

# Efficient Temporal Synopsis of Social Media Streams

## ABSTRACT

Search and summarization of streaming social media, such as Twitter, requires the ongoing analysis of large volumes of data with dynamically changing characteristics. Tweets are short and repetitious - lacking context and structure - making it difficult to generate a coherent synopsis of events within a given time period. Although some established algorithms for frequent itemset analysis might provide an efficient foundation for synopsis generation, the unmodified application of standard methods produces a complex mass of rules, dominated by common language constructs and many trivial variations on topically related results. Moreover, these results are not necessarily specific to events within the time period of interest. To address these problems, we build upon the Linear time Closed itemset Mining (LCM) algorithm, which is particularly suited to the large and sparse vocabulary of tweets. LCM generates only closed itemsets, providing an immediate reduction in the number of trivial results. To reduce the impact of function words and common language constructs, we apply a straightforward filtering step that preserves these terms only when they may form part of a relevant collocation. To further reduce trivial results, we propose a novel strengthening of the closure condition of LCM to retain only those results that exceed a threshold of distinctiveness. Finally, we perform temporal ranking, based on a simple information gain measure, to identify results that are particularly relevant to the time period of interest. We evaluate our work over a collection of Tweets gathered in late 2012, exploring the efficiency and filtering characteristic of each processing step, both individually and collectively. Based on our experience, the resulting synopses from various time periods provide understandable and meaningful pictures of events within those periods, with potential application to tasks such as temporal summarization and query expansion for search.

## Categories and Subject Descriptors

H.4 [Information Systems Applications]: Miscellaneous;

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D.2.8 [Software Engineering]: Metrics—*complexity measures, performance measures*

## General Terms

Theory

## Keywords

ACM proceedings, L<sup>A</sup>T<sub>E</sub>X, text tagging

## 1. INTRODUCTION

The nature of text in social media poses a challenge when applying traditional search and text mining algorithms. Text in social media is usually short, lacks context and structure, and is created at a very high rate. The sheer volume of social media streams mandates the use of efficient algorithms, and the short length of individual documents makes it possible to apply Frequent Itemsets Mining. This family of algorithms is very fast and efficient, however it is not readily suited for application on text. The problem with Frequent Itemsets Mining is the large number of itemsets it produces. Actually it was originally proposed as a preliminary stage to Association Rules Mining, which sifts through the numerous itemsets and produces a smaller number of association rules. The number of itemsets produced can be reduced by setting a higher threshold on the minimum frequency of an itemset. However, in text this threshold has to be kept low because frequencies follow a long tailed Zipfean distribution. Also, the head of the distribution is made up mostly of stop words; a situation that doesn't happen in the domain of market basket data since the packaging of each item is not enlisted in the receipt! Even if a maximum frequency threshold is set, risking to filter out important itemsets, a lot of non-English language constructs will be mined because the proportion of posts in English is much higher than other languages. In this paper we propose methods for adapting Frequent Itemset Mining to be effective on social media text, without degrading its efficiency.

Unlike trending topics<sup>1</sup> [13], the results of Frequent Itemset Mining include itemsets that have high frequency because of sustained relatively high interest as well as a spike of interest. The mined itemsets provide the vocabulary associated with events and can be used in various ways for search and summarization. For example, the collection of mining results of different epochs of time can be used for

<sup>1</sup><http://blog.twitter.com/2010/12/to-trend-or-not-to-trend.html>

temporal query expansion [4]. The mining results of each epoch can be treated as a document, so that the probability of an epoch conditioned on a query can be estimated as the probability of this document on the query using any pseudo-relevance feedback method. Further, terms from individual itemsets relevant to a query (using any relevance ranking) can be used for query expansion, thus acting as precomputed results of a simple version of pseudo-relevance feedback. We also propose a quality ranking scheme that makes the mining results presentable to human users. We don't call this a summary because we don't take into account the quality measures of summaries such as coherence and cohesion. However, the top ranked itemsets cover a variety of open topics, and within one topic different opinions are reported as separate contrastive itemsets.

**Paper outline.** The first three sections provide the necessary background for the rest of the paper. We start by discussing related work in section 2, then we explain Frequent Itemset Mining and the algorithm on which we build our work in sections 3 and 4 respectively. In sections 5, 6 and 7 we propose solutions to different problems faced when applying frequent itemset mining to social media text. We present the outline of those sections graphically using a frequency ordered prefix tree of the itemsets. Each path from root to a leaf in such a tree represents an Itemsets, where each itemset is ordered in non-increasing order of the frequency of items. Itemsets having the same prefix share the nodes representing this prefix. This representation is typical in the literature, and it is actually a very good representation of the Frequent Itemset Mining problem. Figure 1 shows conceptually how section 5, 6 and 7 addresses different parts of the problem. In section 8, we conclude and suggest future directions.

## 2. RELATED WORK

Frequent Itemset Mining is a very large body of work that goes back to the early 90s. We cover the topic only briefly, as the focus of this paper is not Frequent Itemsets Mining but rather its adaptation to social media text. The original Apriori algorithm [1] and algorithms based on it suffer performance degradation and a big increase in memory requirement when the number of distinct items is high. This is caused by the candidate generation bottleneck as explained later. Another famous class of mining algorithms is the FP-Growth [7] based algorithms. FP-Growth skips the candidate generation step, and instead creates a succinct representation of the data as a frequency ordered prefix tree called the FP-tree. An FP-tree imposes the invariant that within a branch the frequency is non-increasing from the root down to the leaves. The memory requirements of FP-Growth algorithm suffers from the sparsity of the data, since the data structure is succinct only if it can find common prefixes within the constraints of its invariant. A less known algorithm that is robust against data sparsity is LCM [16], which we will describe in detail in section 4. We use the implementation submitted to the workshop for Frequent Itemset Mining Implementations (FIMI) '04 [8], which is the workshop's award winner.

The problem that there are too many itemsets, with a lot of noise and redundancy, is addressed either by clustering [18] the itemsets or by mining only itemsets that satisfy a certain condition. The closure condition [14] is a prominent condition upon which many other conditions are based.

Most similar to the strongly closed itemsets we propose are the  $\delta$ -covered [17] and the  $\delta$ -free [3] sets. The  $\delta$ -covered condition is exactly the opposite of the strong closure condition, and it "relaxes the closure condition to further reduce pattern set size" [12]. The  $\delta$ -free condition is like the strongly closed condition but it is used to eliminate different itemsets, since the motivation is to provide a compressed representation of itemsets by sacrificing support information.

The motivation of most of the work in finding representative itemsets is compressing itemsets mined from a general dataset. This leads to decisions different than what are made in order to find the most informative itemsets mined from text. For example, in [11] itemsets are mined from paragraphs of newswire text, and used to determine term weights for query expansion. Improvements in performance have been achieved by using itemsets known to come from a training set of related documents as well as ones from unrelated documents. In [10], similar methods of term weighting were used in pseudo-relevance feedback for Twitter search, and achieved improvements on a weak baseline. **should I include the afore mentioned attempt?** On the other hand, in [19] LCM is also used to mine Twitter posts and then a few itemsets are selected as a "summary" of the data, but they are selected according to their utility in a lossy compression algorithm. The quality of the summary seems to be affected by the choice of utility function, but the only assessment made about this was showing that the actual transactions can be reconstructed from the summary with an accuracy that is expected to be high if the Tweet contains "both recent and frequent keywords". In the experiments, non-negative matrix factorization is used to extract topics from parts of the summary matching a query (world cup) and specific time intervals (before certain matches). It is unclear whether the raw summary would cover a variety of topics, specially non-trending ones, and no ranking scheme was shown to pick interesting topics without specifying a query. Moreover, the data is preprocessed by stemming and stop words removal, which should require a language identification component upstream but it is not discussed.

## 3. FREQUENT ITEMSET MINING

### 3.1 Preliminaries

Traditionally, Frequent Itemset Mining is applied to a *database of transactions* made at a retail store. This terminology is suitable for market basket data and we will stick to it out of convention, even though we are mining text where the terms corpus and document are normally used. Because of the dynamic nature of social media, rather than giving the whole database as input to mining algorithms, the input is an *epoch* of data; data with timestamps within a certain period of time. The epoch's *span* is the length of this period in hours, and the *volume* at this epoch is the number of transactions in the epoch divided by its *span*.

A *frequent itemset* is a set of items that occur together a number of times higher than a given threshold, called the *support* threshold. We also adapt the support threshold to the dynamicity of the *volume* at different times of the day. We define the *minimum support threshold* as the threshold at the hour of the least *volume* during the day. The *minimum support* is supplied as an absolute number,  $a$ , and then converted to a ratio,  $\alpha = \frac{a}{\text{avg}(\text{volume}_{\text{daily}} - \text{minimum})}$ . The actual support threshold used for mining any given epoch is

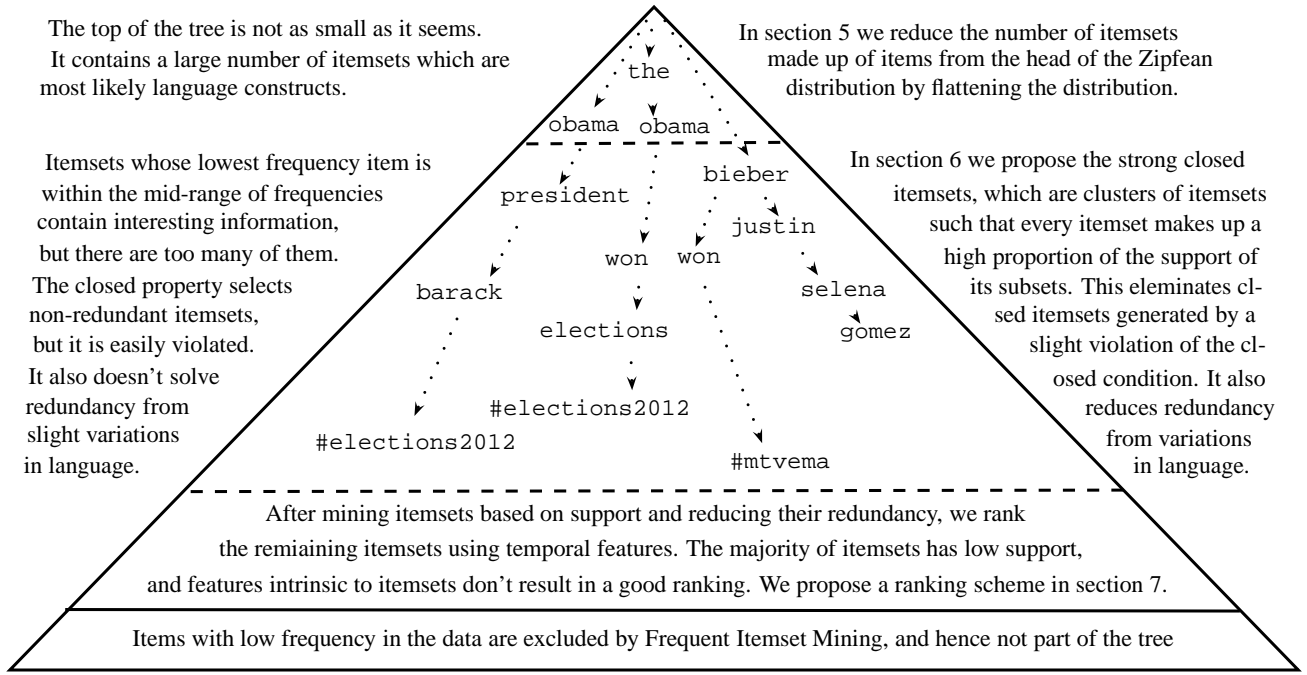


Figure 1: Roadmap of the paper overlaid on an frequency ordered prefix tree

thus  $\alpha$  multiplied by the *epoch's volume*.

We now introduce the notation used in this paper:

- $W = \{w_1, w_2, \dots, w_n\}$ : The set of all items (vocabulary).
- $t_a = \{w_{a1}, w_{a2}, \dots\}$ : A transaction made up of a set of items, not necessarily terms. Each transaction has a sequential id denoted by the subscript letter.
- $E^{span} = \langle t_a, t_b, \dots \rangle$ : An epoch of data of a certain span, such as an hour, made up of a sequence of transactions.
- $s \subset I$ : An itemset, its support is given by  $|T_s|$ .
- $T_s = \{t : s \subseteq t\}$ : All transactions containing itemset  $s$ .
- $|\cdot|$ : Cardinality operator; gives the size of the operand.

### 3.2 Background

The two basic operations of Frequent Itemset Mining algorithms are *Candidate Generation* and *Solution Pruning*. The original Apriori algorithm by Agrawal et al. [1] generates candidates of length  $K$  ( $K$ -itemsets) by merging frequent itemsets of length  $(K-1)$  ( $(K-1)$ -itemsets) that differ in only 1 item, usually the last item given a certain total ordering of items. By using only frequent  $(K-1)$ -itemsets for generating candidate  $K$ -itemsets a lot of possible  $K$ -itemsets are implicitly pruned, based on that all subsets of a frequent itemset has to be frequent (the Apriori property). This still generates a very large number of candidates, specially in early iterations of the algorithm. Consider, for example, the generation of candidate 2-itemsets from a database. This requires producing all unordered pairs of 1-itemsets (terms), after pruning out rare ones with frequency less than the support threshold. In many domains, including text mining, the number of frequent 1-itemsets is large enough to

prohibit generating a number of candidates in the order of this number squared. In text mining, a rather low support threshold has to be used, because the frequency of terms follow a long tailed Zipfean distribution.

### 4. BASIC ALGORITHM

To overcome the bottleneck of *Candidate Generation*, many algorithms are proposed to take hints from the transaction space rather than operating blindly in the item space. Some algorithms do that by traversing a data structure representing the transactions [7], and others do that by generating itemsets having a certain property that helps pruning out more candidates. In this paper we expand on LCM [16], an algorithm based on a property of a class of itemsets called *Closed Itemsets* [14]. A closed itemset contains any item that is present in all the transactions containing this itemset. A formal definition of closed itemsets is given in equation 1:

$$\mathcal{C} = \{s_c : \nexists s_d \text{ where } s_c \subset s_d \text{ and } |T_{s_c}| = |T_{s_d}|\} \quad (1)$$

The properties of Closed Itemsets are the following:

1. Adding any item to a closed itemset reduces its support.
2. A subset of a closed itemset is not necessarily closed, but one or more closed subset must exist for any itemset (formally this could be the empty set, given that any item that appears in all transactions is removed in a preprocessing step).
3. If a closed  $K$ -itemset can be extended any further then one of its supersets will be closed, however not necessarily a  $(K+1)$  superset. Itemsets that cannot be extended any further are called *Maximal Itemsets*, and they are a subclass of closed itemsets.

Besides being much smaller than the solution space of frequent itemsets, the solution space of closed itemsets can be navigated efficiently. By using an arbitrary total ordering of items, any closed itemset can be considered an extension of exactly one of its subsets. Thus, only this subset is extended during candidate generation. All the other subsets do not need to be extended by items that would lead to the longer closed itemset. This is called *Prefix Preserving Closure Extension (PPC-Extension)* and it is proposed and formally proved in [16]. This is achieved by following three rules, which we state after a few definitions to facilitate their statement. First, an item is *larger/smaller* than another item if it comes later/earlier in the total ordering. This terminology comes from that LCM is most efficient if the items are ordered in ascending order of their frequency. Second, the *suffix* of an itemset is one or more items whose removal does not result in an itemset with higher support. Notice that they will necessarily be at the end of the itemset, regardless of the total ordering. Finally, we call the first item added to the suffix of the itemset its *suffix head*. With this terminology, the rules for *PPC-Extension* are:

1. An itemset can be extended only by items *larger* than its *suffix head*. Extending by *smaller* items will lead to closed itemsets already generated.
2. After forming an itemset  $S$ , add to its *suffix* all items whose frequency within  $D_S$  is equal to  $|D_S|$ .
3. If any item in the *suffix* is *smaller* than the suffix head, prune this solution branch. All closed itemsets within this branch have already been generated.

Table 1 is an example of how *PPC-Extension* is used to generate closed itemsets starting from the 1-itemset ‘barack’. The upper table enumerates  $D_{barack}$ . The lower table shows steps of itemsets generation. The current solution along with its frequency is in column 2, solutions marked by an \* are the closed itemsets emitted. All possible extension items and their frequencies are in column 3 with the one being considered bolded. Column 4 is a comment explaining the step. At each step, a pass is done on  $D_{itemset}$  to enumerate and count possible extension items. To enforce a support threshold infrequent extension items are removed, but in this example there isn’t such a threshold. Notice that the number of steps is linear in the number of closed itemsets, and the only additional storage required besides the storage of the documents is that of the possible extension items. Of course this is a simplified example, but it shows in essence how LCM achieves its low run time and memory requirements. We refer the interested reader to [16] for a theoretical proof that the algorithm runs in linear time in the number of closed itemsets, and that this number is quadratic in the number of transactions. Performance on a real data set is shown in section 5. We proceed by describing how to implement this algorithm using an inverted index.

## 4.1 Implementation details

We show in algorithm 1 how to implement LCM and PPO-Extension using an inverted index. The algorithm takes as input an epoch of data and a support threshold as a ratio  $\alpha$ . It outputs the closed itemsets with support more than the threshold. Along with each itemset in the solution, it also outputs the transactions in which it occurs - which is represented as  $\langle items, |T_{itemset}| \rangle$ . The symbol  $\succ$  denotes

that the lefthand side succeeds the righthand side in the total ordering.

The algorithm also lends itself to distributed implementations easily. For example, a Map/Reduce implementation is easy since the only operations are counting (line 14) and projection (line 22). However, the fast execution time and the low memory requirements of the algorithm makes it possible that a distributed implementation will cause an overhead. In the implementation shown, it is not necessary that the index’s tokens list follow the total ordering; all itemsets of length 1 will be considered anyway.

**Input:**  $\alpha$ : Dynamic support ratio

**Data:**  $E$ : Epoch of data

**Result:**  $C$ : Closed itemsets having support  $\alpha$  within  $E$

```

1  $C \leftarrow \{(\emptyset, E)\}$ ; //  $\emptyset$  is a closed itemset!
2  $X \leftarrow$  Inverted index of  $E$ ;
3 foreach  $w \in X.tokens$  do
4    $T_{\{w\}} \leftarrow X.postingsList[w]$ ;
5   if  $|T_{\{w\}}| \geq \alpha \frac{|E|}{E.span}$  then  $LCM(\{w\}, w, T_{\{w\}})$ ;
6 end
7 return  $C$ ;
8 Function  $LCM(s: Current\ itemset, w_{sh}: Suffix\ head,$ 
9  $T_s: Transactions\ (Tweets)\ containing\ s)$  is
10   frequency[1... $w_n$ ]  $\leftarrow 0$ ;
11   suffix  $\leftarrow \{w_{sh}\}$ ;
12   foreach  $t \in T_s$  do
13     foreach  $w \in d$  do
14       frequency[ $w$ ]++;
15       if frequency[ $w$ ] =  $|T_s|$  then suffix.add( $w$ );
16     end
17   end
18   if  $\exists v \in suffix : w_{sh} \succ v$  then return;
19    $C.add(\langle s \cup suffix, T_s \rangle)$ ;
20   foreach  $v \succ w_{sh}$  and  $v \notin suffix$  do
21     if frequency[ $v$ ]  $\geq \alpha \frac{|E|}{E.span}$  then
22        $D \leftarrow T_s \cap v$ ; // Results of query  $s$  AND  $v$ 
23        $LCM(s \cup suffix \cup \{v\}, v, D)$ 
24     end
25   end
26 end

```

Algorithm 1: LCM Frequent Itemsets Mining

## 5. MINING SOCIAL MEDIA

Throughout this paper we use data we have been collecting from the Twitter public stream<sup>2</sup> as of October 1st, 2012. We use only Tweets written in Latin script to facilitate tokenization using white space and other word boundaries. We collected only the Tweet text to avoid reliance on any features specific to a certain social medium, and make the algorithms applicable to other media where text is short such as comments and Facebook or Google+ status updates. The only preprocessing performed was removing duplicate original Tweets (not retweets) using a Bloom filter. This removes spam Tweets sent by BotNets, such as advertisements about Raspberry Ketone diets.

We apply the algorithms to epochs of data, so they are not strictly stream processing algorithms. However, we regard

<sup>2</sup><https://dev.twitter.com/docs/streaming-apis/streams/public>

Doc. Id	Document	Doc. Id	Document
a	barack & mitt	b	brack obama & mitt romney
c	brack obama & romney	d	brack obama

Documents (two per row)

Step	Current Solution	Possible Extension Items	Comments
1	{barack} (4)*	<b>mitt</b> (2), obama (3), romney (2)	Items are ordered lexicographically
2	{barack, mitt} (2)*	<b>obama</b> (1), romney (1)	Extension items reenumerated & counted
3	{barack, mitt, obama} (1)	romney (1)	Rule 2: ‘romney’ appears in all $D_{itemset}$
4	{barack, mitt, obama, romney} (1)*		Rule 2: ‘obama’ is the <i>suffix head</i>
5	{barack} (4)	mitt (2), <b>obama</b> (3), romney (2)	Nothing more to add, back to ‘barack’
6	{barack, obama} (3)*	<del>mitt</del> (1), <b>romney</b> (2)	Rule 1: skipping ‘mitt’, adding ‘romney’
7	{barack, obama, romney} (2)*	<del>mitt</del> (1)	Rule 1: Nothing more to add.
8	{barack} (4)	mitt (2), obama (3), <b>romney</b> (2)	Back to ‘barack’, adding ‘romney’
9	{barack, romney} (2)	<del>mitt</del> (1), obama (2)	Rule 2: add obama to suffix after ‘romney’
10	{barack, romney, obama} (2)	<del>mitt</del> (1)	Rule 3: suffix isn’t ordered, prune solution

Closed itemsets containing ‘barack’

Table 1: Generation of closed itemsets by Prefix Preserving Closure Extension

the process as mining a sliding window that is moved forward by time steps of short span. The time step must be longer than the time needed to mine an epoch of data, and the performance of our algorithms makes it possible to use a time step of a few seconds for epochs up to a day long. Figure 2 shows the runtime of LCM on epochs of increasing length, and we will show in section 6.1 that our extensions don’t degrade the performance. The times reported in figure 2 are averages across all epochs of the specified length in the months of October, November and December, using a time step that is half the epoch length. The variance is very low and the confidence bands are not shown because they appear as dots on top of the bars.

The support threshold used throughout this paper, unless otherwise specified, is  $\alpha = 0.0002$ . This is determined as follows: We picked a topical term that is known to steadily appear with a rather high frequency, and is talked about in all languages; ‘obama’. The maximum likelihood estimate of the probability of the term ‘obama’ within the whole collection of Tweets is 0.0001. Since the average number of Tweets per hour is 100000, so the term ‘obama’ is expected to appear 10 times per hour on average. Thus, we use a minimum support threshold of 10, which translates into  $\alpha = 0.0002$ .

In the rest of this paper we mine epochs of 1 hour span. The reason behind this choice is an observation that the number of closed itemsets mined from epochs of span 1 hour or more, at the same support threshold, remains that same. This indicates that itemsets mined from shorter epochs of social media text are not included in the results of mining longer epochs. Therefore, the epoch span should be minimized. However, when the epoch span is shorter than an hour the frequency required to surpass the support threshold becomes very low, and number of mined itemsets increases because a lot of noise itemsets are mined.

## 5.1 Mining term N-Grams

A large number of itemsets are language constructs that bear no information, such as “such as”. By treating sequential language constructs, and any other multiword expression, as one item we eliminate a large number of such item-

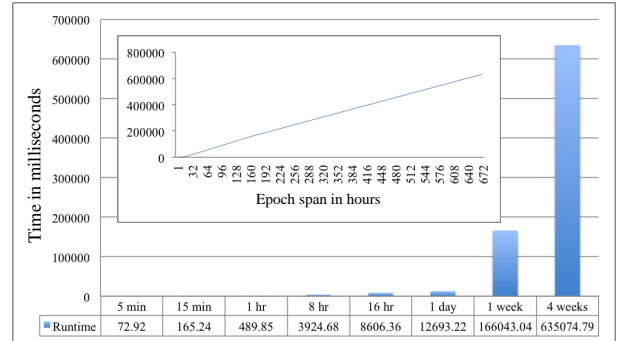


Figure 2: Mean runtime at different epoch spans

sets. Actually, we also eliminate itemsets that are made up of all the different fragments of the language construct along with other items; for example, {we, did, it, #teamobama} can produce 10 other combination of length 2 or more. There are many measures of association that can be used to detect multiword expressions, but each measure is good only under certain conditions and has special properties [15]. We experimented with various measures, and in fact we found out that a very good measure for identifying multiword Named Entities is Yule’s Q; a measure of association and disassociation derived from the odds ratio. However, we have finally found that for the purpose of preprocessing before Frequent Itemsets Mining what works best is tokenizing the documents into term N-Grams with varying N.

We start by tokenizing into unigrams, and counting their frequencies. Then we use the probability of the same word used to determine the support threshold,  $P(\text{‘obama’}) = 0.0001$ , and assume that this is where the head of the Zipfean distribution starts. For each unigram belonging to the head, we create two term bigrams by attaching to it the unigrams before and after it. We repeat this for each  $N \geq 1$  by creating two  $(N+1)$ -Grams for all N-Grams with probabilities above the threshold until there are no more such N-Grams. We do not prevent overlap, because there is no guarantee that the N-Gram created makes any sense. At each N, the probability

threshold is adjusted to account for the increase in the number of tokens and the overall increase in the grand sum of counts, since each high frequency N-Gram is replaced by two lower frequency ones. Let the original probability threshold be  $\eta$ , then the adjusted  $\eta_N$  is:

$$\eta_N = \eta * \frac{\sum_{\{w: w \in W \text{ and } w.length \leq N\}} freq(w)}{\sum_{\{v: v \in W \text{ and } v.length = 1\}} freq(v)} \quad (2)$$

Figure 3 shows the effect of increasing the maximum length of N-Grams from 1 to 5 on the number of tokens, the number of closed itemsets of length more than 1, and the runtime of mining 1 hour epochs of data. The values shown are averages across all 1 hour epochs in the month of November. The value of  $\eta$  used is 0.0001. Figure 3(a) shows that the number of distinct items increases a lot when N moves from 1 to 2, then keeps increasing slightly until it starts decreasing at N=5. The decrease happens because all 4-Grams with probability above the threshold are parts of Tweets from services that use the same text and append a URL, such as Tweets reporting scores from Game Insight<sup>3</sup>. Such Tweets are tokenized into more 4-Grams than 5-Gram, and the 4-Grams appearing in them don't appear elsewhere; thus each two of them are reduced into one 5-Gram. Figure 3(b) shows that the number of itemsets keeps decreasing as expected. Figure 3(c) shows that runtime also decreases as N goes from 1 to 5, since LCM runtime is proportional to the number of closed itemsets, and is not affected by the sparsity of data. The runtimes in this figure are slightly less from those in figure 2 because they don't include the time taken for writing the posting list of each itemset.

After mining term N-Grams we flatten the itemsets to sets of unigrams again. This removes overlap between parts of itemsets making it easier to reason about how they relate to each other. This is also necessary since an itemset will have different N-Gram set representations, and its postings list is the union of those of the different representations.

## 6. STRONGLY CLOSED ITEMSETS

The closed property of an itemset is easily violated by modifying one transaction that contains the itemset and removing one of its items. While an update operation is not supported in the model of frequent itemsets mining, a similar effect happens when people are writing about a certain fine grained topic. For example, on November 6th, 2012 many people were tweeting that “(Barack Obama) was elected (as president of the United States of America)”. If all the Tweets which use the verb phrase “was elected” to report the event also contain the two names (in parenthesis) in full, then there will be one *maximal* itemset with all 12 items. However, either or both the names can be shortened by omitting an underlined part and still convey the meaning. Therefore instead of 1 maximal itemset about the event there will be 8 closed ones. Nevertheless, mining only *maximal* itemsets doesn't solve the problem of redundancy. Consider that it is possible to say that “Obama was elected” as well as “Obama got elected”, resulting in 2 maximal itemsets because of a slight variation in the language. Empirically, this is evident in the relatively small difference between the number of closed itemsets and maximal itemsets of length 2 or more (i.e. excluding itemsets that are merely a frequent

<sup>3</sup><http://www.game-insight.com/>

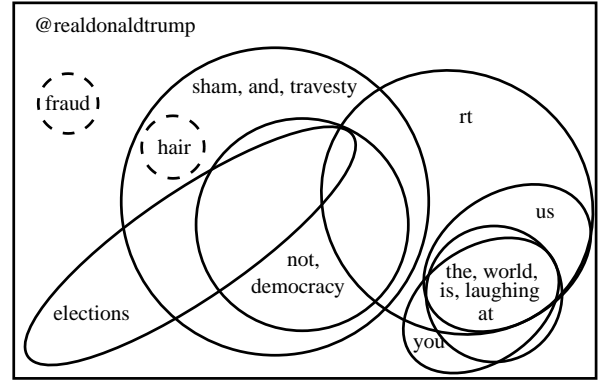


Figure 4: Closed, distinct and strongly closed sets

term); the means across all hour long epochs in our data are 2439.17 closed itemsets and 1831.92 maximal itemsets.

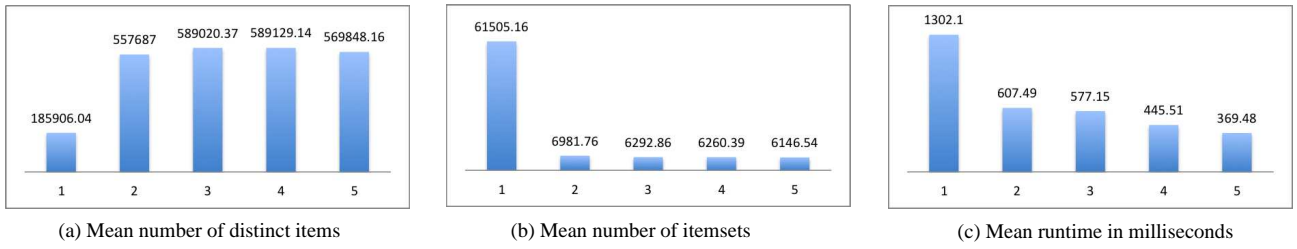
We now define a condition that is not as easily violated as the closed and maximal conditions for selecting itemsets. We define a *distinct* itemset as a closed itemset whose frequency comprises more than a certain proportion of the frequency of its least frequent subset. This condition chooses closed itemsets which substantially violate the closed condition of their subsets. The proportion is a parameter,  $\kappa$ , that controls the selectivity of the *distinctiveness* condition. Formally, the set of *distinct* itemsets,  $\mathcal{D}$ , is defined as follows:

$$\mathcal{D} = \{s : s \in \mathcal{C} \text{ and } \exists s_p \subset s \text{ where } \frac{|T_s|}{|T_{s_p}|} \geq \kappa\} \quad (3)$$

Figure 4 illustrates *distinct* itemsets, in solid lines, and closed itemsets that don't satisfy the distinctiveness condition, in dashed lines. Each area in the figure represents the transactions containing the itemset formed by concatenating the items in all intersecting ellipses. The figure represents itemsets related to Donald Trump's famous Tweets when Obama's victory was clear<sup>4</sup>. The figure also shows the effect of the lack of context and structure in conversations happening on Twitter. Because there was originally no way to refer to a certain Tweet, a Tweet that sparks conversation on Twitter had to be quoted in a retweet along with the retweeter's comment on it. This tradition still continues even though now Tweets can reference each other. Due to the 140 characters length limit of Tweets the quotation usually has to be made as short as possible by selecting only the most discriminative words. This selection is an act of collaborative filtering, but it results in many trivially different subsets from the original Tweet. The additions of retweeters also form many different supersets of the of the original Tweet, some of which represent highly supported opinions.

To remove the redundancy in *distinct* itemsets we merge similar ones into *strongly closed* itemset clusters. The similarity of a *distinct* itemset,  $s_d$ , and another distinct or closed itemset,  $s_c$ , is measured as the overlap of the *transactions* containing both of them with the transactions containing  $s_c$ . A *distinct* itemset is clustered with another itemset if one exists such that the overlap exceeds a similarity threshold, which we take to be  $1 - \kappa$  (the indistinctiveness of  $s_d$  from  $s_c$ ). If more than one satisfies this condition, the *distinct*

<sup>4</sup>[http://www.huffingtonpost.com/2012/11/07/donald-trump-election-revolution\\_n\\_2085864.html](http://www.huffingtonpost.com/2012/11/07/donald-trump-election-revolution_n_2085864.html)



**Figure 3: Effect of the increasing maximum N-Gram length on results of mining of 1hr epochs of data**

itemset is clustered with the one with the highest overlap ratio. When a *distinct* itemset is clustered with an itemset which is already part of a cluster, the *distinct* itemset is added to the existing cluster. Finally, the *strongly closed* itemset is the union of all cluster members, and its support is the size of the union of their postings lists. Formally:

$$S_l = \{w : w \in \bigcup_{s_i \in r_l} s_i \text{ where } r_l \text{ is a cluster}\} \quad (4)$$

To formally define the *strongly closed* itemset clusters, we first introduce the related concept of *association rule confidence*. Confidence is the basic property used for association rules mining, and it is used in the definition of  $\delta$ -free sets [3]. Mining itemsets based on the confidence of rules they induce has long been recognized as a method for finding “interesting patterns” [5], but since this property is not anti-monotone a variation has to be used (for example, *all confidence* [9]). The confidence of an association rule “ $s_j \rightarrow s_i$ ” is defined as:

$$\text{conf}(s_j \rightarrow s_i) = \frac{|T_{s_i} \cap T_{s_j}|}{|T_{s_j}|} \quad (5)$$

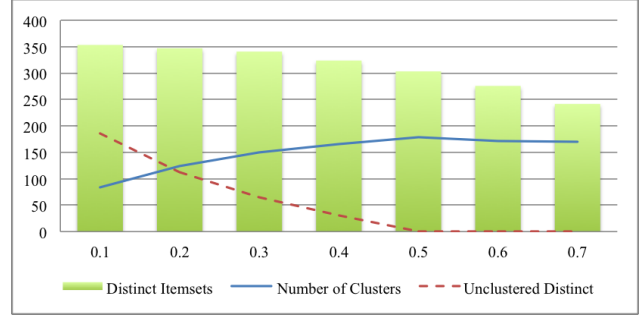
The clustering scheme described selects the cluster that contains the itemset which maximizes the confidence of the rule  $\text{conf}(s_c \rightarrow s_d)$ , with a lower bound on the overlap to maintain distinctiveness. Formally, the clustering result are:

$$\mathcal{R} = \{r : r = \bigcup (s_c, s_d) \text{ where } s_c \in \mathcal{C} \text{ and } s_d \in \mathcal{D} \text{ and } \forall_{(s_c, s_d)} s_c = \text{argmax}_{s_k} \text{conf}(s_k \rightarrow s_d) \text{ and } \text{conf}(s_c \rightarrow s_d) \geq (1 - \kappa)\} \quad (6)$$

This clustering can be implemented efficiently using techniques similar to the ones proposed by Bayardo et al. [2]. The main ideas are to limit the comparisons to a few candidates, and to terminate the comparison early if the similarity threshold will not be met. In our case the postings lists are longer than the itemsets, so we generate candidates for comparison by calculating similarity between itemsets. When calculating the similarity between two postings lists, we can terminate early if the difference exceeds the maximum difference permissible to achieve a similarity of  $1 - \kappa$ , which can be easily derived from equation 5.

Algorithm 2 shows a possible implementation. For each itemset,  $s_i$ , we find the itemsets produced before it and overlapping with it in one or more item. Out of those candidates,  $s_c$ , we find the one that maximizes  $\text{conf}(s_c \rightarrow s_i)$  such that the confidence exceeds  $1 - \kappa$ . Notice that confidence is not a symmetric measure, and we only check the confidence of the rule that the clustering candidate implies the itemset.

Figure 5 show the effect of varying  $\kappa$  on the mean number of *distinct* and *strong closed* itemsets mined from 1 hour long



**Figure 5: Effect of changing  $\kappa$  on mining results**

**Input:**  $\kappa$ : Minimum distinctiveness threshold

**Data:**  $\mathcal{C}$ : Closed Itemsets produced by LCM

**Result:**  $\mathcal{R}$ : Strong closed itemset clusters

```

1 for  $i \leftarrow 2$  to  $|\mathcal{C}|$  do
2    $C \leftarrow \{s_c : s_c \in \mathcal{C} \text{ and } c < i \text{ and } |s_c \cap s_i| > 0\}$ ;
3    $P \leftarrow \{s_p : s_p \in \mathcal{C} \text{ and } p < i \text{ and } s_p \cap s_i = s_i\}$ ;
4    $s_p \leftarrow \text{argmax}_{s_p \in P} \frac{|T_{s_i}|}{|T_{s_p}|}$ ; // Direct parent
5   if  $\frac{|T_{s_i}|}{|T_{s_p}|} < \kappa$  then
6     continue; // Not a distinct itemset
7    $s_m \leftarrow s_i$ ; // Cluster centroid, initially self
8    $\text{maxConf} \leftarrow 0$ ; // Best candidate's score
9   foreach  $s_c \in C$  do
10     $\Delta \leftarrow (1 - (1 - \kappa))|T_{s_c}|$ ; // Maximum difference
11     $\delta \leftarrow \text{difference}(T_{s_c}, T_{s_c \cup s_i}, \Delta)$ ; // Stops early
12    if  $\delta \leq \Delta$  then
13       $\text{conf} \leftarrow \frac{|T_{s_c}| - \delta}{|T_{s_c}|}$ ;
14      if  $\text{conf} > \text{maxConf}$  then
15         $s_m \leftarrow s_c$ ; // Best merge candidate
16         $\text{maxConf} \leftarrow \text{conf}$ ;
17    end
18  end
19 end
20  $\mathcal{R}[s_i] \leftarrow \mathcal{R}[s_m]$ ; // Cluster  $s_i$  with  $s_m$ 
21  $\mathcal{R}[s_m].\text{itemset} \leftarrow \mathcal{R}[s_m].\text{itemset} \cup s_i \cup s_m$ ;
22  $\mathcal{R}[s_m].\text{postingsList} \leftarrow \mathcal{R}[s_m].\text{postingsList} \cup s_i.\text{postingsList} \cup s_m.\text{postingsList}$ ;
23 end
24 return  $\mathcal{R}$ ;

```

**Algorithm 2:** Forming strong closed itemset clusters



epochs. The number of *distinct* itemsets keeps dropping as the distinctiveness threshold increases. On the other hand, the number of *strong closed* clusters formed keeps increasing as the similarity (indistinctiveness) threshold decreases. The dashed line shows that the number of unclustered distinct itemsets reaches zero at  $\kappa = 0.5$ , which explains why the number of clusters changes very slightly after that. We use  $\kappa = 0.25$  for the rest of the paper, which is an arbitrary choice based on the definition not the data. The average number of itemsets (*strongly closed* and unclustered *distinct*) mined from 1 hour epochs at this value of  $\kappa$  is only 224.48 which is about 10% of the number of closed itemsets.

## 6.1 Bounding runtime and memory

We exploit the way PPC-Extension produces itemset to introduce an optimization that improves both runtime, memory requirements and filtering power of the high confidence condition. Unlike the original LCM algorithm, filtering low confidence itemsets requires keeping mining results in memory to calculate the confidence of newly generated ones. Instead of keeping all previous itemsets in memory we keep only a limited number,  $b$  (for buffer size). This obviously improves runtime and memory requirements, and it can also improve the choice of itemsets. If the total ordering of items follows the descending order of items' frequencies, then the call tree of recursive LCM calls will be very similar to the FP-tree. Given the length constraint on Tweets, then the FP-Tree can only be shallow and wide. Therefore, an itemset that has a large number of supersets cannot be appearing with only a limited set of words (topic), but it has to be covering a breadth of possible descendants. Thus this itemset is very likely to be a language construct. Even if it is a language construct that appears with high confidence with some topical words, this high confidence doesn't carry interesting information; for example, {the, of} that appears within {the, united, states, of, america}. Even if it is a topical word with very high frequency, then interesting information will be divided into several topics given its large number of direct descendants. In either case, filtering closed itemsets that have high confidence with such a ubiquitous itemset will not omit any information. If the filtered itemset is part of an informative one then it will be produced when the longer itemset is found to have high confidence with the filtered one.

We use a buffer size  $b = 1000$  in all experiments, which leads to a runtime of about 1 to 3 milliseconds per itemset for filtering out low confidence itemsets and creating strong closed clusters. The number of selected itemsets is also reduced by almost half on average. The difference between the two selections is made up mostly of stop words. A few itemsets including topical words are also filtered, such as the closed itemset {president, Obama} since Obama is an example of a ubiquitous topical word. Such itemsets lack any information on their own and is included in other informative ones. Finally, this enhancement is not necessary for satisfactory performance. With an unlimited buffer, the average runtime is 5.2 seconds for filtering 4265 itemsets (the average number of closed itemsets from an hour long epoch).

Figure 6 shows the total runtime of the LCM algorithm plus the strong closed itemsets clustering (LCM + Clustering) at different epoch spans. The runtime of LCM alone is also plotted for reference. We also plot the performance of another Frequent Itemset Mining algorithm, FP-Zhu [6],

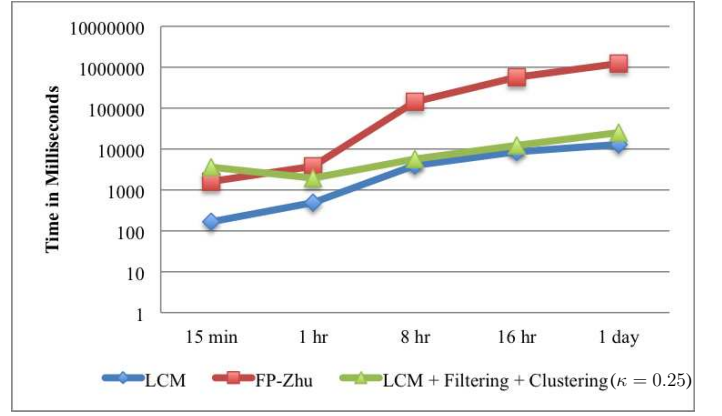


Figure 6: Runtime of itemset mining alone and with filtering and clustering at different epoch spans

which was the runner up at FIMI '04 [8]. We include it to show that our extensions do not degrade the performance of LCM even in the context of competitions. The Y-Axis is in logarithmic scale to keep the scale of the plot suitable for seeing slight differences. The output of LCM is the input to the clustering algorithm, so it is affected by the number of closed itemsets produced. This explains why it takes slightly longer time for clustering results from the 15 minutes epoch and then takes a constant time for epochs of longer span.

## 7. TEMPORAL RANKING

For ranking itemsets according to the importance of each itemset in within the mined epoch, we use itemsets' point-wise KL Divergence from a background model. The point-wise KLD of the probability of an itemset  $s_i$  in the background model  $Q$  from its probability in the current epoch  $P$  can be considered the information gain, IG:

$$\begin{aligned}
 IG(P(s_i), Q(s_i)) &= -(H(P(s_i)) - H(P(s_i), Q(s_i))) \\
 &= \sum P(s_i) \log P(s_i) - \sum P(s_i) \log Q(s_i) = KLD(P(s_i) || Q(s_i))
 \end{aligned} \tag{7}$$

To generalized this to the case of strong closed itemset clusters, we have to take the overlap between the itemsets of the cluster into account when calculating the information gain of the whole cluster. We note that the joint probability of the itemset,  $s_i$ , and its superset,  $s_j$ , is equal to the probability of the superset;  $P(s_i, s_j) = P(s_j)$  and  $Q(s_i, s_j) = Q(s_j)$ . Thus, the information gain of the appearance of an itemset and its superset together is essentially the information gain of the superset;  $IG(P(s_i, s_j), Q(s_i, s_j)) = IG(P(s_j), Q(s_j))$ . Also, their Pointwise Mutual Information is the Self Information of the subset:

$$PMI(s_i, s_j) = \log \frac{P(s_i, s_j)}{P(s_i)P(s_j)} = -\log P(s_i) = I(s_i) \tag{8}$$

Therefore, the information gain of a superset is different from the information gain of its subset by the information gained (or lost) because of the additional items. Thus, the information gain of the appearance of all itemsets in a strong closed itemset cluster,  $S = \{s_1, \dots, s_m\}$ , can be calculated as the information gain of its smallest subset,  $s_{min}$ , plus the differences between the KLDs of member subsets and



the smallest subset. We use the squares of the differences, because the pointwise KLD value can be positive or negative and we only care about the magnitude of the difference. To give the self information of the smallest subset the same influence, we also square it. Thus the information gain of a strong closed itemset is given by:

$$IG^2(P(s_{i1}, \dots, s_{im}), Q(s_{i1}, \dots, s_{im})) = I^2(s_{min}) + \sum_{j=i1..im} (KLD(P(s_j)||Q(s_j)) - KLD(P(s_{min})||Q(s_{min})))^2$$

The formula above can be used directly for ranking clusters, but it will favour largers one. We normalize by the size of the cluster giving our final ranking formula for strong closed itemsets:

$$\overline{IG}(S) = \frac{IG^2(P(s_{i1}, \dots, s_{im}), Q(s_{i1}, \dots, s_{im}))}{m} \quad (9)$$

## 7.1 Empirical Evaluation

For the lack of ground truth data, the performance of the ranking (and the quality of the filtering and clustering on which it is based) is shown by example. The background model we use is the results of mining the previous month, using a *minimum support* value of 1 occurrence (in the least volume hour). Mining the background model at such a low support increases the number of produced itemsets, which is desirable for a background model. All probabilities are smoothed by add-one smoothing.

Tables 2 and 3 show the top 3 itemsets for 1 hour epochs in the days of November 6th and 9th. The itemsets shown are the first appearances of the most interesting itemsets; that is, an hour is shown only if its top 3 feature novel interesting itemsets. The first column is the beginning of the hour, EST time. The second column is the top itemsets. The third column is a commentary to explain the itemsets, but we omit it from table 3 to save space since the itemsets are self explanatory.

On November 6<sup>th</sup>, the US presidential elections took place and we can see how its events unwind from “get out and vote” to the projections and debates, all the way to the “acceptance speech”. Early in the day, itemsets about UEFA Champions football matches and a TV show appear in the top 3 along with itemsets about the still uneventful elections. Actually, the matches keep occupying top positions and timely updates of their scores appear in the top 30 itemsets, until they end and the elections heats up. Short after the results of the elections became clear, news that “weed is legal in Colorado” occupies the top position. This exemplifies the power of social media as a collaborative filter, picking the important news even if it is unexpected. The user centric definition of importance is eminent in the great attention given to the “lady behind Obama with a flag in her hair” during the acceptance speech.

On November 9<sup>th</sup>, no major events were happening but many overlapping minor ones happened. The day starts by news about the end of two careers; the Laker’s “coach Mike Brown” got fired and “CIA director David Petraeus resigns”. The relationship of Justin Bieber also ends as he “broke, up, with, Selena, Gomez” at 22:00. This overlaps with his participation in the MTV Europe Music Awards (MTVEMA), and both topics occupied high rankings. MTVEMA is particularly interesting because votes were solicited from audience through Twitter. This is an example of a topic where people

have different opinions, as the hour 15:00 shows. The top 3 itemsets of this hour are supporting “Katy Perry”, “Justin Bieber” and “Lady Gaga” respectively. They are all reported as separate itemsets, showing how clustering using the postings lists avoid forming incohesive clusters. By the end of the day many congratulations for the Indonesian Hero’s day (“Hari Pahlawan”) appear, and the Turkish commemoration day of Atatürk is also mentioned as the 10th of November starts in these countries. These are examples of itemsets from languages with a relatively low number of users, showing how the absolute popularity of a topic doesn’t affect its rank. If itemsets from only a specific language is desired, language identification can be applied on the itemsets and Tweets from their postings lists. Moving language identification downstream avoids affecting the results of mining because of error in an upstream component.

## 8. CONCLUSION AND FUTURE WORK

We have proposed a method for efficiently creating temporal synopses of social media streams, based on a frequent itemset mining algorithm that is suitable for sparse data, LCM. Our method summarizes an hour long of Twitter data (100,000 Tweets on average) into 224 itemsets in 1945.68 milliseconds on average, and scales well for longer epochs of data. The direct application of LCM on 1 hour long epochs of Twitter data results in 61505.16 closed itemsets and takes 2506.58 milliseconds on average. The improvement is due to using variable length N-grams to mitigate the effect of the skewness of the frequency distribution of unigrams, and strengthening the closure condition such that it selects only itemsets that are distinctively different from its subsets and other itemsets. The distinctiveness between two itemsets is judged based on a parameter  $\kappa$ , which represents the minimum ratio of transactions containing only one of them to the those containing both of them.

We also propose a method for ranking itemsets based on their novelty when compared to mining results from a longer time period. The top 3 itemsets from hours in the elections day and another less eventful day show that the synopses capture important events, and can actually be directly presented to users. We aim to exploit the synopses for temporal query expansion in our future work. We also wish to explore ways to make use of the temporal signal during mining, such as when calculating similarity during clustering.

## 9. ACKNOWLEDGEMENTS

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13:00	0, 1, de, jong	De Jong scores for Ajax
	geordie, shore	Season 5 of the TV series starts
	get, out, and, vote	Still early in the U.S. elections day
17:00	if, obama, wins	Speculations regarding elections
	USERNAME, spots, my, club	Pyramidal marketing scam. Retweeted by people to make money.
	the, polls, close	Polls to start closing at 6 PM
19:00	A partir de que idade você considera alguém velho?	Internet meme from Brazil, discussing when to start considering a person old.
	food, stamps	Discussions pertaining to elections
	linda, mcmahon, senate	Linda McMahon loses CT senate race
20:00	obama, got, this	Announcing states that Obama got projected, winner
	projected, winner	Early projections about who will win
	moving, to, canada	Reaction to projections
21:00	elizabeth, warren	MA senate elections winner
	popular, vote	Comparing Popular vs Electoral votes
	who, is, the, president?	Anticipation for the elections results
22:00	#forward, #obama2012	Obama won, time to move #forward
	my, president, is, still	Some were saying Black (skin colour), others Blue (party colour)
	back, in, office	Obama is back in office
22:30	once, you, go, black	A popular cultural reference.
	@realdonaldtrump, this, elections, ... [his famous Tweet]	Donald Trump is a Republican and he couldn't accept Obama's victory
	concession, speech, write	The losing candidate has to concede before the winner declares victory
23:00	weed, is, legal, in, colorado	This piece of news got the attention as soon as it was announced
	karl, rove	Fox news challenges results from Ohio
	our, country, is, trouble	Dramatic reactions to elections
23:30	acceptance, speech, wrote	Obama wrote his speech,
	give, his, speech	and is going to deliver it
	on, cnn	on CNN
00:30	behind, obama, with, flag, in, her [hair]	During the speech, attention on social media goes to a woman appearing behind Obama on TV!
	the, best, is, yet, to, come	Quote from Obama's speech
	flag, that, weaves	Play on words about the "flag in hair"

Table 2: Top 3 itemsets for hours in November 6th

12:00	#iwillneverunderstand, why
	breaking, news, head, coach, mike, brown, have, fired
	USERNAME, spots, available, my, club
14:00	cia, director, david, petraeus, resigns
	você, acha, que
	#tvoh [the voice of Holland], babette
15:00	#emawinkaty, i, think, katy, perry, will, be, the, big,
	#mtvema, winner, tweet, your, pick, at, URL
	#emawinbieber, i, think, justin, bieber, ... [same as above]
22:00	#emawingaga, ... [same as above]
	justin, bieber, and, selenagomez, broke, up
	selamat, hari, pahlawan
	aniyorus, kemal, mustafa [ataturk]

Table 3: Top 3 itemsets for hours in November 9th

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