 12, Table 3].

We compare Sadeghi’s list of references with two other taxonomies with mostly overlapping sources [78, 89], and arrived at the following list of artefacts: *ASM* [49], *Addicted* [80], *Amandroid* [97], *ApkCombiner* [60], *AppAudit* [100], *AppCaulk* [82], *AppCracker* [17], *AppFence* [51], *AppGuard* [8], *AppProfiler* [79], *AppRay* [91], *AppSealer* [102], *Aquifer* [72], *AuthDroid* [96], *Bagheri* [11], *Bartel* [12], *Bartsch* [13], *Bifocals* [23], *Buhov* [16], *Buzzer* [20], *CMA* [85], *COVERT* [9], *CRPE* [25], *CoChecker* [28], *ComDroid* [24], *ConDroid* [83], *ContentScope* [54], *Cooley* [26], *Copes* [80], *CredMiner* [107], *CryptoLint* [34], *Dare (Enck)*¹, *Desnos* [30], *DexDiff* [67], *DroidAlarm* [106], *DroidCIA* [22], *DroidChecker* [21], *DroidGuard* [10], *DroidRay* [105], *DroidSearch* [76], *FineDroid* [103], *FlowDroid* [6], *Gallo* [42], *Geneiatakis* [43], *GrabNRun* [36], *Harehunter* [1], *HornDroid* [18], *IC3 (Epicc)* [73], *IPCInspection* [40], *IVDroid* [37], *IccTA* [59], *Juxtapp* [48], *KLD* [84], *Kantola* [56], *Lintent* [15], *Lu* [63], *MalloDroid* [35], *Matsumoto* [65], *Mutchler* [71], *NoFrack* [44], *NoInjection* [55], *Onwuzurike* [74], *PCLeaks* [58], *PaddyFrog* [98], *PatchDroid* [69], *PermCheckTool* [95], *PermissionFlow* [81], *Poeplau* [75], *Pscout* [7], *QUIRE* [31], *Ren* [52], *SADroid* [47], *SCanDroid* [41], *SEFA* [99], *SMVHunter* [88], *STAMBA* [14], *Scoria* [93], *SecUp* [101], *Smith* [87], *Stowaway* [39], *Supor* [53], *Tongxingli* [62], *Vecchiato* [94], *VetDroid* [104], *WeChecker* [29], *Woodpecker* [46], *Zuo* [108]. Note that in cases where the approach or tool has no proper name, we use the first author’s last name instead.

In a next step, we further categorized and filtered the tools. First, we tried to obtain all artefacts by employing the following strategy:

1. Review the paper, look for links or directions on how to obtain the artefact,
2. look for the artefact on author’s or research group’s websites,
3. search online with contemporary search engines for the artefact, and
4. contact the authors and inquire whether the tool is available or can be made available to us (at most two requests by email within three weeks).

By reviewing the corresponding paper for each tool, we identified 13 tools that cannot be executed because they are formal models or frameworks only and thus are not suitable to be included in the benchmark. For 14 tools, we could obtain the artefact by reviewing the paper or using contemporary search engines.

¹<http://siis.cse.psu.edu/ded/>

For the remaining 61 tools, we have contacted the researchers to request access to their artefacts. Out of those 61 requests, 49 remained unanswered and four researchers refused to give us access due to the commercial usage of their tool.

This results in the following list of 22 available and obtainable tools: *Amandroid*, *APK-Combiner*, *AppCaulk*, *AppGuard*, *Aquifer*, *ASM*, *ConDroid*, *COVERT*, *CRePE*, *Dare (Enck)*, *DexDiff*, *IC3 (Epicc)*, *FlowDroid*, *Geneiatakis*, *Grab n’Run*, *HornDroid*, *IccTA*, *Lintent*, *MalloDroid*, *NoFrack*, *PScout*, and *SCanDroid*.

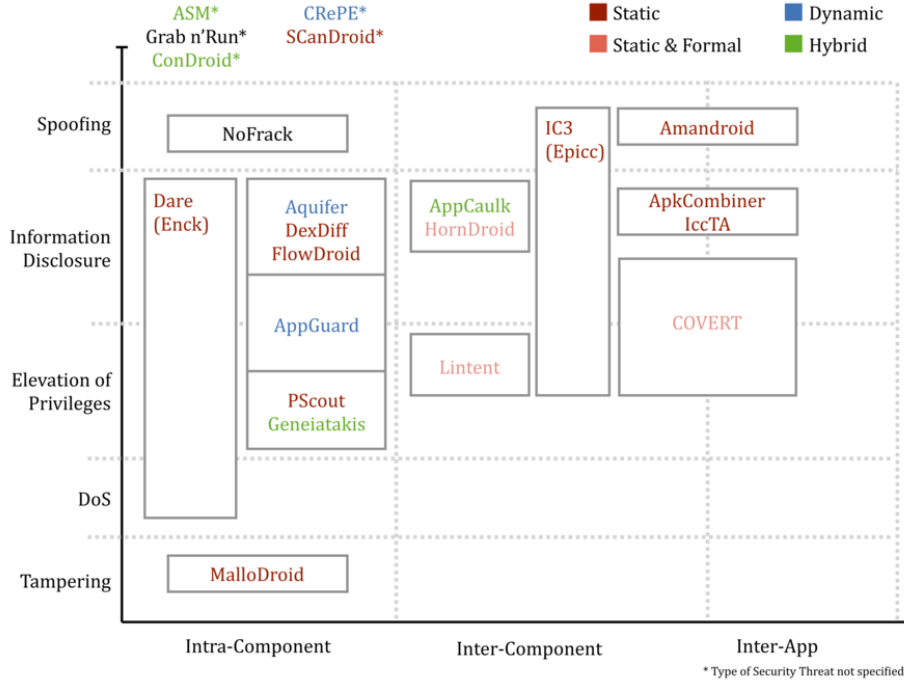


Figure 3.1: Overview of tool categorization. Tools marked with an asterisk are not classified in Sadeghi’s work.

Figure 3.1 provides an overview of the obtainable tools categorized according to the type of security threat, type of approach and scope of the analysis. In order to compare the tools, they need to focus on the same type of vulnerability. As can be seen in Fig. 3.1, the domain of information disclosure or rather data leak vulnerabilities holds most tools with 11 approaches compared to escalation of privileges with just 7 tools. That is why we selected data leak vulnerabilities as the domain for this work.

From the artefacts available, we further filtered out 17 tools for the following reasons:

- **Dynamic or hybrid approach:** In this benchmark, we focus solely on static approaches for comparability reasons. Dynamic or hybrid approaches include *AppCaulk*, *AppGuard*, *Aquifer*, *ASM*, *ConDroid*, *CRePE*, *Geneiatakis*, *Grab n’Run*.
- **Not in the selected domain:** Tools not in the information disclosure domain include *Amandroid*, *Lintent*, *MalloDroid*, *NoFrack*, *PScout* and *SCanDroid*.
- **Analysis enabler:** *ApkCombiner* is used to combine several APKs into one APK to enable inter-app communication analysis. *Dare (Enck)* is used to decompile Android applications from the installation image to source code on which other analysis tools can perform. Neither of the two tools are reporting any data leaks and thus are excluded from the benchmarking.
- **Setup failed:** We failed to setup *DexDiff*. This is mainly due to poor documentation and lack of support by the responsible researcher.

Out of the initial list of 87 papers and their corresponding artefacts, we could obtain 22 tools. We selected data leak detection as the target domain and therefore selected five tools from these 22 that are suitable for benchmarking.

3.4 Selected Tools

In this benchmark, we will focus on the following five tools: *COVERT*, *FlowDroid*, *HornDroid*, *IccTA* and *IC3*. Table 3.1 provides an overview of the five selected tools and their main characteristics.

FlowDroid is a taint analysis based on the *Soot* and *Heros* frameworks. The authors report high rates of precision and recall. Especially its call graph construction procedure is popular, since it helps increasing the context-, flow- and object-sensitivity and is therefore used by other tools such as *IccTA*, *COVERT*, or *IC3*.

To model application lifecycle, *FlowDroid* uses a list of source and sinks. To generate such a list, the researches have created the *SuSi* tool [77], which we also applied in this benchmark. It takes an *android.jar* as an input and produces a list of all sources and sinks found using machine learning algorithms. Along with the list of sources and sinks, *FlowDroid* also uses the

Name	COVERT	FlowDroid	Horn-Droid	IccTA	IC3
Type	static, formal	static	static, formal	static	static
Scope	Inter-App	Inter-Component	Inter-Component	Inter-Component	Inter-Component
Artefact	Manifest, Bytecode	Manifest, Layout, Bytecode (DEX)	Bytecode	Manifest, Layout, Bytecode (DEX)	Manifest, Bytecode
Sensitivity	context, value	object, context, flow, field	object (p), context, value, flow (p), field (p)	object, context, flow, field	context, flow, field
Sources and sinks	yes	yes	yes	yes	yes

(p) partial

Table 3.1: Overview of main tool characteristics

manifest file, dex files and layout files to perform the analysis. In order to ensure flow-sensitive information flow analysis, it creates a dummy main method emulating each possible interleaving of the callbacks during the application lifecycle.

In addition to static analysis, *HornDroid* is also using a formal approach. It translates the application into so called horn clauses, which are used to abstract the program semantics. These horn clauses are then discharged using an off-the-shelf theorem prover (Z3)². This formal proof of soundness helps excluding unreachable program parts. *HornDroid* uses Dalvik VM bytecode as a base artefact. It is value- and context-sensitive, and partially flow-, field-, and object-sensitive. So, it can skip unreachable program parts, approximate runtime values, compute different static approximations upon different method calls, take the order of statements into account, and approximate static fields. *HornDroid* uses a list of sources and sinks supplied by the *DroidSafe* tool. As a known limitation, *HornDroid* does not find any implicit information flows and might produce false positives for heap abstractions, exceptions and inter-app communication.

IccTA is an inter-component communication based taint analysis tool. However, the approach is generic and promises to be applicable to any data-

²<https://github.com/Z3Prover/z3>

flow analysis. It reports precise tracks of information flow from source points to sinks. First, *IccTA* decompiles the Android app in the form of Dalvik bytecode into Jimple using the *Dexpler* decompiler. Jimple is an internal representation of the *Soot* framework. *IccTA* then extracts all ICC methods using *Epicc*. Then it uses *IC3* to parse the URIs in order to support content provider related ICC methods. The target components are retrieved, based on the manifest file and bytecode. It then connects components to enable a data-flow analysis between the different components, on which an intra-component analysis is performed. For that purpose, *IccTA* uses a modified version of *FlowDroid*. This yields in a control flow graph of the whole application. The tainted paths or leaks are stored in a database to reuse results. Similarly to *FlowDroid*, *IccTA* is a context-, object-, flow-, and field-sensitive analysis.

IC3 detects inter-component communication with a special focus on inferring values of complex objects with multiple fields such as intents or URIs. First, it decompiles the Android source code to Java bytecode using *Dare*. Then it generates an Inter-procedural Control Flow Graph (ICFG) using the *FlowDroid* call graph construction procedure. This ICFG and the so-called COAL (constant propagation language)³ specifications describing the structure of the composite objects are passed to a COAL Solver. This COAL Solver uses the *Heros IDE* to resolve the Inter-procedural Distributive Environment Problems (IDE). In a last step, it performs argument value analyses such as String analysis to determine the values of a function's argument. *IC3* is a context-, flow- and field-sensitive analysis and replaces the previous tool *Enck*.

COVERT is also a static and formal security analysis. However, its main focus is inter-app communication and escalation of privileges. It tries to find unsafe combinations of apps exploiting these vulnerabilities. In order to also cover the topic of information disclosure, it was extended with a static information flow analysis, performed by *FlowDroid*. *COVERT* consists of a model extractor and a formal analyzer. The model extractor takes the applications manifest file and byte code and specifies a first model of the potential inter-process communications. In a second step, a reachability analysis is performed to determine what permission required functionalities are being exposed by unguarded execution paths. This results in an extended manifest file for each application. Those extended manifest files are transformed into so called Alloy modules using the Alloy language (formal modeling language) and then fed to the formal compositional *Alloy Analyzer*. The analyzer proofs which combinations of applications are safe and for which combinations the

³<http://siis.cse.psu.edu/coal/index.html#banner>

unguarded execution paths can be exploited. The benefit of *COVERT* is the inter-app communication, which is not possible with most static analysis tools. *COVERT* is a value-, context-and flow-sensitive analysis. *COVERT* also provides a desktop client; however, we did not manage to set it up and execute it. According to the researchers, the desktop client is outdated and not supported any more. Instead, one should only use the command line to perform the analysis.

4

Benchmark Implementation

In this chapter, we explain our benchmarking concept and present our Java tool to perform the analysis on all selected tools. Our benchmarking implementation merges all output files and provides a homogeneous overview of the reported vulnerabilities. It allows us to automate the analysis, which is necessary for larger scale analyses where manual inspection is infeasible.

4.1 Benchmarking Concept

This benchmark compares the tools' performances in practice. To do so, we need to run the analyses on a common dataset of applications. However, the size and structure of the dataset of applications has implications on what information we can deduce from the analysis. A smaller dataset of applications allows for a more thorough analysis with manual validations of the results. A larger dataset will improve the validity of the results, but will make manual investigations of the results much more time consuming.

Synthetic applications with seeded vulnerabilities are useful for determining precision and recall; real-world applications, on the other hand, provide a realistic setting at the expense of not being able to determine those metrics. The two targets complement each other and we use both in our evaluation.

Therefore, this work is built on two pillars. The first pillar is a smaller scale qualitative study, which is done by running the selected tools on a small scale dataset with known data leak vulnerabilities. Based on the analyses

results, we can draw conclusions on the quality of the results and formulate hypotheses on how the tools might perform in a bigger scale. In this first pillar, the focus lies on true/false positives and negatives and related metrics such as precision and recall.

The second pillar is a large scale quantitative study, for which we run the tools on a set of real-world applications from the *F-Droid* app repository. This gives us insights on the tools' performance on real-world applications and helps us to validate the hypotheses. Here, the number of reported vulnerabilities and the overlap with other tools play a key role. Moreover, this analysis provides insights into the general state of real-world applications concerning data leak vulnerabilities.

4.2 Benchmarking Java Implementation

Besides being useful for summarising the tools' outputs and thus helping us to compare the results, there are other potential use cases. Especially developers, who want to analyse their applications before publication can profit from our implementation. The original reports generated by the tools are often hard to find (e.g., in a nested folder or printed to standard output and error streams). Furthermore, they are generally cryptic and over-populated with additional information such as call paths or register values. An example of such a report produced by the *IccTA* tool can be seen in Listing 4.1. It shows the report for a single vulnerability.

Listing 4.1: Example excerpt of console output by *IccTA*

```
Found a flow to sink virtualinvoke $r3.<org.cert.sendsms.MainActivity: void startActivityResult(android.content.Intent,int)>($r2, 0) on line 26, from the following sources: - $r6 = virtualinvoke $r5.<android.telephony.TelephonyManager: java.lang.String getDeviceId()>() (in <org.cert.sendsms.Button1Listener: void onClick(android.view.View)>) on Path [$r6 = virtualinvoke $r5.<android.telephony.TelephonyManager: java.lang.String getDeviceId()>(), virtualinvoke $r2.<android.content.Intent: android.content.Intent.putExtra(java.lang.String,java.lang.String)>("secret", $r6), virtualinvoke $r3.<org.cert.sendsms.MainActivity: void startActivityResult(android.content.Intent,int)>($r2, 0)] - $r4 = virtualinvoke $r3.<org.cert.sendsms.MainActivity: java.lang.Object getSystemService(java.lang.String)>("phone") (in <org.cert.sendsms.Button1Listener: void
```

```
id onClick(android.view.View)>) on Path [$r4 = virtualinvoke
$r3.<org.cert.sendsms.MainActivity: java.lang.Object getSystemService(java.lang.String)>("phone"), $r5 = (android.telephony.TelephonyManager) $r4, $r6 = virtualinvoke $r5.<android.telephony.TelephonyManager: java.lang.String getDeviceId()>(), virtualinvoke $r2.<android.content.Intent: android.content.Intent.putExtra(java.lang.String,java.lang.String)>("secret", $r6), virtualinvoke $r3.<org.cert.sendsms.MainActivity: void startActivityForResult(android.content.Intent,int)>($r2, 0)]
```

As a result, the exact location of the leaky code is hard to make out within the generated reports. Our wrapper, on the other hand, can present the results in a standardised way, focusing solely on the location of the data leak within the code. Thus, we only present the name of the class and method containing the leaky line of code and the exact sink (method call) that led to the data leak. We abstain from showing additional information such as register values or call paths.

Since our approach can easily be extended with additional analysis tools, it could make a wide range of program analysis more accessible. In addition, there are currently only very few tools covering multiple security domains. Consequently, the end user would need to identify several analysis tools covering several Android security domains to get a complete picture of the security state of the under-lying application. Contrarily, our Java tool can be extended with tools covering other security domains and thus simplify the user's life by providing a single interface. The user will not have to search for available tools and can launch all provided tools with just one click receiving all results in a standardised and understandable way.

Listing 4.2 shows a sample output of our Java tool for the same vulnerability presented in listing 4.1.

```
Class Name:      org.cert.sendsms.Button1Listener
Method Name:     onClick(android.view.View) void
Sink Method:     startActivityForResult(Intent,int) void
Found by Tool:  [iccta, horndroid]
```

Listing 4.2: Example Java tool output

The benchmarking implementation consists of runners and parsers. Each tool artefact included in the benchmark must extend and implement these interfaces. A tool runner is a simple Java class that specifies the tools main characteristics such as commands used to run it in the shell, location of the results files and how to get these results. Furthermore, a tool parser needs

to be implemented that parses generated output and creates reports in the form mentioned above (class name, method name, sink). So, to include a new tool in the benchmark, one should follow these steps:

1. Obtain the artefact and adapt the tool to run from the command line (which is already the case for most tools).
2. If the tools do not already create a separate results file, then adapt the command to store the results in such a file.
3. Extend the abstract `Tool` class by specifying the command used to trigger the analysis and the location of the results file.
4. Extend the abstract `Parser` class that obtains the relevant data.
5. Provide a path to the application (.apk file), that should be analyzed.

The benchmark is implemented in Java and available from the supplementary website¹. While we do not redistribute the integrated tools' artefacts, we do provide detailed instructions on how to set them up.

Launching our benchmark tool will sequentially copy and paste the command used to trigger the analysis for each tool to the shell and execute it. Afterwards, all results files are gathered and parsed for class name, method name and sink method. This results in a list of all leaks detected by the tools in a standardised format. However, we want a summarised list, so that if two tools detect the same vulnerability, we do not have several entries in our list. Therefore, we are first sorting the list and then summarising entries that have the same class name, method name and sink method by adding the tool name to the indicated list of tools that detected the vulnerability. Finally, we are storing the summarised and standardised list of vulnerabilities in a text file and in a csv (comma separated values) file. If the analysis fails for one of the tools, the others are still performing the analysis and the results summary will simply miss the results of the failed tool.

The analysis of real-world applications is time consuming and may take hours or even days to complete. In our benchmark, the user can set a time out for the analysis being aborted after a certain number of minutes and then proceeding with the next tools. If a tool runs longer than the specified timeout, it will be aborted.

¹<https://github.com/tiimoS/droid-Security-Thesis>

5

Experimental Setup

To benchmark our five selected tools—*IccTA*, *IC3 (Epicc)*, *HornDroid*, *FlowDroid* and *COVERT*—we have integrated them in our Java tool. In this chapter, we describe the experimental setup in terms of specific configurations and database selection used in the evaluation.

5.1 Database Selection

As many others [32, 68, 78, 2, 92, 19, 57, 5, 61, 97, 38, 70] in recent research into Android security analysis, we have selected *DroidBench*¹ as a test suite for our small scale qualitative study. *DroidBench* consists of 119 synthetic applications. They specifically target different data leak vulnerabilities with a total of 125 leaks, which are described and indicated directly in the source code. These synthetic applications usually only consist of two to three classes and are therefore very small in size. The leak indication within the source code allows us to review the generated reports and identify the true and false positives and negatives. Based on these metrics, we can calculate precision and recall.

For the large scale quantitative study, we have selected *F-Droid*². *F-Droid* is an online application store for open-source software. However, due to the long run time of the analyses with larger applications (which often exceeds

¹<https://github.com/secure-software-engineering/droidbench>

²<https://f-droid.org>

three hours, depending on the tool), we only analyse a random sample of 250 applications. Compared to the *DroidBench* applications, we have no information about possible data leaks, since they are not indicated in the source code.

With these two databases of applications at hand, we perform the analysis as follows:

1. Run all tools on the small *DroidBench* dataset and manually inspect the applications and results for the number of true and false positives and negatives.
2. Run our benchmarking implementation on the same *DroidBench* dataset and compare the summarised results with the results from the first run. We evaluate the quality of our parsers and results summarising implementations in Section 7.1.
3. Run the benchmarking implementation on the large *F-Droid* dataset and check the number of reported vulnerabilities and overlaps among the tools.

5.2 Tools Configuration

For a fair and unbiased comparison, one needs to make sure that the tools are running with similar configurations and are based on similar resources. We have made out the following resources that the selected tools commonly use:

- A list of sources and sinks,
- a list of Android callback methods to simulate the Android lifecycle,
- an Android SDK (`android.jar`), and
- a specific version of the *Apktool*³ used to reverse engineer APK files.

5.2.1 List of Sources and Sinks

As already elaborated in Section 2.2.1, sources and sinks play a key role in data flow analysis. All the tools included in the benchmark, except for *IC3*, are providing their own list of sources and sinks. These lists vary heavily in length, namely from a few hundred entries to several thousand entries.

³<https://ibotpeaches.github.io/Apktool/>

COVERT is using an additional filter list to parse the sources and sinks list. Mostly, these lists of sources and sinks are based on specific Android OS versions. Hence, one has to provide their own list or use a tool such as *SuSi* [77] to generate such a list. Listing 5.1 shows two typical entries from a list of sources and sinks.

```
<android.telephony.TelephonyManager: java.lang.String getDeviceId(int)> -> _SOURCE_
<android.os.Handler: boolean sendMessage(android.os.Message)> ->
    _SINK_
```

Listing 5.1: Example sources and sinks list entry

5.2.2 Original Configurations

Table 5.1 gives an overview of the different tools' original configurations.

The number of callback methods is similar for all tools. The biggest differences in the configurations is the number of entries in the list of sources and sinks. With over forty thousand entries, *HornDroid* has by far the most exhaustive list. The list stems from the *DroidSafe*⁴ tool. *HornDroid* is followed by *COVERT* with just a bit over twenty five thousand entries. The other tools only have a few hundred entries in their lists of sources and sinks. For *IC3* there is no list provided.

The list of Android callback methods used by all tools except *IC3* is required to model the Android lifecycle during the analysis. The number of entries is similar for all tools.

FlowDroid and *IC3* do not provide an *android.jar* file for the analysis. Hence, the user will have to provide the path to a valid Android SDK in the command used to trigger the analysis. *COVERT* and *IccTA*, on the other hand, are shipped with different versions of the *android.jar* that are provided together with their artefact. *IccTA* uses Android 4.3 (Jelly Bean) from 2013. For *COVERT*, the version of the Android file could not be identified. But based on the year of the paper publication, it cannot be newer than Android 5.1 (Lollipop) from 2015.

COVERT and *HornDroid* are using the *Apktool* to decompile the Android applications.

⁴<https://github.com/MIT-PAC/droidsafe-src>

Name	Flow-Droid	Horn-Droid	COVERT	IC3	IccTA
Sources and sinks	355	40 986	26 591	N/A	160
Callbacks	182	183	185	N/A	182
Android version	user	N/A	≤ 5.1	user	4.3
Apktool	N/A	2.3.0	1.5.2	N/A	N/A

Table 5.1: Overview of tool configurations. The value "user" indicates that the user has to provide the path to the artefact in the running command.

5.2.3 Shared Configurations

For the benchmark we use Android 6.0 (Marshmallow, API Level 23). Later versions such as Android Nougat or Oreo have caused *COVERT* to crash during the analysis. To identify sources and sinks for this Android version, we use the *SuSi* tool. It takes the *android.jar* as an input and based on machine learning algorithms generates a list of sources and sinks. For Android 6.0 (API Level 23), the list contains 4.050 entries.

In this list, there are entries that show additional tags after the *SOURCE* indication (see Listing 5.2). These tags have caused crashes with the *HornDroid* tool.

```
<android.telephony.TelephonyManager: android.telephony.CellLocation
  getLocation()> android.permission.ACCESS_FINE_LOCATION
  android.permission.ACCESS_COARSE_LOCATION -> _SOURCE_|LOCATION
  INFORMATION
```

Listing 5.2: Problematic source entry for HornDroid

Therefore, we have built an additional pre-processor for *HornDroid* that parses the list of sources and sinks and eliminates all these additional tags, resulting in a special list just for *HornDroid*.

For the list of Android callback methods, there are to the best of our knowledge no tools available to generate such a list. Hence, we have merged the existing lists from *FlowDroid*, *IccTA*, *COVERT*, and *HornDroid* to get a base list for our shared configurations.

For the *Apktool*, we are using the latest version 2.3.1.

Configuring the tools in such a way allows us to compare their performance solely based on the approaches and algorithms. However, it is important to note that although it is possible to configure the tools as described, the out-of-the-box experience may differ considerably. End-users that are

unaware of the internals of such tools may thus never be able to properly configure tools and miss reports. In Section 6.2.3, we evaluate the differences in running the tools with original or with shared configurations.

5.2.4 Switching Between Configurations

In the benchmark, we want to easily change between shared and original configurations. For this purpose, we have created soft links for all shared resources and have created two shell scripts. Before running the benchmark implementation, the user can decide whether to run it in original or shared configuration by simply running the corresponding shell script. The scripts will then change the soft links to either point to shared resources or to the resources that the tools were provided with. This allows us to simplify the comparison between shared and original configurations.

5.3 Evaluation Process

The following section presents the evaluation process for the two datasets of applications used in this benchmark.

5.3.1 DroidBench Evaluation Process

For the small scale qualitative analysis on the *DroidBench* applications, we run our benchmark twice. Once with shared configuration and once with original configuration. After the first test run, we manually review the obtained reports. For each *DroidBench* application, we proceed as follows:

1. Create a list of all known vulnerabilities for all *DroidBench* applications by scanning the source code of each application for the *DroidBench* author's sink indication.
2. Compare the list of known vulnerabilities with the reports generated by the tools and record which tools found the leak (true positives) and which ones did not (false negative).
3. For each additional report, record the false positives and true negatives.

After processing 125 DroidBench vulnerabilities, we have a complete list of all known vulnerabilities in the *DroidBench* test suite and for each tool a list of true positives, true negatives, false positives, and false negatives. This allows us to report on each tool's precision and recall.

Precision expresses the amount of reported real vulnerabilities amongst all reports, while recall is the ratio between reported vulnerabilities and all existing vulnerabilities. Thus, we can compare the performance of our five tools, which is not feasible by simply reviewing existing literature. Furthermore, we apply the *McNemar's Test* [66] which can be used since our tools run on the same configurations and the same dataset. *McNemar's Test* is a statistical test for determining whether two tools are likely to report similar issues.

In the second test run, we use the initial, original configuration. Then we calculate precision and recall and compare the results with the shared configurations. Thus, conclusions on the effect of the change in configurations can be made.

5.3.2 F-Droid Evaluation Process

For the large scale quantitative analysis, we run the benchmarking implementation on the *F-droid* dataset. Since we do not identify the true and false positives and negatives, we can only report on the number of reported vulnerabilities and the number of overlaps among the tools. Furthermore, as the applications analysed in this step are real-world applications from the *F-Droid* app repository, we can draw conclusions on the security of those applications. Finally, we evaluate the hypotheses stated in section Section 3.1.

6

Results

In this chapter, we report on the qualitative and quantitative evaluations. We compare the tools' performance in terms of precision, recall and accuracy and interpret the results from the McNemar's test. Furthermore, we investigate the differences of the results when running the tools with shared configurations or with the original out-of-the-box configurations.

In the second part of this chapter, we interpret our findings from the larger scale quantitative analysis and evaluate our hypotheses from Section 3.1.

6.1 Small Scale Qualitative Analysis

For the small scale qualitative analysis, we evaluate and interpret the results gathered from the *DroidBench* dataset analysis. We first look at the results of the analysis with shared configurations. Hence, we focus on the individual tool's performance and then on the combined performance, where we compare the performance of our benchmark implementation with the individual tools. In the end, we compare the tools pairwise.

In Section 6.2.3, we also present the results for the analysis with original configurations and try to interpret potential differences to the results with shared configurations.

6.1.1 Individual Performance

The manual analysis of the *DroidBench* test suite resulted in a list of 125 vulnerabilities that analysis tools ought to detect. We ran the five tools—*COVERT*, *FlowDroid*, *HornDroid*, *IccTA*, and *IC3*—on the *DroidBench* dataset and added reported leaks that are not documented in *DroidBench* (which are potential false positives). This resulted in a list of 269 distinct reports, including the aforementioned 125 real vulnerabilities. We reviewed the list of reports and obtained a Boolean value indicating whether *DroidBench* considers the finding a vulnerability or not. For each of the 269 reports, we indicate whether a particular tool has reported this vulnerability or not.

This allows us to easily make out the number of true positives (TP) and true negatives (TN), as well as false positives (FP) and false negatives (FN) for each tool.

We consider a reported vulnerability that is not described in *DroidBench* a false positive. We further manually checked all 144 potential false positives and verified them not to be true positives. We note that the tools only report on potential vulnerabilities and thus do not report negatives explicitly.

Table 6.1 summarises the results.

Metric	<i>Flow-Droid</i>	<i>Horn-Droid</i>	<i>COVERT</i>	<i>IC3</i>	<i>IccTA</i>
TP	99	99	8	4	97
FP	54	87	3	37	59
TN	90	57	141	107	85
FN	26	26	117	121	28

Table 6.1: Raw counts of true/false positives/negatives.

As can be seen in Table 6.1, *IC3* reported the least true positives and therefore has a poor performance on the *DroidBench* dataset compared to the other tools. Out of 41 reported data leaks, only 4 reports were actually true positives. The remaining 37 reports were all false positives. Hence, only close to 10% of the reported vulnerabilities were in fact true. This makes it the only tool reporting more false positives than true positives. We further observe that the tool only reports leaks in certain *DroidBench* data leak categories such as *inter-component communication*, *inter-app communication* and *emulator detection*. Since most of the other tools are reporting leaks for all categories, it becomes clear why *IC3* underperforms in comparison.

With just 11 reports, *COVERT* reported the least findings of all tools. Out of those 11 reports, 8 reports were in fact true positives. With slightly

over 70% of the reported vulnerabilities being true, *COVERT* is already a strong improvement compared to *IC3* in terms of credibility of the results. Similar to *IC3*, *COVERT* does not report leaks for all categories of data leak vulnerabilities present in *DroidBench*. As we will observe later on in this chapter, it only reports data leaks in 2 out of 13 categories. We further note that *COVERT* and *IC3* overall produce far less reports than the other tools, even in the categories where all tools are reporting data leaks.

For *FlowDroid* and *IccTA*, we observe similarly high numbers of true positives and false positives. Both tools report more than 12 times more true positives than *COVERT* and *IC3*. However, they also report more false positives.

HornDroid and *FlowDroid* report most true positives. Nonetheless, *HornDroid* also reports most false positives. As a known limitation, *HornDroid* might produce false positives for heap abstractions, exceptions and inter-app communication, which might be an explanation for the higher number of false positives.

Precision

To evaluate the quality of the reports, we calculate the precision, recall, and accuracy of each tool.

The precision of a tool is the ratio of correctly reported vulnerabilities to all reported vulnerabilities $((TP)/(TP + FP))$. Thus, it helps us to answer the question of how many of the reported vulnerabilities are in fact real data leaks.

The closer precision, recall, and accuracy are to 1, the better is the quality of the results. Table 6.2 shows the values for each tool.

Metric	<i>Flow-Droid</i>	<i>Horn-Droid</i>	<i>COVERT</i>	<i>IC3</i>	<i>IccTA</i>
Recall	0.792	0.792	0.064	0.032	0.776
Precision	0.647	0.532	0.727	0.098	0.622
Accuracy	0.703	0.580	0.554	0.413	0.677

Table 6.2: Collected metrics for each tool as observed on *DroidBench* vulnerabilities. Bold values indicate maxima for the respective metric.

In terms of precision, *COVERT* has the best performance with over 72%. Accordingly, most of its reported vulnerabilities are actually real data leaks. One reason for this good performance can be that *COVERT* creates a formal

model of the security specifications of each application before starting the analysis. This allows the tool to find critical security properties that ought to be analysed.

FlowDroid and *IccTA* have a slightly poorer precision with over 62%, but are still on a decent level. One of the reasons for the good performance in terms of precision is their aim for a flow-sensitive analysis. Flow-sensitivity allows the tool to take the order of statements into account during the analysis. Asynchronous callbacks in Android applications is a major challenge for this case. *IccTA* and *FlowDroid* are both using control flow graphs of callbacks implemented in the application to make out the program paths of the application. Calls to Android APIs can, however, have multiple callbacks with multiple possible sequences of such callbacks depending on the call site. To handle these cases, both tools create a *dummy main method* emulating the application lifecycle with all possible callbacks. This is crucial to ensure flow-sensitivity and high precision. However, the construction of this dummy main method is very complex and mis-constructions can lead to missing data leaks. Nonetheless, both tools seem to overcome this difficulty with their dummy main methods.

HornDroid and especially *IC3* both show high numbers of false positives compared to the number of true positives. Consequently, both tools have a poor precision. *HornDroid* is only partially flow-sensitive for register values, callbacks and heap locations. It might produce more false positives due to problems with heap abstractions, exceptions and inter-app communications leading to poorer precision.

Recall

Recall is the ratio between true positives and all reported true positives and false negatives ($(TP)/(TP + FN)$). Hence, recall tells us how many of all vulnerabilities labelled as real data leaks are actually discovered by the tool. The recall is of particular interest, since it only includes *DroidBench* vulnerabilities in the calculation (true positives and false negatives).

In terms of recall, both *FlowDroid* and *HornDroid* perform equally well on the dataset with close to 80% recall. *IccTA* performs similarly, with a recall of slightly over 77%. Again, we state that this is mostly due to the flow-, field-, context-, and object-sensitive nature of the tools allowing them to take different program paths into account and abstract different runtime values.

COVERT and *IC3* both underperform with a low recall of just 6.4% and 3.2%. Thus, these two tools miss most of the vulnerabilities present in the dataset. *COVERT* focuses mainly on inter-app communication and

escalation of privileges and *IC3* is focussing mainly on abstracting complex field values. Both tools are lacking the focus on other categories of data leak vulnerabilities and therefore show poorer performance there.

Accuracy

Accuracy is the ratio of correctly reported findings to the total number of reported findings $((TP + TN)/(TP + TN + FP + FN))$, *i.e.*, how many of the reported findings are actually correct.

We observe that both *FlowDroid* and *IccTA* show a good accuracy with more than 67%, followed by *HornDroid* and *COVERT* with an accuracy of about 55% each.

IC3 underperforms again with a precision of just 9.8% and an accuracy of 41%.

Overall Performance

Looking at the overall performance, both *COVERT* and *IC3* underperform and cannot compete with the other tools, with *IC3* having the worst performance of the selected tools. *COVERT* shows the best precision, but the poor recall rate of 6.4% and the accuracy of just over 55% somewhat spoil the effect. Both tools lack a broader focus on other categories of data leak vulnerabilities and especially lack the ability to model the applications lifecycle and handle Android callback methods.

In contrast, *FlowDroid*, *IccTA*, and *HornDroid* are detecting most of the real data leaks in the *DroidBench* dataset with a relatively good accuracy and precision. *FlowDroid* delivers the best performance, followed shortly by *IccTA* and then *HornDroid*, which suffers from the high number of false positives.

We state that these good results are mainly due to the flow-, field, object-, and context-sensitive nature of the tool and their abilities to model the application lifecycle and tackle the challenge of asynchronous Android callbacks.

6.2 Category Performance

DroidBench is presenting its synthetic applications in categories reflecting the type of data leak that they target. These categories range from inter-component communication to general Java types and even to implicit flows. Table 6.3 shows the recall for all available categories. This allows us to discuss on which type of data leaks the individual tools perform best and on which they face problems.

Category	<i>Flow-Droid</i>	<i>Horn-Droid</i>	<i>COVERT</i>	<i>IC3</i>	<i>IccTA</i>
Aliasing (0)	0	0	0	0	0
Android Specific (11)	0.909	0.727	0	0	0.909
Arrays / Lists (4)	1	1	0	0	0.667
Callbacks (21)	0.762	0.905	0	0	0.810
Emulator (23)	0.870	0.957	0	0.043	0.870
Field-/obj.-sensit. (3)	1	0.667	0	0	1
General Java (20)	0.750	0.800	0	0	0.800
Implicit Flows (9)	0	0	0	0	0
Inter-app comm. (8)	0.750	0.500	0.375	0.250	0.750
ICC (18)	0.944	0.833	0.278	0.056	0.889
Lifecycle (1)	1	1	0	0	1
Reflection (4)	0.750	1	0	0	0.750
Threading (5)	1	1	0	0	0.600

Table 6.3: Recall for each tool as observed on *DroidBench* for each category. Bold values indicate maxima for the respective metric and the value in parentheses indicates the number of data leaks present in this category. ICC stands for Inter-component communication.

We can observe that *COVERT* and *IC3* only report vulnerabilities in the field of inter-component and inter-app communication, with the latter also reporting some vulnerabilities for emulator detection. We can see that both tools do not report any leaks for the callback or lifecycle category. As discussed before, this has an impact on the overall precision of the tools. Taking that into account, it becomes apparent why the overall performance is poor compared to the other tools. Even in the two aforementioned categories, the two tools are still only detecting below 40% of the present data leaks.

Initially, *COVERT* was built only for detecting escalation of privileges with a special focus on inter-app communication. Later, it was extended with *FlowDroid* to also use taint analysis. The more surprising is the different behaviour from the standard *FlowDroid* tool. Even for inter-component and inter-app communication, *COVERT* recalls much less vulnerabilities than its competitor *FlowDroid*. Therefore, there must be internal algorithms in *COVERT* influencing and filtering the results.

We further note that none of the tools is able to detect data leaks for implicit information flows. *HornDroid* mentions this as a known limitation for its analysis. On the other hand, *FlowDroid* claims to support the detection of

implicit flows¹. However, it only comes optional with the command addition *-implicit*, which was not used in our tests since we wanted to benchmark the tools in their out-of-the-box configuration. We executed additional test runs for *FlowDroid* with this additional option but could not observe any changes in the result. *COVERT*, *IC3*, and *IccTA* are supporting implicit intents, but are not mentioning any additional support for implicit information flows in general.

FlowDroid, *HornDroid*, and *IccTA* show good recalls over all categories. All of them are reporting the highest recall in at least four categories each. Nonetheless, *HornDroid* recalls only 50% of inter-app communication vulnerabilities and *IccTA* only 66.7% for arrays and lists. Apart from these two dips, the three tools in general recall more than 75% of all present data leaks over all categories.

6.2.1 Combined Performance

The idea of our benchmarking implementation is to facilitate and automate the analysis of several tools, which enhances the reproducibility of our results and improve scalability. However, we are also interested in whether our own implementation can leverage the different base approaches and improve the recall. Hence, we are interested in knowing how many vulnerabilities are reported by at least one tool.

We observe that, when simply aggregating the results of all tools, this is the case for 113 of the 125 real *DroidBench* data leaks—which means a recall of 90.4%—topping the best individual recall by 11.2%. This is surprisingly high, considering the variety of applications and types of data leak vulnerabilities in *DroidBench*. Clearly, this comes at the cost of more false positives, since they would also be increasing in the combined approach.

Nonetheless, a future tool could consider how many of the approaches are reporting the same leak. If there are multiple, possibly fundamentally different approaches that are reporting the same leak, then we could report this leak with a higher confidence than if just a single tool would have reported it. Table 6.4 shows the probability of a true positive, if two specific tools are both reporting the same potential vulnerability. On the diagonal, we can see the individual probability—so, how likely the tool reports a true data leak. For *HornDroid*, for example, this would mean that 52.9% of all reported vulnerabilities are actually real data leaks.

We note that cases where both tools report a negative are not considered,

¹<https://blogs.uni-paderborn.de/sse/2013/10/01/flowdroid-implicit-flows/>

since our *DroidBench* test suite and the tools themselves are focussing only on positives.

	<i>Flow-Droid</i>	<i>Horn-Droid</i>	<i>COVERT</i>	<i>IC3</i>	<i>IccTA</i>
<i>FlowDroid</i>	0.643	0.731	0.700	0.231	0.631
<i>HornDroid</i>		0.529	0.700	0	0.700
<i>COVERT</i>			0.727	0	0.700
<i>IC3</i>				0.098	0.286
<i>IccTA</i>					0.618

Table 6.4: Probabilities of a real data leak when two analyses report the same potential vulnerability. The bold value indicates the maxima. Values on the diagonal are the base values of a single tool.

If we look at vulnerabilities that were reported by *HornDroid* and *IccTA*, then we can say that 70% of the reported vulnerabilities are real data leaks. In fact, we can leverage the individual probabilities by a certain combination with another tool in nearly all cases. Only for *COVERT* we can see a slight decrease in the probability when looking at combined results. We also observe that combinations with *IC3* are problematic, since its poor performance has a negative impact on the combined performance. Nonetheless, we can increase *IC3*'s performance when looking at the combined results with *FlowDroid* or *IccTA*.

We note that the best combination of tools is *FlowDroid* and *HornDroid*. Hence, when both *FlowDroid* and *HornDroid* are reporting the same vulnerability, then it is a real data leak with a probability of 73.1%. By including *IccTA* as well, we observe a real data leak with a probability of 71.9%. However, this comes at the cost of a lower recall. If we only look at the data leaks that all three tools report, then the recall is just 65.6%.

We recognise the importance of still listing all reported vulnerabilities to the user to ensure a high recall whilst classifying the findings depending on the tool combination to indicate the confidence level for each finding.

6.2.2 Pairwise Comparison

Here, we answer the question of how similar every two tools are in the detection behaviour compared to each other. To obtain such a measure expressing similarities, we apply McNemar's Test [66]. To do so, we obtain the following four numbers for each pair:

- The number of times both tools report correctly (TP+TN)
- The number of times both tools report incorrectly (FP+FN)
- The number of times one tool reports correctly and the other does not
- The number of times the other tool reports correctly and the first one does not

. Table 6.5 summarises the collected data.

		<i>HornDroid</i>		<i>COVERT</i>		<i>IC3</i>		<i>IccTA</i>	
		T	F	T	F	T	F	T	F
<i>FlowDroid</i>	T	121	68	97	92	66	123	179	10
	F	33	45	52	28	45	35	3	77
<i>HornDroid</i>	T			62	92	18	136	117	37
	F			87	26	93	20	65	48
<i>COVERT</i>	T					104	45	92	57
	F					7	113	90	30
<i>IC3</i>	T							62	49
	F							120	38

Table 6.5: Pairwise comparison of tools. For each pair, we list four numbers, denoting the various combinations of correct and incorrect classification. T denotes a true positive or negative, F denotes a false positive or negative. For example, *IC3* and *IccTA* both classify 62 times correctly, while in 120 cases, *IC3* classifies incorrectly and *IccTA* correctly.

We apply McNemar’s Test by determining if there is a statistically significant difference between two tools. This is done by calculating the χ^2 value for each pair of tools:

$$\chi^2 = \frac{(|n_{01} - n_{10}| - 1)^2}{n_{01} + n_{10}}$$

In the formula above, n_{01} and n_{10} indicate the number where one tool reports correctly and the other one does not and vice versa. Under the null hypothesis, the two tools perform equally well. We choose a significance level of $\chi^2_{1,0.01} = 6.635$, which corresponds to a confidence interval of 99%. Higher values indicate a statistically significant difference between the performance of the pair of tools. Lower values indicate that the null hypothesis holds with a probability of at least 99% and that the tools perform similarly.

Table 6.6 summarises the results for the McNemar’s Test, with bold values indicating pairs of tools with statistically significant differences in their performance.

	<i>Horn-Droid</i>	<i>COVERT</i>	<i>IC3</i>	<i>IccTA</i>
<i>FlowDroid</i>	9.94	10.56	35.29	2.77
<i>HornDroid</i>		0.2	8.53	6.01
<i>COVERT</i>			26.33	6.97
<i>IC3</i>				28.99

Table 6.6: χ^2 -values from McNemar’s test for each pair of tools. Bold values indicate a statistically significant difference.

Based on the results from the McNemar’s Test, we note that most pairs of tools show a statistically significant difference in their performance. Only for three pairs of tools we can say that the null hypothesis holds and the tools perform similarly. This is the case for the pairs *COVERT* and *HornDroid*, *HornDroid* and *IccTA* and last *IccTA* and *FlowDroid*.

For the pair, *HornDroid* and *COVERT*, we can see that the tools are not statistically significantly different and that the null hypothesis holds. This is surprising, considering the vastly different sets of reported vulnerabilities, especially considering the different number of reported findings.

Even though the McNemar’s test is suited to compare classifiers in general, it may not be best suited for this case since it only considers the disagreements, which are evidently close to each other in that case. However, the absolute number of disagreements is very high with 179 ($87 + 92$). In this case, the test fails to identify the significant differences.

Looking at the presented metrics in this chapter so far, we can postulate that *FlowDroid* is performing best based on precision, recall and accuracy, followed closely by *IccTA* and *HornDroid*. Both *FlowDroid* and *HornDroid* can be considered statistically significantly similar to *IccTA* in terms of their performance. Therefore, we would expect a similar behaviour for real-world applications.

COVERT is a precise tool, but lacks a good recall. *IC3* cannot compete with the other tools and shows a poor performance among all metrics. Accordingly, it can be considered statistically significantly different to all the tools included in this benchmark.

We showed that our benchmarking implementation can leverage the different base approaches to achieve higher recall. However, this comes at the

cost of more false positives. Nonetheless, we can present combinations of tools that, whilst agreeing on vulnerabilities, increase the confidence in the results and probability of true data leaks. Such a pair is *FlowDroid* and *HornDroid*. Therefore, we can overcome the disadvantage of presenting more false positives to the user by ranking the reports depending on the level of agreement between combinations of tools.

6.2.3 Differences to Original Configurations

While we based this benchmark mainly on running the tools with similar configurations, we should also consider examining them in their out-of-the-box configuration. In the end, this is the setup that the users unfamiliar with the internals will be using.

In this section, we are interested in the knowing if there is a difference in running the tools with shared or original configurations. Table 6.7 summarises the findings.

Metric	<i>Flow-Droid</i>	<i>Horn-Droid</i>	<i>COVERT</i>	<i>IC3</i>	<i>IccTA</i>
Reports	-0.355	-0.091	+0.364	0	-0.208
True Positives	-0.356	+0.077	0	0	-0.293
False Positives	-0.352	-0.271	+0.800	0	-0.097

Table 6.7: Calculated relative change in selected metrics from shared configurations to original configurations. Bold values indicate a positive effect on the quality of results.

We observe that there are indeed differences in running the tools with shared or with original configurations. We note that *FlowDroid* reports 35.5% less vulnerabilities with the original configurations. However, the ratio of true and false positives remains approximately the same with a slightly negative impact on precision and recall. It misses more reports, but at the same time also reports similarly less false positives. With original configurations, *FlowDroid*'s list of sources and sinks only holds 355 entries as with shared configuration, the list has 4.050 entries. This can explain the higher number of reported leaks with shared configurations, since more data flows between sources and sinks can be analysed.

HornDroid also reports slightly less vulnerabilities. Nonetheless, whilst it reports around 27.1% less false positives, the number of true positives

increases. This has a positive effect on the tool’s recall and precision. Consequently, the tool achieves better results in the original configurations.

COVERT, on the other hand, reports more false positives at a similar number of true positives and therefore makes its performance worse. This again is surprising, since it uses *FlowDroid* for its taint analysis, which behaved differently.

For *IC3*, there is no change visible in running it with shared or original configurations. This is mainly because most resources must be provided by the user in the running command. In those cases, we used the same resources as arguments in both runs.

IccTA behaves similar to *FlowDroid* with an overall decreasing effect. The number of findings decreases by 20.8%. Precision and recall for *IccTA* weaken, since the declining number of true positives is relatively bigger than the one of false positives.

In Section 5.2.2 we discussed the highly diverging number of entries in the list of sources and sinks. With shared configurations, the number of entries in the list of sources and sinks increased by factor 11 for *FlowDroid* and by factor 25 for *IccTA*. However, we do not observe a similarly strong change for the number of findings or the number of true and false positives. This leads us to the assumption that both lists of sources and sinks are already pretty complete in terms of potentially leaky sinks and source—in spite of the low number of entries.

On the other hand, the initial list was reduced by 84% for *COVERT* and for *HornDroid* by even more than 90%. *COVERT* reported almost the same number of vulnerabilities with the reduced list, but with more false positives in shared configurations. Surprisingly, *HornDroid* reported more vulnerabilities with the reduced list, but with a slight decline in true positives and more false positives.

Overall the number sources and sinks does have an impact on the number of findings and the quality of the results. However, a longer list does not necessarily lead to better results, as observed with *COVERT*. All tools using lists of sources and sinks reacted to a change in the number of entries. Nonetheless, there might be additional internal factors such as special filters, mechanisms or algorithms that boost the effect.

6.3 Large Scale Qualitative Analysis

In this section, we discuss the results for the large scale analysis on applications from the *F-Droid* store. In total, we have analysed 250 randomly

selected open source applications. Since we cannot identify true and false positives and negatives for every application, the metrics precision, recall, accuracy, and McNemar’s Test cannot be applied. Therefore, we base our evaluations on the following metrics:

- **Number of reported leaks** to compare the tools amongst each other and with the observed behaviour on the *DroidBench* dataset.
- **Number of overlaps** to validate the results from the McNemar’s test.
- **Prevalence of sink methods** to give indications on which data leaks are common with real-world applications.
- **Number of time outs** to draw conclusions on the performance within the given time frame.

Each metric will be discussed in a separate section. We must note here that *HornDroid* had to be excluded from the *F-Droid* analysis. For a sample set of 50 applications, *HornDroid* produced a timeout (set to one hour) for every application. Even with a higher timeout of three hours, *HornDroid* did not finish the analysis a single time. For this reason, we have excluded *HornDroid* from the *F-Droid* analyses, since it does not produce reports within a reasonable amount of time.

6.3.1 Number of Reported Leaks

In the *DroidBench* analysis, we have already observed big differences in the number of reported vulnerabilities between the tools. Especially, *COVERT* and *IC3* reported only very few vulnerabilities, whereas *FlowDroid* and *IccTA* produced far more reports. Here, we are interested in whether this trend can also be observed in the *F-Droid* analysis. Table 6.8 shows the results for the number of reported leaks.

Metric	<i>Flow-Droid</i>	<i>COVERT</i>	<i>IC3</i>	<i>IccTA</i>
Leaks	1735	1	3353	0
Average leaks	9.971	0	14.088	0.004

Table 6.8: Number reported leaks for each tool and average number of leaks per application.

In total, the tools produced 5052 distinct reports. We observe that only *FlowDroid* and *IC3* are frequently reporting vulnerabilities with an average of close to 10 to 14 reports per application. We note that *IC3* reports close to twice as many vulnerabilities as *FlowDroid*. However, we have little confidence in the reports from *IC3* considering its poor performance on the *DroidBench* test suite. Only close to 10% of all reported leaks were in fact true positives there. We predict a similar behaviour on the *F-Droid* test suite and suspect that most of *IC3*'s reports are false positives.

A possible root cause for the high number of false positives of *IC3* can be the value-insensitive nature of the analysis. Value-sensitivity allows the analysis to approximate runtime values and skip unreachable program branches [86]. Because *IC3* is context-sensitive, the tool can compute different static approximations upon different method calls, resulting in more program branches to be checked during the analysis. However, since the tool is value-insensitive, it cannot skip unreachable program branches resulting in the detection of non-existent data flows and therefore more false positives. Nonetheless, verifying this is beyond the scope of this work and remains as future work.

FlowDroid, on the other hand, showed a good performance on the *DroidBench* test suite. We therefore suspect that the tool detected most of the existing vulnerabilities in the real-world applications. Especially, since the tool is both context- and value-sensitive, it does not face similar issues to *IC3* and reports less false positives.

COVERT only reports one vulnerability confirming the behaviour observed on *DroidBench*. The tool focusses on the detection of inter-app communication vulnerabilities and especially on the detection of unsafe combinations of applications that could lead to security exploits. Since we analysed the applications one by one, this could have an impact on internal mechanisms of the tool resulting in less reports.

IccTA cannot continue its good performance from the *DroidBench* test suite and does not report any vulnerabilities. The major issue here is, that the tool only completed the analysis for 13 out of 250 applications. For the remaining 237 applications, the analysis was aborted due to internal exceptions or due to time outs.

Given the poor results of *IccTA* we cannot confirm the statistical significant similarities observed in *DroidBench* between *FlowDroid* and *IccTA*.

6.3.2 Number of Overlaps

In the *DroidBench* dataset we observed a lot of overlaps between the tools, especially between *FlowDroid* and *IccTA*.

Table 6.9 summarises the results for the *F-Droid* dataset.

	<i>FlowDroid</i>	<i>COVERT</i>	<i>IC3</i>	<i>IccTA</i>
<i>FlowDroid</i>	1.735	0	37	0
<i>COVERT</i>		1	0	0
<i>IC3</i>			3.353	0
<i>IccTA</i>				0

Table 6.9: Number of overlaps for all pair of tools.

We only observe overlaps between *FlowDroid* and *IC3*, which can be expected given the low number of reports for the other tools. We note that over all 5.052 reported vulnerabilities, there are only 37 overlaps. So, in only 0.7% of the cases, there is an agreement between the tools. This is very surprising, given the different situation in the *DroidBench* analysis. There, about 59% of all reported leaks showed overlaps between the tools. For the *DroidBench* test suite we have calculated the probability of a real data leak when two particular tools report the same vulnerability. For the pair of *FlowDroid* and *IC3* this probability was 23.1%. For the 37 overlaps, this would mean that less than 9 reports are in fact true positives.

We can state that the tools' agreement for real-world applications diverge and the reports are much more heterogeneous than with the synthetic *DroidBench* applications.

One reason might be that tool developers used *DroidBench* to evaluate their tools and therefore optimize them to achieve high precision and recall on this dataset. Whereas for real-world applications, there are other factors impacting the tools' performance such as obfuscation or the size of the application. Thus, a good performance on the *DroidBench* dataset does not necessarily yield in good results for real-world applications. This confirms the necessity of evaluations not only on synthetic applications, but also on real-world applications.

6.3.3 Prevalence of Sink Methods

In this section, we are interested in knowing the prevalence of particular data leaks in real-world applications. The analysis on the *F-Droid* dataset allows us to observe increasing numbers of occurrences of certain sink methods. This could give us an indication what developers should try to avoid whilst developing their application. Out of the 5.043 reports, we can see only 265 distinct sink methods. Table 6.10 shows the five most common sinks.

We note that the five most reported vulnerabilities together make up

Sink Method	<i>Total</i>	<i>Percentage</i>
<code>startActivity(Intent)</code>	968	0.192
<code>getStringExtra(String)</code>	337	0.067
<code>startActivityForResult(Intent,int)</code>	274	0.054
<code>getString(String)</code>	233	0.046
<code>startService(Intent)</code>	203	0.040

Table 6.10: Top 5 of most reported sink methods and absolute as well as relative number of occurrences.

39.9% of all reported leaks. The vulnerability that was reported most is the call to the `startActivity(android.content.Intent)` method. The `Intent` passed as argument describes an action to be performed and carries necessary data that might be required. An activity presents a single screen within the application. So, the `startActivity(android.content.Intent)` method will start a new instance of an `Activity` and perform the defined action with the passed data. If this data contains sensitive information, it will leave the current component. In that case, a data leak occurs. This sink method can be categorized as an inter-component communication (ICC) data leak.

Also the methods `startActivityForResult(android.content.Intent,int)` and `startService(android.content.Intent)` pass an `Intent` to another component and belong to the same category. In fact, such ICC API calls starting another service or activity make up around 28.9% of all reported vulnerabilities. We state that they are the most prevalent type of data leak vulnerabilities in our set of real-world applications.

The security concern with ICC is that we do not have control on what happens to the sensitive information passed to the other component. The `Intent` message could be intercepted during message passing or the receiving component could for example send out the sensitive information from the `Intent` to a phone number as an SMS [64]. However, these kinds of sink methods can only be considered data leaks, when the `Intents` are actually containing sensitive information. Verifying this is beyond the scope of this work and remains as future work.

The remaining two sink methods in our top five are also related to `Intent` handling. The `getStringExtra(java.lang.String)` method and `getString(java.lang.String)` method can be used to extract information from the `Intent` such as the `String extra`, which could contain for example the device UID. However, given our definition of sink methods and data leaks from Section 2.2.1, we cannot consider reading such data from `Intents` as a sink

method leaking sensitive information. Consequently, we suspect all these cases to be false positives. In fact, these method calls were only reported by *IC3*, where we already commented on the poor confidence level in the results.

We also note the high prevalence of log activities among all reports. In summary, there are 341 log activity sink methods; making these kinds of sink methods the actual second most prevalent of all reports with a total share of 6.8%. An example of such an API call is the `Log.d(java.lang.String, java.lang.String)` method. Here, a `DEBUG` message is sent containing a piece of potentially sensitive information as a `String` text to the log files. Log files are shared resources in Android. Therefore, they could be accessed by other malicious applications to obtain the sensitive information². Prior to Android 4.0, any application with the `READ_LOGS` could access the logs of any other application using the command-line tool `LogCat`³. Since Android 4.1, an application can only access its own logs. However, using a PC, one can still obtain log output from other applications⁴.

To mitigate this issue, Google advises programmers to limit log usage and to use debug flags or custom `Log` classes⁵.

6.3.4 Time Outs

From the 250 analysed applications, there are cases where one or several of the tools was aborted due to time outs or exceptions.

Table 6.11 summarises these cases.

Metric	<i>Flow-Droid</i>	<i>COVERT</i>	<i>IC3</i>	<i>IccTA</i>
Time outs	76	0	9	118
Exceptions	0	0	3	119
Completed	174	250	238	13

Table 6.11: Number of time outs, exceptions and completed analysis for each tool.

We observe that *COVERT* and *IC3* nearly always finished the analysis without reaching the time out or throwing an exception. Therefore, the tools

²<https://blogs.uni-paderborn.de/sse/2013/05/17/privacy-threatened-by-logging/>

³<https://developer.android.com/studio/command-line/logcat>

⁴<https://wiki.sei.cmu.edu/confluence/display/android/DRD04-J.+Do+not+log+sensitive+information>

⁵<https://developer.android.com/training/articles/security-tips>

can be considered time efficient. For *FlowDroid* we can see that in about 30% of the cases a time out was reached. *IccTA* produced most time outs and is therefore time intensive compared to the other tools. In fact, the analysis completed only in 5.3% of all cases for *IccTA*. In 47.6% of all cases, an exception was thrown such as a `RuntimeException`. These exceptions must have internal root causes in the *IccTA* tool. We suspect that larger applications or code obfuscation cause such issues with *IccTA*. However, verifying this is beyond the scope of this work and remains as future work.

6.4 Hypotheses Evaluation

In this section, we evaluate the hypotheses from Section 3.1.

H1 *The more tools report the same vulnerability, the more likely is a true positive.*

In Section 6.2.1 we showed that when a particular combination of tools report the same vulnerability, then the probability of it being a true data leak increase. The best combination is the pair *FlowDroid* and *HornDroid*.

H2 *The less tools report a certain vulnerability, the more likely is a false positive.*

This can also be observed in Section 6.2.1. However, it depends on the underlying tool reporting the vulnerability. For tools with a poor performance such as *IC3* it is certainly true. There are also cases where a tool covers edge cases ignored by other tools. Therefore, we cannot say that this hypothesis is true in general.

H3 *If a tool reports much more data leaks than all the other tools, then it is likely to report more false positives than the others.*

This is the case for *IC3* on the *F-Droid* data set. However, verifying the reports from being true or false positives is beyond the scope of this work and remains as future work.

7

Threats to Validity

In this chapter, we point out threats that might impact the validity of the results. We validate our benchmarking implementation and assess impacts of other threats.

7.1 Benchmarking Implementation Validation

Our implementation may contain bugs that directly influence the results. Possible areas of bugs that could impact the results are during the analysis triggering, parsing of the results, or writing to the summary files. Especially parsing the results files from the different tools is not trivial, since the results usually do not have a standardised format and differ in the amount of information they show. To mitigate this threat, we implemented unit tests during the development to ensure that produced results are correctly parsed. Furthermore, we verified the generated output manually by comparing the results from the individual tools with our summarised reports and with the manually gathered data from *DroidBench*.

Our validation shows that about 90% of all *DroidBench* reports were correctly parsed for class name, method name, and sink method. In the remaining reports, one of three issues impacts the validity of the reports:

1. ***FlowDroid* reports are missing information.** In some cases, *FlowDroid* shows the class and method name of the source instead of the ones where the sink occurs. In fact, this is an issue of *FlowDroid* itself.

Hence, we are only able to identify it by manually checking within the application source code whether the indicated class and method contain the reported sink method. However, our parser cannot automatically detect if the reported class and method name belong to the sink or the source. Since we manually analysed the *DroidBench* reports, we were able to identify these cases and correct them.

Consequently, this issue can only impact the *F-Droid* results in that some of the *FlowDroid* reports show the wrong class or method name, but the correct sink method. In the *DroidBench* dataset, we observed this *FlowDroid* issue in 13 out of 141 reports; so, in about 9.2% of all *FlowDroid* reports.

2. ***IccTA* reports are parsed for the wrong class and method name.** *IccTA* vulnerability descriptions are extensive. Depending on the vulnerability, they can contain additional information such as special register values or path descriptions. Class and method name indicating where the sink method occurs can therefore change the location within the description. We have identified seven different possible locations and created the corresponding unit tests for these cases. We have invested a considerable amount of work to optimise the *IccTA* parser and are now able to identify the correct class and method in six out of these seven test cases.

In addition, we manually verified our parser on the *DroidBench* dataset. It correctly parsed 64% of all *IccTA* reports. In 52 cases, the parser was returning the correct sink, but showed an issue with the class or method name. In 42 cases, the class name was correct, but the method name was incorrect. In 8 cases, both—i.e., class name and method name—were incorrect. We note that the reported sink method was in all cases correctly parsed.

To mitigate this threat, a future project could modify the original *IccTA* tool to return the class name, method name and sink method using API calls. With this approach, parsing the results file would not be required anymore.

3. ***COVERT* is not reporting the sink method.** This tool only reports the class and method name where the sink occurs, missing the sink method name in the report. However, *COVERT* indicates the sink type such as `SEND_SMS`. Based on these sink types, we could manually identify if a sink method matching the sink type was present in the reported class and method. Nonetheless, our benchmarking tool

cannot do this automatically. Consequently, the *F-Droid* analysis never shows an overlap of a *COVERT* report with another tool’s report, since it misses the sink method that would be required to identify such an overlap. For the *DroidBench* analysis however, we were able to overcome this issue by manually checking whether the sink type can be referenced with a sink method in the indicated class and method.

7.2 DroidBench Validation

We assume that the synthetic *DroidBench* applications do not contain other vulnerabilities than the ones that are indicated by the authors of this test suite. By manually analysing potential false positives reported by the tools, we could mitigate this threat. In fact, none of the reported false positives turned out to be a true positive. Hence, there was no vulnerability missing in the test suite.

However, it is still possible that our manual check misclassified such true or false positives due to lack of expertise. Due to the small size of the programs in *DroidBench*, we are, however, confident that these vulnerabilities would have been detected by the community, the original authors, or us.

DroidBench and *FlowDroid* originate from the same research group. Hence, a selection bias that favours *FlowDroid* is possible. Nonetheless, other tools such as *HornDroid* or *IccTA* also used *DroidBench* for their evaluation. To mitigate this threat, other sets of applications with known vulnerabilities could be used in the future such as the one provided by Mitra *et al.* [68]. Creating and maintaining such a set is out of scope for this paper and remains as future work.

7.3 Misconfiguration

We have used the tools included in this benchmark as distributed. Furthermore, we have followed the provided instructions to set up and run the tools. Nonetheless, it is possible that we have misconfigured some of them. We mitigate this threat by only making minimal changes to shared configurations (as described in Chapter 5).

Our choice to normalise configurations and run the tools with shared configurations may impact the results. We mitigate this threat by also running the tools in the original configurations and by pointing out differences in the results (Section 6.2.3). Nevertheless, we argue that the threat is minimal and that using the same configuration for the tools is a sensible choice.

8

Conclusions and Future Work

8.1 Summary

In this work, we investigate to what degree tools assessing data leak vulnerabilities in Android applications are available, and how they perform in practice compared to each other. We report the process of elimination during the tool selection process, and how, from an initial list of 87 vulnerability detection tools, we arrive at 5 available tools focusing on data leak vulnerability detection.

We present our Java implementation of an automated benchmark to automatically execute the tools and to consolidate the results. To benchmark the tools, we evaluate them in a small scale qualitative study on the *DroidBench* dataset with synthetical applications with known data leak vulnerabilities. In addition, we perform a large scale quantitative study on the *F-Droid* dataset of real-world applications.

We then present and discuss the results, where we report on several aspects of the tools. Within the context of the small scale qualitative analysis, we discuss values such as precision, recall and accuracy for each tool. We further perform a pairwise comparison of the tools using the McNemar's Test [66]. Moreover, we evaluate differences stemming from running the tools with shared or with original configurations.

For the large scale quantitative analysis, we use real-world applications. Using real applications makes it considerably harder to determine precision

and recall of the tools; we would need to manually determine whether a report is a false positive, which is a challenging task, even for experts in the domain. Hence, we discuss values such as number of reported vulnerabilities or the number of overlaps. In addition, we evaluate the prevalence of certain sink methods in real-world applications.

We show that our benchmarking implementation can leverage different base approaches to achieve higher recall rates. It improves the state of the art in that it provides an analysis of five tools in the same domain. To ensure a fair comparison, we configured the tools to use the same configurations for the analysis and apply them to the same targets.

8.2 Conclusions

In this work, we arrive at the following main conclusions:

- **Tool availability is poor.**

Even though there exist hundreds of tools in the domain of Android Security, we observe that only little of them are available for a specific sub-domain such as data leak vulnerability detection. We further note, that the obtainable tools are often out of date and face issues when being used with newer Android versions.

- **There are huge differences in how the tools perform in practice.**

We observe that, even though tools show similar structures on paper, they perform quite differently when being compared in practice. Consequently, we see considerable differences in key metrics such as precision, recall or accuracy.

- **Different lists of sources and sinks have an impact on recall and precision.**

The tools' original configurations and especially the number of entries in the list of sources and sinks is tremendously different. However, a long list is no guarantee for detecting more vulnerabilities as observed with *FlowDroid*, *IccTA*, and *HornDroid*. Changing the list of entries impacts both precision and recall for each tool differently.

- **Tool evaluation on real-world applications is important.**

A good performance on the *DroidBench* test suite with synthetic applications is no guarantee that the tools perform equally well for real-world applications. *HornDroid* and *IccTA* show a good performance

on the *DroidBench* test suite, but face major issues with real-world applications such as time outs or exceptions. We note the importance in evaluating tools not only on synthetical applications with known vulnerabilities, but also on real-world applications. Hence, large applications or obfuscation as seen in such applications can impact the tools performance considerably and should be considered during the tool's evaluation.

- ***FlowDroid* has the best performance.**

In comparison to the other tools *FlowDroid* shows the best overall performance. It detects most vulnerabilities on the *DroidBench* test suite with a good precision. It is the only tool reporting credible results for real-world applications and can in most cases complete its analysis within the time limit of one hour.

- **Leveraging different base approaches can lead to higher recall.**

We showed that our own benchmarking implementation can leverage the different base approaches to achieve higher recall. However, it comes at the cost of more false positives. We propose to rank reports depending on the level of agreement between the tools to overcome this threat.

- **The most prevalent sinks in real-world applications are log activities.**

The analysis on real-world applications shows that log activities such as `Log.i` make up a considerable amount of all reported vulnerabilities. A future work could further investigate the usage of log activities in applications and to what extent they pose a security threat.

8.3 Future Work

Future work could include adding additional tools to the benchmark, either in the same domain or in other domains such as escalation of privileges. Nowadays, a developer wanting to evaluate the overall security of an application is required to use several tools from different security domains. By including other domains in the benchmark, we could overcome this issue and deliver an overall security evaluation for an application in just a few steps. Furthermore, the benchmarking implementation could be extended to show the confidence level in each finding, depending on which tools reported it and if there are overlaps between the tools.

Case studies on larger applications such as *Facebook* or *Candy Crush* could be considered to further evaluate the tools' performance on larger applications.

Besides, future efforts should validate *HornDroid* for real-world applications with higher time outs. This would allow us to evaluate whether *HornDroid* does in general not work for such applications or if it is just very time intensive.

Finally, another direction we could explore is to include the benchmark suite in development tools. Here, the idea is to run analyses on demand, for example, on the program that is currently open in an integrated development environment (IDE). We then can provide the reports directly to the programmer, who can benefit from data from multiple tools, while having to deal with a single, easy to understand interface and report format.

9

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Thanks to some extra efforts on all sides, we were able to submit an additional paper on the topic of "Benchmarking Android Data Leak Detection Tools" for the International Symposium on Engineering Secure Software and Systems (ESSoS 18)¹, which was accepted as an idea paper [27].

I would also like to thank the experts who were involved in the validation either of the aforementioned research paper or this thesis here. Special thanks go to Dr. Mohammed Ghafari for his valuable inputs during the project.

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¹<https://distrinet.cs.kuleuven.be/events/essos/2018/index.html>



Anleitung zu wissenschaftlichen Arbeiten

The “Anleitung zu wissenschaftlichen Arbeiten” consists of:

- The ESSoS ’18 paper “Idea: Benchmarking Android Data Leak Detection Tools” (C. Corrodi, T. Spring, M. Ghafari, O. Nierstrasz) [27], and
- the instructions to setting up the benchmark in the following pages.

A.1 Tool Setup

Our benchmarking implementation is available on GitHub¹. This manual will lead you through the project and tool setup and explains how to use our benchmarking implementation to reproduce the results.

The benchmark can be setup in four steps:

1. Setup static analysis tools,
2. run `gradle shadowJar` to generate build,
3. Adapt the properties file and change the paths to the tools,
4. Run the tool with `java -jar benchmarking.jar`

We start with the setup of the tools included in the benchmark—*COVERT*, *FlowDroid*, *HornDroid*, *IC3*, and *IccTA*.

The GitHub page contains the following folders that we used for running the analysis. One can either use our proposed folder structure or adapt the properties file with the modified paths to the required artefacts.

We propose the following folderds:

- **src**

Contains the source code to our own benchmarking implementation to run all tools and collect and summarise the results.

- **tools**

For each tool included in the benchmark, there exists a separate folder. The tools' artefacts and resources will need to be placed there. During the analysis, the tools will create results files in those folders, that will be used by our benchmarking implementation.

The tools folder also contains a *commonConfig* folder containing all shared resources used for the analysis. To run the tools with shared configurations, you have to provide the following resources in the *commonConfig* folder:

- ***android.jar*** We used API level 23 for this work. There are, however, good collections for different API levels available online².

¹<https://github.com/tiimoS/droid-Security-Thesis/tree/essos>

²<https://github.com/Sable/android-platforms>

- ***SourcesAndSinks.txt*** Text file containing the list of sources and sinks to be checked for data flows during the analysis. We recommend to use the *SuSi*³ tool to obtain such a list based on the selected *android.jar*.
- ***AndroidCallbacks.txt*** Text file containing the list of callback methods. We recommend to use an existing list for example the one provided by *FlowDroid*.
- ***apktool.jar*** We recommend to use the latest version of the *Apktool*⁴.

- **results**

The summarised and grouped findings from our benchmarking implementation will be stored here; including a summary text file report, a summary csv file report and a csv file showing which tools timed out and which ones completed the analysis.

- **apksToTest**

All applications that should be analysed need to be put in this folder.

A.1.1 *COVERT* Setup

Clone the benchmarking project from GitHub and navigate to the *tools/covert* folder inside the project. Then follow the steps below to set up *COVERT*.

1. Obtain *COVERT*

Note, that the folder already contains certain files and folders. These should not be changed for the benchmark to work. In order for *COVERT* to run, you need to obtain the following artefacts and store them in the *covert* folder. Make sure to rename the files as indicated or adapt the running command later on.

First, you have to obtain the *COVERT* back-end⁵. Then unpack the folder and extract the following files and folders to the *covert* folder in the benchmarking project directory.

- ***covert.bat***
- ***covert.sh***

³<https://blogs.uni-paderborn.de/sse/tools/susi/>

⁴<https://ibotpeaches.github.io/Apktool/>

⁵<https://www.ics.uci.edu/~seal/projects/covert/>

- ***appRepo*** (folder)
- ***resources*** (folder)

Copy and paste the following resources into the *configCustom* folder inside the *covert* directory. We denote the root directory of the project as \sim :

- `~/resources/AndroidPlatforms/android-8/android.jar`
- `~/resources/Covert/resources/apktool/apktool.jar`
- `~/resources/FlowDroid/resources/AndroidCallbacks.txt`
- `~/resources/FlowDroid/resources/SourcesAndSinks.txt`

These are the configurations that we will change to run the tool with shared configurations. However, we copy the artefacts to the *custom-Config* folder to allow the user to easily switch between original and shared configurations.

2. **Running Command** Go to the *covert* directory and make sure that all required files and artefacts are present. Copy a sample application into the *app_repo* folder. Then run the following command to start the analysis.

```
./covert.sh <APPLICATION_NAME>.apk
```

Instead of a single application, you could also analyse multiple applications with *COVERT*. To do so, you have to create an additional folder inside the *app_repo* folder. Instead of the application name, you then pass the name of the newly created folder in the running command.

Make sure to run the tool on an sample application and verify that the results are as expected. We recommend to use the SendSMS.apk provided by *DroidBench*⁶ for a test run. *COVERT* should detect a data leak for the `sendTextMessage()` method. The source code for the application is available on *DroidBench*.

3. **Results** The results of the analysis are located in the *app_repo* folder. A new folder with the same name as the analysed application should have been created during the analysis. This folder contains the results file in a *.xml* file.

⁶<https://github.com/secure-software-engineering/DroidBench>

***FlowDroid* Setup**

Clone the benchmarking project from GitHub and navigate to the *tools/flowDroid* folder. Then follow the steps below to set up the *FlowDroid* tool.

1. Obtain *FlowDroid*

Note, that the folder already contains certain files and folders. These should not be changed for the benchmark to work. In order for *FlowDroid* to run, you to obtain the following artefacts and store them in the *flowDroid* folder. Make sure to rename the files as indicated or adapt the running command later on.

The following additional libraries are required:

- ***soot-trunk.jar***
- ***soot-infoflow.jar***
- ***soot-infoflow-android.jar***
- ***slf4j-simple-1.7.5.jar*** (libraries for logging)
- ***slf4j-simple-1.7.5.jar*** (libraries for logging)
- ***axml-2.0.jar*** (Android XML parser library)
- ***android.jar*** (Android SDK): For the analysis we use API Level 23

Furthermore, you need to obtain the following configuration files and store them in the same folder as the artefacts above: Make sure to also copy and paste the *SourcesAndSinks.txt* and *AndroidCallbacks.txt* file in the *configCustom* folder.

- ***EasyTaintWrapperSource.txt*** (taint wrapper)
- ***AndroidCallbacks.txt*** (Android callbacks)
- ***SourcesAndSinks.txt*** (sources and sinks)

Make sure to also copy and paste the *SourcesAndSinks.txt* and *AndroidCallbacks.txt* file in the *configCustom* folder. This is required to later create soft links that allow switching between running the tools with original or shared configurations.

2. **Running Command** Go to the *flowDroid* directory and make sure that all required files and artefacts are present. Then run the following command to start the analysis.

```
java -Xmx4g -cp soot-trunk.jar:soot-infoflow.jar:soot-infoflow-android.jar:slf4j-api-1.7.5.jar:slf4j-simple-1.7.5.jar:axml-2.0.jar soot.jimple.infoflow.android.TestApps.Test <PATH_TO_APPLICATION> ./android.jar > flowdroid_results.txt
```

For the benchmark, we put applications that should be analysed in the *apksToTest* folder. The corresponding folder path would be `../../apksToTest/<APPLICATION_NAME>.apk`.

Make sure to run the tool on a sample application and verify that the results are as expected. We recommend to use the *SharedPreferences1.apk* provided by *DroidBench* for a test run. *FlowDroid* should detect all data leaks as indicated in the applications source code that is available on *DroidBench*.

3. **Results** The results of the analysis are stored in the *flowdroid_results.txt* file. They also get printed to the console.

***HornDroid* Setup**

For *HornDroid* the setup is straightforward and well documented on its project page.

1. **Obtain *HornDroid***

First, clone the *HornDroid* GitHub project⁷ and build it using `mvn clean package`. A new folder *target* is created. Move the content from this folder to the *horndroid* folder inside the benchmarking directory.

In addition, copy and paste the following files to the *horndroid/config-Custom* folder:

- ***apktool.jar***
- ***Callbacks.txt***
- ***SourcesAndSinks.txt***

2. **Running Command** Go to the *horndroid* directory and make sure that all required files and artefacts are present. Then run the following command to start the analysis.

⁷<https://github.com/ylya/horndroid.git>

```
java -jar fshorndroid-0.0.1.jar / ./apktool.jar <PATH_TO_APPLICATION>
```

Make sure to run the tool on an sample application and verify that the results are as expected. We recommend to use the `SharedPreferences1.apk` provided by *DroidBench* for a test run. *HornDroid* should detect all data leaks as indicated in the applications source code that is available on *DroidBench*.

3. **Results** The results of the analysis are located in the *OUTPUT.report* folder as a *JSON* file. They are also printed to the console.

IC3 Setup

The setup of *IC3* also requires the installation of the *Dare* tool. It is used to decompile Android applications from the installation image to source code on which *IC3* can perform.

1. Obtain IC3

We start by setting up the *Dare* tool. First, go to the Dare installation page⁸ and download the latest version. After the download, unzip the installation package and create a folder named *output*. Move the whole content of the installation package to the *tools_helper/dare* folder.

Now, you can start with the setup of the *IC3* tool. Begin by cloning the project from GitHub⁹. Then go the *IC3* directory and build the tool with the following command:

```
git clone https://github.com/siis/ic3
cd ic3
mvn clean package -P standalone
```

Move the content from the newly created *target* folder to the *ic3* folder inside the benchmarking directory. There are already two folders present—*dareOutput* and *ic3ouput*. These folders are used to store the results of *Dare* and *IC3* respectively.

2. **Running Command** Before starting the analysis, we have to run the *Dare* tool on the application to decompile it. To do so, run the following command from within the *Dare* directory:

⁸<http://siis.cse.psu.edu/dare/installation.html>

⁹<https://github.com/siis/ic3>


```
./dare -d <PATH_TO_IC3_DAREOUTPUT> <PATH_TO_APPLICATION>
```

Afterwards, the folder `~/ic3/dareOutput` should contain the decompiled files. Now we are all set to run the *IC3* tool with the following command from within the *ic3* folder:

```
./runIC3.sh <PATH_TO_APPLICATION>
```

Make sure to run the tool on an sample application and verify that the results are as expected. We recommend to use the `StartActivityForResult1.apk` provided by *DroidBench* for a test run. *IC3* should detect one data leaks at the `startActivityForResult()` method.

3. **Results** The results of the analysis are located in the *ic3output* folder as a text file.

IccTA Setup

The setup of *IccTA* is the most time intensive one. It requires to build the tool yourself and set up a *mySQL* database to store the intermediate results. Furthermore, it uses several frameworks that you can build yourself. However, there are already built versions available online that we use for this benchmark.

1. **Obtain *IccTA*** Begin by importing all the following projects to Eclipse or another IDE:
 - *jasmin*
 - *heros*
 - *soot*
 - *soot-infoflow*
 - *soot-infoflow-android*
 - *soot-infoflow-android-iccta*

Then change the build path of *soot-infoflow-android-iccta* to include the above projects and then build each project.

Afterwards, create a new folder called *output_iccta* inside the *soot-infoflow-android-iccta* directory.

2. **Database Setup** To store intermediate results during the analysis, *IccTA* uses a *mySQL* database. Start *mySQL* and create the following database:

```
mysql -u root -p -e 'create database cc?;'
mysql -u root -p cc < res/schema;
```

As a database name, you have to use *cc* since it is hardcoded in the *IC3* tool provided. When you are done with importing the *IccTA* *schema*, you have to adapt the database properties of the tool. The following files need to be updated with the correct username and password to the database:

- `~/iccProvider/ic3/runIC3.sh`
- `~/iccProvider/ic3/runIC3Wrapper.sh`
- `~/res/jdbc.xml`
- `~/res/iccta.properties`
- `~/src/soot/jimple/infoflow/android/iccta/util/Constants`

3. **Running Command** Before starting the analysis, we have to run the *IC3* tool on the application. To do so, run the following command from within the `~/iccProvider/ic3` directory:

```
./runIC3.sh <PATH_TO_APPLICATION>
```

Now, go to the *release* directory and run the following command to start *IccTA*:

```
java -jar IccTA.jar <PATH_TO_APPLICATION> <PATH_TO_ANDROID_SDK> \
newline -iccProvider ../iccProvider/ic3
```

Make sure to run the tool on a sample application and verify that the results are as expected. We recommend to use the `SharedPreferences1.apk` provided by *DroidBench* for a test run. *IccTA* should detect all data leaks as indicated in the applications source code that is available on *DroidBench*.

4. **Results** The results of the analysis are located in the `output_iccta` folder as a text file.

A.2 Benchmarking Setup

In order for our benchmarking implementation to work, you have to make sure that all tools are working as explained in Appendix A.1.

A.2.1 Build benchmark

After setting up the tools, you have to build the benchmark as a `jar` file containing all required dependencies. This can be done using our *gradle script*. Note that these are only dependencies such as to *apache commons* or *junit* and not to the tools themselves.

Run the following command to build the benchmark:

```
gradle shadowJar
```

This generates an executable jar file containing all relevant dependencies.

A.2.2 Adapt Configurations

In case you have setup the tools in the same folder structure as presented in this manual, then you can skip this section. In case you have changed the location or name of certain artefacts, then you can easily adapt the paths used to these benchmarks in our `data_leak_detection.properties.defaults` file.

Furthermore, you can exclude certain tools from the benchmark by simply changing the `enabled` flag in the properties file to `false`.

A.2.3 Run Benchmark

Put the APK file of the application you would like to analyse in the *apksToTest* folder. Then run the following command to start the benchmark with the enabled tools:

```
java -jar benchmarking.jar
```

The results of the analysis can be found in the *results* folder. It contains the individual tools' results as well as the summarised reports. The *timedOut* folder contains a list of all applications that were analysed with an indication for each tool if the analysis was completed (indicated as 0), interrupted with and exception (indicated as 1), or timed out (indicated as 2).

A.3 Concluding remarks

By following this manual, the benchmark can be set up and presented results reproduced. The majority of work stems from setting up the tools; we do not ship them with our benchmark. Once the tools are up and running, running the benchmark is trivial.

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