

TSP (graph)

Input: a complete weight graph

output: Hamiltonian cycle with the minimum weight and the maximum weight

Let mst be getMST(graph)

Let dfs-path be getDFSPath(mst)

Let Hamiltonian-cycle be getHamiltonianCycle(dfs-path)

Let weight be getWeight(Hamiltonian-cycle)

Return weight, Hamiltonian cycle

get MST (graph) :

input : a complete weighted graph

output : a minimum spanning tree

Let S be an empty set

Let Q be an empty priority queue

Let mst-edges be an empty list

Add root node "0" to queue

while queue is not empty, do:

 let current be node with min weight

 if current not in S :

 Add current to S

 Add current to mst-edge

 for all neighbor of current:

 if neighbor not in S

 Append neighbor to queue.

return \leftarrow make it a traversable tree
 to Tree(mst-edges)

get DFS Path (mst, root)

inputs: minimum spanning tree, a root node
output: Depth First search path

Let path be an empty list

let neighbors be neighbors of root

if neighbors is empty:

return

for all neighbor in neighbors:

append to path getDFSPath(mst, neighbor)

append root to path

return path

get Hamiltonian Cycle (dfs-path)

input: depth first search path of a tree

output: Hamiltonian cycle represented by
a list of vertex

Let result be an empty list

for all vertex in dfs-path:

if vertex is not in empty list :

append vertex to result

append first node dfs-path [0] to result

return result

get Hamiltonian Weight (cycle)

input: List of vertex representing
Hamiltonian cycle

output: Total weight of cycle

Let $a = 0$

Let $b = 1$

Let total_weight be 0

while b less than length of cycle

Let weight be weight between vertex

at index a and index b

total_weight = total_weight + weight

$a += 1$

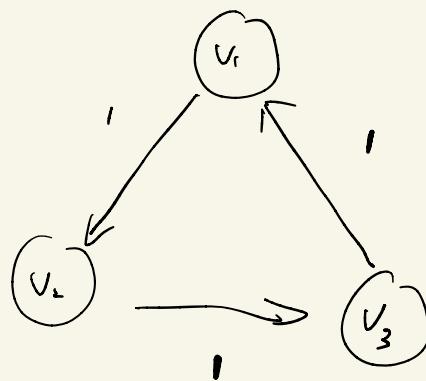
$b += 1$

return weight +

Proof:

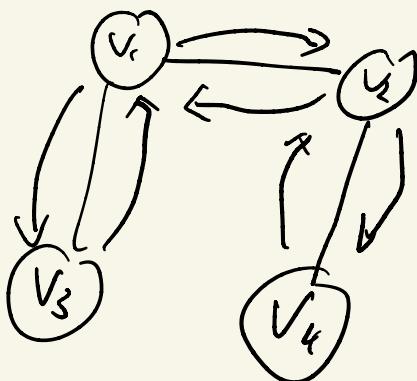
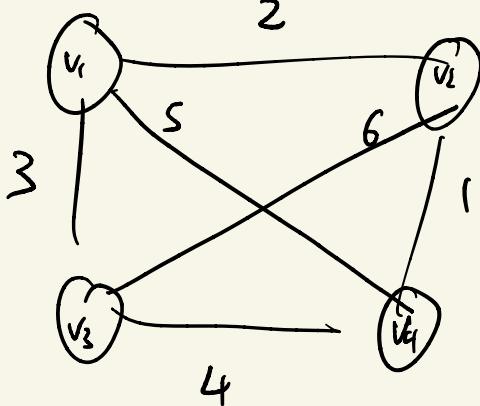
Lemma: There exist a complete, metrically weighted graph $G = (V, E)$ such that your approximation always finds a Hamiltonian cycle of min weights.

Example:



TSP will get a min MST with $\{(V_1, V_2), (V_2, V_3)\}$ and the dfs will be $[V_1, V_2, V_1, V_2, V_1]$ and path is $[V_1, V_2, V_3]$ so weight is 3 always.

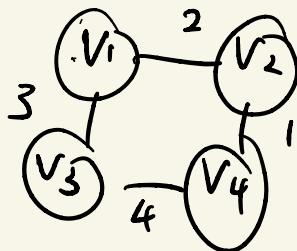
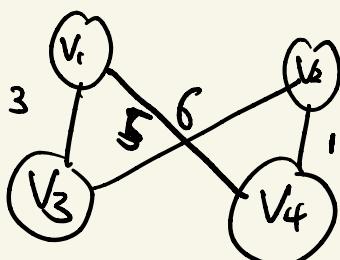
Lemma: There exist a complete, metrically weighted graph $G = (V, E)$ such that your approximation may find a Hamiltonian cycle of greater than minimum weight.



$$dPS = [v_1, v_3, v_1, v_2, v_1, v_4, v_2, v_1]$$

Hamiltonian path:

$$[v_1, v_3, v_2, v_4]$$



TPS returns

$$1 + 3 + 5 + 6 = \boxed{15}$$

$$1 + 2 + 3 + 4 = \boxed{10}$$

Optimal is