# OS 2018

Homework 4: Memory Allocator Implementation

(Due date 01/03 23:59:59)



## **Objective**

- Understand how malloc() and free() work
- Understand how to manage Heap



# Requirements



## Requirements

- 1. Implement a **Memory Allocator Library** for user application (*slides 4-16*)
- 2. Write a **User Application** to test the memory allocator library (*slides 17,18*)
- 3. Please follow the **Input/output format** (*slides 19,20*)



## 1. Memory Allocator Library Requirements

- The library must provide the following **3 functions**, and you should follow the format in next slide:
  - 1. hw\_malloc() (slides 6-13)
  - 2. hw free() (*slides 14-16*)
  - 3. hw\_get\_start\_brk() (heap 性忧化量)
- Use chunk (slides 7,8), bin (slides 9,10) and sbrk() to manage heap
- Use chunk (*slides 7,8*) to manage every mmap-allocated memory



## **Functions format**

### void \*hw\_malloc(size\_t bytes)

• bytes: the required memory size in bytes

• Return =  $\begin{cases} \text{the valid virtual address} \\ \text{(starting address of the data part)} \end{cases}, if success \\ \text{NULL} , otherwise \end{cases}$ 

#### 2. int hw free(void \*mem)

- mem: starting address of the data part
- Return =  $\begin{cases} 1, & if success \\ 0, & otherwise \end{cases}$

## 3. void \*hw\_get\_start\_brk()

• Return the starting address of the heap



## hw\_malloc() requirements

- Use mmap\_threshold to decide the memory allocate method. If the allocated size(data size + chunk header size) > mmap\_threshold, use mmap allocation method; else use Heap allocation method.
  - mmap threshold is **initial**: **32 KiB** (32 \* 1024)

#### mmap allocation method:

- Use **mmap()** system call to allocate the space
- Use **chunk** (slides 7,8) to manage the allocated space

The allocate size = request size (data size) + chunk header size

• Use **mmap\_alloc\_list** (slide 11) to manage allocated mmap chunks

#### **Heap allocation method:**

- Use **chunk** (slides 7,8) and **bin** (slides 9,10) to manage heap
- Should follow the rules of **Heap initialization** (*slide 12*)
- Should follow the rules of **Split** (*slide 13*)
- The allocated size (data size + chunk header size) should be the best fit size



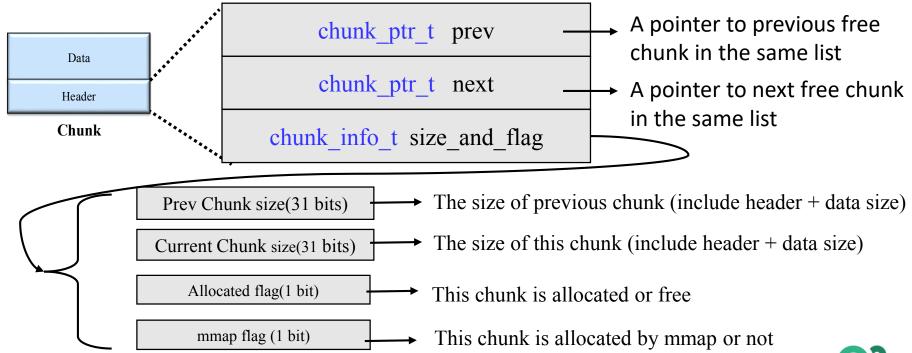
## Chunk requirements

- The continuous heap space is split into chunk(s) for management
- Each chunk contains two parts, header and data (*in next slide*)
  - Header (lower address)
  - Data (higher address), the actual memory space return to caller



## Chunk header format

- Chuck header (24 bytes)
  - There are 3 members in the header
    - chunk\_ptr\_t, chunk\_size\_t, and chunk\_sizeandflag\_t can be defined by yourself, but each of them should be 8 bytes
    - <a href="mailto:chunk\_info\_t">chunk\_info\_t</a> should include 4 information (Prev Chunk size, Current Chunk size, Allocates flag and mmap flag)



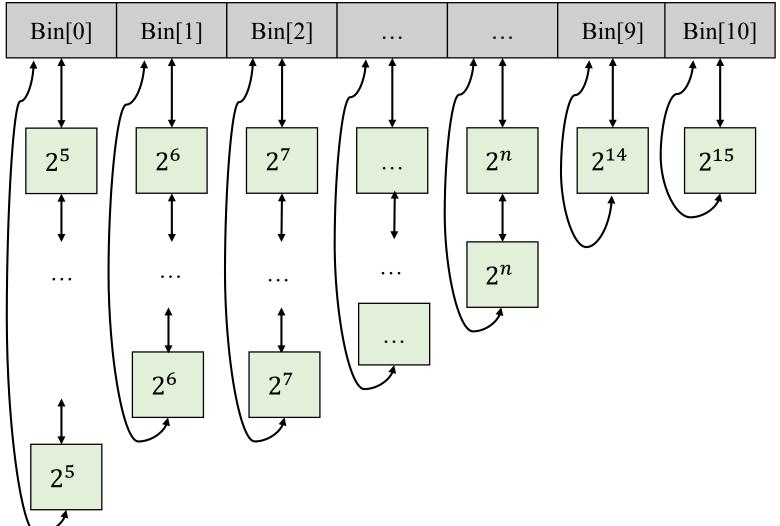


## Bins requirements

- Bins are use to manage free chunks of Heap's segment.
- A bin is a circular doubly-linked list of free chunk(s) (next slide)
- Add the chunk to rear of the bin.
- You should manage 11 bins
  - bin[0]-bin[10] hold chunks with fixed size (next slide)
  - Every chunk size should be the nth power of 2 (n is a number of 5 to 15)
- Use the best fit size to select a chunk during memory allocation
  - If there are multiple chunks with the same size, select the one with the lowest address

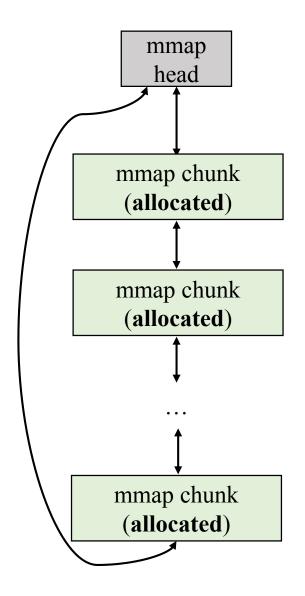


# Bin example





# mmap\_alloc\_list

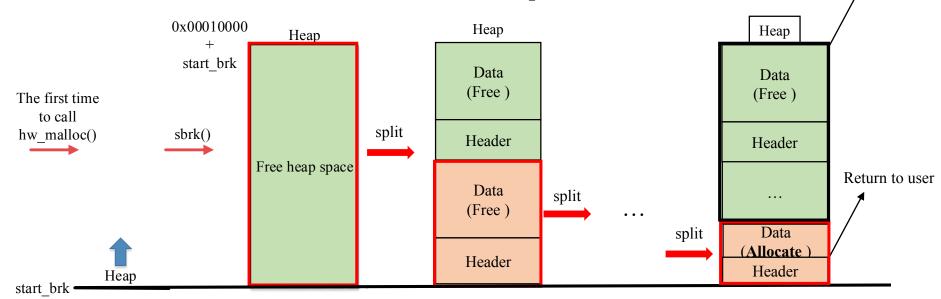


- mmap\_alloc\_list is a circular doublylinked list of mmap allocated chunk(s)
- mmap chunk header is same as slide 8
   (use chunk ptr t)
- mmap\_alloc\_list should be ordered by size (ascending)
- If there is/are multiple chunk(s) of the same size, add new chunk after it/them.



# Heap initialization & First-time Heap allocation

- Initialize the Heap before its first use :
  - Use sbrk() to allocate a **64KiB** heap
- After the heap initialization, Split (*next slide*) until the allocated size match the best fit of the nth power of 2
  - an allocated chunk (lower address), returned to the caller
  - Other free chunks should insert into specific bins



Heap initialization

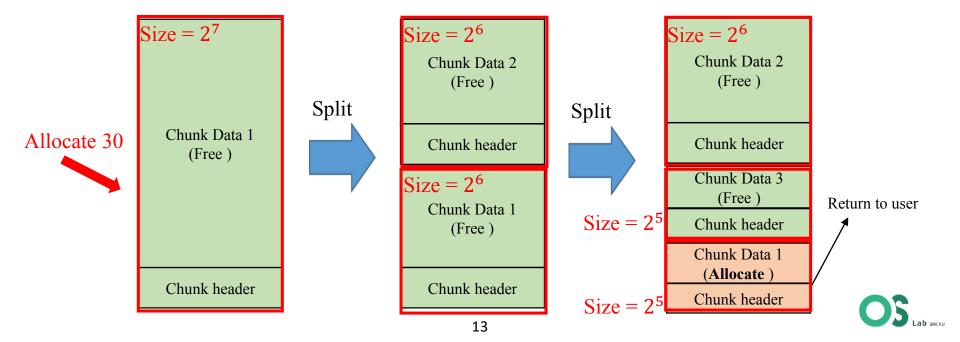


Insert into

specific bins

## **Split**

- When Heap allocation method is used, *split* may be performed:
  - If chunk size is too large (not the best fit for the allocation size), they must be split into two equal-sized chunks
  - Should split the lowest address chunk
  - Chunk size should always be the nth power of the 2
  - Must be split until it reach the best fit for the allocation size
- Example: (hw\_malloc(b), allocation size = 30, Chunk1 = 2<sup>7</sup>)



## hw\_free() requirements

• Use Chunk header information (*slides 7, 8*) to check the address was allocated by mmap or Heap allocation method.

## mmap free method:

- Use **munmap()** system call to free the space
- The chunk header should be free too.
- This free chunk do not need to be added into bin.
   (bin is not used in mmap allocation method)

## Heap free method:

- Use bin (*slides 9,10*) to manage free chunks
- Should follow the rules of **Merge** (*slides 15,16*)



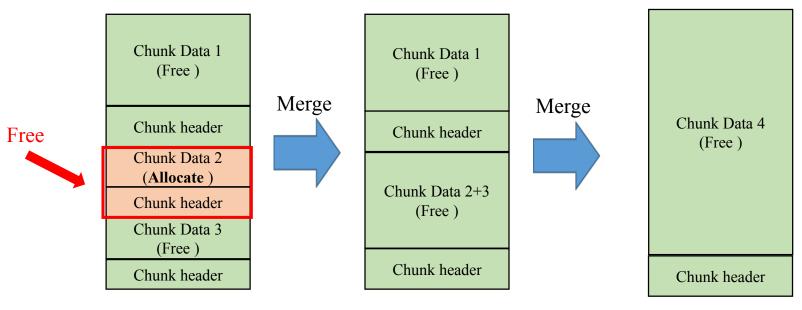
## Merge

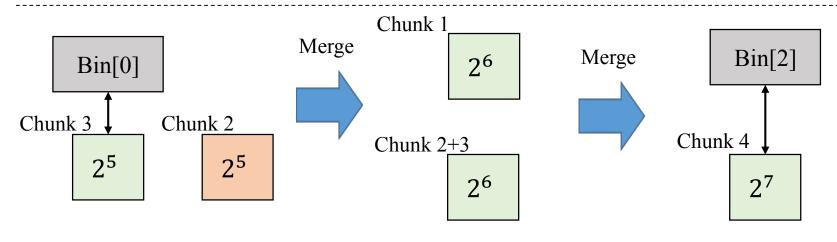
- When Heap free method is called, adjacent free chunks must be merged into one if they have the same size:
  - Merge operations should be repeated until the two adjacent free chunks do not have the same size



## Merge Example

• Example: (Chunk 1 size  $=2^6$ , Chunk 2 size = Chunk 3 size  $=2^5$ )







# 2. User Application Requirements

Write a user application to test the memory allocator library

- Should receive 4 kinds of commands:
  - 1. alloc N
  - 2. free ADDR
  - 3. print BIN
  - 4. print mmap\_alloc\_list
- Continuously receive commands from stdin until *EOF* (*Ctrl+D*)
  - Should successfully run

    "cat testfile.txt | hw4\_mm\_test > outputfile.txt"



## **Commands format**

#### 1. alloc N

- Call hw malloc(N) to allocate N bytes of data memory
- Print relative data address (i.e., offset between start\_brk and the address returned by hw\_malloc())

#### 2. free ADDR

- Call hw\_free() to free the memory at (start\_brk + ADDR)
- Print either "success" or "fail"

## 3. print bin[i]

- Print relative data address and size information of a given bin
  - bin[i] can be bin[0], bin[1] ..... bin[10]

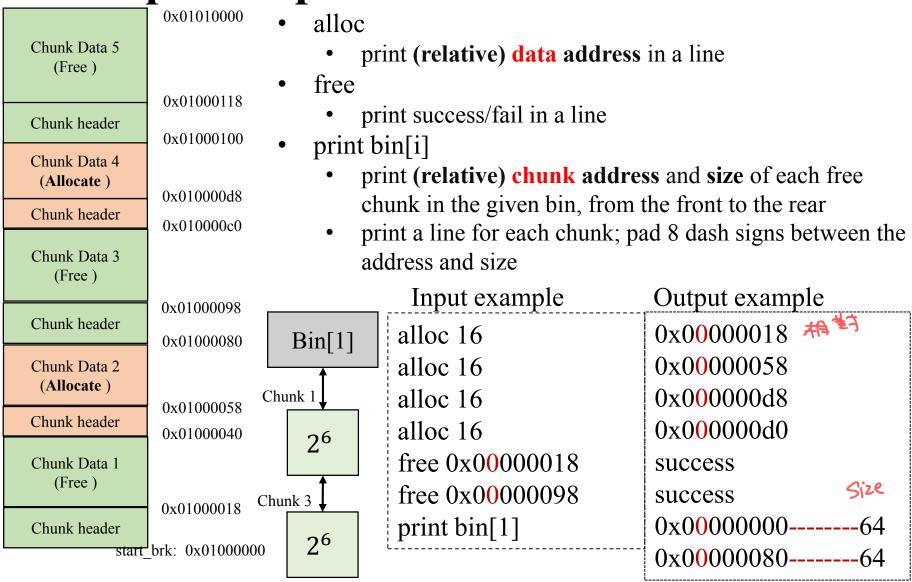
## 4. print mmap\_alloc\_list

Print data address and size information.

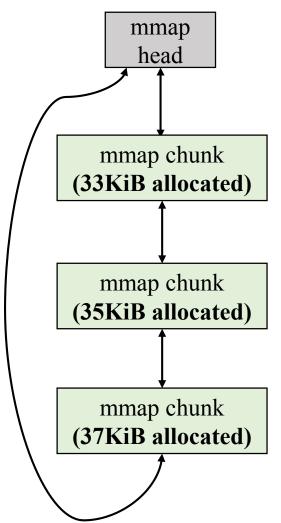




## 3. Input/output format



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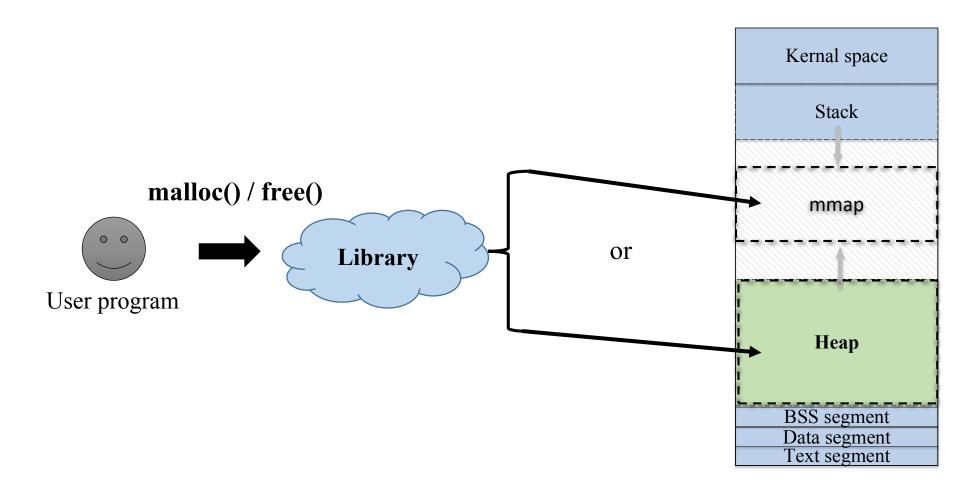
- print mmap alloc list
  - print chunk address and size of each free chunk in the list, from the front to the rear
  - print a line for each chunk; pad 8 dash signs between the address and size

Output example	_
0xfdcf00018	
0xfecff0f018	
0xffcff0f018	
0xfdcf0000033792	
0xffcff0f00035840	
0xfecff0f00037888	

Output example



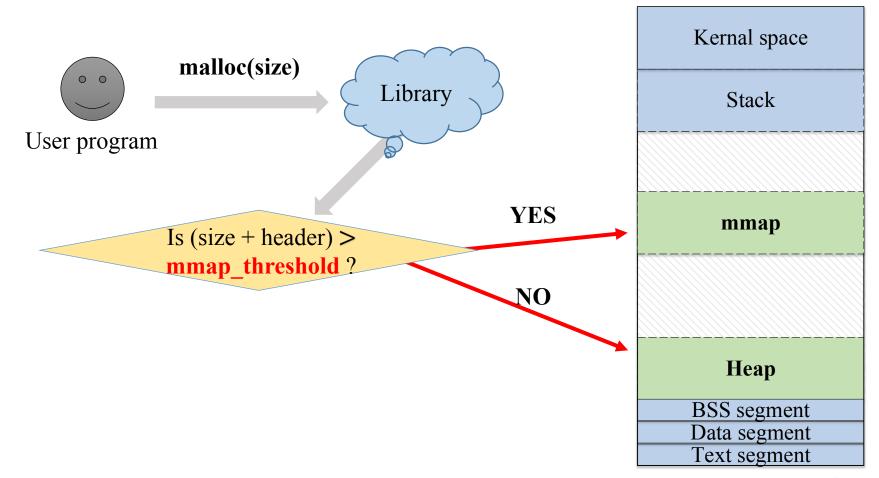
## Concepts (malloc / free)





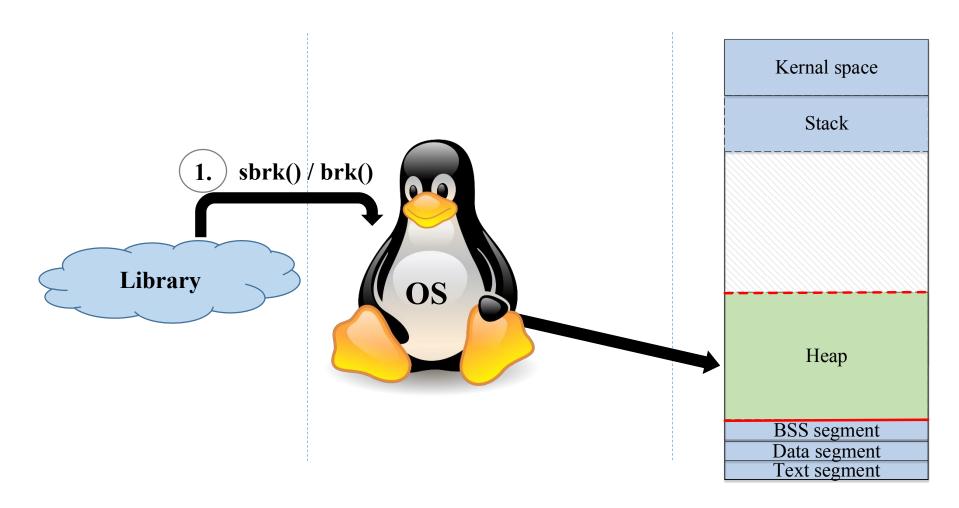
## Concepts (mmap\_threshold)

Use mmap\_threshold to decide the memory allocate method





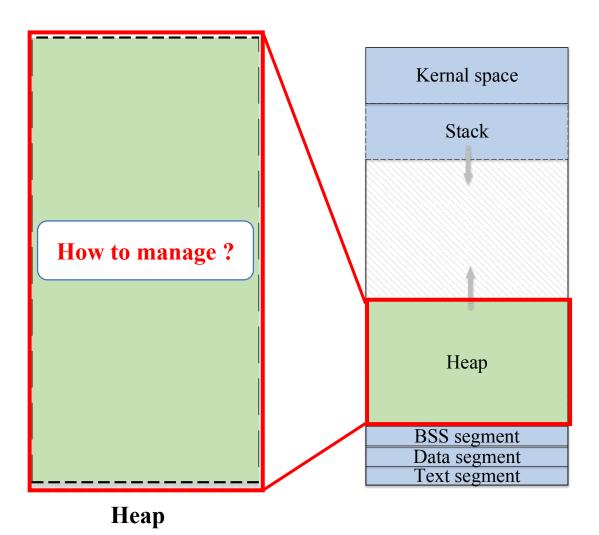
# **Concepts (Heap initialization)**





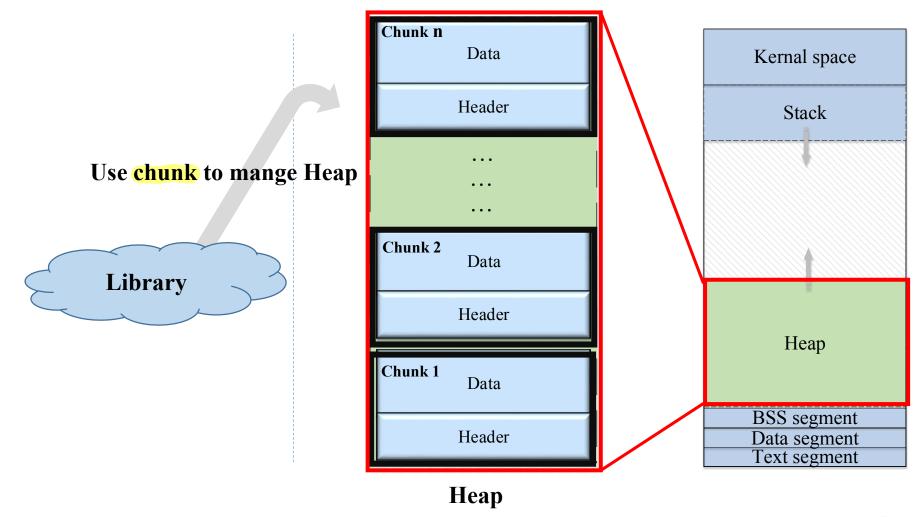
# Concepts (heap)

Library

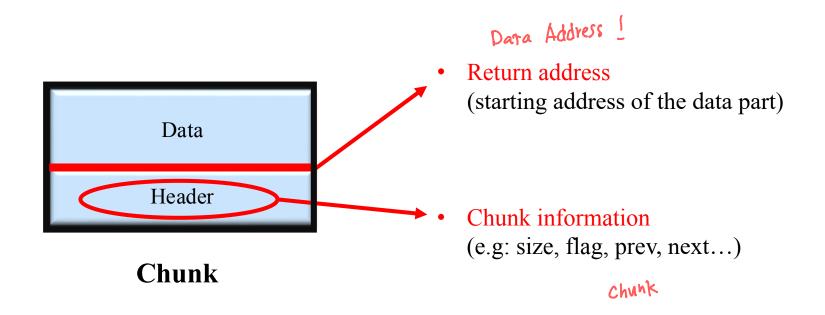




# **Concepts (heap segment)**

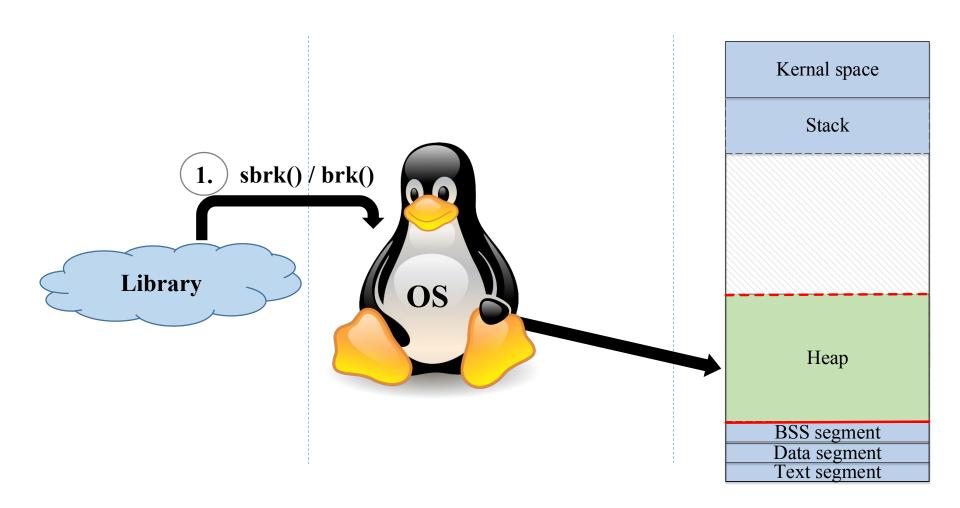


## Concepts (chunk)

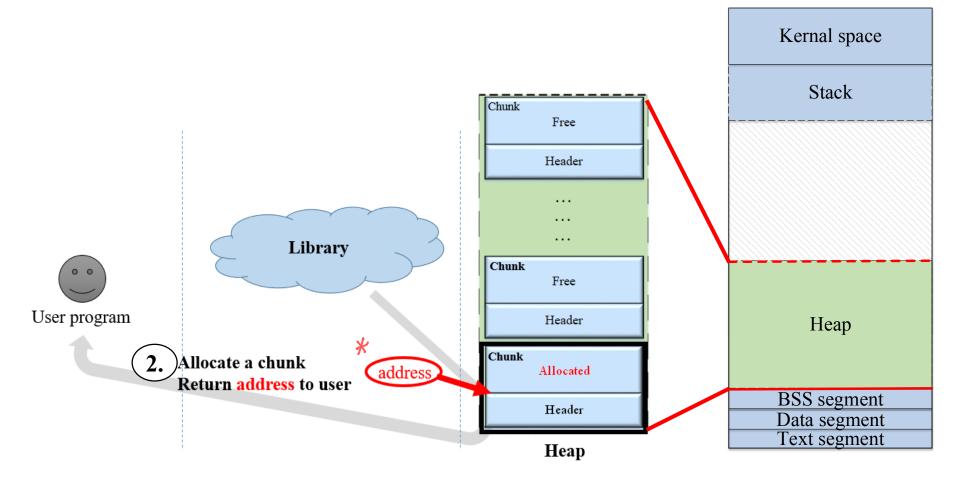




# **Concepts (Heap initialization)**

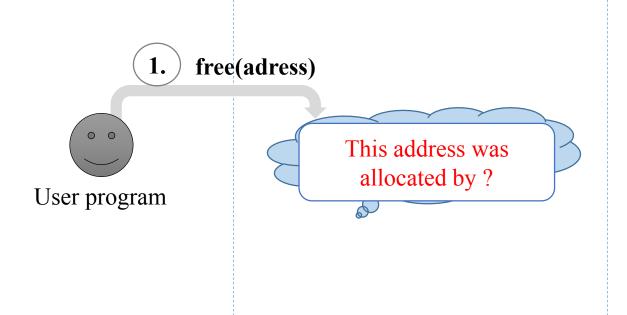


## **Concepts (heap segment)**





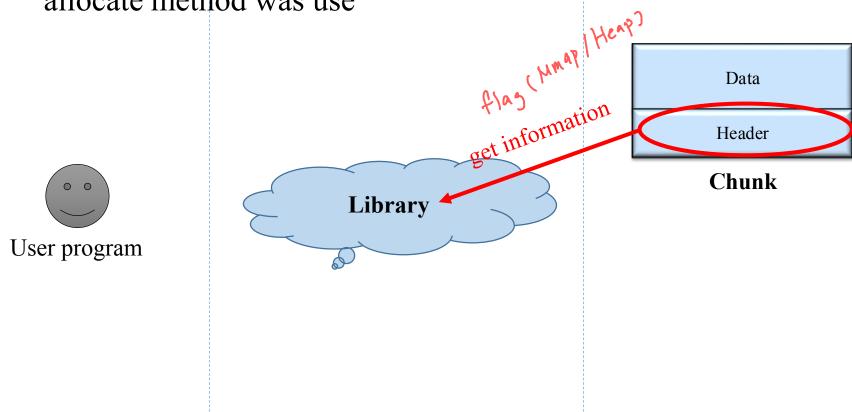
# **Concepts (free)**





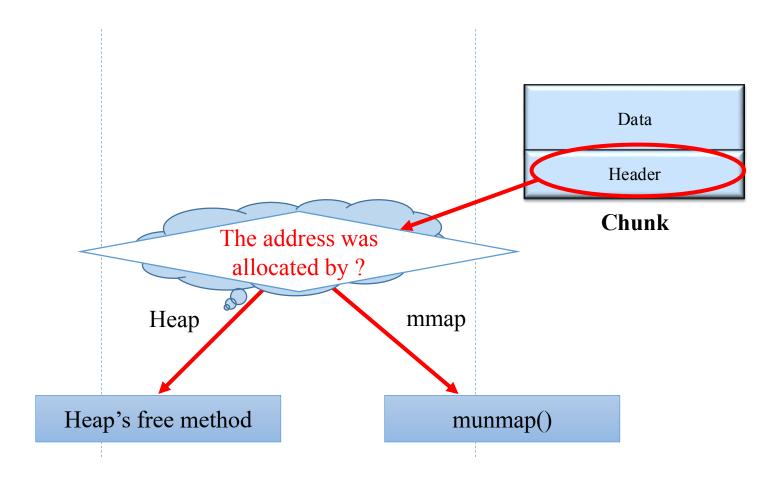
## **Concepts** (free)

• Get information from Chunk header to know which memory allocate method was use



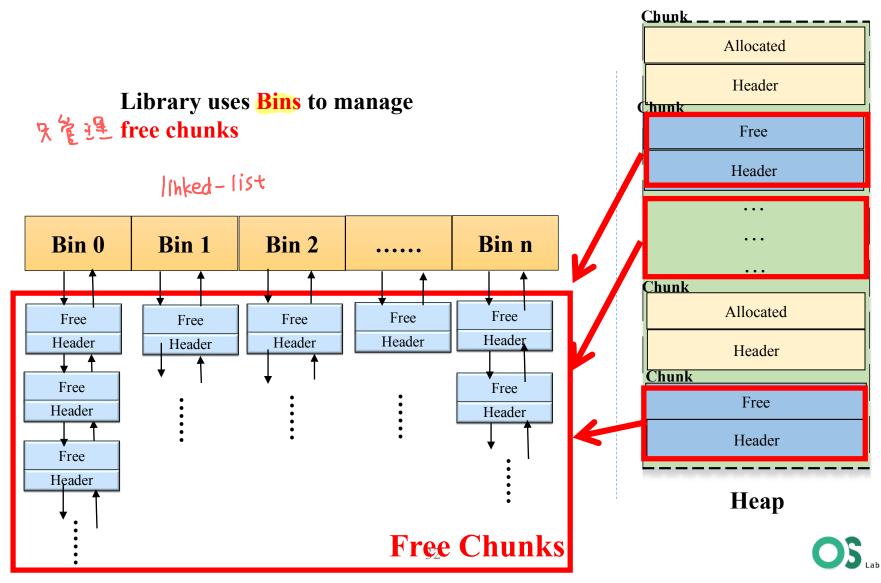


# **Concepts (free)**

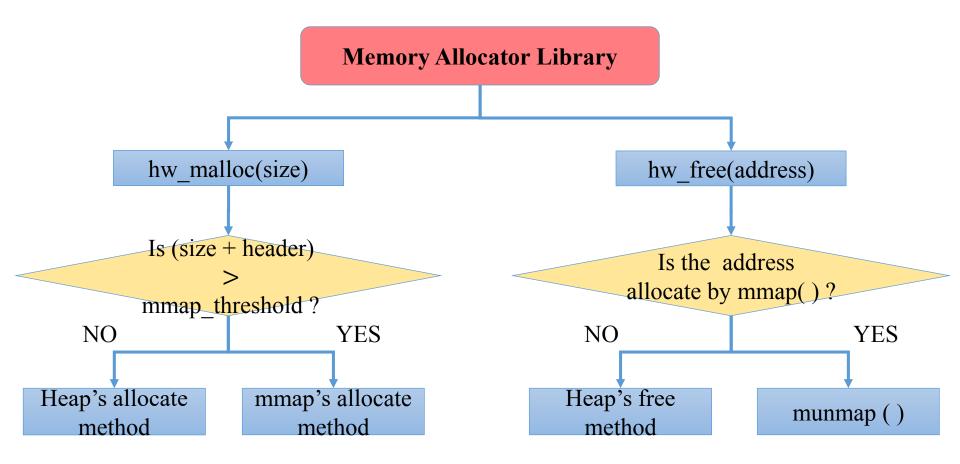




# Concepts (Heap's free method)



## **Concepts (Architecture)**





## References

- sbrk()
  - Linux man page
- Streams, pipes, and redirects
  - IBM

