

Math 578A: Part 2, Homework 2

Due on May 1 Midnight

Total points: 80 points

Points will be given for correctness and clarity. You may submit a scan/photo of your work, however if it is difficult to read due to handwriting, etc., points may be deducted.

Question 1

20 points

Eulerian graphs

- (a) [5 points] Write pseudo-code to determine if a graph has an Eulerian cycle.
- (b) [15 points] Provide an algorithm to find an Eulerian cycle in an undirected graph in $O(N)$ time.

Question 2

20 points

Assembly

- (a) [10 points] An experiment attempts to provide the k -mer spectrum of a genome that does not have any repeats of length $k - 1$, in either the forward or reverse direction, however the experiment randomly provides either a k -mer, or its reverse complement, for the entire spectrum (e.g. there are $n - k + 1$ k -mers). Give an algorithm for constructing a genome.
- (b) [10 points] Can one use a 12-inch-long wire to form a cube (each of the 12 cube edges is 1-inch long). If not, what is the smallest number of cuts one must make to form this cube?

Question 3

20 points

Approximation ratios

One of the first greedy algorithms you learned (but maybe didn't know it) is to make change. As input, you are given a list of coin values in decreasing value $C = c_1, c_2, \dots, c_N$, and a value v , and you want to determine n the fewest coins to use to total v . For example, in US currency, $C = 25, 10, 5, 1$, and to give 0.82 you would give 3 c_1 , 0 c_2 , 1 c_3 , 2 c_4 , $n = 6$. A greedy algorithm to find the number of coins to use is:

- (a) [5 points] Provide a list of coin values and a total value for which GreedyChange incorrectly calculates the minimal number of coins to return as change.
- (b) [5 points] Assuming there isn't a coin that exactly matches the return value, what is the minimal value for GreedyChange?
- (c) [10 points] What is the approximation ratio for GreedyChange? Note this allows for any set of coin values C .

Question 4

20 points

Phylogenetics

Design a backtracking procedure to reconstruct the optimal assignment of characters in the Sankoff algorithm for the Weighted Small Parsimony problem.

Algorithm 1 Greedy change

```
1: procedure GREEDYCHANGE( $C, v$ )  
2:    $i = 1, n = 1$   
3:   while  $v > 0$  do  
4:      $n_i \leftarrow \lfloor v/c_i \rfloor$   
5:      $v \leftarrow v - n_i * c_i$   
6:      $n \leftarrow n + n_i$   
7:      $i = i + 1$   
8:   end while  
9:   return  $n$   
10: end procedure
```
