

## 1 System Theory

### 1.1 Nonlinear Systems

$$\begin{aligned}\dot{x} &= g(x, u) & y &= h(x, u) \\ \dot{x}_s &= g(x_s, u_s) = 0 & y_s &= h(x_s, u_s) \\ A^c &= \left. \frac{\partial g}{\partial x^T} \right|_{x=x_s, u=u_s} & B^c &= \left. \frac{\partial g}{\partial u^T} \right|_{x=x_s, u=u_s} \\ C^c &= \left. \frac{\partial h}{\partial x^T} \right|_{x=x_s, u=u_s} & D^c &= \left. \frac{\partial h}{\partial u^T} \right|_{x=x_s, u=u_s}\end{aligned}$$

### 1.2 Linear Systems

#### Continuous

$$\begin{aligned}\dot{x}(t) &= A^c x(t) + B^c u(t) \\ x(t) &= e^{A^c(t-t_0)} x_0 + \int_{t_0}^t e^{A^c(t-\tau)} B^c u(\tau) d\tau \\ e^{A^c t} &= \sum_{n=0}^{\infty} \frac{(A^c t)^n}{n!}\end{aligned}$$

#### Discrete

$$\begin{aligned}x_{k+1} &= A x_k + B u_k \\ y_k &= C x_k + D u_k \\ x_{k+N} &= A^N x_k + \sum_{i=0}^{N-1} A^i B u_{k+N-1-i}\end{aligned}$$

**Forward Euler**  $A = I + T_s A^c, B = T_s B^c, C = C^c, D = D^c$

$$\begin{aligned}x_{k+1} &= x_k + T_s g^c(x_k, u_k) = g(x_k, u_k) \\ y_k &= h^c(x_k, u_k) = h(x_k, u_k)\end{aligned}$$

**Exact discretization** (assume constant  $u(t)$  during  $T_s$ )

$$\begin{aligned}A &= e^{A^c T_s}, B = \int_0^{T_s} e^{A^c(T_s-\tau')} B^c d\tau \\ B &= (A^c)^{-1} (A - I) B^c, \text{ if } A^c \text{ invertible}\end{aligned}$$

### 1.3 Lyapunov Stability

System is stable in the sense of Lyapunov iff it stays in any arbitrarily small neighborhood of the origin when it is disturbed slightly.

**Lyapunov stable** iff  $\forall \epsilon > 0 \exists \delta(\epsilon)$  s.t.  $\|x_0\| < \delta(\epsilon) \Rightarrow \|x_k\| < \epsilon, \forall k \geq 0$

**asymptotically stable** in  $\Omega \subseteq \mathbb{R}^n$  if Lyapunov stable and *attractive*  $\lim_{k \rightarrow \infty} x_k = 0, \forall x_0 \in \Omega$ .

**Lyapunov Function**  $V: \mathbb{R}^n \rightarrow \mathbb{R}$  continuous at the origin, finite  $\forall x \in \Omega$ ,  $V(0) = 0$  and  $V(x) > 0, \forall x \in \Omega \setminus \{0\}$

$$V(g(x)) - V(x) \leq -\alpha(x), \forall x \in \Omega \setminus \{0\}$$

where  $\alpha: \mathbb{R}^n \rightarrow \mathbb{R}$  is continuous positive definite, equilibrium at  $x = 0$  and  $\Omega \subset \mathbb{R}^n$  closed and bounded set containing the origin.

**Lyapunov Theorem** If system admits Lyapunov function  $V(x)$ , then  $x = 0$  is **asymptotically stable** in  $\Omega$  (sufficient but not necessary).

If additionally  $\|x\| \rightarrow \infty \Rightarrow V(x) \rightarrow \infty$  **globally asymptotically stable**. To check if  $V(x) = x^T P x$  is valid Lyapunov function of system

$x_{k+1} = A x_k$  check if  $(A^T P A - P)$  has neg. eigen values.

In other words: Iff eigenvalues of  $A$  inside unit circle (i.e. stable) then  $\exists$  unique  $P > 0$  that solves  $A_{cl}^T P A_{cl} - P = -Q, Q > 0$  and  $V(x) = x^T P x$  is a lyapunov function.

### 1.4 Observability $\Rightarrow$ Detectability, Controllability $\Rightarrow$ Stabilizability

**(A, C) observable** if  $\text{rank}(O) = n$  (full col. rank) for

$$O = \begin{bmatrix} C^T \\ (CA)^T \\ \vdots \\ (CA^{n-1})^T \end{bmatrix} \text{ or } \text{rank} \begin{bmatrix} A - \lambda I \\ C \end{bmatrix} = n \quad \forall \lambda_i \text{ of } A \text{ (PBH-test).}$$

**(A, C) detectable** iff  $\text{rank} \begin{bmatrix} A - \lambda I \\ C \end{bmatrix} = n \quad \forall \text{unstable } |\lambda_i| \geq 1 \text{ of } A$ .

**(A, B) controllable** if  $\text{rank } C = n, C = [B \ AB \ \dots \ A^{n-1}B]$  or if  $\text{rank}([\lambda_j I - A \ B]) = n \quad \forall \lambda_i \text{ of } A$  (PBH-test).  
Intuition: Can reach any state in (at most)  $n$  steps.

**(A, B) stabilizable** if  $\text{rank}[\lambda_j I - A \ B] = n \quad \forall \text{unstable } |\lambda_i| \geq 1 \text{ of } A$ .  
Intuition: Can reach origin in (at most)  $n$  steps.

## 2 Unconstrained Control

### 2.1 Block Approach (used also for $\bar{w}$ substitution)

$$\begin{bmatrix} x_0 \\ x_1 \\ \vdots \\ x_N \end{bmatrix} = \begin{bmatrix} I \\ A \\ \vdots \\ A^N \end{bmatrix} x(0) + \begin{bmatrix} 0 & 0 & \dots & 0 \\ B & 0 & \dots & 0 \\ AB & B & \dots & 0 \\ \vdots & \vdots & \ddots & \vdots \\ A^{N-1}B & \dots & AB & B \end{bmatrix} \begin{bmatrix} u_0 \\ u_1 \\ \vdots \\ u_{N-1} \end{bmatrix}$$

$$\begin{aligned}x &= S^x \cdot x(0) + S^u \cdot u & \text{size}(S^x) &= [n_{\text{states}} \cdot (N+1), N] \\ & & \text{size}(S^u) &= [n_{\text{states}} \cdot (N+1), n_{\text{states}}]\end{aligned}$$

$$\bar{Q} = \text{diag}(Q, \dots, Q, P) \quad \text{size}(\bar{Q}) = [n_{\text{states}} \cdot (N+1), n_{\text{states}} \cdot (N+1)]$$

$$\bar{R} = \text{diag}(R, \dots, R) \quad \text{size}(\bar{R}) = [n_{\text{input}} \cdot N, n_{\text{input}} \cdot N]$$

$$H = S^{uT} \bar{Q} S^u + R \quad F = S^{xT} \bar{Q} S^x$$

$$Y = S^{xT} \bar{Q} S^x$$

#### Optimal cost and control

$$J^*(x_0) = -x_0^T F H F^T x_0 + x_0^T Y x_0$$

$$u^*(x_0) = -H^{-1} F^T x_0 = -\left(S^{uT} \bar{Q} S^u + R\right)^{-1} S^{uT} \bar{Q} S^x x_0$$

### 2.2 Recursive Approach

$$J_k^*(x_k) = \min_{u_k} I(x_k, u_k) + J_{k+1}(x_{k+1})$$

Is a feedback controller as opposed to the Batch Approach. For LQR solve via Riccati Difference Equation (RDE).

$$F_k = -(B^T P_{k+1} B + R)^{-1} B^T P_{k+1} A$$

$$P_k = A^T P_{k+1} A + Q - A^T P_{k+1} B (B^T P_{k+1} B + R)^{-1} B^T P_{k+1} A$$

$$u_k^* = F_k x_k \quad J_k^*(x_k) = x_k^T P_k x_k \quad P_N = P$$

For unconstrained Infinite Horizon Problem, substituting  $P_\infty = P_k = P_{k+1}$  into RDE gives DARE. Uniquely solvable, iff  $(A, B)$  stabilizable and  $(A, G)$  detectable, where  $G G^T = Q$ . Follows from closed-loop system  $x_{k+1} = (A + B F_k) x_k$

## 3 (Convex) Optimization

**General Problem**  $\min_{x \in \text{dom}(f)} f(x)$  s. t.  $g_i(x) \leq 0$  and  $h_j(x) = 0$ .

**Norm**  $f(x): \mathbb{R}^n \rightarrow \mathbb{R}$

$$f(x) = 0 \Rightarrow x = 0, \quad f(x) \geq 0$$

$$f(\alpha \cdot x) = |\alpha| \cdot f(x) \quad \text{for scalar } \alpha$$

$$f(x+y) \leq f(x) + f(y) \quad \forall x, y \in \mathbb{R}^n$$

### 3.1 Convexity

**Convex set**  $\mathcal{X}$  iff  $\forall \lambda \in [0, 1] \forall x, y \in \mathcal{X} \quad \lambda x + (1 - \lambda)y \in \mathcal{X}$ . Intersection preserves convexity, union does not.

**Affine set**  $\mathcal{X} = \{x \in \mathbb{R}^n | Ax = b\}$  for some  $A, b$

**Subspace** is affine set through origin, i.e.  $b = 0$ , aka Nullspace of  $A$ .

**Hyperplane**  $\mathcal{X} = \{x \in \mathbb{R}^n | a^T x = b\}$  for some  $a, b$ .

**Halfspace**  $\mathcal{X} = \{x \in \mathbb{R}^n | a^T x \leq b\}$  for some  $a, b$ .

**Polyhedron**  $\mathcal{P} = \{x | a_i^T x \leq b_i, i = 1, \dots, n\} = \{x | Ax \leq b\}$

**Cone**  $\mathcal{X}$  if for all  $x \in \mathcal{X}$ , and for all  $\theta > 0, \theta x \in \mathcal{X}$ .

**Ellipsoid**  $\mathcal{E} = \{x | (x - x_c)^T A^{-1} (x - x_c) \leq 1\}$ ,  $x_c$  center point.

**Convex function**  $f: \text{dom}(f) \rightarrow \mathbb{R}$  is convex iff  $\text{dom}(f)$  is convex and  $f(\lambda x + (1 - \lambda)y) \leq \lambda f(x) + (1 - \lambda)f(y), \forall \lambda \in (0, 1), \forall x, y \in \text{dom}(f)$ .

**Norm ball** is convex (for all norms).

**Epigraph set**  $f: \text{dom}(f) \rightarrow \mathbb{R}$  is the set

$$\text{epi}(f) := \left\{ \begin{bmatrix} x \\ t \end{bmatrix} \mid x \in \text{dom}(f), f(x) \leq t \right\} \subseteq \text{dom}(f) \times \mathbb{R}$$

**Level set**  $L_a$  of a function  $f$  for value  $a$  is the set of all  $x \in \text{dom}(f)$  for which  $f(x) = a$ :  $L_a = \{x | x \in \text{dom}(f), f(x) = a\}$ .

**Sublevel set**  $C_a$  is defined by  $C_a = \{x | x \in \text{dom}(f), f(x) \leq a\}$ .

### 3.2 Linear Programming (LP)

**Problem statement**  $\min c^T x$  such that  $Gx \leq h$  and  $Ax = b$ .

$$\begin{aligned}\text{Norm } l_\infty \quad \min_{x, t} \|x\|_\infty &= \min_{x \in \mathbb{R}^n} [\max\{x_1, \dots, x_n, -x_1, \dots, -x_n\}]: \\ &\text{subject to } x_i \leq t, -x_i \leq t, & Fx &\leq g \\ \Leftrightarrow \min_{x, t} t &\text{subject to } -1t \leq x \leq 1t, & Fx &\leq g.\end{aligned}$$

$$\begin{aligned}\text{Norm } l_1 \quad \min_{x, t} \|x\|_1 &= \min_x [\sum_{i=1}^m \max\{x_i, -x_i\}]: \\ &\text{subject to } x_i \leq t_i, -x_i \leq t_i, & Fx &\leq g \\ \Leftrightarrow \min_t 1^T t &\text{subject to } -t \leq x \leq t, & Fx &\leq g.\end{aligned}$$

Note that for  $\dim x = 1$ ,  $l_1$  and  $l_\infty$  are the same. Note also that  $t$  is scalar for norm  $l_\infty$  and a vector in norm  $l_1$ .

#### Piecewise Affine

$$\begin{aligned}\min_x \left[ \max_{i=1, \dots, m} \{c_i^T x + d_i\} \right] &\text{ s.t. } Gx \leq h \\ \Leftrightarrow \min_{x, t} t &\text{ s.t. } c_i^T x + d_i \leq t, Gx \leq h\end{aligned}$$

### 3.3 Duality

#### Lagrangian Dual Function

$$L(x, \lambda, \nu) = f(x) + \sum_{i=1}^m \lambda_i g_i(x) + \sum_{i=1}^p \nu_i h_i(x)$$

$$d(\lambda, \nu) = \inf_{x \in \mathcal{X}} L(x, \lambda, \nu) \quad \text{i.e. } \nabla_x L(x, \lambda, \nu) = 0$$

**Dual Problem (always convex)**  $\max_{\lambda, \nu} d(\lambda, \nu)$  s. t.  $\lambda \geq 0$ .

Optimal value is lower bound for primal:  $d^* \leq p^*$ .

If primal convex, **Slater condition** (strict feasibility) implies *strong duality*:  $\{x | Ax = b, f_i(x) < 0, \} \neq \emptyset \Rightarrow d^* = p^*$

**Karush-Kuhn-Tucker (KKT) Conditions** are necessary for optimality (and sufficient if primal convex).

$$\begin{array}{lll} \text{Primal Feasibility} & f_i(x^*) \leq 0 & i = 1, \dots, m \\ & h_i(x^*) = 0 & i = 1, \dots, p \\ \text{Dual Feasibility} & \lambda^* \geq 0 & \\ \text{Complementary slackness} & \lambda_i^* \cdot f_i(x^*) = 0 & i = 1, \dots, m \\ \text{Stationarity} & \nabla_x L(x^*, \lambda^*, \nu^*) = 0 & \end{array}$$

## Dual of LP

$$\begin{aligned} \min_x c^T x \quad \text{s.t. } Ax = b, Cx \leq e \\ \Leftrightarrow \max_{\lambda, \nu} -b^T \nu - e^T \lambda \quad \text{s.t. } A^T \nu + C^T \lambda + c = 0, \lambda \geq 0 \end{aligned}$$

## Dual of QP

$$\begin{aligned} \min_x \frac{1}{2} x^T Q x + c^T x \quad \text{s.t. } Cx \leq e \\ \Leftrightarrow \min_{\lambda, \nu} \frac{1}{2} \lambda^T C Q^{-1} C^T \lambda + (C Q^{-1} c + e)^T \lambda + \frac{1}{2} c^T Q^{-1} c \\ \text{s.t. } Qx + \nu + c^T \lambda = 0, \lambda \geq 0 \end{aligned}$$

## 4 Constrained Finite Time Optimal Control (CFTOC)

### 4.1 MPC with linear cost

$$J(x_0, u) = \|P x_N\|_p + \sum_{i=0}^{N-1} \|Q x_i\|_p + \|R u_i\|_p.$$

The CFTOC problem can be formulated as an  $\infty$ -norm LP problem as shown below.

$$\begin{aligned} \min_z \quad & \epsilon_0^x + \dots + \epsilon_N^x + \epsilon_0^u + \dots + \epsilon_{N-1}^u \\ \text{s.t.} \quad & -\mathbf{1}_n \epsilon_i^x \leq \pm Q \left[ A^i x_0 + \sum_{j=0}^{i-1} A^j B u_{i-1-j} \right] \\ & -\mathbf{1}_r \epsilon_N^x \leq \pm P \left[ A^N x_0 + \sum_{j=0}^{N-1} A^j B u_{N-1-j} \right] \\ & -\mathbf{1}_m \epsilon_N^u \leq \pm R u_i \\ & x_i = A^i x_0 + \sum_{j=0}^{i-1} A^j B u_{i-1-j} \in \mathcal{X} \\ & x_N = A^N x_0 + \sum_{j=0}^{N-1} A^j B u_{N-1-j} \in \mathcal{X} \end{aligned}$$

$$u_i \in \mathcal{U}$$

Converting to LP form:

$$\begin{aligned} \min_z \quad & c^T z \\ \text{s.t.} \quad & \bar{G} z \leq \bar{w} + \bar{s} x_k \\ & z = [\epsilon_0^x \dots \epsilon_N^x \quad \epsilon_0^u \dots \epsilon_{N-1}^u \quad u_0^T \dots u_{N-1}^T] \\ & c = [1 \dots 1 \quad 1 \dots 1 \quad 0 \dots 0] \\ & \bar{G} = \begin{bmatrix} G_\epsilon & G_u \\ 0 & G \end{bmatrix}, \quad \bar{w} = \begin{bmatrix} w_\epsilon \\ w \end{bmatrix}, \quad \bar{s} = \begin{bmatrix} s_\epsilon \\ s \end{bmatrix} \end{aligned}$$

Where  $G$  is the normal problem constraints and  $[G_\epsilon G_u]$  form the constraints of the newly introduced variable  $\epsilon$  as given in the first 3

constraints in the section above. For example, we require:

$$\begin{aligned} -\epsilon_i^u &\leq u_i \leq \epsilon_i^u \\ -\epsilon_0^x &\leq A x_0 + B u_0 \leq \epsilon_0^x \\ -\epsilon_1^x &\leq A^2 x_0 + B u_1 + A B u_0 \leq \epsilon_1^x \end{aligned}$$

### 4.2 QP with substitution (see also Batch approach)

$$J^*(x_k) = \min_u \begin{bmatrix} u^T & x_k^T \end{bmatrix} \begin{bmatrix} H & F^T \\ F & Y \end{bmatrix} \begin{bmatrix} u \\ x_k \end{bmatrix}$$

$$\text{s. t. } G u \leq w + E x_k$$

Latter gives three sets (same for without substitution)

$$\mathcal{X} = \{x | A_x x \leq b_x\}$$

$$\mathcal{U} = \{u | A_u u \leq b_u\}$$

$$\mathcal{X}_f = \{x | A_f x \leq b_f\}$$

State equations are in cost matrix, usually in the form:

$$A_x = \begin{bmatrix} 1 \\ -1 \end{bmatrix}, b_x = \begin{bmatrix} b_{\max} \\ -b_{\min} \end{bmatrix}$$

$$G = \begin{bmatrix} A_u & 0 & \dots & 0 \\ 0 & A_u & \dots & 0 \\ \vdots & \vdots & \ddots & \vdots \\ 0 & 0 & \dots & A_u \\ A_x B & 0 & \dots & 0 \\ A_x A B & A_x B & \dots & 0 \\ \vdots & \vdots & \ddots & \vdots \\ A_x A^{N-2} B & A_x A^{N-3} B & \dots & 0 \\ A_f A^{N-1} B & A_f A^{N-2} B & \dots & A_f B \end{bmatrix} \quad E = \begin{bmatrix} 0 \\ 0 \\ \vdots \\ 0 \\ -A_x \\ -A_x A \\ -A_x A^2 \\ \vdots \\ -A_x A^{N-1} \\ -A_f A^N \end{bmatrix} \quad w = \begin{bmatrix} b_u \\ b_u \\ \vdots \\ b_u \\ b_x \\ b_x \\ b_x \\ \vdots \\ b_x \\ b_f \end{bmatrix}$$

### 4.3 QP without substitution

State equations represented in equality constraints ( $k$  fixed, usually  $k = 0$ ).

$$J^*(x_k) = \min_z \begin{bmatrix} z^T & x_k^T \end{bmatrix} \begin{bmatrix} \bar{H} & 0 \\ 0 & Q \end{bmatrix} \begin{bmatrix} z \\ x_k \end{bmatrix}$$

$$\text{s.t. } G z \leq w + E x_k$$

$$G_{\text{eq}} z = E_{\text{eq}} x_k, \quad \text{system dynamics}$$

$$\bar{H} = \text{diag}(Q, \dots, Q, P, R, \dots, R)$$

$$z = \begin{bmatrix} x_1 \\ x_N \\ u_0 \\ u_{N-1} \end{bmatrix} \quad G_{\text{eq}} = \begin{bmatrix} I & I & -B & -B \\ -A & I & -A & I \\ & & & -B \end{bmatrix} \quad E_{\text{eq}} = \begin{bmatrix} A \\ 0 \\ \vdots \\ 0 \end{bmatrix}$$

$$w = \begin{bmatrix} b_x \\ \vdots \\ b_f \\ b_u \end{bmatrix} \quad G = \begin{bmatrix} 0 & A_x & & \\ & & A_x & \\ & & & A_d \\ & & A_d & A_d \end{bmatrix} \quad E = \begin{bmatrix} -A_x^T \\ 0 \\ \vdots \\ 0 \end{bmatrix}$$

### 4.4 Invariance

**Pos. Invariant set**  $O$  iff  $x_k \in O \Rightarrow x_{k+1} = g(x_k) \in O \forall k$ .

**Max. Pos. Invariant set**  $O_\infty \subset \mathcal{X}$  iff  $0 \in O_\infty$ ,  $O_\infty$  invariant and contains all invariant sets  $O$  with  $0 \in O$ .

**Pre-Set**  $\text{pre}(S) := \{x | g(x) \in S\} = \{x | A x \in S\}$

Linear systems:  $S = \{x | F x \leq f\} \Rightarrow \text{pre}(S) = \{x | F A x \leq f\}$ .

Note:  $O$  invariant  $\Leftrightarrow O \subseteq \text{pre}(O) \Leftrightarrow \text{pre}(O) \cap O = O$ .

**Calculate max. invariant set** by  $\Omega_{i+1} \leftarrow \text{pre}(\Omega_i) \cap \Omega_i$ , terminating when  $\Omega_{i+1} = \Omega_i$ , starting with  $\Omega_0 = \mathcal{X}$ .

## 4.5 Stability and Feasability

**Main Idea** Choose  $\mathcal{X}_f$  and  $P$  to mimic infinite horizon. LQR control law  $\kappa(x) = F_\infty x$  from solving DARE. Set terminal weight  $P = P_\infty$ , terminal set  $\mathcal{X}_f$  as maximal invariant set:

$$\begin{aligned} x_{k+1} = A x_k + B F_\infty x_k \in \mathcal{X}_f \quad \forall x_k \in \mathcal{X}_f \quad \text{terminal set invariant} \\ \mathcal{X}_f \subseteq \mathcal{X}, \quad F_\infty x_k \in \mathcal{U} \quad \forall x_k \in \mathcal{X}_f \quad \text{constraint satisfied} \end{aligned}$$

### Result

1. Positive stage cost function,
2. invariant terminal set by construction and
3. Terminal cost is Lyapunov function with  $x_{k+1}^T P x_{k+1} - x_k^T P x_k = -x_k^T (Q + F_\infty^T R F_\infty) x_k$

Extension to non-linear (time-invariant) MPC possible since terminal set and cost do not rely on linearity.

## 5 Practical Issues

### 5.1 MPC for tracking

Target steady-state conditions  $x_s = A x_s + B u_s$  and  $y_s = C x_s = r$  and constraints give:

$$\min_{x_s, u_s} u_s^T R u_s \quad \text{subj. to} \quad \begin{bmatrix} I - A & -B \\ C & 0 \end{bmatrix} \begin{bmatrix} x_s \\ u_s \end{bmatrix} = \begin{bmatrix} 0 \\ r \end{bmatrix}, x_s \in \mathcal{X}, u_s \in \mathcal{U}$$

Usually assume  $x_s, u_s$  unique and feasible. If no solution exists, compute closest steady-state  $\min(C x_s - r)^T Q (C x_s - r)$  s. t.

$$x_s = A x_s + B u_s.$$

MPC problem to drive  $y \rightarrow r$  is:

$$\min_u \|y_N - C x_s\|_{P_y}^2 + \sum_{i=0}^{N-1} \|y_i - C x_s\|_{Q_y}^2 + \|u_i - u_s\|_R^2$$

### 5.2 Delta formulation

Reference  $r$ ,  $\Delta x_k = x_k - x_s$ ,  $\Delta u_k = u_k - u_s$ :

$$\min V_f(\Delta x_N) + \sum_{i=0}^{N-1} \Delta x_i^T Q \Delta x_i + \Delta u_i^T R \Delta u_i$$

$$\text{s.t. } \Delta x_0 = \Delta x_k$$

$$\Delta x_{k+1} = A \Delta x_k + B \Delta u_k$$

$$H_x x \leq k_x \Rightarrow H_x \Delta x \leq k_x - H_x x_s$$

$$H_u u \leq k_u \Rightarrow H_u \Delta u \leq k_u - H_u u_s$$

$$\Delta x_N \in \mathcal{X}_f \quad \text{adjusted accordingly, shift (and scaled)}$$

$$x_s \oplus \mathcal{X}_f \subseteq \mathcal{X}$$

$$K \Delta x + u_s \in \mathcal{U}$$

Control given by  $u_0^* = \Delta u_0^* + u_s$ .

### 5.3 Offset free tracking

$$x_{k+1} = A x_k + B u_k + B_d d_k$$

$$d_{k+1} = d_k$$

$$y_k = C x_k + C_d d_k$$

$$\begin{bmatrix} I - A & -B \\ C & 0 \end{bmatrix} \begin{bmatrix} x_s \\ u_s \end{bmatrix} = \begin{bmatrix} B_d \hat{d} \\ r - C_d \hat{d} \end{bmatrix}$$

Choice of  $B_d, C_d$  requires that  $(A, C)$  is observable and  $\begin{bmatrix} A - I & B_d \\ C & C_d \end{bmatrix}$

has full  $(n_x + n_d)$  column rank (i.e.  $\det \neq 0$ ). Intuition: for fixed  $y_s$  at steady-state,  $d_s$  is uniquely determined.

If plant has no integrator we can choose  $B_d = 0$  since  $\det(A - I) \neq 0$ .

$$\begin{bmatrix} \hat{x}_{k+1} \\ \hat{d}_{k+1} \end{bmatrix} = \begin{bmatrix} \mathbf{A} & \mathbf{B}_d \\ \mathbf{0} & \mathbf{I} \end{bmatrix} \begin{bmatrix} \hat{x}_k \\ \hat{d}_k \end{bmatrix} + \begin{bmatrix} \mathbf{B} \\ \mathbf{0} \end{bmatrix} u_k + \begin{bmatrix} \mathbf{L}_x \\ \mathbf{L}_d \end{bmatrix} \left( -y_k^m + \mathbf{C}\hat{x}_k + \mathbf{C}_d\hat{d}_k \right)$$

where  $y_k^m$  measured output; choose  $\begin{bmatrix} \mathbf{L}_x \\ \mathbf{L}_d \end{bmatrix}$  s.t. error dynamics stable and converge to zero.

If 1) number of dist. = number of outputs, 2) target steady-state problem feasible and no constraints active at steady-state, 3) closed-loop system converges, then target achieved without offset. Extend *Delta formulation* from above with

$$\Delta x_{k+1} = \mathbf{A}\Delta x_k + \mathbf{B}\Delta u_k + \mathbf{B}_d\Delta d_k$$

$$\Delta d_{k+1} = \Delta d_k$$

Algorithm becomes:

1. Estimate state and disturbance  $\hat{x}, \hat{d}$ ,
2. Obtain  $(x_s, u_s)$  target condition,
3. Solve MPC problem (adapted Delta formulation)

**Theorem** Case  $n_d = n_y$  and RHC is recursively feasible and unconstrained for  $k \geq j$  for some  $j \in \mathbb{N}$  and closed-loop converges, it converges to reference, i.e.  $y_k^m \rightarrow r$ .

#### 5.4 Soft-constraints via slack variables

$$\min_x f(z) + l_\epsilon(\epsilon) \quad \text{s.t. } g(z) \leq \epsilon, \epsilon \geq 0$$

Requirement: Softened problem has same minimiser as original problem if feasible.

Quadratic error function  $l_\epsilon(\epsilon) = v\epsilon + w\epsilon^2$ ,  $w > 0$  gives smoothness, choose  $v > \lambda^* \geq 0$  for exact penalty (above requirement fulfilled).

#### 5.5 Move Blocking

Main idea to set a number of inputs as the same,  $u_2 = u_3 = \dots = u_N$ , to reduce computational burden, at the slight cost of sub-optimality.

### 6 Robust MPC

**Enforcing terminal constraints** by robust invariance:

$$x \in O^W \Rightarrow g(x, w) \in \Omega^W \quad \forall w \in W$$

$$\text{pre}^W(\Omega) = \{x | g(x, w) \in \Omega \quad \forall w \in W\}$$

**Enforcing sequential constraints** for uncertain system  $\phi$ :

$$\phi_i(x_0, u, w) = \left\{ x_i + \sum_{j=0}^{i-1} \mathbf{A}^j w_j \mid w \in \mathcal{W}^i \right\} \subseteq \mathcal{X}$$

$$\phi_N(x_0, u, w) \in \mathcal{X}_f \quad \text{as well}$$

The uncertain system evolves with the summation of all the disturbances up to time  $i$ , hence we have to restrict the open-loop (determine control before disturbance is measured):

$$\mathbf{A}_x x \leq b_x \text{ becomes } \mathbf{A}_x x_i + \mathbf{A}_x \sum_{j=0}^{i-1} \mathbf{A}^j w_j \leq b_x :$$

$$x_i \in \mathcal{X} \ominus (\mathcal{W} \oplus \mathbf{A}\mathcal{W} \oplus \dots \oplus \mathbf{A}^{i-1}\mathcal{W})$$

$$= \left( \bigoplus_{j=0}^{i-1} \mathbf{A}^j \mathcal{W} \right) = [\mathbf{A}^0 \quad \dots \quad \mathbf{A}^{i-1}] \mathcal{W}^i$$

For example: Robust invariant set calculation of  $x_{k+1} = 0.5x_k + w_k$  under  $-10 \leq x \leq 10$  and  $-1 \leq w \leq 1$ .

$$\Omega_0 = [-10, 10]$$

$$\begin{aligned} \text{pre}^W(\Omega_0) &= \{x | -10 \leq 0.5x + w \leq 10 \text{ for } -1 \leq w \leq 1\} \\ &= \{x | -20 - 2w \leq x \leq 20 + 2w \text{ for } -1 \leq w \leq 1\} \\ &= \{x | -18 \leq x \leq 18\} \end{aligned}$$

$$\Omega_1 = [-10, 10] \cap [-18, 18] = [-10, 10] = \mathcal{O}_\infty^W$$

For example: Terminal set calculation of  $x_{k+1} = w_k$ ,  $-1 \leq w \leq 1$ ,  $N=2$ :

$$\mathcal{X}_f^W = \mathcal{X}_f \ominus \left( \bigoplus_{j=0}^1 \mathbf{A}^j \mathcal{W} \right) = \mathcal{X}_f \ominus 2\mathcal{W} = [-10, 10] \ominus [-2, 2] = [-8, 8]$$

**Tube-MPC** We want nominal system  $z_k = \mathbf{A}z_k + \mathbf{B}v_k$  with “tracking” controller  $u_k = \mathbf{K}(x_k - z_k) + v_k$  i.e. closed-loop,  $\mathbf{K}$  found offline.

Step 1: Compute the minimal robust invariant set  $\mathcal{E} = \bigoplus_{j=1}^\infty \mathbf{A}_{cl}^j \mathcal{W}$ .

Step 2: Shrink Constraints:

$$\begin{aligned} \{z_i\} \oplus \mathcal{E} \subseteq \mathcal{X} &\Rightarrow \{z_i\} \in \mathcal{X} \ominus \mathcal{E} \\ u_i \in \mathbf{K}\mathcal{E} \oplus \{v_i\} \subset \mathcal{U} &\Rightarrow \{v_i\} \in \mathcal{U} \ominus \mathbf{K}\mathcal{E} \\ z_n \in \mathcal{X}_f \ominus \mathcal{E} \end{aligned}$$

Also check that the set  $\mathcal{X}_f$  is invariant for the nominal system with tightened constraints:  $(\mathbf{A} + \mathbf{B}\mathbf{K})\mathcal{X}_f \subseteq \mathcal{X}_f$ , and that it satisfies the constraints:  $\mathcal{X}_f \subseteq \mathcal{X} \ominus \mathcal{E}$  and  $\mathbf{K}\mathcal{X}_f \subseteq \mathcal{U} \ominus \mathbf{K}\mathcal{E}$ .

### 7 Explicit MPC

$z^*(x_k)$  is continuous and polyhedral piecewise affine over feasible set.

#### 7.1 Quadratic Cost

$J^*(x_k)$  is continuous, convex and polyhedral piecewise quadratic. Using the formulation with substitution,

$$J(x_k) = \min z^T H z - x_k^T (Y - F H^{-1} F^T) x_k$$

$$\text{s.t. } G z \leq w + S x_k$$

$$z(x_k) = U + H^{-1} F^T x_k$$

$$S = E + G H^{-1} F^T$$

$$U^* = z^*(x_k) - H^{-1} F^T x_k$$

The first solution gives  $u^*(x_k) = \kappa(x_k)$ , which is continuous and piecewise affine on polyhedra  $\kappa(x) = F_j x + g_j$ .

#### 7.2 $1/\infty$ -norm

$J^*(x_k)$  is continuous, convex and polyhedral piecewise affine. Optimal solution:  $u_0^* = [0 \dots 0 \quad I_m \quad 0 \dots 0] z^*(x_k)$ , and is in the same form as the QP case above.

#### 7.3 Explicit Example

1. Write out KKT conditions and Lagrangian.
2. Determine infeasible regions from primal feasibility constraints. For example,  $x_1 < 10$ .
3. From primal and dual feasibility, and complementary slackness conditions, list out all cases that can occur.

$$\begin{aligned} \lambda_1 &= 0 & \lambda_1 &\geq 0 \\ g_1(x) &< 0 & g_1(x) &= 0 \end{aligned}$$

4. Solve for each case:  $z^*(x_1, x_2)$  and  $J^*(x_1, x_2)$ , listing the active constraints, and range of validity.

### 8 Hybrid MPC

#### 8.1 Piecewise Affine (PWA) Systems

Affine dynamics and output for each region:

$$\begin{cases} x_{k+1} = \mathbf{A}^i x_k + \mathbf{B}^i u_k + \mathbf{f}^i \\ y_k = \mathbf{C}^i x_k + \mathbf{D}^i u_k + g_i \end{cases} \text{ if } x_t \in \mathcal{X}_j$$

Polyhedral partition of the  $(x, u)$ -space:

$$\{\mathcal{X}_i\}_{j=1}^s = \{x, u | H_j x + J_j u \leq K_j\}$$

#### 8.2 Mixed Logical Dynamical Hybrid Model (MLD)

**Idea** associate boolean to binary:  $p_j \Leftrightarrow \delta_i = 1, \neg p_j \Leftrightarrow \delta_i = 0$ .

**Goal** Given a boolean formula  $F(p_1, \dots, p_n)$  define polyhedral set  $P$  s.t. set of binary values  $\{\delta_1, \dots, \delta_n\}$  satisfies Boolean formula  $F$  in  $P$   $F(p_1, \dots, p_n) \Leftrightarrow \mathbf{A}\delta \leq b, \delta \in \{0, 1\}^n$ .

#### 8.3 Analytical Approach

1. Transform into **Conjunctive Normal Form (CNF)**, i.e.

$$F(p_1, \dots, p_n) = \bigvee_m \left[ \bigwedge_j p_j \right].$$

2. Translate CNF into algebraic inequalities.

#### Translate logic rules to Linear Integer Inequalities

AND	$p_1 \wedge p_2$	$\delta_1 \geq 1, \delta_1 \geq 1$ also $\delta_1 + \delta_2 \geq 2$
OR	$p_1 \vee p_2$	$\delta_1 + \delta_2 \geq 1$
NOT	$\neg p_1$	$1 - \delta_1 \geq 1$ also $\delta_1 = 0$
XOR	$p_1 + p_2$	$\delta_1 + \delta_2 = 1$
IMPLY	$p_1 \rightarrow p_2$	$\delta_1 - \delta_2 \leq 0$
IFF	$p_1 \leftrightarrow p_2$	$\delta_1 - \delta_2 = 0$
ASSIGN	$p_3 \leftrightarrow p_1 \wedge p_2$	$\delta_1 + (1 - \delta_3) \geq 1$ <b>and</b> $\delta_2 + (1 - \delta_3) \geq 1$ <b>and</b> $(1 - \delta_1) + (1 - \delta_2) + \delta_3 \geq 1$
<b>CNF-Clause 0</b>	$p_1 \vee p_2 \vee p_3$	$\delta_1 + \delta_2 + \delta_3 \geq 1$
<b>CNF-Clause 1</b>	$\neg p_1 \vee p_2 \vee p_3$	$\delta_1 - \delta_2 - \delta_3 \leq 0$
<b>CNF-Clause 2</b>	$\neg p_1 \vee \neg p_2 \vee p_3$	$\delta_1 + \delta_2 - \delta_3 \leq 1$
<b>CNF-Clause 3</b>	$\neg p_1 \vee \neg p_2 \vee \neg p_3$	$\delta_1 + \delta_2 + \delta_3 \leq 2$

<b>CNF-Clause 0</b>	$p_1 \vee p_2 \vee p_3$	$\delta_1 + \delta_2 + \delta_3 \geq 1$
<b>CNF-Clause 1</b>	$\neg p_1 \vee p_2 \vee p_3$	$\delta_1 - \delta_2 - \delta_3 \leq 0$
<b>CNF-Clause 2</b>	$\neg p_1 \vee \neg p_2 \vee p_3$	$\delta_1 + \delta_2 - \delta_3 \leq 1$
<b>CNF-Clause 3</b>	$\neg p_1 \vee \neg p_2 \vee \neg p_3$	$\delta_1 + \delta_2 + \delta_3 \leq 2$

#### Logic Equality Rules (for Jenwei)

$$\begin{aligned} \neg(\mathbf{A} \wedge \mathbf{B}) &= \neg \mathbf{A} \vee \neg \mathbf{B} \\ \mathbf{A} \wedge (\mathbf{B} \vee \mathbf{C}) &= (\mathbf{A} \wedge \mathbf{B}) \vee (\mathbf{A} \wedge \mathbf{C}) \\ \mathbf{A} \vee (\mathbf{B} \wedge \mathbf{C}) &= (\mathbf{A} \vee \mathbf{B}) \wedge (\mathbf{A} \vee \mathbf{C}) \end{aligned}$$

#### 8.3.1 Translate continuous and logical components into Linear Mixed-Integer Relations

Event generator:  $\delta_e(k) = f_{EG}(x_c(k), u_c(k), t)$ .

Consider:  $p \Leftrightarrow a^T x \leq b, \mathcal{X} = \{x | a^T x - b \in [m, M]\}$ .

Translated to linear inequalities:  $m\delta < a^T x - b \leq M(1 - \delta)$ , where  $[m, M]$  are lower and upper bounds.

#### Representing Switched Affine Dynamics as "IF-THEN-ELSE" relations

IF  $p$  THEN  $z_k = a_1^T x_k + b_1$  else  $z_k = a_2^T x_k + b_2 \Leftrightarrow$

$$(m_2 - M_1)\delta + z_k \leq a_2^T x_k + b_2 \leq -(m_1 - M_2)\delta + z_k$$

$$(m_1 - M_2)(1 - \delta) + z_k \leq a_1^T x_k + b_1 \leq -(m_1 - M_2)(1 - \delta) + z_k$$

This results in a linear MLD model

$$x_{k+1} = \mathbf{A}x_k + \mathbf{B}_1 u_k + \mathbf{B}_2 \delta_k + \mathbf{B}_3 z_k$$

$$y_k = \mathbf{C}x_k + \mathbf{D}_1 u_k + \mathbf{D}_2 \delta_k + \mathbf{D}_3 z_k$$

$$E_2 \delta_k + E_3 z_k \leq E_4 x_k + E_1 u_k + E_5$$

where the last equation describes the relationship between the continuous and integer variables. Physical constraints on cont.

variables:  $\{C\} = \left\{ \begin{bmatrix} x_c \\ u_c \end{bmatrix} \in \mathcal{R}^{n_c+m_c} \mid Fx_c + Gu_c \leq H \right\}$

## 8.4 CFTOC for Hybrid Systems

$$J^*(x) = \min_U l_N(x_N) + \sum_{k=0}^{N-1} l(x_k, u_k, \delta_k, z_k)$$

$$\text{s.t. } x_{k+1} = Ax_k + B_1 u_k + B_2 \delta_k + B_3 z_k$$

$$E_2 \delta_k + E_3 z_k \leq E_4 x_k + E_1 u_k + E_5$$

$$x_N \in \mathcal{X}_f, x_0 = x(0)$$

## 8.5 MILP/MLQP

$$\min c_c z_c + c_b z_b + d \quad \text{OR} \quad [z_c z_b] H [z_c z_b] + q [z_c z_b] + d$$

$$\text{s.t. } G_c z_c + G_b z_b \leq W$$

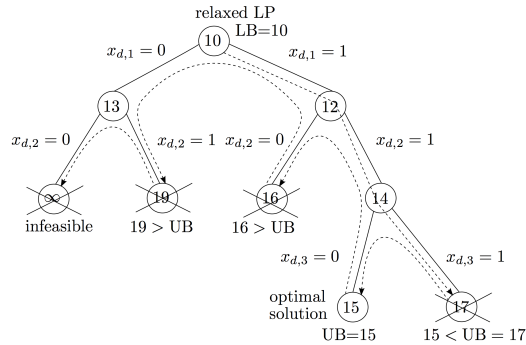
$$z_c \in R^{s_c}, z_b \in \{0, 1\}^{s_b}$$

Explicit solution is a time varying fb law for both problems:

$$u_k^*(x_k) = F_k^j x_k + G_k^j \text{ if } x_k \in \mathcal{R}_k^j.$$

**Brute force:** enumerating all the  $2^{s_b}$  integer values of the variable  $z_b$  and solve the corresponding problem.

**Branch and Bound:** relaxation of binaries:  $\{0, 1\} \rightarrow [0, 1]$ . Lower bound on the optimal solution of the modified problem is found. Any feasible solution to original problem is upper bound on optimal cost.



## 9 Numerical Optimization – Iterative Methods

### 9.1 Gradient descent

$x_{i+1} = x_i - h_i \nabla f(x_i)$  with step-size  $h_i = \frac{1}{L}$  for  $L$ -smooth  $f(x)$ :

$$\exists L : \|\nabla f(x) - \nabla f(y)\| \leq L \|x - y\| \quad \forall x, y \in \mathbb{R}^n$$

$\Leftrightarrow \nabla f$  is Lipschitz continuous

$\Leftrightarrow f$  can be upperbounded by a quadratic function:

$$f(x) \leq f(y) + \nabla f(y)^T (x - y) + 0.5L \|x - y\|^2 \quad \forall x, y \in \mathbb{R}^n$$

### 9.2 Newton's Method

$$x_{i+1} = x_i + h_i \Delta x_{nt} := x_i - h_i (\nabla^2 f(x_i))^{-1} \nabla f(x_i)$$

Line search problem: choose  $h_i > 0$  s.t.  $f(x_i + h_i \Delta x_{nt}) \leq f(x_i)$ .

Either compute exact and best  $h_i$  using:

$$h_i^* = \operatorname{argmin} x_i + h_i \Delta x_{nt}$$

Or use the backtracking search method:

For  $\alpha \in (0, 0.5)$  and  $\beta \in (0, 1)$ :

Initialise  $h_i = 1$ ;

while  $f(x_i + h_i \Delta x_{nt}) > f(x_i) + \alpha h_i \nabla f(x_i)^T \Delta x_{nt}$  do  $h_i \leftarrow \beta h_i$

For given equality constraint  $Ax = b$  solve:

$$\begin{bmatrix} \nabla^2 f(x_i) & A^T \\ A & 0 \end{bmatrix} \begin{bmatrix} \Delta x \\ \lambda \end{bmatrix} = \begin{bmatrix} -\nabla f(x_i) \\ 0 \end{bmatrix}$$

### 9.3 Constrained optimization with $g_i(x) \leq 0$

**Gradient method**  $x_{i+1} = \pi_Q(x_i - h_i \nabla f(x_i))$  where  $\pi_Q$  is a projection  
 $\pi_q = \arg \min_x \frac{1}{2} \|x - y\|_2^2$ . Projection can be solved directly if simple enough, else solve the dual.

### 9.4 Interior-Point methods

Assumptions  $f(x^*) < \infty, \tilde{x} \in \operatorname{dom}(f)$ .

**Barrier method**  $\min f(x) + \kappa \phi(x)$ . Approximate  $\phi$  using diff'able log barrier (instead of indicator function):

$$\phi(x) = \sum_{i=1}^m I_{-}(g_i(x)) = - \sum_{i=1}^m \log(-g_i(x))$$

$$\lim_{\kappa \rightarrow 0} x^*(\kappa) = x^*$$

Analytic center:  $\arg \min_x \phi(x)$ , central path  $\{x^*(\kappa) \mid \kappa > 0\}$ .

### Path following method

1. Centering  $x^*(\kappa) = \arg \min_x f(x) + \kappa \phi(x)$  with newton's method:

$$1.1. \Delta x_{nt} = [\nabla^2 f(x) + \kappa \nabla^2 \phi(x)]^{-1} (-\nabla f(x) - \kappa \nabla \phi(x)).$$

1.2. Line search:

$$\text{retain feasibility: } \arg \max_{h>0} \{h | g_i(x + h \Delta x) < 0\}$$

$$\text{Find } h^* = \arg \min_{h \in [0, h_{\max}]} \{f(x + h \Delta x) + \kappa \phi(x + h \Delta x)\}$$

2. Update step  $x_i = x^*(\kappa_i)$

3. Stop if  $m \kappa_i \leq \epsilon$

4. Decrease  $\kappa_{i+1} = \kappa_i / \mu, \mu > 1$ .

### Centering step with equality constraints

$$\begin{bmatrix} \nabla^2 f(x) + \kappa \nabla^2 \phi(x) & c^T \\ c & 0 \end{bmatrix} \begin{bmatrix} \Delta x_{nt} \\ \nu \end{bmatrix} = - \begin{bmatrix} \nabla f(x) + \kappa \nabla \phi(x) \\ 0 \end{bmatrix}$$

### Relaxed KKT

$$Cx^* = d \quad g_i(x^*) + s_i^* = 0$$

$$\nabla f(x^*) + \sum_{i=1}^m \lambda_i^* \nabla g_i(x^*) + c^T \nu = 0 \quad \lambda_i^* = \kappa \frac{\partial \phi}{\partial g_i} = -\frac{\kappa}{g_i}$$

$$\lambda_i^* g_i(x^*) = -\kappa \quad \lambda_i^*, s_i^* \geq 0$$

### Primal Dual Search Direction Computation

$$\begin{bmatrix} H(x, \lambda) & c^T & G(x)^T & 0 \\ c & 0 & 0 & 0 \\ G(x) & 0 & 0 & I \\ 0 & 0 & S & \Lambda \end{bmatrix} \begin{bmatrix} \Delta x \\ \Delta u \\ \Delta \lambda \\ \Delta s \end{bmatrix} = - \begin{bmatrix} \nabla f(x) + c^T \nu + G(x)^T \lambda \\ cx - d \\ g(x) + s \\ s\lambda - \nu \end{bmatrix}$$

$S = \operatorname{diag}(s_1, \dots, s_m), \Lambda = \operatorname{diag}(\lambda_1, \dots, \lambda_m)$  and

$\nu$  is a vector for choosing centering parameters.