

Consider a permutation test (the ideas here apply broadly to all sampling-based tests of independence or conditional independence, including marginal permutation test, conditional permutation test, marginal randomization test, and conditional randomization test). Let  $B$  be the number of resamples. Let  $T_1, \dots, T_B$  be the ordered resampled test statistics. Let  $T^*$  be the test statistic computed on the raw data. Let  $I_i$  be the indicator that  $T^*$  is less than  $T_i$ , i.e.

$$I_i = \mathbb{I}(T^* \leq T_i).$$

We have the basic fact under the null hypothesis that

$$\mathbb{E}(I_i) = \mathbb{P}(T^* \leq T_i) = i/B.$$

We want to combine  $I_1, \dots, I_B$  in such a way so as to produce an e-value  $E$  (i.e., a non-negative variable  $E$  such that  $\mathbb{E}(E) \leq 1$ ). To test the hypothesis, we check if  $E > \beta$  for some threshold  $\beta$ .

We call  $f : \mathbb{R}^B \rightarrow \mathbb{R}^{\geq 0}$  a “combining function” if  $f$  combines  $I_1, \dots, I_B$  in such a way so as to produce a valid E-value, i.e.,

$$f(I_1, \dots, I_B) \geq 0; \quad \mathbb{E}[f(I_1, \dots, I_B)] \leq 1.$$

Let  $\mathcal{F}$  be the class of combining functions. An interesting subset of  $\mathcal{F}$  is the set of linear combining functions. A function  $g \in \mathcal{F}$  is a linear combining function if

$$g(I_1, \dots, I_B) = c + \sum_{i=1}^B a_i I_i$$

for given scalars  $a_1, \dots, a_B$ , and  $c$ . The linearity of expectation implies that

$$\mathbb{E}[g(I_1, \dots, I_B)] = c + \sum_{i=1}^B a_i \mathbb{E}(I_i) = c + \frac{1}{B} \sum_{i=1}^B (a_i)(i).$$

Therefore, under the stricter requirement that  $\mathbb{E}(E) = 1$ , we have the constraints

$$\text{i } \sum_{i=1}^B (a_i)i = B(1 - c)$$

$$\text{ii } c + \sum_{i=1}^B a_i I_i \geq 0.$$

We give a couple examples of linear combining functions.

**Example 1.** Let  $c = 0$  and  $a_i = 1/i$  for all  $i$ . Then

$$\sum_{i=1}^B (1/i)i = B(1 - 0),$$

confirming constraint (i). Next,

$$0 + \sum_{i=1}^B \frac{I_i}{i} \geq 0,$$

confirming constraint (ii). Therefore,

$$E = \sum_{i=1}^B \frac{I_i}{i}$$

is a valid  $e$ -value. This is a left-tailed test: if  $T^*$  is small, then  $E$  is big, leading us to reject the null.

**Example 2.** The standard permutation test is a special case of the proposed framework. Let  $c = 3/2 + 1/(2B)$  and  $a_i = -1/B$  for all  $a_i$ . Verifying condition (i),

$$\begin{aligned} \sum_{i=1}^B (a_i)i &= -\frac{1}{B} \sum_{i=1}^B i = -\frac{1}{B} \left( \frac{B^2 + B}{2} \right) = -(1/2)B - 1/2 \\ &= B(1 - (3/2) - 1/(2B)) = B(1 - c). \end{aligned}$$

Next, verifying condition (ii),

$$(3/2) + 1/(2B) - (1/B) \sum_{i=1}^B I_i \geq (3/2) + 1/(2B) - 1 \geq 0.$$

Therefore,

$$E = 3/2 + 1/(2B) - \frac{1}{B} \sum_{i=1}^B I_i$$

is a valid  $e$ -value. Note that  $p_B := \frac{1}{B} \sum_{i=1}^B I_i$  is simply the  $p$ -value corresponding to the right-sided permutation test. For given  $\alpha \in (0, 1)$ , we reject the null hypothesis if and only if  $p_B < \alpha$ , which is equivalent to rejecting the null hypothesis if and only if

$$E > 3/2 + 1/(2B) - \alpha.$$

Therefore, the  $p$ -permutation test is a special case of the  $e$ -permutation test. We typically think of  $e$ -values as controlling FDR, but here  $e$ -values control type-I error.

We can extend this idea to control family-wise error rate.

**Example 3.** Suppose we are using CRT or CPT to test conditional independence. If we have misspecified the model for  $X|Z$ , then the resampled test statistics  $T_1, \dots, T_B$  will be off. In particular, qq-plots often reveal that in such settings the tail of the distribution is incorrect, leading to  $p$ -value inflation. We can use the proposed framework to downweight re-sampled test statistics in the tail of the distribution, possibly leading to better calibration. For example, consider again example 1, where previously we set  $a_i = 1/i$  for all  $i$ . Suppose that we use (e.g.) a Gaussian kernel to assign the weights. Let

$$\hat{\mu} = \frac{1}{B} \sum_{i=1}^B T_i$$

be the average of the resampled test statistics. Next, let

$$K(t) = \frac{1}{\sqrt{2\pi}} e^{-(1/2)t^2}$$

be the Gaussian kernel. Define  $a_i$  as

$$a_i := (B/i) \frac{K(|T_i - \hat{\mu}|/h)}{\sum_{i=1}^B K(|T_i - \hat{\mu}|/h)}.$$

Again, set  $c = 0$ . We see that

$$\sum_{i=1}^B (a_i) i = \frac{\sum_{i=1}^B K(|T_i - \hat{\mu}|/h)}{\sum_{i=1}^B K(|T_i - \hat{\mu}|/h)} = B = B(1 - c),$$

verifying the first condition. Moreover,  $A_i \geq 0$  for all  $i$ , satisfying the second condition. It follows that

$$E := \sum_{i=1}^B a_i I_i$$

is a valid  $e$ -value. This  $e$ -value is less sensitive to the tails of the sampling distribution than the  $e$ -value from example 1.

**Advantages to this approach.** The proposed framework has several advantages (I think). First, because the outputs are  $e$ -values, they can be combined under arbitrary dependence to control FDR using the  $e$ -BH procedure. Second, the framework is flexible: we can choose among many different “combining functions” (including the easy-to-use linear variants) to achieve different objectives. Certain combining functions might have robustness properties, for example. Third, we easily can combine  $e$ -values across data splits, enabling us to run multiple train-test splits without worrying about combining the resulting  $p$ -values.

**Open questions.** What is the “best” combining function that we could use?

**Additional properties of the  $I_i$ s.** Recall that  $I_i = \mathbb{I}(T^* \leq T_i)$ . We state several additional properties of the  $I_i$ s. Let  $k_1, k_2, \dots, k_p \in \mathbb{N}$  be indexes such that

$$1 \leq k_1 \leq k_2 \leq \dots \leq k_p \leq B.$$

Furthermore, let  $r_1, r_2, \dots, r_p \in \mathbb{N}$  be non-negative integers. We have that

$$I_{k_1}^{r_1} I_{k_2}^{r_2} \dots I_{k_p}^{r_p} = I_{k_1}.$$

The reason is as follows. First, because  $I_{k_i}$  is Bernoulli,

$$I_{k_1}^{r_1} I_{k_2}^{r_2} \dots I_{k_p}^{r_p} = I_{k_1} I_{k_2} \dots I_{k_p}.$$

Next, we take cases on  $T^*$ . First, suppose that  $T^* \leq k_1$ . Then  $T^* \leq k_2, \dots, k_p$ . Therefore,

$$I_{k_1} I_{k_2} \dots I_{k_p} = 1 \cdot 1 \dots 1 = I_{k_1}.$$

Next, suppose that  $T^* > k_1$ . Then  $I_{k_1} = 0$ , and so

$$I_{k_1} I_{k_2} \dots I_{k_p} = 0 \cdot I_{k_2} \dots I_{k_p} = 0 = I_{k_1}.$$

This property is useful because it allows us to more easily evaluate expressions of the form  $(\sum_{i=1}^n I_i)^r$  for  $r \in \mathbb{N}$ . As a warmup, we consider  $r = 1$ :

$$\sum_{i=1}^B I_i = a_i^{(1)} I_i,$$

where  $a_1^{(1)} = 1$ . Next, consider  $r = 2$ :

$$\begin{aligned} \left( \sum_{i=1}^B I_i \right)^2 &= \sum_{i=1}^B \sum_{j=1}^B I_i I_j = \sum_{i=1}^B I_i^2 + 2 \sum_{i=1}^{B-1} \sum_{j=i+1}^B I_i I_j = \sum_{i=1}^B I_i + 2 \sum_{i=1}^{B-1} \sum_{j=i+1}^B I_i \\ &= \sum_{i=1}^B I_i + \sum_{i=1}^{B-1} 2(B-i) I_i = \sum_{i=1}^B I_i + \sum_{i=1}^B 2(B-i) I_i = \sum_{i=1}^B (2B-2i+1) I_i = \sum_{i=1}^B a_i^{(2)} I_i, \end{aligned}$$

where  $a_i^{(2)} = 2B - 2i + 1$ .

Consider the multinomial theorem:

$$\left( \sum_{i=1}^B I_i \right)^r = \sum_{k_1+k_2+\dots+k_B=r} \binom{r}{k_1, k_2, \dots, k_B} \prod_{i=1}^B I_i^{k_i},$$

where

$$\binom{r}{k_1, k_2, \dots, k_B} = \frac{r!}{k_1! k_2! \dots k_B!}.$$

Let

$$C(r, B) := \left\{ (k_1, k_2, \dots, k_B) \in \{1, \dots, r\}^B : \sum_{i=1}^B k_i = r \right\},$$

i.e., the set of tuples of length  $B$  of integers from 1 to  $r$  such that the elements of the tuple sum to  $r$ . Let  $\tau : C(r, B) \rightarrow B$  be defined by

$$\tau(k_1, \dots, k_B) = \min\{i : \{1, \dots, B\} : k_i \geq 1\},$$

i.e.,  $\tau$  is the minimal nonzero index of a given configuration. The inverse  $\tau^{-1}$  of  $\tau$  is the set of configurations with a given minimal nonzero index. That is, for  $s \in \{1, \dots, B\}$ ,

$$\tau^{-1}(s) = \left\{ (0, \dots, 0, k_s, k_{s+1}, \dots, k_B) : \sum_{i=s}^B k_i = r \right\}.$$

Therefore, the multinomial theorem reduces to

$$\begin{aligned} \left( \sum_{i=1}^B I_i \right)^r &= \sum_{(k_1, \dots, k_B): k_1 + \dots + k_B = r} \binom{r}{k_1, k_2, \dots, k_B} I_{\tau(k_1, \dots, k_B)} \\ &= \sum_{s=1}^B \sum_{(0, \dots, 0, k_s, \dots, k_B): k_s \geq 1, k_s + \dots + k_B = r} \binom{r}{k_1, k_2, \dots, k_B} I_s. \end{aligned}$$

We consider the above sum:

$$\begin{aligned} &\sum_{(0, \dots, 0, k_s, \dots, k_B): k_s \geq 1, k_s + \dots + k_B = r} \binom{r}{k_1, k_2, \dots, k_B} \\ &= \sum_{(k_1, \dots, k_{B-s+1}): k_1 \geq 1, k_1 + \dots + k_{B-s+1} = r} \binom{r}{k_1, k_2, \dots, k_B} \\ &= \sum_{j=1}^r \sum_{l_1 + \dots + l_{B-s} = r-j} \binom{r}{j, l_1, \dots, l_{B-s}} \\ &= \sum_{j=1}^r \sum_{l_1 + \dots + l_{B-s} = r-j} \frac{r!}{j! l_1! \dots l_{B-s}!} \\ &= \sum_{j=1}^r \sum_{l_1 + \dots + l_{B-s} = r-j} \frac{r(r-1) \dots (r-j+1)(r-j)!}{j! l_1! \dots l_{B-s}!} \\ &= \sum_{j=1}^r \sum_{l_1 + \dots + l_{B-s} = r-j} \frac{r!}{(r-j)! j!} \binom{r-j}{l_1, \dots, l_{B-s}} \\ &= \sum_{j=1}^r \frac{r!}{(r-j)! j!} \sum_{l_1, \dots, l_{B-s}} \binom{r-j}{l_1, \dots, l_{B-s}} \\ &= \sum_{j=1}^r \binom{r}{j} (B-s)^{r-j} = \sum_{j=1}^r \binom{r}{j} (B-s)^{r-j} 1^j \\ &= \sum_{j=0}^r \binom{r}{j} (B-s)^{r-j} 1^j - 1(B-s)^r = (B-s+1)^r - (B-s)^r := M(B, r, s). \end{aligned}$$

Therefore,

$$\left( \sum_{i=1}^B I_i \right)^r = \sum_{s=1}^B M(B, r, s) I_s.$$

The number  $M(B, r, s)$  can be quite large. However, normalizing the sum by  $B$  neutralizes this problem:

$$\begin{aligned} \left( \frac{1}{B} \sum_{i=1}^B I_i \right)^r &= (1/B)^r \sum_{i=1}^B [(B-i+1)^r - (B-i)^r] I_i \\ &= \sum_{i=1}^B [(1-i/B + 1/B)^r - (1-i/B)^r] I_i. \end{aligned}$$

All terms are easy to compute.

## 1 Checking correctness

Let us check this formula for specific values of  $r$ . First, setting  $r = 1$ , we have  $M(B, 1, s) = (B - s + 1) - (B - s) = 1$ . Therefore,

$$\left( \sum_{i=1}^B I_i \right)^1 = \sum_{i=1}^B I_i.$$

Next, setting  $r = 2$ , we have

$$M(B, 2, s) = (B + s - 1)^2 - (B - s)^2 = 2B - 2s + 1.$$

Therefore,

$$\left( \sum_{i=1}^B I_i \right)^2 = \sum_{i=1}^B [2B - 2i + 1] I_i,$$

which matches what we have above.

## 2 Polynomials

We easily can compute a polynomial of  $\sum_{i=1}^B I_i$ . We have that

$$\begin{aligned}
\left(\sum_{i=1}^B I_i - x_0\right)^l &= \sum_{k=0}^l \binom{l}{k} \left(\sum_{i=1}^B I_i\right)^k x_0^{l-k} \\
&= x_0^l + \sum_{k=1}^l \binom{l}{k} \left(\sum_{i=1}^B I_i\right)^k x_0^{l-k} \\
&= x_0^l + \sum_{k=1}^l \binom{l}{k} \left[ \sum_{i=1}^B M(B, k, i) I_i \right] x_0^{l-k} \\
&= x_0^l + \sum_{k=1}^l \sum_{i=1}^B x_0^{l-k} \binom{l}{k} M(B, k, i) I_i \\
&= x_0^l + \sum_{i=1}^B \left[ \sum_{k=1}^l x_0^{l-k} \binom{l}{k} M(B, k, i) \right] I_i
\end{aligned}$$

Finally, for an integer  $r$  and polynomial coefficients  $a_0, a_1, \dots, a_r \in \mathbb{R}$ , we have that

$$\begin{aligned}
\sum_{l=0}^r a_l (S - x_0)^l &= a_0 + \sum_{l=1}^r a_l (S - x_0)^l \\
&= a_0 + \sum_{l=1}^r a_l \left[ x_0^l + \sum_{i=1}^B \left[ \sum_{k=1}^l x_0^{l-k} \binom{l}{k} M(B, k, i) \right] I_i \right] \\
&= \left( a_0 + \sum_{l=1}^r a_l x_0^l \right) + \sum_{i=1}^B \left[ \sum_{l=1}^r \sum_{k=1}^l a_l x_0^{l-k} \binom{l}{k} M(B, k, i) \right] I_i.
\end{aligned}$$

Typically, we will have  $a_l = f^{(l)}(x_0)/l!$ , where  $f$  is a  $p$ -to- $e$  calibrator. This follows from the following  $r$ th order Taylor expansion:

$$f\left(\sum_{i=1}^B I_i\right) \approx \left(a_0 + \sum_{l=1}^r a_l x_0^l\right) + \sum_{i=1}^B \left[ \sum_{l=1}^r \sum_{k=1}^l \frac{f^{(l)}(x_0) x_0^{l-k} \binom{l}{k} M(B, k, i)}{l!} \right] I_i.$$

Of course, the above is a linear combination of the  $I_i$ s. We define

$$W(f, x_0, r, B, i) = \sum_{l=1}^r \sum_{k=1}^l \frac{f^{(l)}(x_0) x_0^{l-k} \binom{l}{k} M(B, k, i)}{l!}.$$



Therefore, we can write

$$f\left(\sum_{i=1}^B I_i\right) \approx \left(a_0 + \sum_{l=1}^r a_l x_0^l\right) + \sum_{i=1}^B W(f, x_0, r, B, i) T_i.$$

In practice, we could compute the  $W$  terms ahead of running the algorithm. We even could save these terms in within the package for common choices of  $f$ ,  $x_0$ ,  $r$ , and  $B$  (e.g.,  $f(p) = 2p^{-1}$ ,  $x_0 = 1/2$ ,  $r = 20$ ,  $B = 1000$ ).

### 3 $p$ -to- $e$ calibrators

A decreasing function  $f : [0, 1] \rightarrow [0, \infty]$  is a  $p$ -to- $e$  calibrator if and only if  $\int_0^1 f \leq 1$ . It is admissible if and only if  $f$  is upper semicontinuous,  $f(0) = \infty$ , and  $\int_0^1 f = 1$ . An example of a  $p$ -to- $e$  calibrator is

$$f_\kappa(p) = \kappa p^{\kappa-1}.$$

We derive the Taylor series of  $f$ . First, we compute the derivatives of  $f$ . The first derivative is

$$f_\kappa^{(1)}(p) = \kappa(\kappa - 1)p^{\kappa-2}.$$

Meanwhile, the second derivative is

$$f_\kappa^{(2)}(p) = \kappa(\kappa - 1)(\kappa - 2)p^{\kappa-3}.$$

In general, we have that

$$f_\kappa^{(m)}(p) = \prod_{j=0}^{m-1} (\kappa - j) p^{\kappa-m-1}.$$

Therefore, the Taylor series of  $f$  at  $a$  is

$$f_\kappa(p) = \sum_{m=0}^{\infty} \frac{f_\kappa^{(m)}(a)}{m!} (p - a)^m.$$