

Unit 2

Word and Morphology

Finite State Machine, Morphology, Word Construction

Natural Language Processing (NLP)

MDS 555



Objective

- Regular Expressions
- Formal Language Processing
 - Finite state machine
 - Finite state transducers
- Further Study
 - Chapter 2 , 3 of Text book



Regular Expressions (RE)

- First developed by Kleene (1956)
- language for specifying text search strings
- The Regular expression languages used for searching texts in UNIX (vi, **Perl**, **Emacs**, **grep**)
- RE features exist in the various Web search engines.



Regular Expressions (RE)

- A **string** is a sequence of symbols
 - for the purpose of most text based search techniques, a string is any sequence of alphanumeric characters (letters, numbers, spaces, tabs, and punctuation).
 - Regular expression search requires a **pattern** that we want to search for, and a **corpus**



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RE: Basic Patterns

- Regular expressions are case sensitive

RE	Example Patterns Matched
/woodchucks/	"interesting links to <u>woodchucks</u> and lemurs"
/a/	" <u>M</u> ary Ann stopped by Mona's"
/Claire_says,/	" "Dagmar, my gift please," <u>Claire</u> says,"
/DOROTHY/	"SURRENDER <u>DOROTHY</u> "
/!/	"You've left the burglar behind again!" said Nori

- Disjunction

RE	Match	Example Patterns
/[wW]oodchuck/	Woodchuck or woodchuck	" <u>Woodchuck</u> "
/[abc]/	'a', 'b', or 'c'	"In uomini, in soldat <u>i</u> "
/[1234567890]/	any digit	"plenty of <u>7</u> to 5"



RE: Basic Patterns

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RE	Match	Example Patterns Matched
/ [A-Z] /	an uppercase letter	“we should call it ‘ <u>D</u> renched Blossoms”
/ [a-z] /	a lowercase letter	“ <u>m</u> y beans were impatient to be hoed!”
/ [0-9] /	a single digit	“Chapter <u>1</u> : Down the Rabbit Hole”

- caret ^ for negation

RE	Match (single characters)	Example Patterns Matched
[^A-Z]	not an uppercase letter	“Oyfn pripetchik”
[^Ss]	neither ‘S’ nor ‘s’	“ <u>I</u> have no exquisite reason for’t”
[^\.]	not a period	“ <u>o</u> ur resident Djinn”
[e^]	either ‘e’ or ‘^’	“look up <u>^</u> now”
a^b	the pattern ‘a^b’	“look up <u>a^b</u> now”



the period (/./), a wildcard expression

- One very important special character is the period (/./), a wildcard expression that matches any single character (except a carriage return):

RE	Match	Example Patterns
/beg.n/	any character between <i>beg</i> and <i>n</i>	<u>begin</u> , <u>beg'n</u> , <u>begun</u>



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Finite State Machine

- Finite State Automata
- It is a computation model that can be implemented with hardware or software and can be used to **simulate sequential logic** and some computer programs.
- It has fixed set of possible states, **a set of inputs that change the state** and **set of possible outputs**.



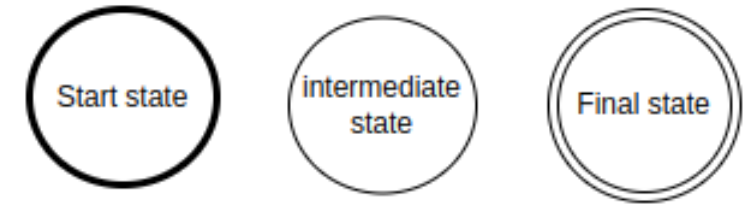
Finite State Machine

- Finite state automata generate regular languages.
- Finite state machines can be used to model problems in many fields including
 - mathematics, artificial intelligence, games, and linguistics.



State transition diagram

- States
 - Start State – Circle with bold border
 - State – intermediate states
 - Final State – double border circle
- Transition is shows by arrow

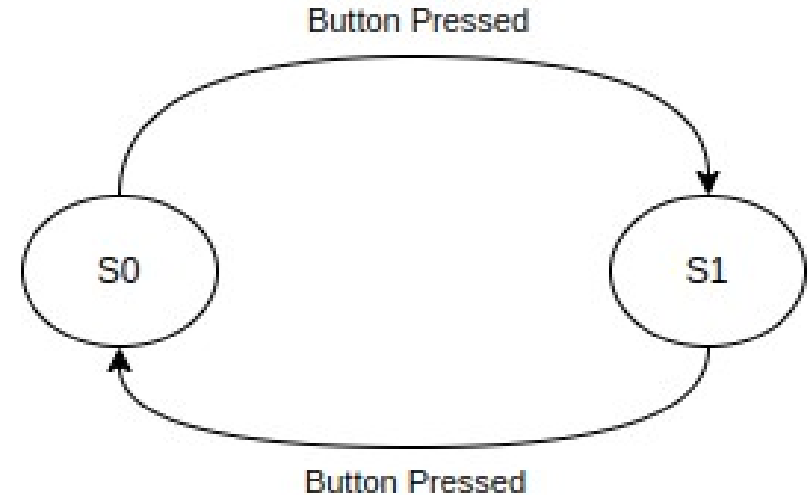


State transition diagram

- Lets consider pen as a machine

PEN: has push button on top and writing NIB on the bottom

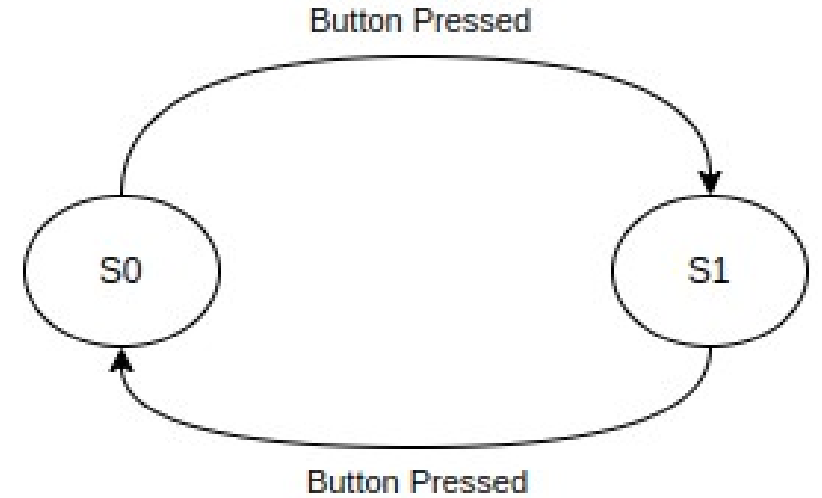
- S0 : NIB Retracted
- S1 : NIB Extended



State transition table

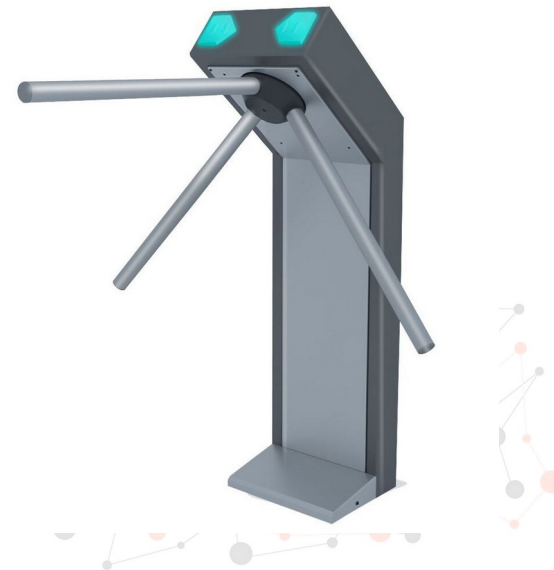
- Table with
 - All possible input
 - Current State
 - Output or Next state after input is applied

Input	Current State	Next State
Button pressed	NIB retracted	NIB extended
Button pressed	NIB extended	NIB retracted



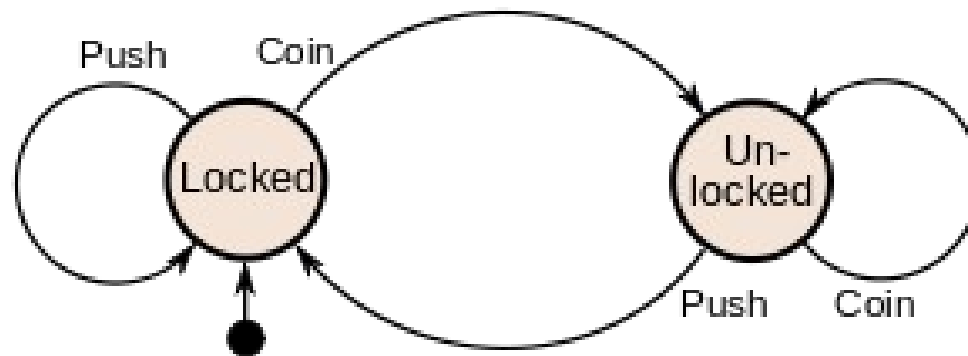
FSM – Example (Turnstile)

- Inserting a coin into a turnstile will unlock it, and after the **turnstile** has been pushed, it locks again. Inserting a coin into an unlocked turnstile, or pushing against a locked turnstile will not change its state



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Deterministic Finite State Machine (DFA):

- In a DFA, **each state has a well-defined transition** for every possible input.
- The transition from one state to another is uniquely determined by the current state and the input.
- DFAs are often used in scenarios where the system's behavior is **straightforward** and **deterministic**.



DFA : Formal definition

A deterministic finite automaton (DFA) is described by a five-element tuple: $(Q, \Sigma, \delta, q_0, F)$

Q = a finite set of states

Σ = a finite, nonempty input alphabet

δ = a series of transition functions

q_0 = the starting state

F = the set of accepting states

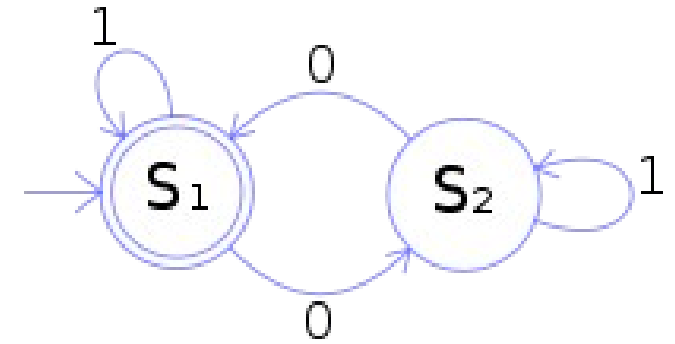
There must be exactly one transition function for every input symbol in Σ from each state.



DFA

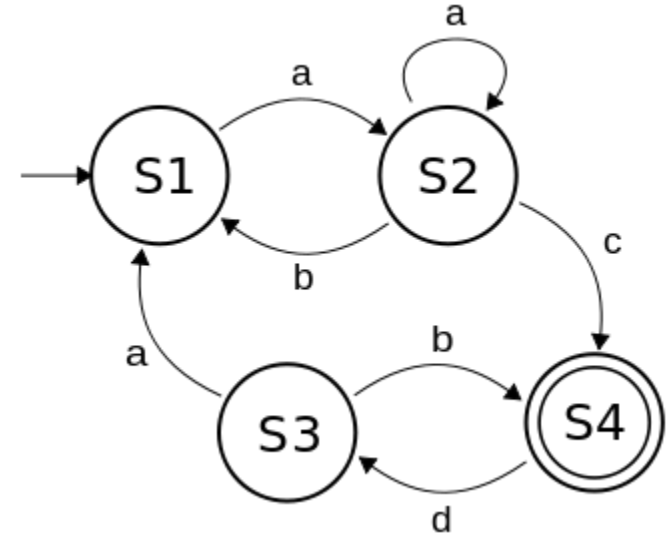
- $Q = \{s_1, s_2\}$
- $\Sigma = \{0, 1\}$
- $q_0 = s_1$
- $F = s_1$
- The following table describes δ :

current state	input symbol	new state
s_1	1	s_1
s_1	0	s_2
s_2	1	s_2
s_2	0	s_1



DFA

- What string cannot be generated by the finite state machine below?
 - abacdaac
 - abac
 - aaaaac
 - aaaacd



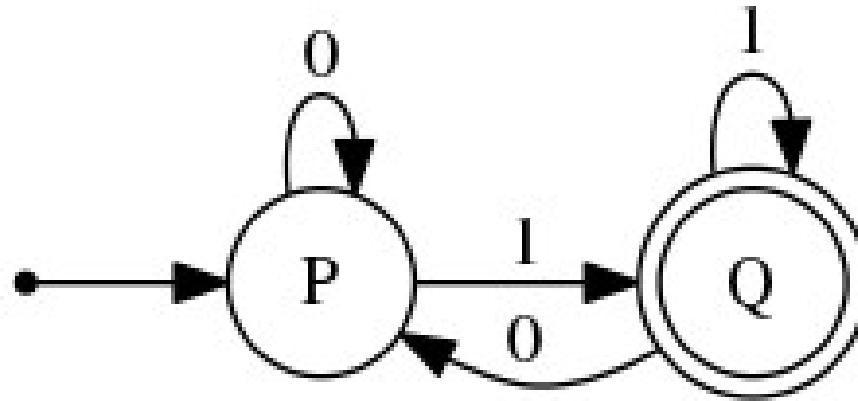
DFA

- Draw a diagram for a DFA that recognizes the following language:
 - The language of all strings that end with a 1.



DFA

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 - The language of all strings that end with a 1.



Non-Deterministic Finite State Machine (NFA)

- In an NFA, there can be **multiple possible transitions** for a given input in a given state.
- NFAs are used when the system's behavior is more complex and might have **multiple valid paths**.



NFA – Formal Defination

Similar to a DFA, a nondeterministic finite automaton (NDFA or NFA) is described by a five-element tuple: $(Q, \Sigma, \delta, q_0, F)$

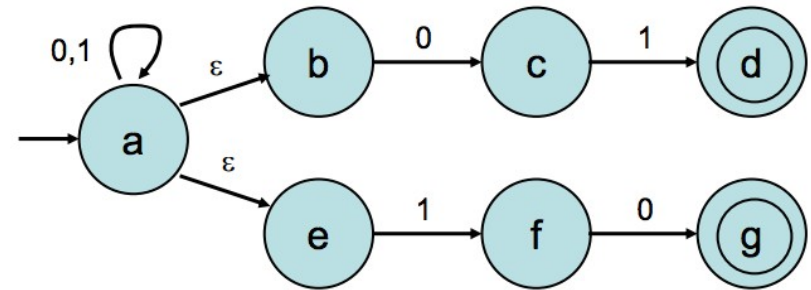
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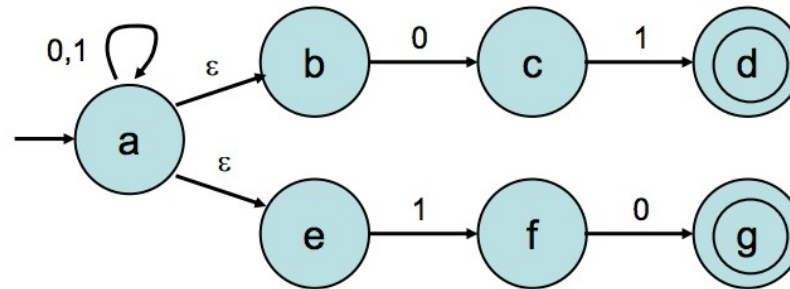
NFA – Formal Defination

- Unlike DFAs, NDFAs are not required to have transition functions for every symbol in Σ , and there can be **multiple transition functions** in the same state for the same symbol.
- Additionally, NDFAs can use **null transitions**, which are indicated by ϵ .
- Null transitions allow the machine to jump from one state to another without having to read a symbol.
- An NDFA accepts a string x if there exists a path that is compatible with that string that ends in an accept state.



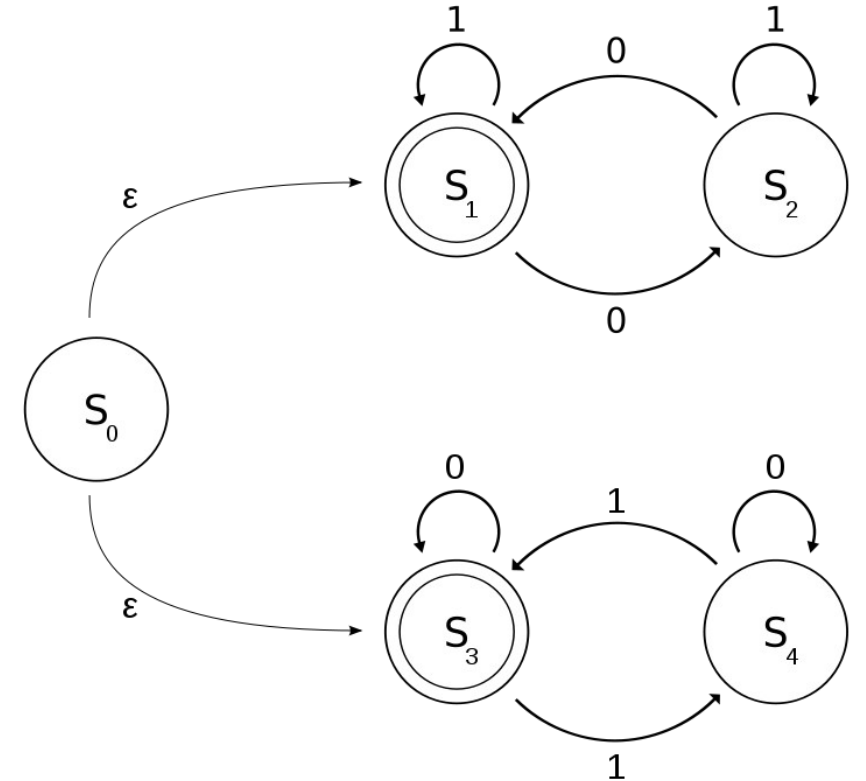
NDFA

- The NDFA that recognizes strings that end in “10” and strings that end in “01.”



NDFA

- Which string cannot be generated by the finite state machine below?
 - 1
 - 01001
 - 1011101
 - 1000
 - 0



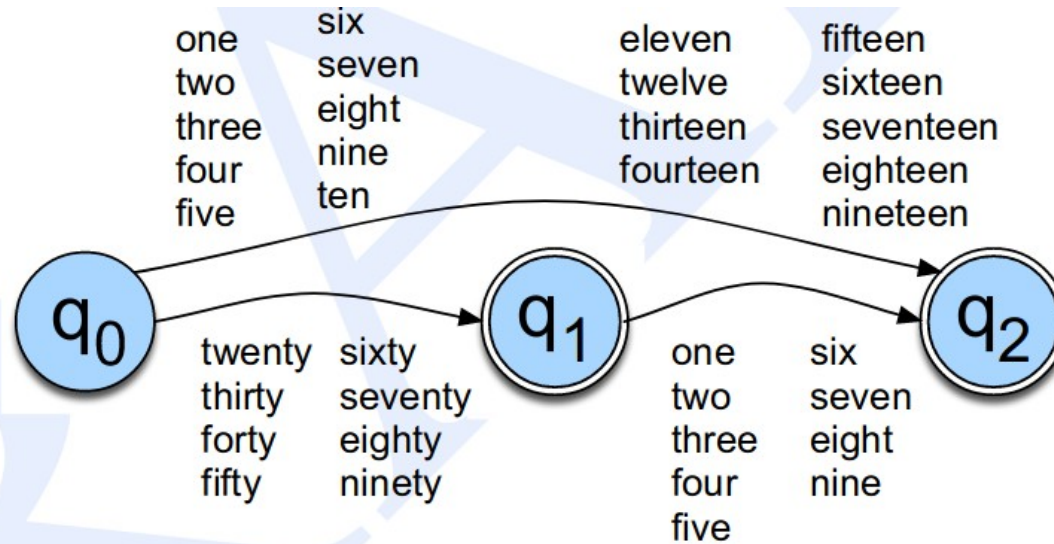
FSA - Example

- Make FSA for checking if given text represent english number or not?



FSA - Example

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An FSA for the words for English numbers 1–99



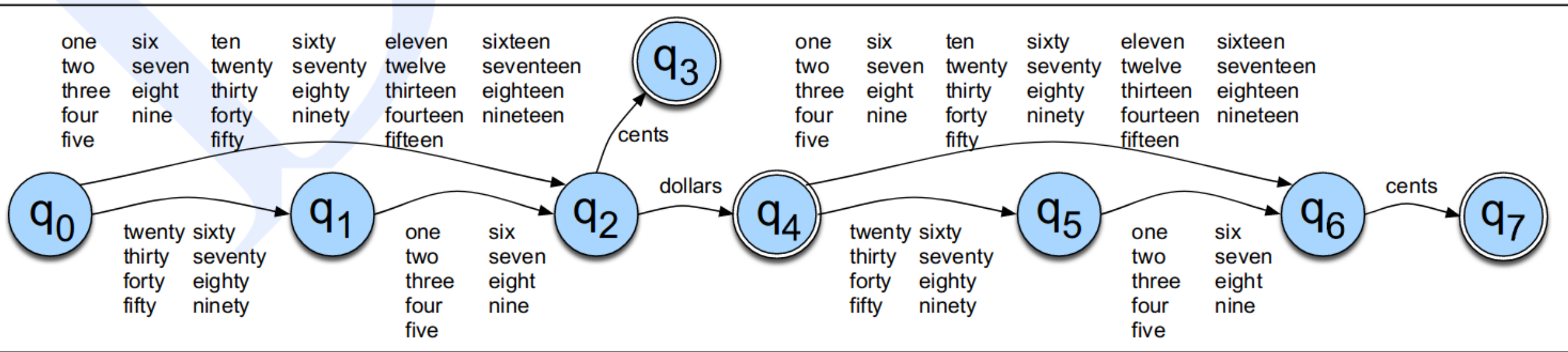
FSA - Example

- Make FSA to check if “two dollars four cents” is valid language or not.



FSA - Example

- Make FSA to check if “two dollars four cents” is valid language or not.



FSA for the simple dollars and cents



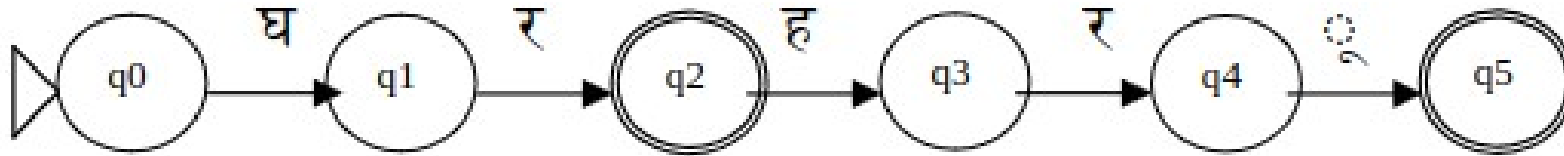
FSM – language processing

- Construct a FSM To validate the regular expression
 - Prefix/suffix detection
 - String end with “ing”



FSM – language processing

- For illustration, an automaton that accepts a string from the Nepali language घर 'house' and घरहरू 'houses' is visualized in Figure below



- This FSA accepts घर 'house' and घरहरू 'houses' because the inputs lead to final states. No other strings are accepted by this FSA.



NLP using Finite State Automata / Machine

- FSA/FSM, are useful in **classical, rule-based** NLP tasks that involve **pattern recognition, regular languages, or finite context constraints**.
- FSAs are not expressive enough for full natural language understanding,
- Still powerful tools in specific types of NLP tasks, particularly those **that don't require deep hierarchical or long-distance dependencies**.



FSA Use - Morphological Analysis

- Breaking words into morphemes (roots, suffixes, etc.).
- Finite State Transducers (FSTs, an extension of FSAs) are used for mapping surface forms to lemmas and morphological features.
- Example: walked → walk + PAST



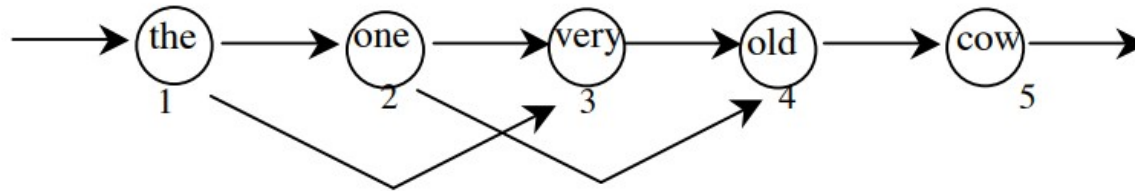
Weighted Finite State Machines (WFSM)

- In many applications, we are not so interested in whether a sequence is a valid sentence in a language
- For instance, speech recognition engines typically use a model of **word transition probabilities** (e.g., how often a word **w** follows a word **w'**) to drive the recognition process.
 - Adding such a model on top of straight acoustic recognition can double the word accuracy rate.
- We can model this data with a variant of an FSM called a **weighted finite state machine (WFSM)**.
- A WFSM is just like an FSM expect that it has numbers on the transitions.



Weighted Finite State Machines (WFSM)

- We can generate
 - ... the very old cow
 - ... the one very old cow
 - ... the one old cow

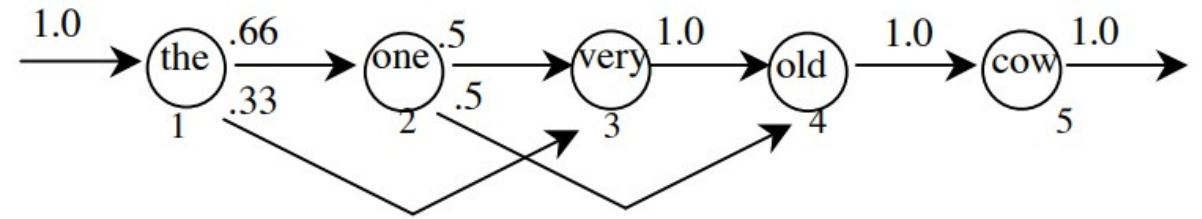


Which language is most likely?



Weighted Finite State Machines (WFSM)

- We can generate
 - ... the very old cow
 - ... the one very old cow
 - ... the one old cow



$$\begin{aligned} P(\text{The one old cow}) &= P(\text{transition 1 to 2}) * P(\text{transition 2 to 3}) * P(\text{transition 3 to 5}) \\ &= 0.66 * 0.5 * 1.0 = 0.33 \end{aligned}$$



Markov Chain

- A deterministic probabilistic finite state machine, as previous one is, is often called a **Markov Chain**.
- It has the nice properties of DFSA in that we can very quickly determine whether a sequence is accepted by the machine.
- In addition, **we can compute the probability of a sequence by simply multiplying the probabilities** on each of the transitions together.



Thank you

